

A Constraint-Based Algorithm for Contract-Signing Protocols

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Joint work with Ralf Küsters

Overview

- Introduction / Motivation
- Formal model / Problem
- Constraint solving / Algorithm
- Conclusion

Contract signing

Example: Alice wants to buy a house from Bob

- 1) Parties agree upon contractual text
- 2) Both parties sign a copy
- 3) Parties exchange signed copies simultaneously



Alice

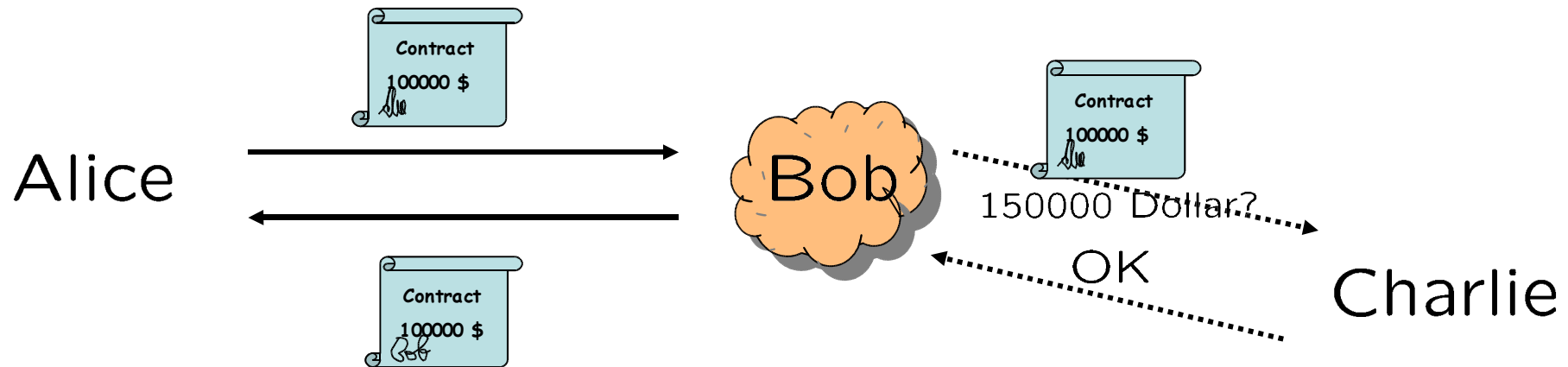
Bob



Contract signing over a network

How to do the exchange of signed copies?

Naïve approach

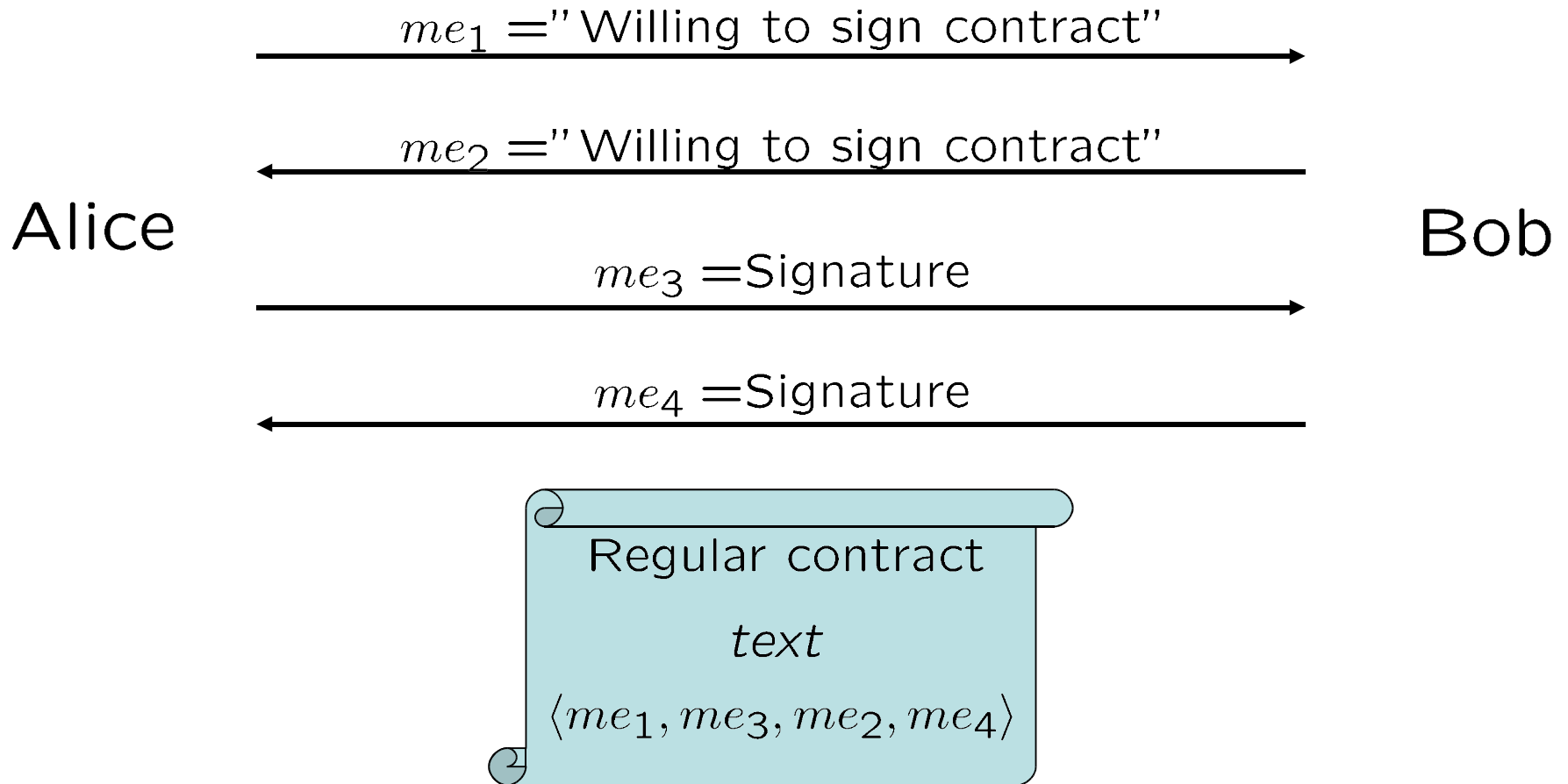


A contract-signing protocol is **balanced for Alice** if at no stage of the protocol execution Bob has both

- a strategy to obtain a valid contract, and
- a strategy to prevent Alice from getting a valid contract.

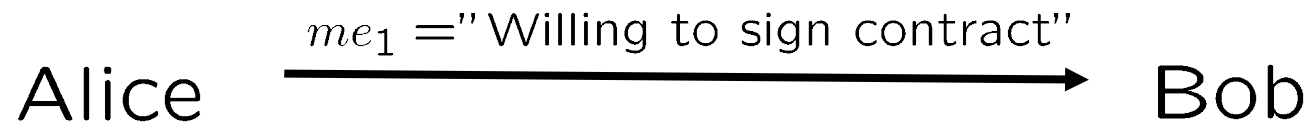
ASW-protocol

ASW-protocol: Optimistic two party contract-signing protocol



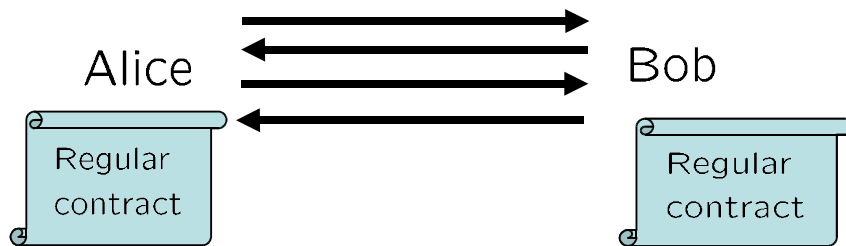
ASW-protocol

Abort subprotocol

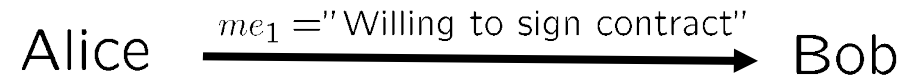


ASW-protocol

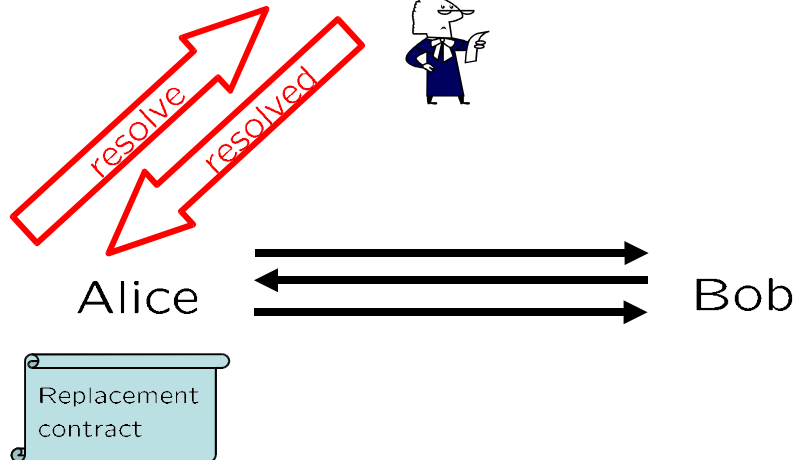
Main protocol



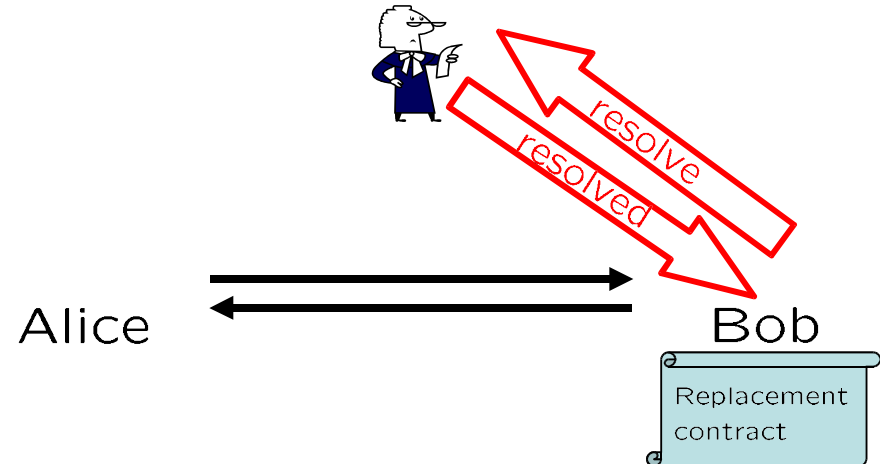
Abort subprotocol



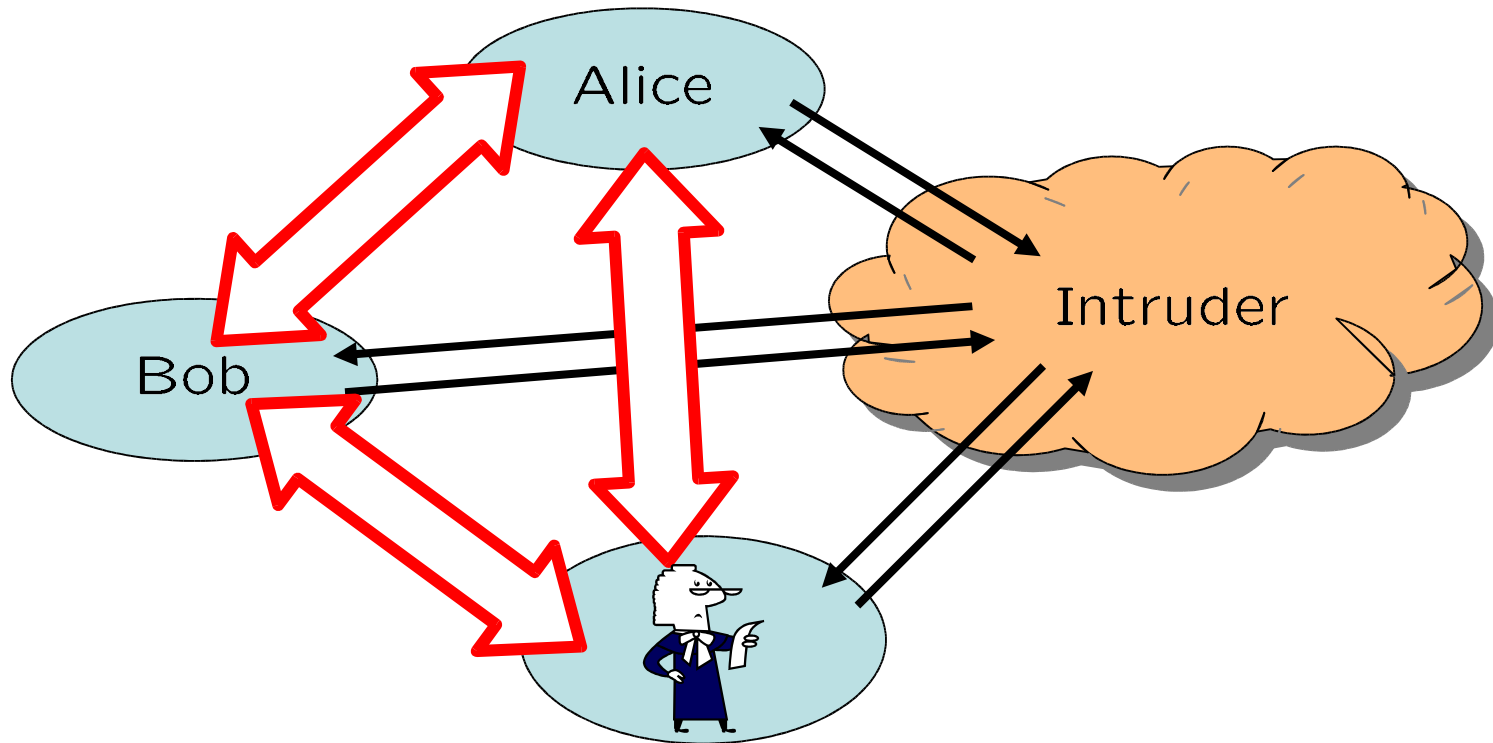
Resolve Subprotocol (Alice)



Resolve Subprotocol (Bob)



Communication model (extended Dolev-Yao)



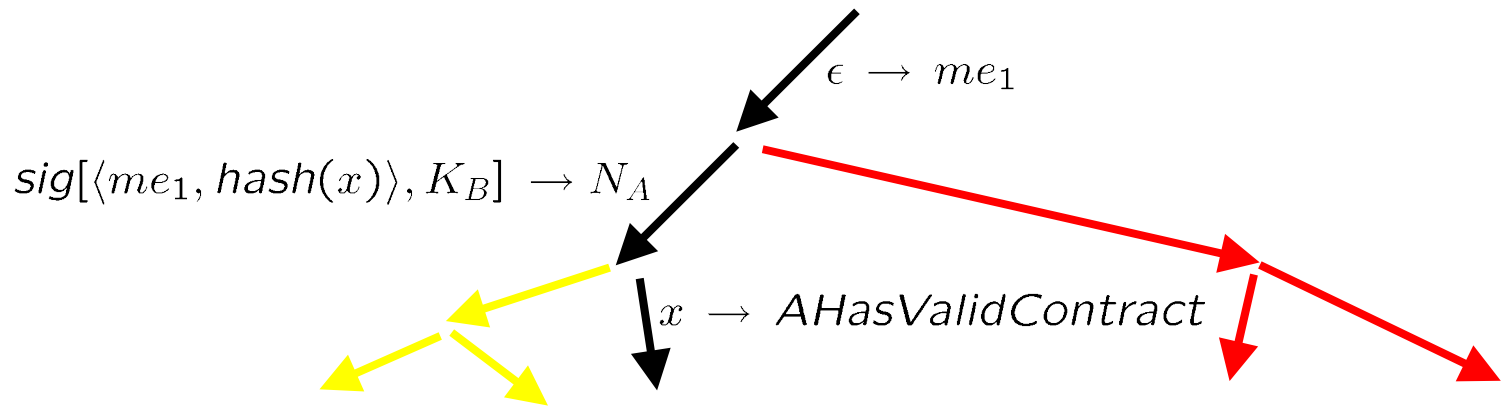
- The network is the intruder
- Secure channels

Finite-session model of a model
proposed by Scedrov et al. [2001]

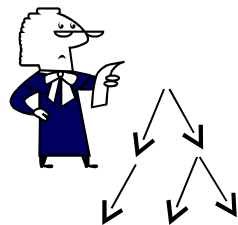
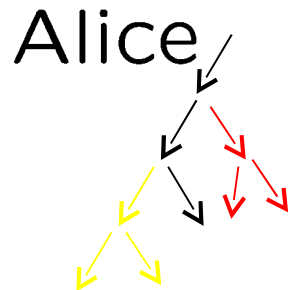
Formal model

Participants are rule trees

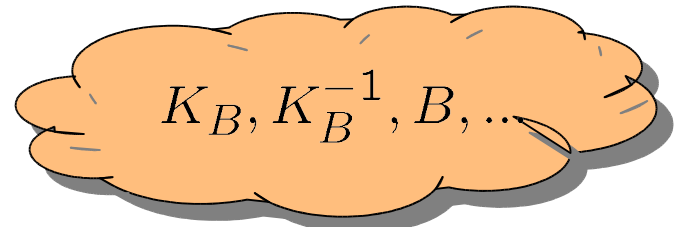
Alice (ASW-protocol) as rule tree Π_A



A protocol P consists of
a family of participants + initial intruder knowledge



$$P = (\{\Pi_i\}_i, K)$$

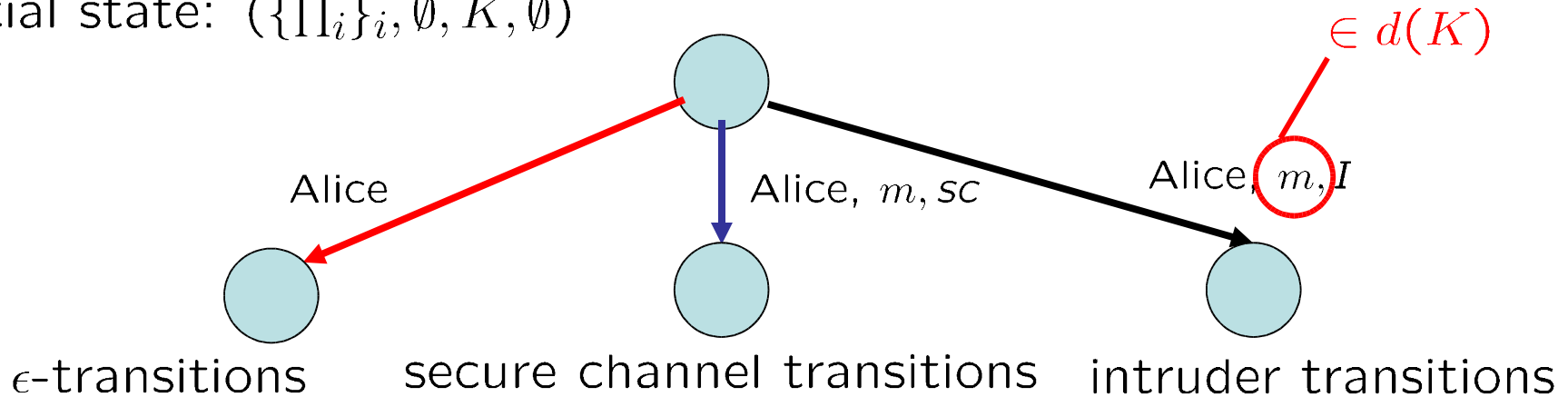


Protocol States

Given a protocol $P = (\{\Pi_i\}_i, K)$.

State of protocol P : $(\{\Pi'_i\}_i, \sigma, K', S)$

Initial state: $(\{\Pi_i\}_i, \emptyset, K, \emptyset)$

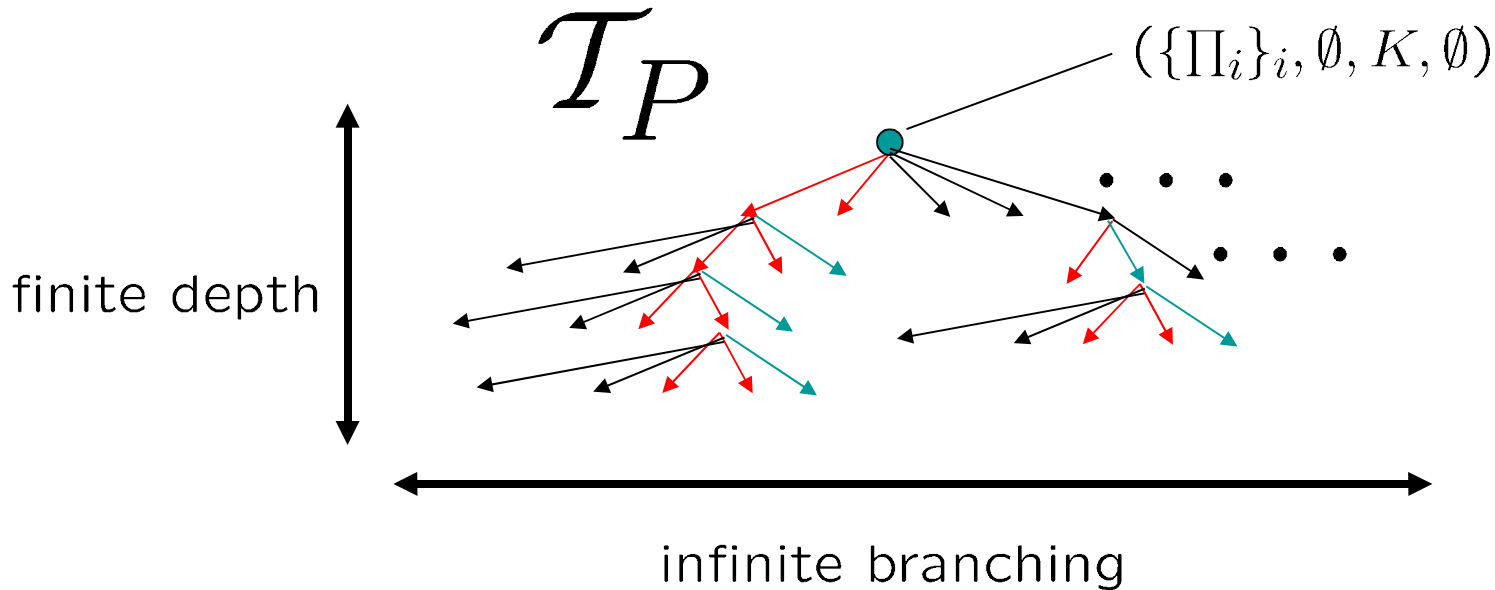


$d(K)$: set of messages constructable by the intruder

Transition tree

Given a protocol $P = (\{\Pi_i\}_i, K)$.

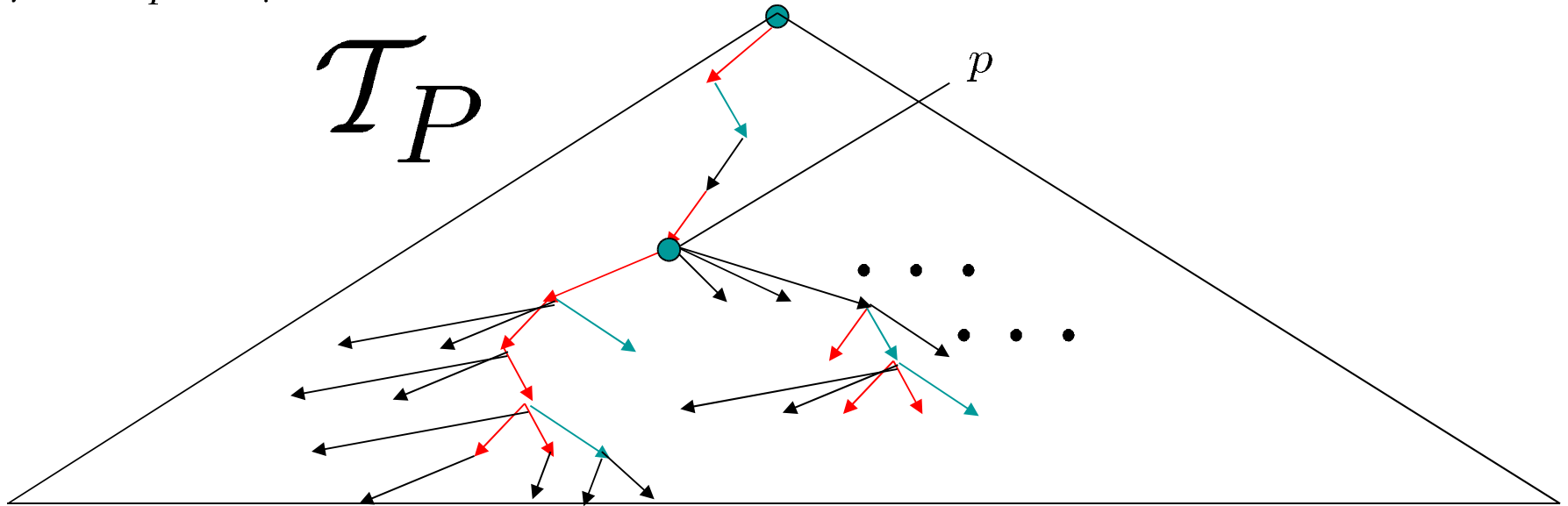
What is the execution of protocol P ?



The set of runs can be thought of as a tree rooted at $(\{\Pi_i\}_i, \emptyset, K, \emptyset)$

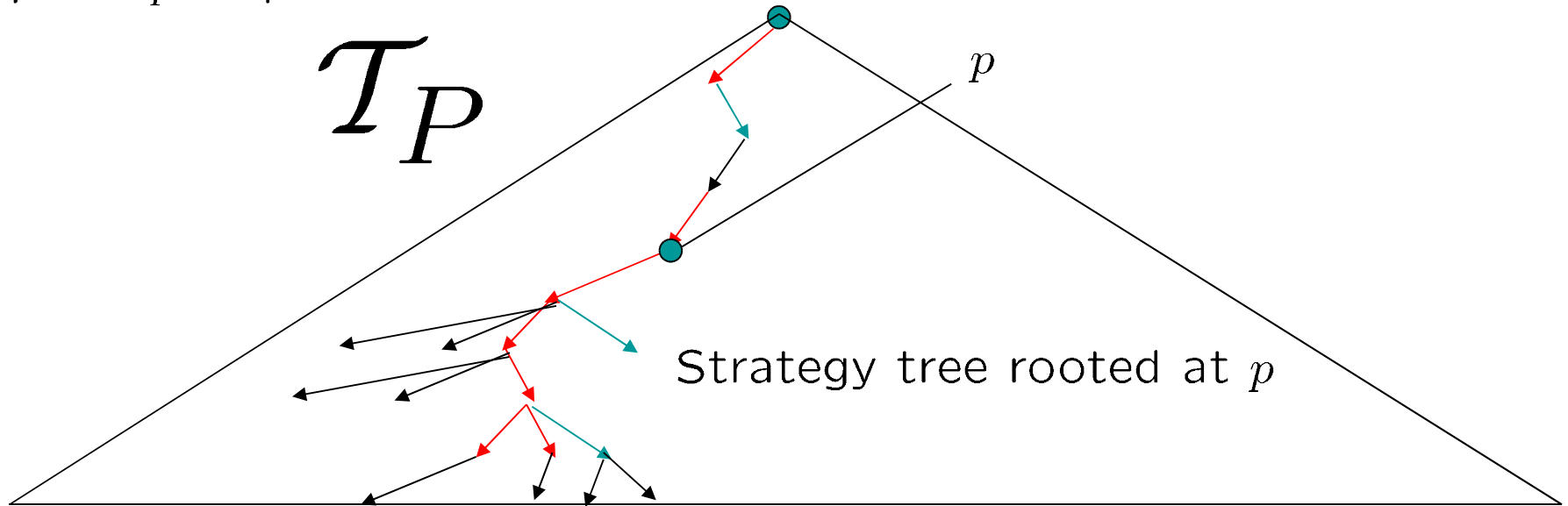
Intruder Strategies (1)

Balance talks about intruder strategies.
What is an intruder strategy at some
point p of protocol execution?



Intruder Strategies (2)

Balance talks about intruder strategies.
What is an intruder strategy at some
point p of protocol execution?



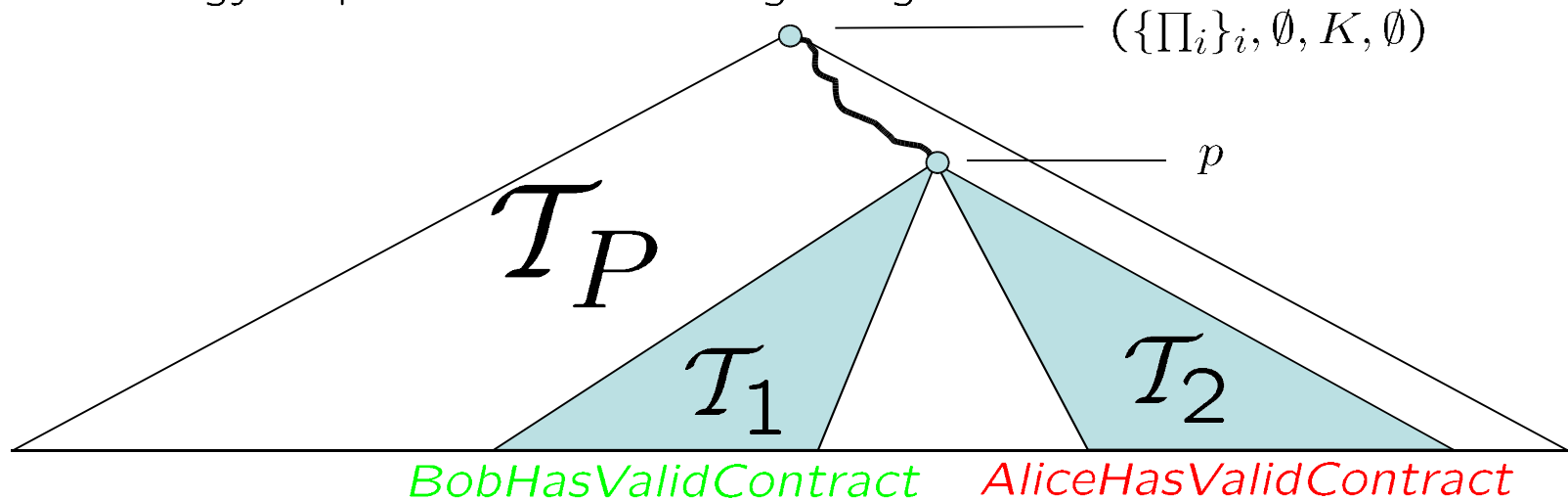
- Intruder may choose not to take some intruder transitions
- ε -transitions and secure channel transitions are not under control of the intruder

Strategy properties

How can we formulate game-theoretic security requirements?

A contract-signing protocol is **unbalanced for Alice** if there exists a point in the protocol execution where Bob has both

- a strategy to obtain a valid contract, and
- a strategy to prevent Alice from getting a valid contract.



Problem:

Given a protocol P and a strategy property C
Is there a state $p \in \mathcal{T}_P$ that satisfies C ?

Problem is decidable K., Küsters [2005]

Constraint solving

A **constraint** is of the form

$$m : T$$

where m is a term and T is a finite set of terms

A **constraint system** is a sequence of constraints

$$\begin{array}{l} m_1 : T_1 \\ m_2 : T_2 \\ \vdots \\ m_k : T_k \end{array}$$

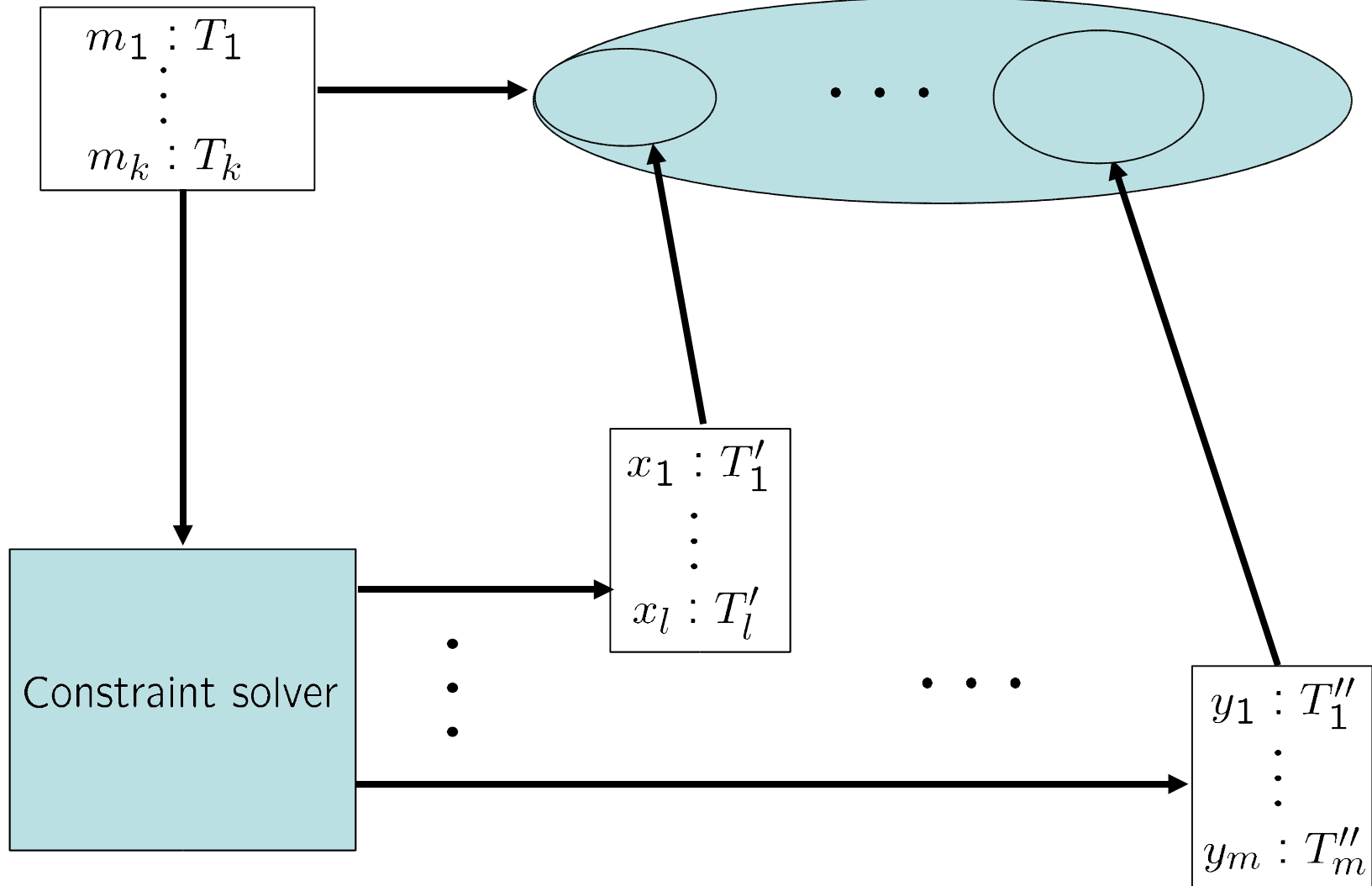
A **solution** of a constraint system \mathbf{C} is a substitution σ of the variables in \mathbf{C} by messages such that

$$\begin{array}{l} \sigma(m_1) \in d(\sigma(T_1)) \\ \sigma(m_2) \in d(\sigma(T_2)) \\ \vdots \\ \sigma(m_k) \in d(\sigma(T_k)) \end{array}$$

Constraint solver (1)

Constraint system \mathcal{S}

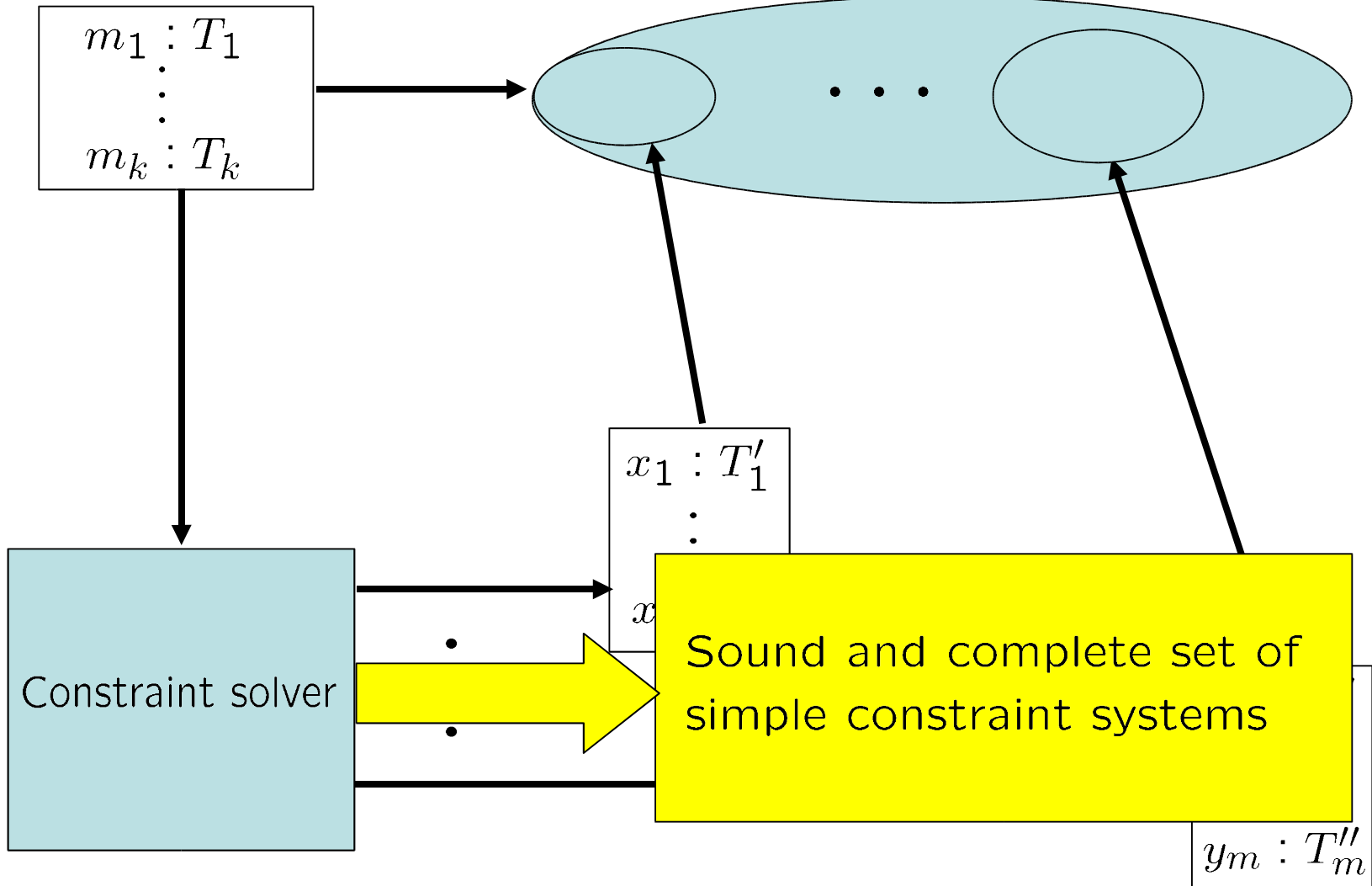
Set of solutions for \mathcal{S}



Constraint solver (2)

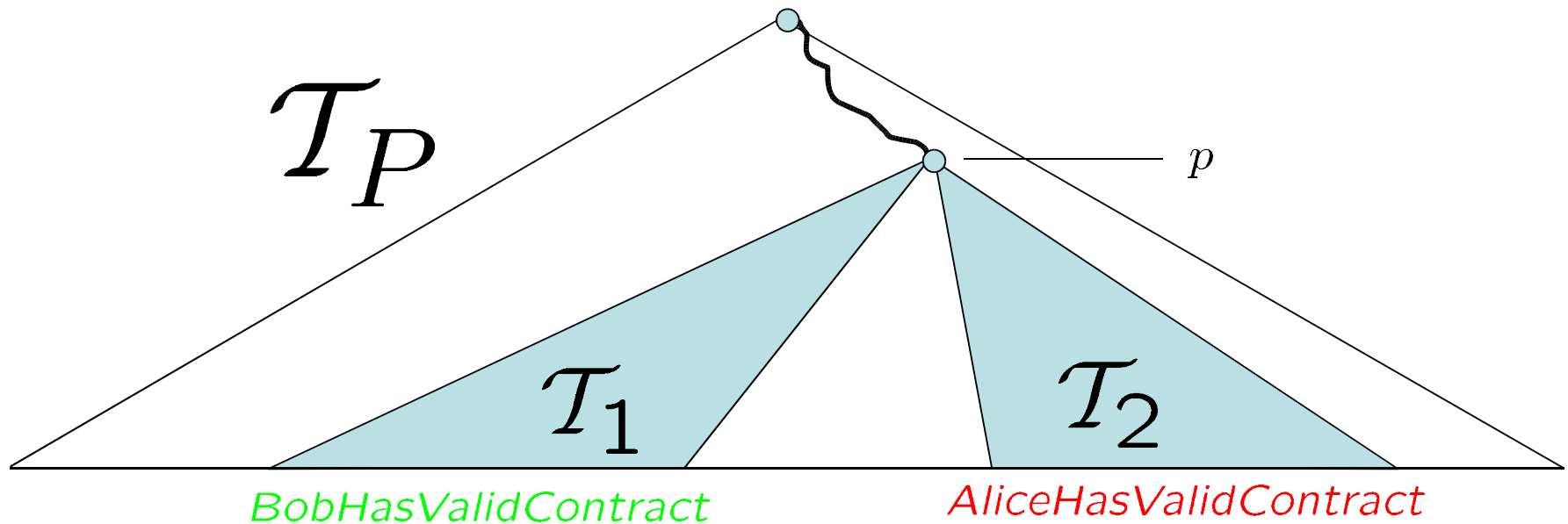
Constraint system \mathcal{S}

Set of solutions for \mathcal{S}



Algorithm

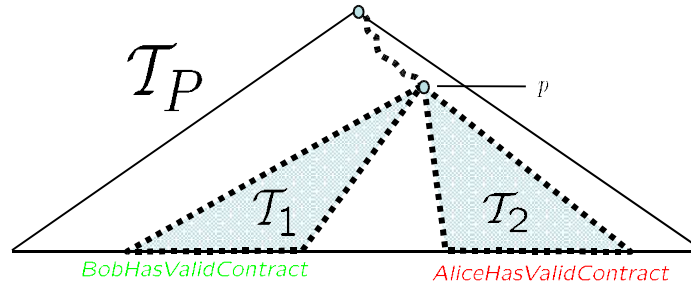
Given: Protocol P and a strategy property C
Is there a state $p \in \mathcal{T}_P$ that satisfies C ?



Algorithm

Given: Protocol P and a strategy property C
Is there a state $p \in \mathcal{T}_P$ that satisfies C ?

Guess
symbolic
attack



Construct
and solve
constraint
system

$m_1 : T_1$
 $m_2 : T_2$

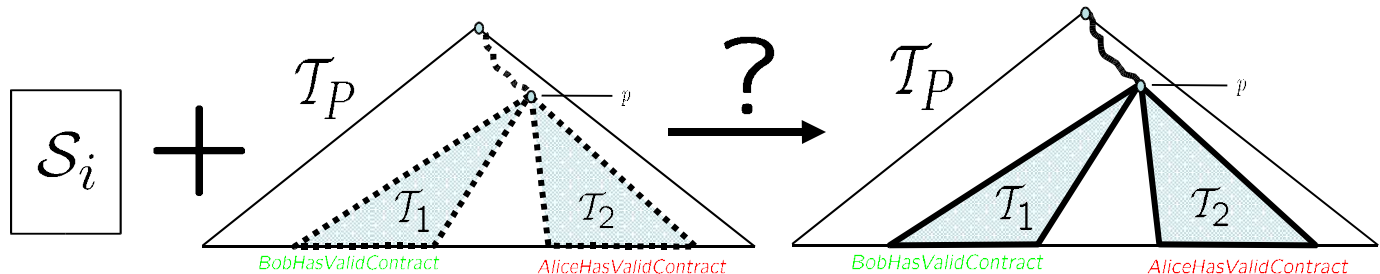
 $m_{17} : T_{17}$

Constraint
solver

sound and complete

\mathcal{S}_1 \mathcal{S}_2 \dots \mathcal{S}_b

Check
solutions



Related work

- Contract-signing protocols
Asokan, Shoup, and Waidner [1998]
Garay, Jacobsson, and MacKenzie [1999]
• • •
- Finite state analysis of contract-signing protocols
Mitchell, Shmatikov [2001]
Kremer, Raskin [2002]
• • •
- Infinite state analysis
Chadha, Kanovich, and Scedrov [2001]
• • •
- Constraint solving
Millen, Shmatikov [2001]
• • •

Conclusions

- Studied game-theoretic properties of infinite transition graphs induced by cryptographic protocols.
- Showed that balance and related game-theoretic properties are decidable using constraint solving algorithm.
- Future work: implementation, complexity analysis