



# **On the Decidability of the Reachability Problem for Planar Differential Inclusions**

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## Overview of the presentation

- Introduction
- Simple Planar Differential Inclusion System (SPDI)
  - Reachability Problem
    - \* Difficulties
    - \* Our solution
  - Reachability Algorithm (Example)
- Conclusions

## Introduction

- Reachability for Hybrid Systems
  - Undecidable in most cases
  - Approximation methods
    - \* Polyhedral approximations
    - \* Ellipsoidal approximations
    - \* Level-set approximations
  - Exact decision methods for sub-classes
    - \* Finite bisimulation (TA, Initialized Rectangular Automata, etc.)
    - \* Quantifier elimination
    - \* “Geometric” methods (only on the plane)

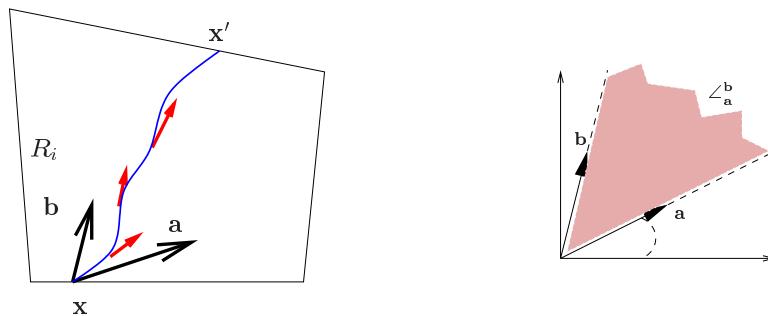
## Introduction (Cont.)

- “Geometric” methods - Ideas
  - Based on: Topology + Dynamical Systems
  - Require: Explicit solutions of Differential Equations
  - Take advantage of:
    - \* Topology of the plane
    - \* Explicit formula for Successors
    - \* Acceleration of cycles

# SPDI: Simple Planar Differential Inclusion System

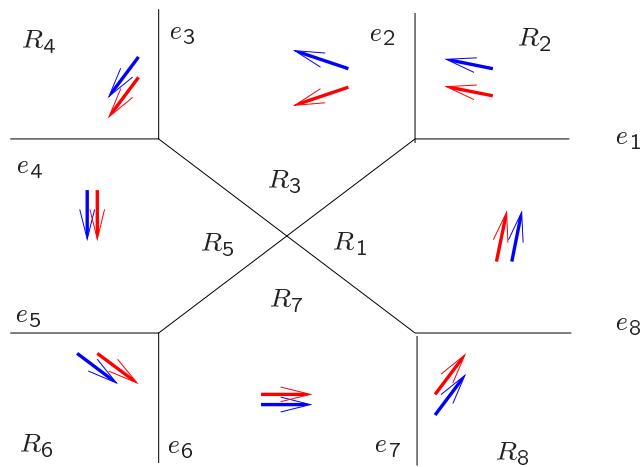
- SPDI:
  - A partition of the plane into convex polygonal regions
  - A constant differential inclusion for each region

$$\dot{x} \in \angle_a^b \text{ if } x \in R_i$$



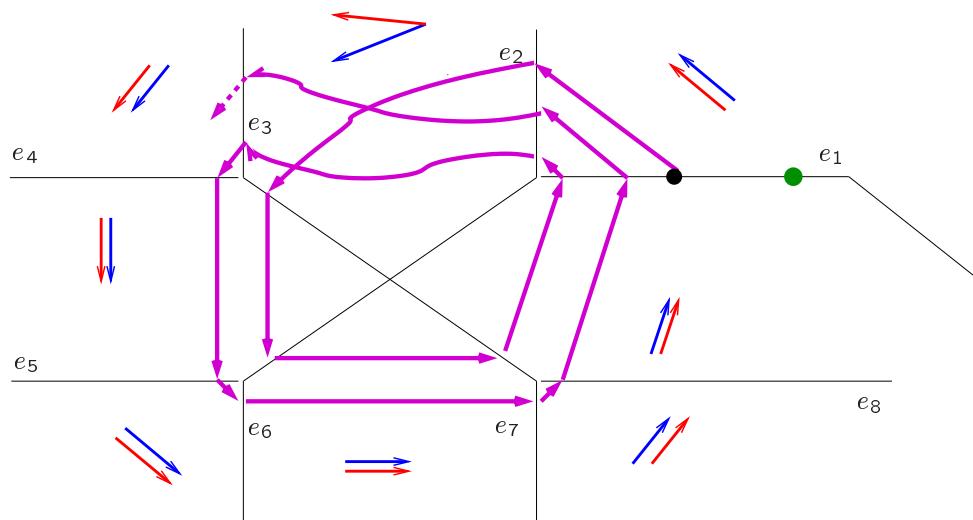
# SPDI: Simple Planar Differential Inclusion System

Example: Swimmer trying to escape from a whirlpool



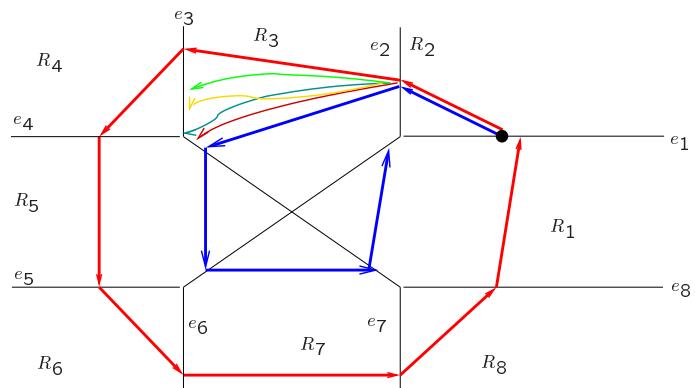
## Reachability Problem

Given  $x, x' \in \mathbb{R}^2$ , is there a trajectory  $\xi$  and  $t \geq 0$  such that  $\xi(0) = x$  and  $\xi(t) = x'$ ?



## Reachability Difficulties

Infinitely many trajectories (even locally)

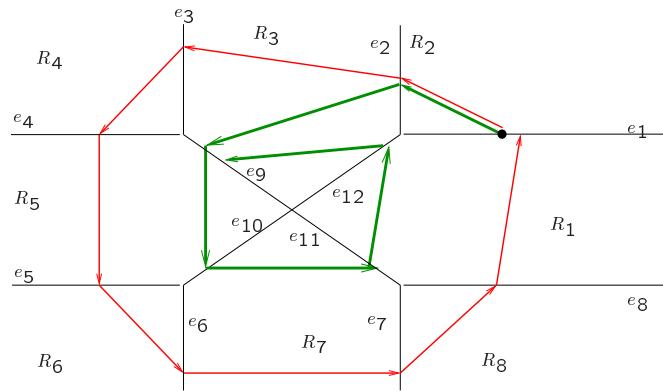


## Reachability Difficulties

Infinitely many qualitative behaviors:

$$\xi_1 = e_1 e_2 e_9 e_{10} e_{11} e_{12} e_9 \cdots e_{12} \cdots$$

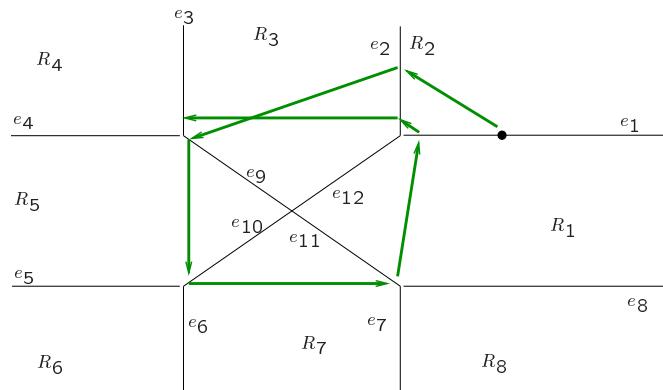
$$\xi_2 = e_1 e_2 \cdots e_8 e_1 \cdots$$



# Reachability Difficulties

Self-crossin s:

$$\xi_3 = e_1 \ e_2 e_9 e_{10} e_{11} e_1 \ e_2 e_3 \dots$$



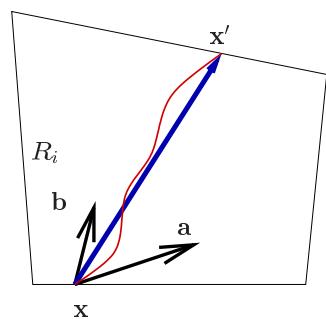
## Reachability (Solution)

Our approach:

- Simplification of trajectories
- Enumeration of qualitative behavior of trajectories
- Analysis of each qualitative behavior (computing successors)

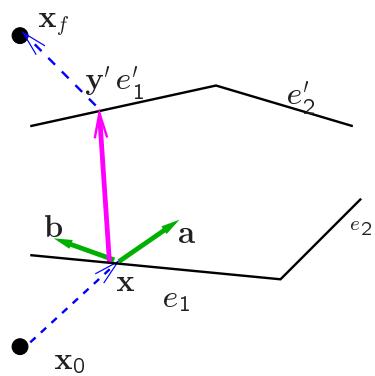
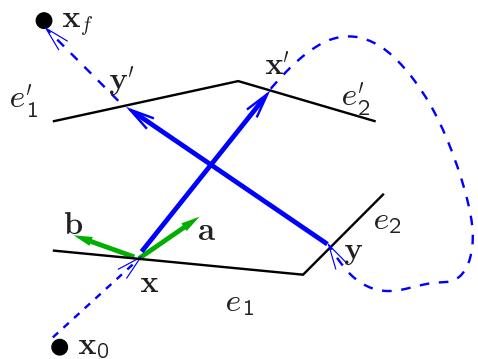
# Simplification of Trajectories

Strainin :



## Simplification of Trajectories (Cont.)

Removing self-crossings:

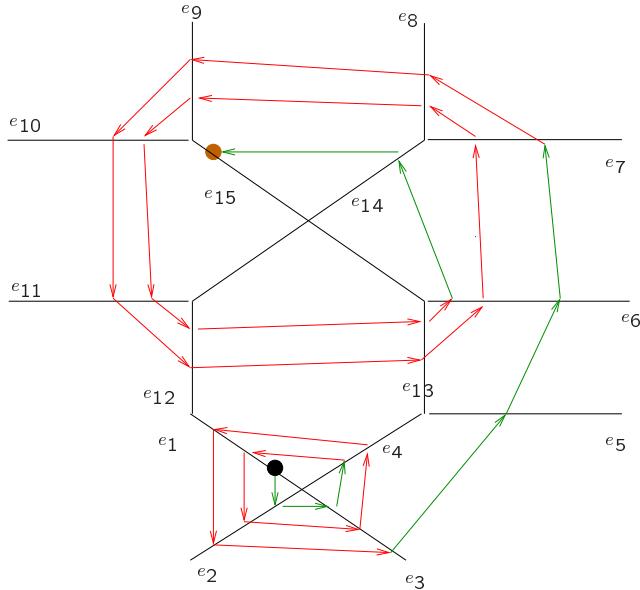


# Qualitative Behavior (of Simplified Trajectories)

- Classification Theorem: Any edge  $e$  is of nature.  $\sigma$  belongs to a type

$$type(\sigma) = r_1(s_1)^*r_2(s_2)^*\dots r_n(s_n)^*r_{n+1}$$

- Example:



with  $\sigma =$

$$e_1 e_2 e_3 \ (e_4 e_1 e_2 e_3)^2 e_5 e_6 \ (e_7 \dots e_{13} e_6)^2 \ e_{14} e_{15}$$

## Qualitative Behavior (of Simplified Trajectories)

- Properties:
  - $r_i$  is a seq. of pairwise different edges;
  - $s_i$  is a simple cycle;
  - $r_i$  and  $r_j$  are disjoint
  - $s_i$  and  $s_j$  are different
- Prop. The set of type of signatures is finite

## Preliminaries - for computing Successors

- Affine function (AF):  
 $f(x) = ax + b$  with  $a > 0$
- Affine multi-valued function (AMVF):  
 $\tilde{F}(x) = \langle f_1(x), f_2(x) \rangle$
- Truncated affine multi-valued function (TAMVF):  
 $F(x) = \tilde{F}(x) \cap \langle L, U \rangle$

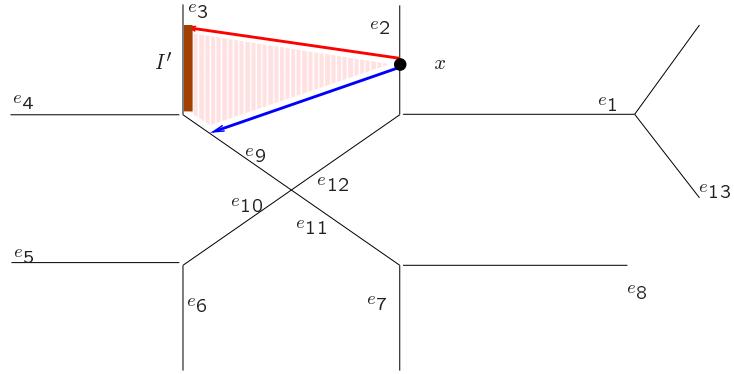
**Lemma:** AF, AMVF and TAMVF are closed under composition.

**Lemma:** Fixpoint equations  $F(I) = I$  can be explicitly solved (without iteration)\*.

\*Non trivial for T MVF.

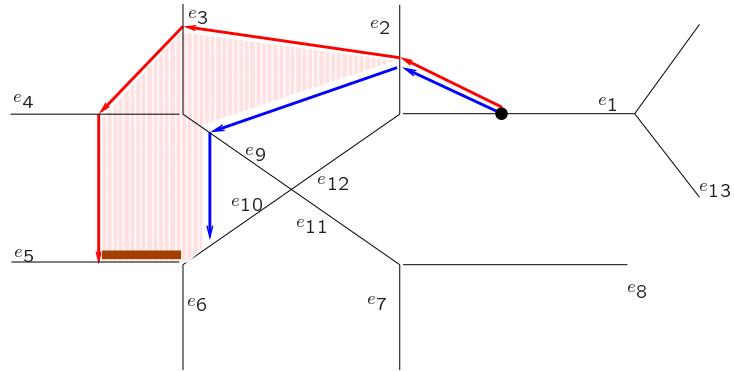
## Successors (by $\sigma$ )

- One step ( $\sigma = e_2e_3$ )



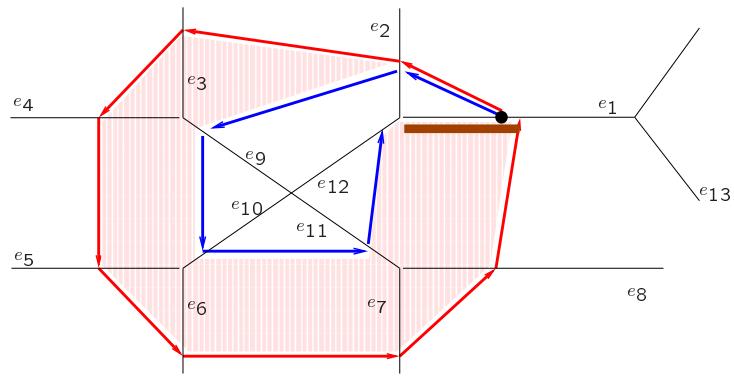
$$I' = \text{Succ}_{e_1e_2}(x) = [f_b(x), f_a(x)] = F(x)$$

- Several steps ( $\sigma = e_1e_2e_3e_4e_5$ )

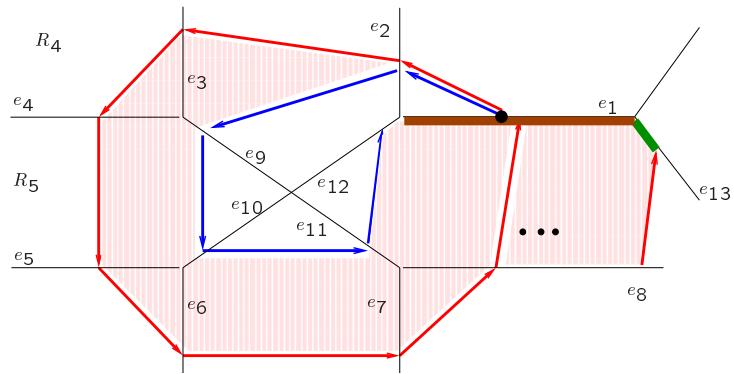


## Successors (by $\sigma$ )

- One cycle ( $\sigma = s = e_1 e_2 \cdots e_8 e_1$ )



- One cycle iterated: *solution of fixpoint equation (acceleration)* ( $\sigma = (s)^* e_{13}$ )



## Reachability Algorithm

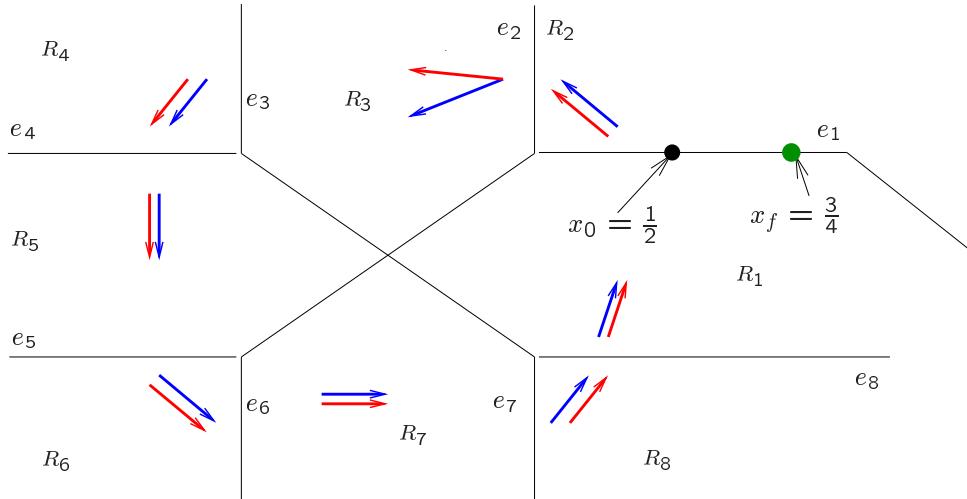
- We have that

$$\begin{aligned}x_f \in \text{Reach}(x_0) \\ \text{iff} \\ \text{for some } \sigma, \quad x_f \in \text{Succ}_\sigma(x_0)\end{aligned}$$

- **Algorithm:** Let  $x_0 \in e_0$  and  $x_f \in e_f$ .  
For each type  $\text{type}(\sigma) = r_1(s_1)^* \dots (s_n)^* r_{n+1}$ ,  
with  $e_0 = \text{first}(r_1)$  and  $e_f = \text{last}(r_{n+1})$   
check whether

$$x_f \in \text{Succ}_\sigma(x_0)$$

## Reachability Algorithm (Example)

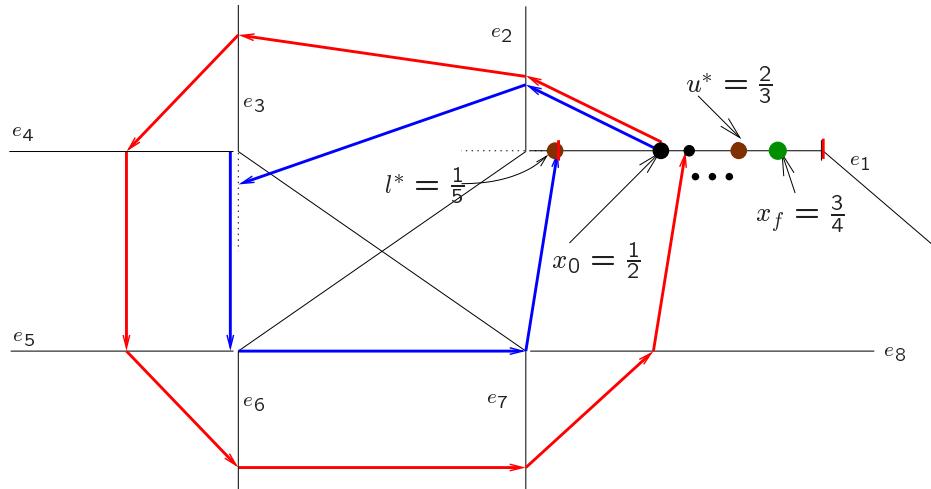


- Type of si nature:  $\sigma = (e_1 \cdots e_8)^*$
- Successor for the loop  $s = e_1 \dots e_8$ :

$$\text{Succ}_{e_1 \dots e_8}(l, u) = [\frac{l}{2} - \frac{1}{10}, \frac{u}{2} + \frac{1}{3}] \cap (\frac{1}{5}, 1)$$

## Reachability Algorithm (Example)

- Fixpoint equation:  $\text{Succ}_{e_1 \dots e_8}(I^*) = I^*$
- Solution:  $I^* = [l^*, u^*] = [\frac{1}{5}, \frac{2}{3}]$
- Hence:  $\text{Succ}_{e_1 \dots e_8}(x_0) \subseteq [\frac{1}{5}, \frac{2}{3}]$



Conclusion:  $x_f \notin I^*$  ( $\frac{3}{4} \notin [\frac{1}{5}, \frac{2}{3}]$ ).  
Hence,  $x_f \notin \text{Reach}(x_0)$

**Note:** the solution was found without iterating  
(*acceleration*).

## Conclusions (Achievements)

- Reachability is decidable for SPDI
- First application of “eometric” methods to planar non-deterministic systems
  - For deterministic planar systems:
    - \* PCD: Piece-wise Constant Derivatives (Maler & Pnueli - CAV'93)
    - \* Planar Piece-wise Hamiltonian Systems (Ācerāns & J. Vīksna - HS III)

## Conclusions (Future Works)

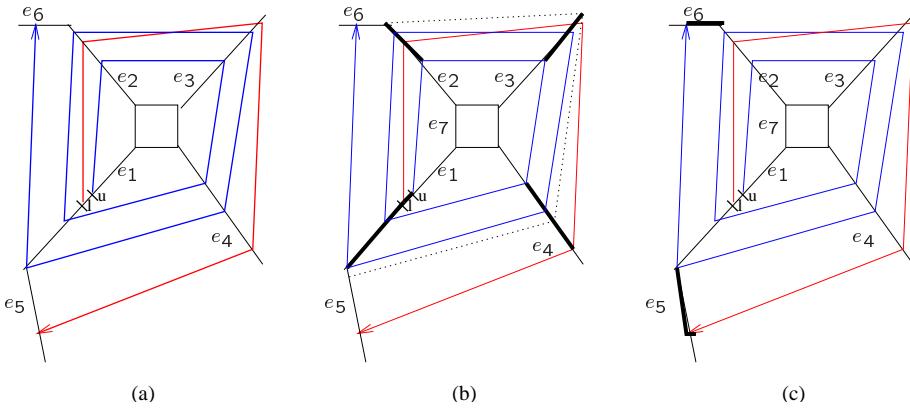
- Complexity analysis
- Extensions:
  - Hierarchical SPDI's
  - 2-dim surfaces (ex.: torus)
  - Finding limit cycles
- Applications:
  - Approximation of non-linear differential equations

## “Types” of cycles

1. **STAY** The cycle is not abandoned by any of the two trajectories:  $L \leq l^* \leq u^* \leq U$ .
2. **DIE** The right trajectory exits the cycle through the left (consequently the left one also exits):  $u^* < L \vee l^* > U$ .
3. **EXIT-BOTH** Both trajectories exit the cycle (the left one through the left and the right one through the right):  
$$l^* < L \wedge u^* > U.$$
4. **EXIT-LEFT** The leftmost trajectory exits the cycle but not the other:  
$$l^* < L \leq u^* \leq U.$$
 Dual: **EXIT-RIGHT**.

## “Types” of cycles

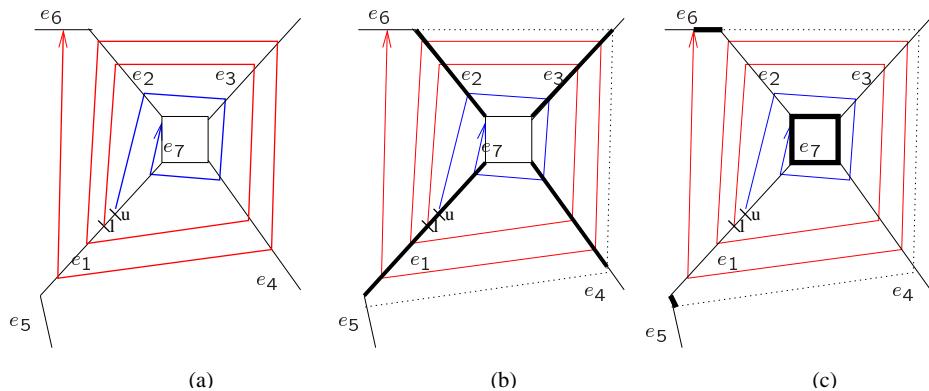
- Die



- (a) Both trajectories leave the cycle  $(e_1, e_2, e_3, e_4)^*$  through the left, whereas the right one tends to the limit  $u^*$ ;
- (b) Reachable points on the cycle (in bold);
- (c) Possible continuation after leaving the cycle (in bold).

## “Types” of cycles

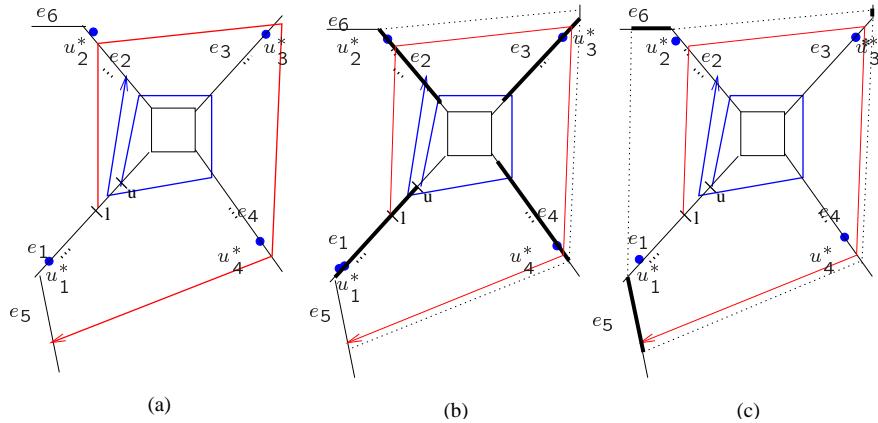
- Exit Both



- (a) Both trajectories leave the cycle  $(e_1, e_2, e_3, e_4)^*$ ;
- (b) Reachable points on the cycle (in bold);
- (c) Possible continuation after leaving the cycle (in bold).

## “Types” of cycles

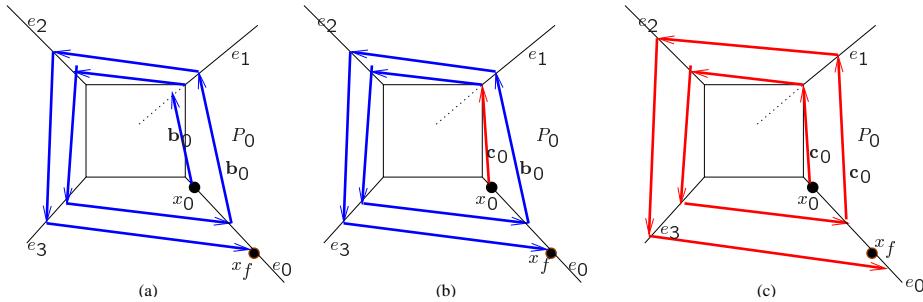
- Exit Left



- (a) The left trajectory leave the cycle  $(e_1, e_2, e_3, e_4)^*$  through the left;
- (b) Reachable points on the cycle (in bold);
- (c) Possible continuation after leaving the cycle (in bold).

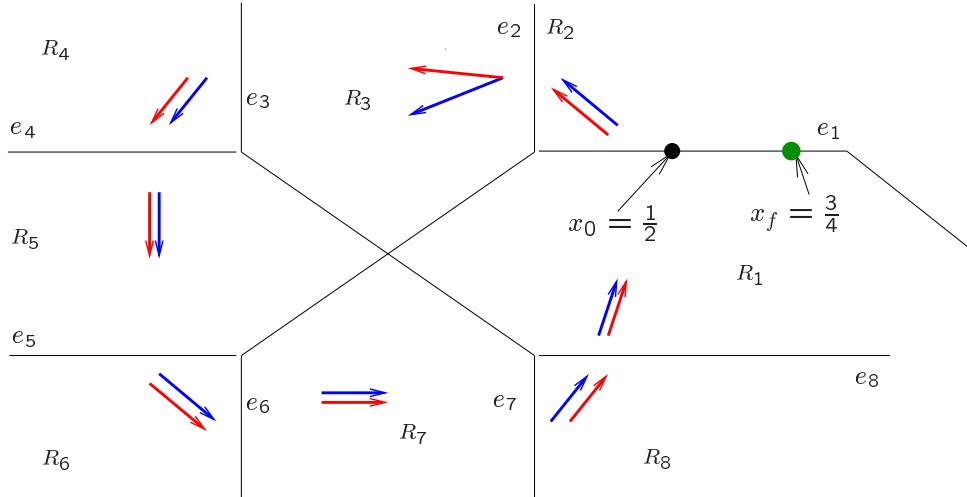
## Parametric Synthesis Problem

- Decidability of reachability for SPDI does not imply decidability for the parametric case.



- (a)  $x_f$  is reachable from  $x_0$  following the leftmost (truncated) trajectory;
- (b) Parameter must be “at least”  $c_0$ ;
- (c) Takin  $c_0$  makes  $x_f$  unreachable.

## Reachability Algorithm (Example)

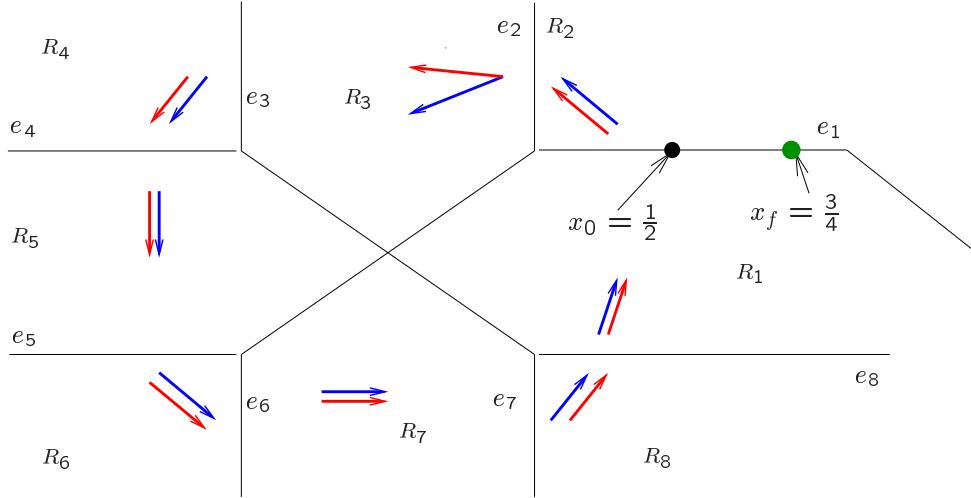


- Dynamics:

- $a_1 = b_1 = (1, 5),$
- $a_2 = b_2 = (-1, \frac{1}{2}),$
- $a_3 = (-1, \frac{11}{60})$  and  $b_3 = (-1, -\frac{1}{4}),$
- $a_4 = b_4 = (-1, -1),$
- $a_5 = b_5 = (0, -1),$
- $a_6 = b_6 = (1, -1),$
- $a_7 = b_7 = (1, 0),$
- $a_8 = b_8 = (1, 1).$

- Type of signature:  $\sigma = (e_1 \cdots e_8)^*$

## Reachability Algorithm (Example)



- Successor functions are:

$$\text{Succ}_{e_1 e_2}(x) = \left[ \frac{x}{2}, \frac{x}{2} \right] \cap (0, 1)$$

$$\text{Succ}_{e_2 e_3}(x) = \left[ x - \frac{3}{10}, x + \frac{2}{15} \right] \cap (0, 1)$$

$$\text{Succ}_{e_i e_{i+1}}(x) = [x, x] \cap (0, 1), \text{ for all } i \in 3..7$$

$$\text{Succ}_{e_8 e_1}(x) = \left[ x + \frac{1}{5}, x + \frac{1}{5} \right] \cap (0, 1)$$

- Successor for the loop  $s = e_1 \dots e_8$ :

$$\text{Succ}_{e_1 \dots e_8}(l, u) = \left[ \frac{l}{2} - \frac{1}{10}, \frac{u}{2} + \frac{1}{3} \right] \cap \left( \frac{1}{5}, 1 \right)$$