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Graphics Hardware

UlfAssarsson

Graphics hardware – why?

• Often said to be "100x" faster than CPU.

- Reason: Simple to parallelize triangle rendering :
 - over individual triangles, pixels, (even over x,y,z,w, and r,g,b,a)
 - Hardware fixed functions: clipping, rasterizer, texture filtering, fragment-merge, ...

• Current hardware:

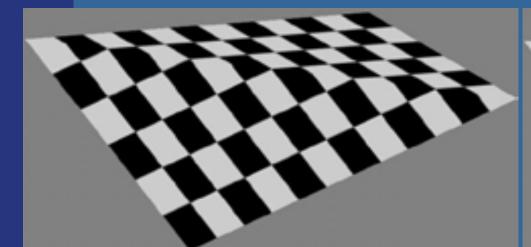
- Triangle rasterization with programmable shading.
- Massive parallel general-purpose computations:
 - CUDA/OpenCL/Compute Shaders (~4000 ALUs)
- AI computations:
 - ~500 tensor cores, each performing a 4x4-matrix mul+add.
- GPU Ray tracing:
 - NVIDIA RTX (via OptiX, Vulcan, Microsoft DXR api)
 - Although, can write your own GPU ray-tracer (e.g., CUDA or shader based)

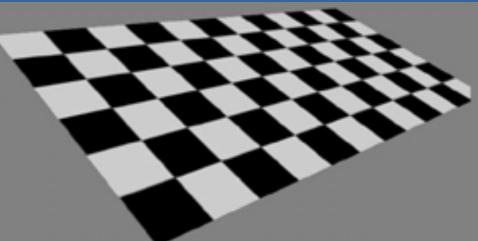
Perspective-correct interpolation of texture coordinates (and actually all screen-space-interpolated pervertex data)



Perspective-correct texturing

How is texture coordinates interpolated over a triangle?Linearly?





Linear interpolation

Perspective-correct interpolation

- Perspective-correct interpolation gives foreshortening effect!
- Hardware does this for you, but you need to understand this anyway!

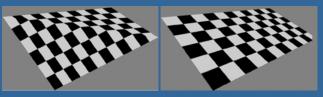
Recall the following

Vertices are projected onto screen by non-linear transform. Hence, tex coords cannot be linearly interpolated in screen space (just like a 3Dposition cannot be).

- Perspective projection introduces a non-linear transform by the homogenization step:
 - Projection: $\mathbf{p} = \mathbf{M}\mathbf{v}$
 - After projection p_w is not 1!
 - Homogenization: $(p_x/p_w, p_y/p_w, p_z/p_w, 1)$
 - Gives (x, y, z, 1), where x, y are the screen-space coordinates and z is depth

$$\mathbf{p} = \mathbf{M}\mathbf{v} = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & -1/d & 0 \end{pmatrix} \begin{pmatrix} v_x \\ v_y \\ v_z \\ 1 \end{pmatrix} = \begin{pmatrix} v_x \\ v_y \\ v_z \\ -v_z/d \end{pmatrix}$$

Perspective-correct interpolation



- Linear interpolation in screen space does not work for u,v
- Why:
 - We have applied a non-linear transform to each vertex position (x/w, y/w, z/w, w/w).
 - Non-linear due to 1/w factor from the homogenisation
 - Surprisingly, we can screen-space interpolate any vertex attribute a/w (including 1/w) perspective correctly.
 - For a proof, see Jim Blinn,"W Pleasure, W Fun", IEEE Computer Graphics and Applications, p78-82, May/June 1998

• Solution:

 Interpolate (u/w, v/w, 1/w), from each vertex, where w is from homogeneous coordinate (x,y,z,w). (Screen-space coord is (x/w, y/w, z/w, 1))

• Then at each pixel, get u_i,v_i as:

$$- w_i = 1 / (1/w)$$

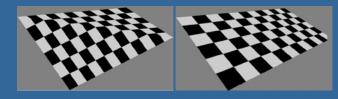
$$- u_i = (u/w)_i * w_i$$

 $- v_i = (v/w)_i * w_i$

 $(U_2/W_2, V_2/W_2, 1/W_2)$ $(u/w)_i$, $(v/w)_i$, $(1/w)_i$ $(u_1/w_1, v_1/w_1, 1/w_1)$ (40 mo, Vomo, 7/mo)

Shading is automatically interpolated this way too (though, not as annoying as textures). Perspective correct interpolation nowadays handled automatically by the GPU.

Perspective-correct interpolation



- Why we can screen-space interpolate an attribute a/w perspective correctly:
 - Let (x,y,z,w) be the vertex' homogeneous coordinate (after mult. with the modelViewProjectionMatrix).
 - Then (x/w, y/w, z/w, 1) is its screen-space position.
 - In screen-space, we will linearly interpolate between the vertices $\left(\frac{x_0}{w_0}, \frac{y_0}{w_0}, \frac{z_0}{w_0}\right), \left(\frac{x_1}{w_1}, \frac{y_1}{w_1}, \frac{z_1}{w_1}\right), \left(\frac{x_2}{w_2}, \frac{y_2}{w_2}, \frac{z_2}{w_2}\right).$
 - And we know we could transform back an interpolated position $\left(\frac{x_i}{w_i}, \frac{y_i}{w_i}, \frac{z_i}{w_i}\right)$ to homogeneous space if we had w_i to muliply with.
 - Hence, this interpolated position is perspectively-correct interpolated regarding
 - the screen-space z-position $\left(\frac{z_i}{w}\right)$ (of course x, and y as well). And z is not special.
 - We see that we can interpolate any value a/w correctly, if we have w_i . So, we need w_i perspective-correctly interpolated.
 - We cannot interpolate the special case w/w since that =1.To get w_i , we let a=1 to linearly interpolate 1/w. Then, we get $w_i=1/(1/w)_i$.

Overview of GPU architecture

-History / evolution

- GPU design: Several **cores** consisting of many **ALU**s (NVIDIA terminology: **Streaming Multiprocessors (SMMs)** of many **cores**
- GPU vs CPU

Take-away: bandwidth (cost of memory accesses) is a major problem

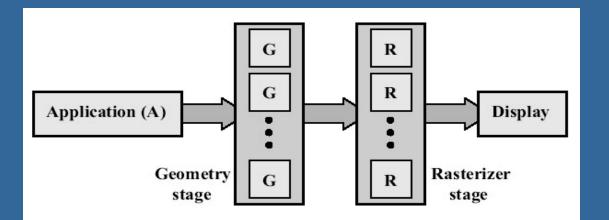
Background: Graphics hardware architectures

- Evolution of graphics hardware has started from the end of the pipeline
 - Rasterizer was put into hardware first (most performance to gain from this)
 - Then the geometry stage
 - Application will not be put into GPU hardware (?)
- Two major ways of getting better performance:
 - Pipelining
 - Parallellization
 - Combinations of these are often used

Parallellism

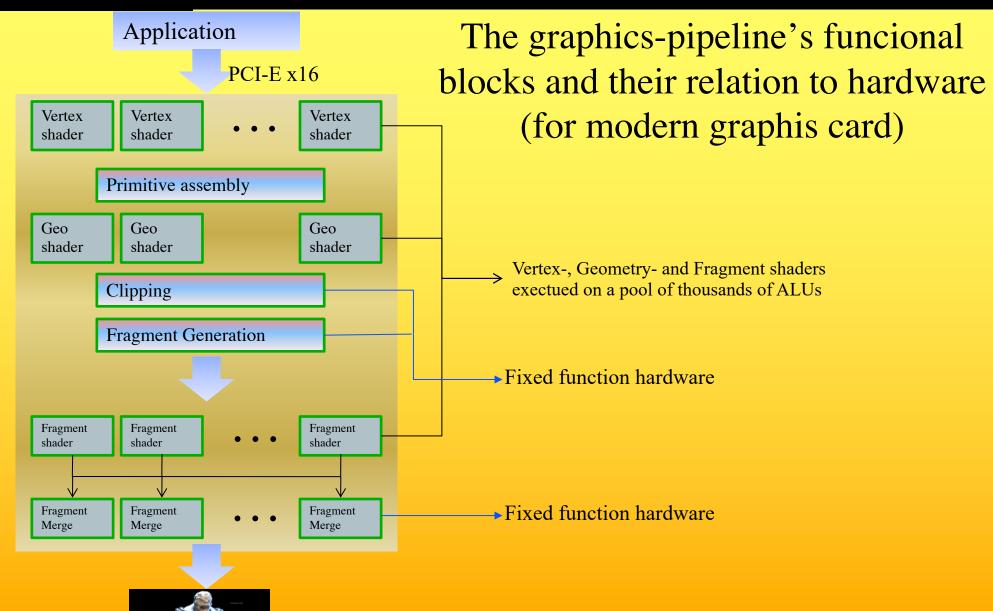
- "Simple" idea: compute n results in parallel, then combine results
- Not always simple
 - Try to parallelize a sorting algorithm...
 - But vertices are independent of each other, and also pixels, so simpler for graphics hardware

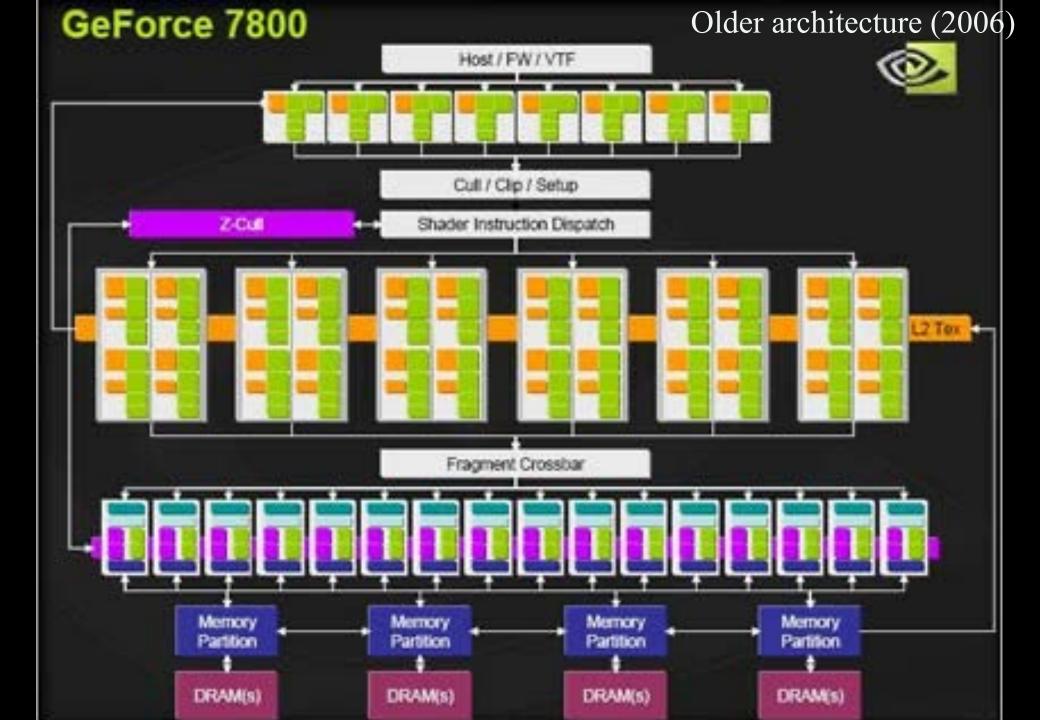
• Can parallellize both geometry and rasterizer stage:



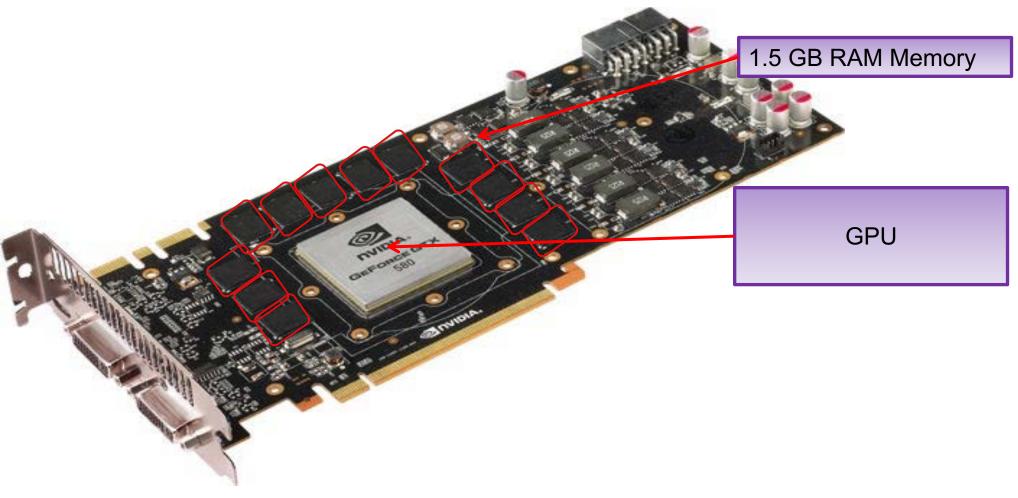
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Graphics Processing Unit - GPU



NVIDIA Geforce GTX 580

NVIDIA Maxwell (GTX 980)

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2014

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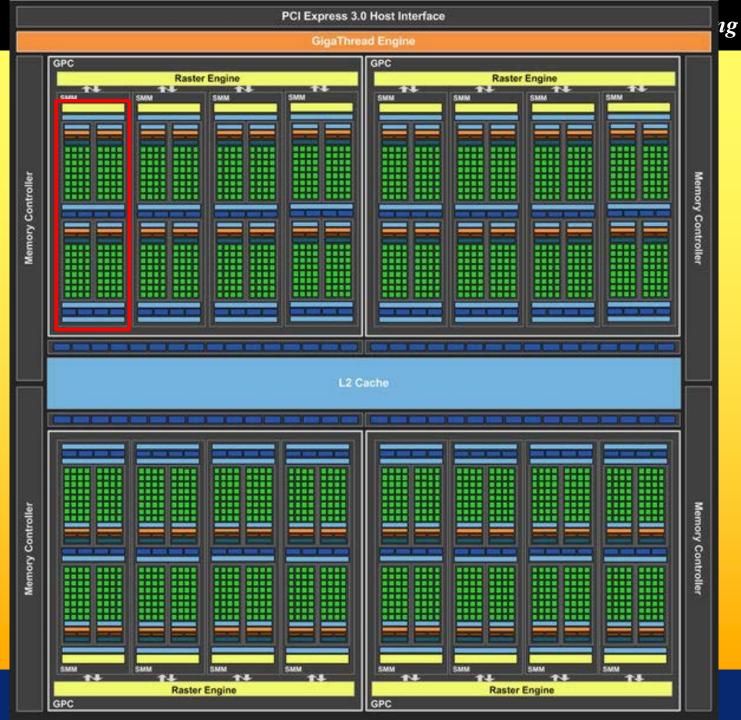
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16 Cores ("SMM") 2MB L2 cache 64 output pixels / clock (i.e., 64 ROPs) 2048 ALUs ("cores") ~6 Tflops

Each Core:

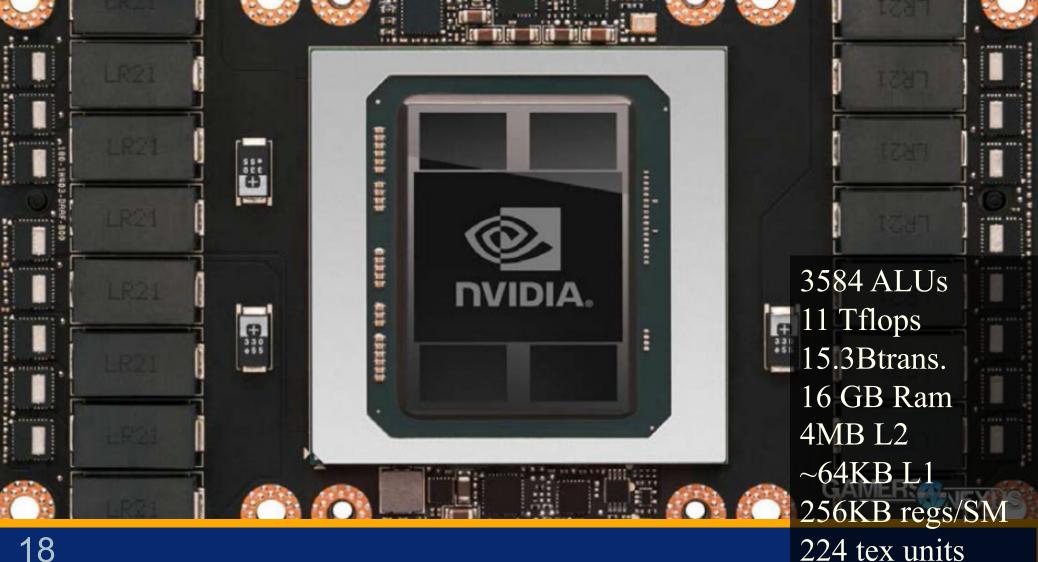
- 128 ALUs
- 96KB L1 cache
- 8 TexUnits
- 32 Load/Store units for access to global memory



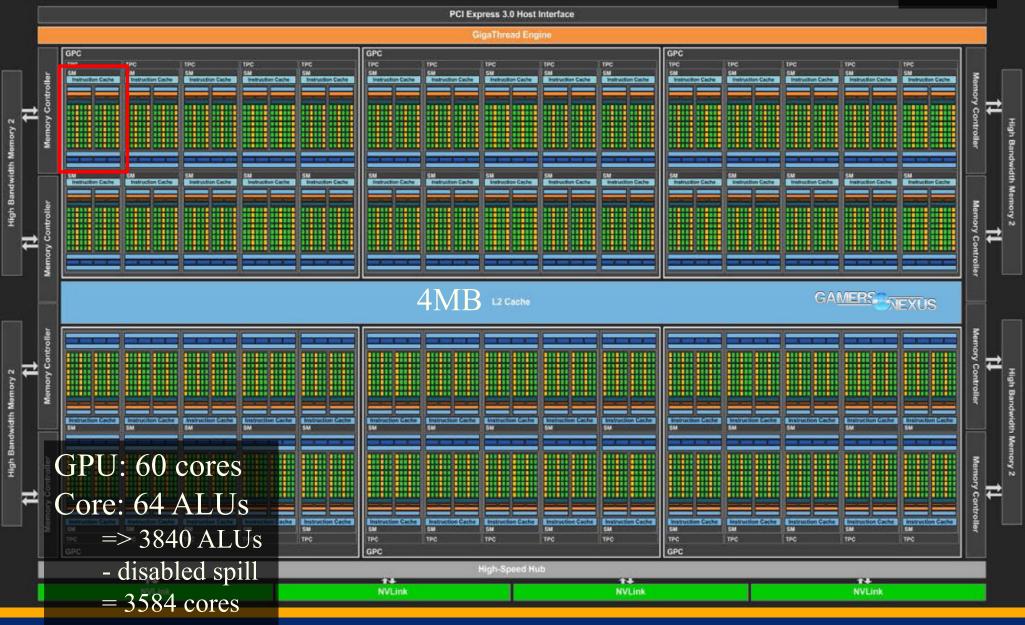
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NVIDIA Pascal GP100 (GTX 1080 / Titan X)

2016



Department of Computer 2016



SM

GAMER: NEXUS

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NVIDIA Volta GV100

2018 (Dec. 2017)



NVIDIA Volta GV100

2018



NVIDIA Turing TU102

2018

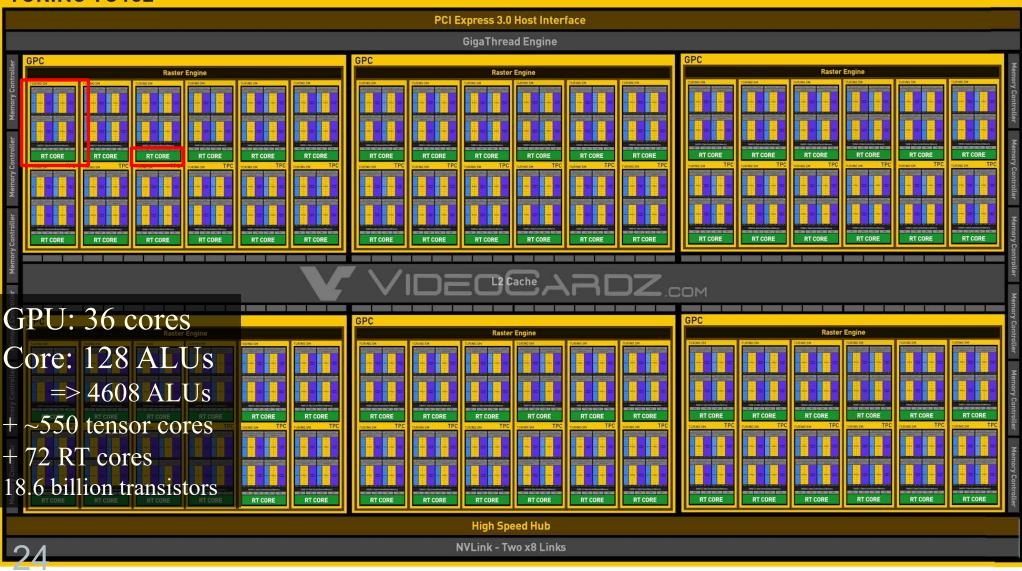
h to H

GPU: 36 cores Core: 128 ALUs => 4608 ALUs + ~550 tensor cores + 72 RT cores 18.6 billion transistors

NVIDIA Turing TU102

2018

TURING TU102



NVIDIA Ampere

GPU: 82 cores Core: 128 ALUs => 10496 ALUs + ~328 tensor cores + 82 RT cores 28.3 billion transistors



2020

NVIDL

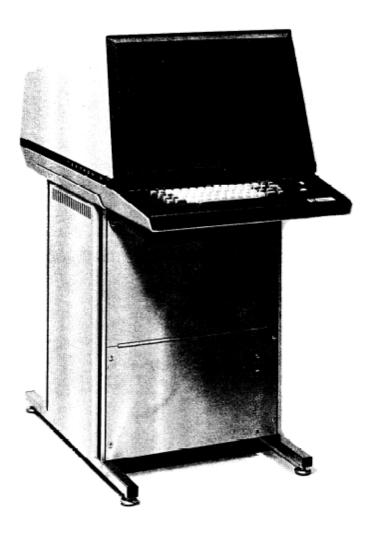


GPC GPC BM2 GPU: 82 cores Core: 128 ALUs SM \Rightarrow 10496 ALUs ~128KB L1\$ $+ \sim 328$ tensor cores SM SM TPC GPC +82 RT cores 28.3 billion transistors **NVLink NVLink** 26

Graphics Hardware History

Direct View Storage Tube:

- Created by Tektronix (early 70's)
 - -First with "frame buffer" (moveto/lineto)
 - -Did not require constant refresh
 - -Standard interface to computers
 - Allowed for standard software
 - Plot3D in Fortran
 - -Relatively inexpensive
 - Opened door to use of computer graphics for CAD community
 - 4096 * 4096 addressable points (4096 * 3120 viewable).



Graphics Hardware History - functionality

- 80's:
 - linear interpolation of color over a scanline
 - Vector graphics
- 91' Super Nintendo, Neo Geo,
 - Rasterization of 1 single 3D rectangle per frame (FZero)
- 95-96': Playstation 1, 3dfx Voodoo 1
 - Rasterization of whole triangles (Voodoo 2, 1998)
- 99' Geforce (256)
 - Transforms and Lighting (geometry stage)
- 02' 3DLabs WildCat Viper, P10
 - Pixel shaders, integers,
- 02' ATI Radion 9700, GeforceFX
 - Vertex shaders and **Pixel shaders** with floats
- 06' Geforce 8800
 - Geometry shaders, integers and floats, logical operations
- Then: More general multiprocessor systems, higher SIMD-width, more cores
- 09' Tesselation Shaders (Direct3D '09, OpenGL '10)
- 17' Tensor cores
- 18' RT cores, Mesh Shaders





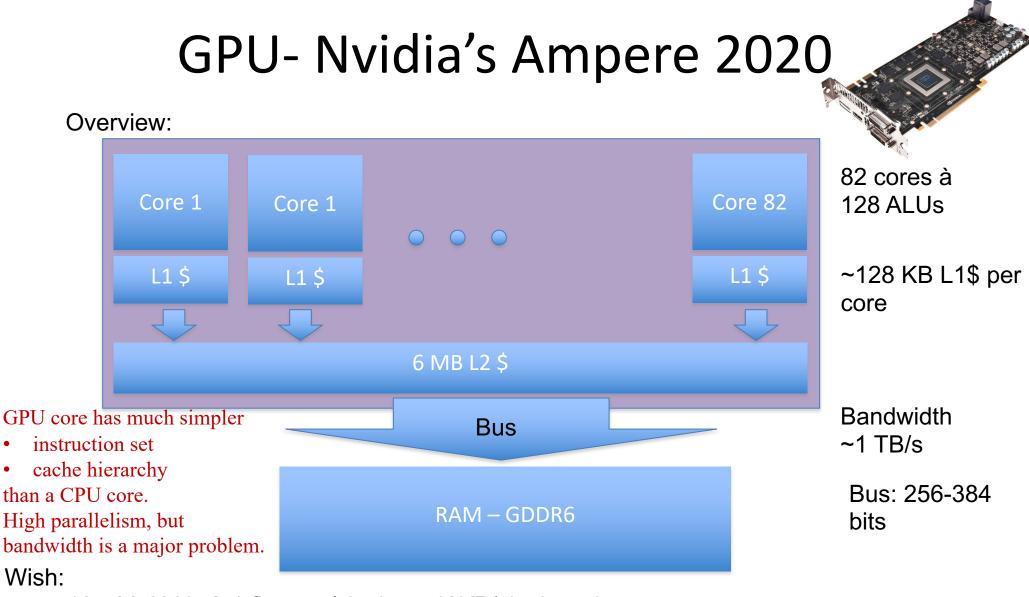
Graphics Hardware History - specs

2001 • In GeForce3: 600-800 pipeline stages! 57 million transistors

– First Pentium IV: 20 stages, 42 million transistors,

• Evolution of cards:

- 2004 X800 165M transistors
- 2005 X1800 320M trans, 625 MHz, 750 Mhz mem, 10Gpixels/s, 1.25G verts/s
- 2004 GeForce 6800: 222 M transistors, 400 MHz, 400 MHz core/550 MHz mem
- 2005 GeForce 7800: 302M trans, 13Gpix/s, 1.1Gverts/s, bw 54GB/s, 430 MHz core,mem 650MHz(1.3GHz)
- 2006 GeForce 8800: 681M trans, 39.2Gpix/s, 10.6Gverts/s, bw:103.7 GB/s, 612 MHz core (1500 for shaders), 1080 MHz mem (effective 2160 MHz), GDDR3
- 2008 Geforce 280 GTX: 1.4G trans, 65nm, 602/1296 MHz core, 1107(*2)MHz mem, 142GB/s, 48Gtex/s
- 2007 ATI Radeon HD 5870: 2.15G trans, 153GB/s, 40nm, 850 MHz, GDDR5, 256bit mem bus,
- 2010 Geforce GTX480: 3Gtrans, 700/1401 MHz core, Mem (1.848G(*2)GHz), 177.4GB/s, 384bit mem bus, 40Gtexels/s
- 2011 GXT580: 3Gtrans, 772/1544, Mem: 2004/4008 MHz, 192.4GB/s, GDDR5, 384bit mem bus, 49.4 Gtex/s
- 2012 GTX680: 3.5Gtrans (7.1 for Tesla), 1006/1058, 192.2GB/s, 6GHz GDDR5, 256-bit mem bus.
- 2013 GTX780: 7.1G, core clock: 837MHz, 336 GB/s, Mem clock: 6GHz GDDR5, 384-bit mem bus
- 2014 GTX980: 7.1G?, core clock: ~1200MHz, 224GB/s, Mem clock: 7GHz GDDR5, 256-bit mem bus
- 2015 GTX Titan X: 8Gtrans, core clock: ~1000MHz, 336GB/s, Mem clock: 7GHz GDDR5, 384-bit mem bus
- 2016 Titan X: 12/15Gtrans, core clock: ~1500MHz, <u>480GB/s, Mem clock: 10Gbps GDDR5X</u>, 4096-HBM2
- 2018 Nvidia Volta: 21.1Gtrans, core clock: ~1500MHz, 900GB/s, Mem: 4096-bit HBM2, (or GDDR6)
- 2020 Nvidia Ampere: 54 Gtrans, ~1500MHz, 1500GB/s, Mem: 4096-bit HBM2, (or 900BG/s GDDR6)
 Lesson learned: #trans doubles ~per 2 years. Core clock increases slowly. Mem clock –increases with new technology DDR2, DDR3, GDDR5/6, HBM2 and with more memory busses (à 64-bit). Now stacked.
 - We want as fast memory as possible! Why?
 - Parallelization can cover for slow core clock. Parallelization more energy efficient than high clock frequency; power consumption proportional to freq².
 - Memory transfers often the bottleneck

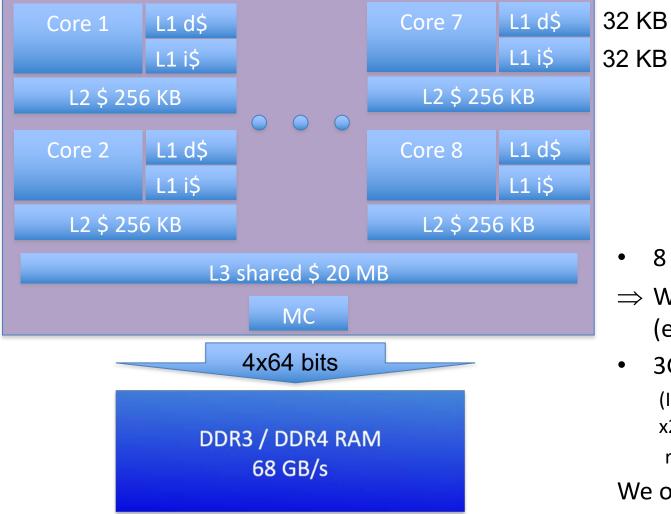


- ~10.500 ALUs à 1 float.op/clock => 42KB/clock cycle
- ~1.7GHz core clock => 71 TB/s request

We have ~1TB/s. Hence, would need to do ~70 computations between each RAM–read/write. Ameliorated by L1\$ + L2\$ + latency hiding (warp switching) but still a main problem!

Intel's Sandybridge / Haswell / Broadwell

CPU – 2014-2016



- 21500 vs 768 GB/s ≈ 30x diff.
- You could say bandwith is 2 orders of magnitude more important on GPU than CPU, due to parallelism.

1 – 8 cores à 8 SIMD floats

- 8 cores à 8 floats
- ⇒ We want 256 bytes/clock (e.g. from RAM).
- 3GHz CPU => 768 GByte/s.
 (In addition both for GPU & CPU x2, since:

r1 = r2 + r3;)

We only have 30-68 GB/s.

Solved by \$-hierarchy +

registers + thread switching

Memory bandwidth usage is huge!!

- On top of that bandwith usage is never 100%.
- However, there are many techniques to reduce bandwith usage:
 - Texture caching with prefetching
 - Texture compression
 - Hierarchical Z-occlusion testing
 - E.g., for every 8x8 pixel block of frame buffer, store its z_{min} , z_{max} .
 - If triangle is behind pixel block, skip rasterize it.
 - If triangle is in front, skip accessing 8x8 individual z-values.



Taxonomy of hardware design

for how to resynchronize (sort) parallelized work.

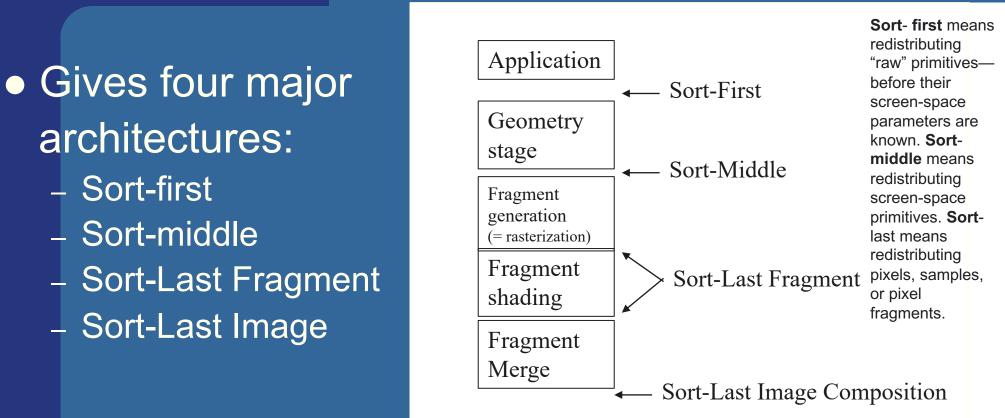
Outputs to frame buffers must respect incoming triangle order.

Take-aways: Sort-first, Sort-middle, Sort-Last Fragment, Sort-Last Image

Taxonomy of Hardware

- We can do many computations in parallel:
 - Pixel shading, vertex shading, geometry shading
- But result on screen must be as if each triangle were rendered one by one in their incoming order (according to OpenGL spec)
 - I.e., for every pixel, the rasterized fragments must be merged to the buffers in the original input triangle order
 - E.g., for blending/transparency, (z-culling + stencil test)
- Hence, results need to be sorted somewhere before reaching the screen...

Taxonomy of hardware Need to sort the results of the parallelization



 Will describe these briefly. Sort-last fragment (and sort middle) are most common in
 commercial hardware

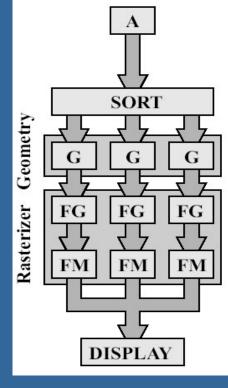
Sorting/dividing work to parallel execution units.

Sort-First

- Sorts primitives before geometry stage
 - Screen in divided into large regions
 - Blocks or scanlines
 - A separate pipeline is responsible for each region (or many)
- Not explored much at all, since:
 - Poor load balancing if uneven triangle distribution between regions.
 - Vertex shader can change triangle position

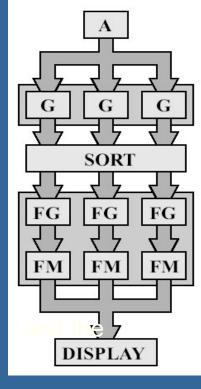


- A fragment is all the generated information for a pixel on a triangle
- FG is Fragment Generation (finds which pixels are inside triangle)
- FM is Fragment Merge (merges the created fragments with various buffers (Z, color))



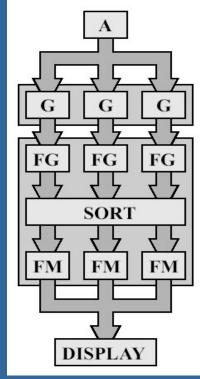
Sort-Middle

- Sorts between G and R
- Pretty natural, since after G, we know the screen-space positions of the triangles
- Older/cheaper hardware uses this
 - Examples include InfiniteReality (from SGI) KYRO architecture (from Imagination)
- Spread work arbitrarily among G's
- Then depending on screen-space position, sort to different R's
 - Screen can be split into "tiles". For example:
 - Rectangular blocks (8x8 pixels)
 - Every n scanlines
- The R is responsible for rendering inside tile
- Bads (same as Sort-First):
 - A triangle can be sent to many FG's depending on overlap (over tiles)
 - May give poor load balancing if triangles are unevenly distributed over the screen tiles



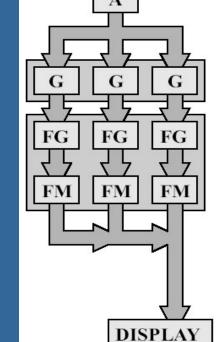
Sort-Last Fragment

- Sorts betwen FG and FM
- Most graphics cards use this.
 - Each pixel block responsible for sorting its fragments.
 - No need for redundant rasterization
 - \Rightarrow Hence cheaper to have small pixel blocks
 - \Rightarrow Better for load balancy
- Again spread work among G's
- The generated work is sent to FG's
- Then sort fragments to FM's
 - An FM is responsible for a tile of pixels
- A triangle fragment is only sent to one FG, so this avoids doing the same work twice
- (Bad: many more fragments to sort than triangles)



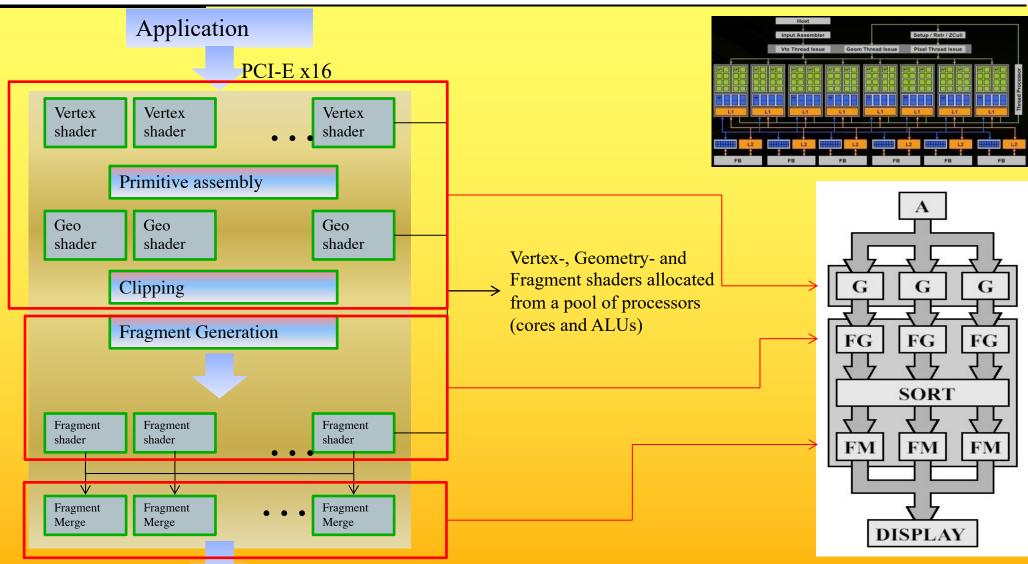
Sort-Last Image

- Sorts after entire pipeline
- So each FG & FM has a separate frame buffer for entire screen (Z and color)
 - Typically: one whole graphics card per pipeline.



- After all primitives have been sent to the pipeline, the z-buffers and color buffers are merged into one color buffer
- Can be seen as a set of independent pipelines
- Huge memory requirements!
- Used in research, but not much commerically.
- Problematic for transparency.

Functional layout of the graphics pipeline and relation to a graphics card:





The history implies the future

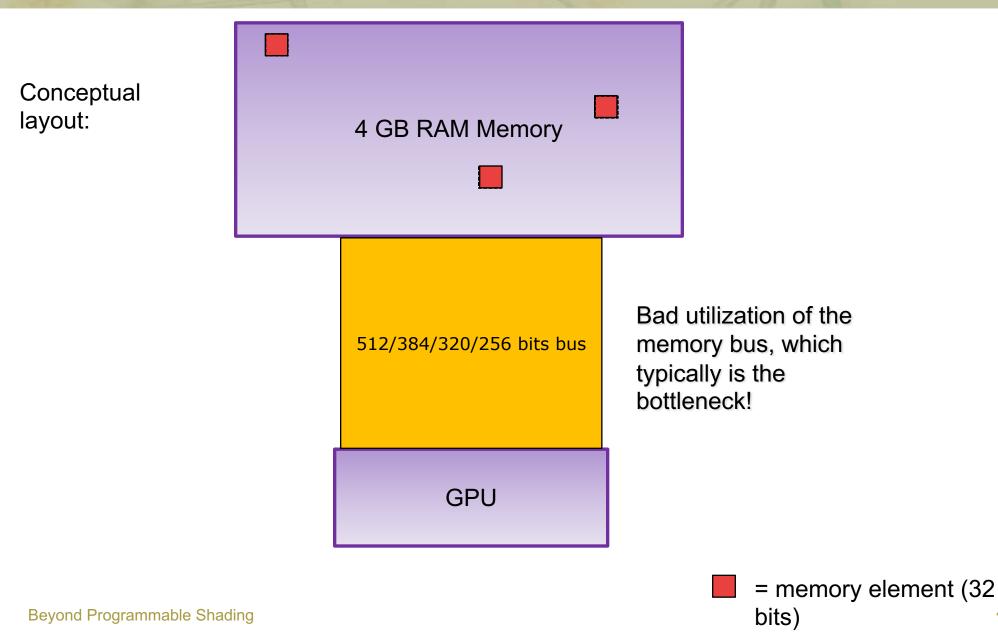
- Cell 2005, Sony Playstation 3
 - 8 cores à 4-float SIMD, 256KB L2 cache/core, 3.2 GHz
- NVIDIA 8800 GTX Nov 2006
 - 16 cores à 8-float SIMD (GTX 280 30 cores à 8, june '08)
 - 16 KB L1 cache, 64KB L2 cache
 - 1.2-1.625 GHz
- NVIDIA Fermi GF100 2010, (GF110 2011)
 - 16 cores à 2x16-float SIMD (1x16 double SIMD)
 - 16/48 KB L1 cache, 768 KB L2 cache
- NVIDIA Kepler 2012 16 cores à 2x3x16=96 float SIMD
- NVIDIA Kepler 2013 16 cores à 2x6x16=192 float SIMD
- NVIDIA Titan X 2016 60 cores à 2x4x8=64 float SIMD
- NVIDIA Volta 2018 84 cores à 64 float SIMD + tensor cores (16-bit matrix mul+add)
 NVIDIA Turing 2018 36 cores à 128 float SIMD + ~550 tensor cores (16-bit matrix mul+add) + 72 RT cores
- NVIDIA Ampere 2020 82 cores à 128 ALUs + ~328 tensor cores + 82 RT cores

If we have time...

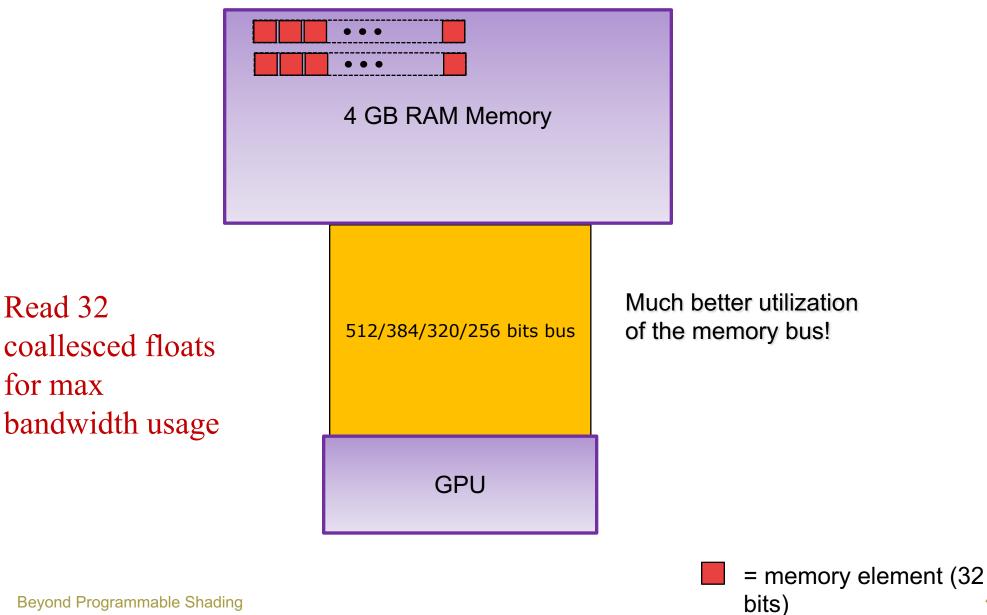
How create efficient GPU programs?

Answer: coallesced memory accesses

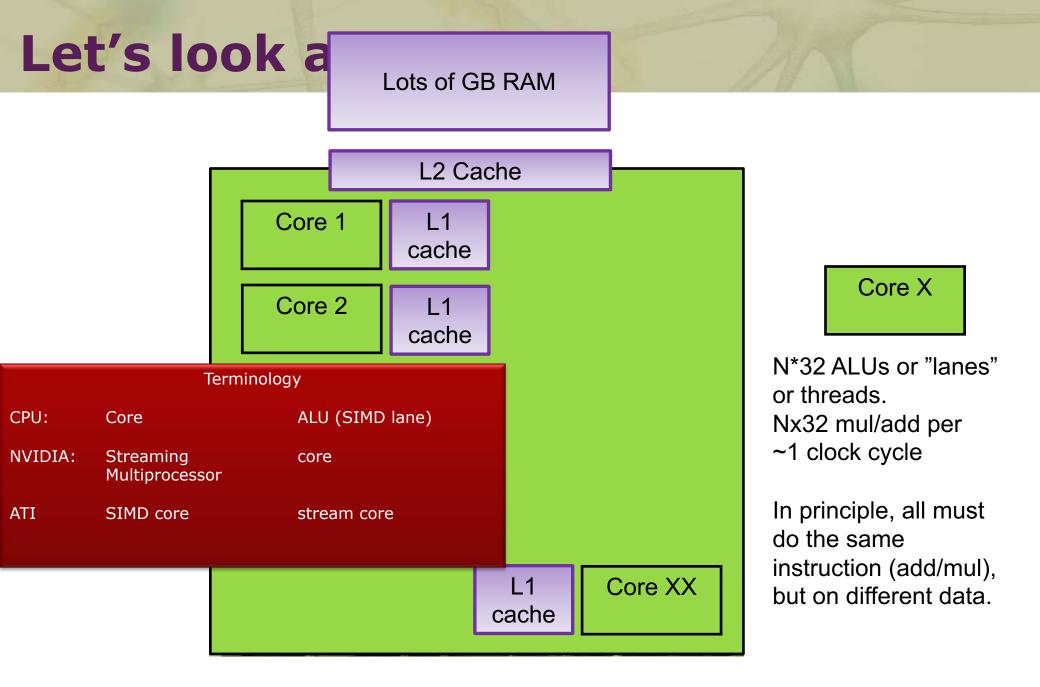
Graphics Processing Unit - GPU



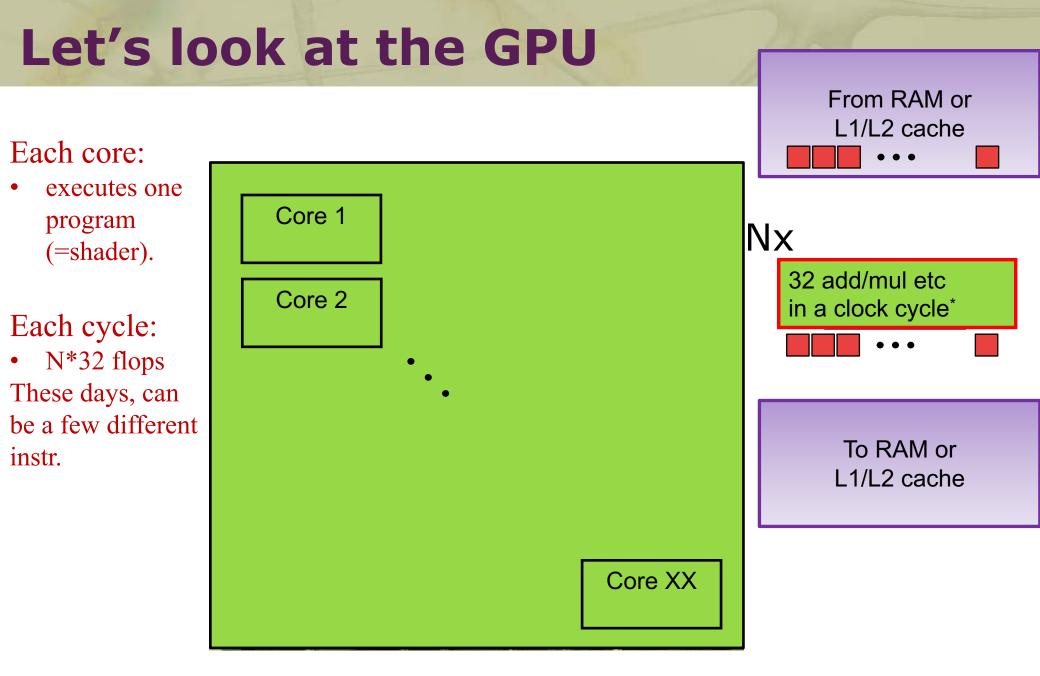
Graphics Processing Unit - GPU



Beyond Programmable Shading



Beyond Programmable Shading



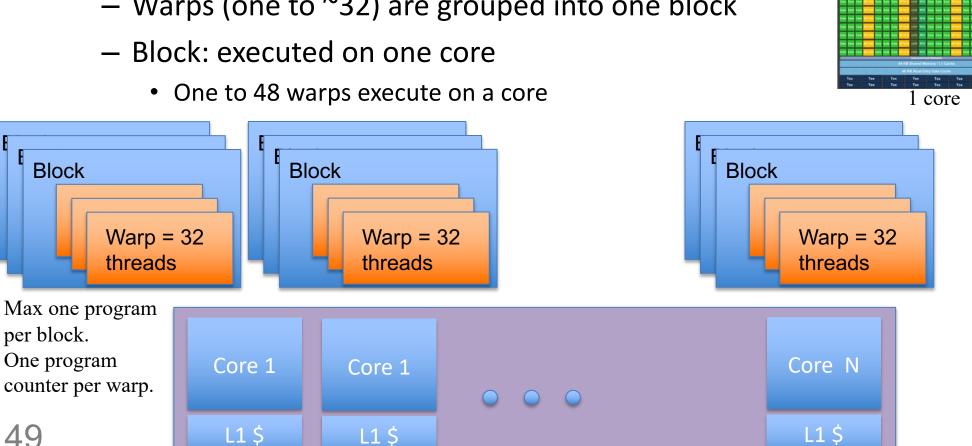
Low level APIs for GPU programming

• CUDA

- C++ compiler
- Works best for NVIDIA GPUs
- CUDA SDK
 - Numerous examples and documentation (most for single GPU)
 - Has most functionality
- OpenCL
 - C compiler
 - Platform independent
 - AMD
 - NVIDIA
 - Less control/functionality than CUDA
- Compute Shaders (DirectX, OpenGL).

CUDA

- A kernel (=CUDA program) is executed by 100:s-1M:s threads
 - A "warp" = 32 threads, one thread per ALU
 - Warps (one to ~32) are grouped into one block



Read whole cache blocks (128 bytes)

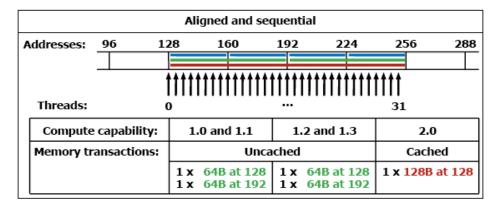
• Global mem accesses.

• One transaction:

Bandwidth to GPU RAM is the most precious resource, so two transactions is often bad.

Fermi:

• Two transactions:



Aligned and non-sequential									
Addresses:	96	128	160	192	224	256	288		
Threads:		0				31			
Compute capability:		ility: 1	.0 and 1.1 1.2 and 1.3		2.0				
Memory tr	ansacti	ions:	Uncached			Cached			
			32B at 128		4B at 128	1 x 1288	3 at 128		
		8 X	32B at 160 32B at 192		4B at 192				
			32B at 224						

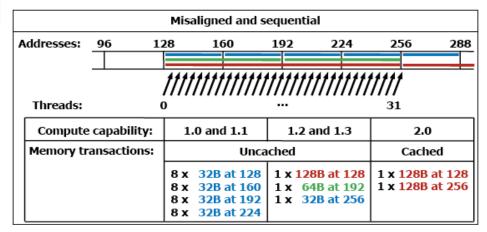


Figure G-1. Examples of Global Memory Accesses by a Warp, 4-Byte Word per Thread, and Associated Memory Transactions Based on Compute Capability

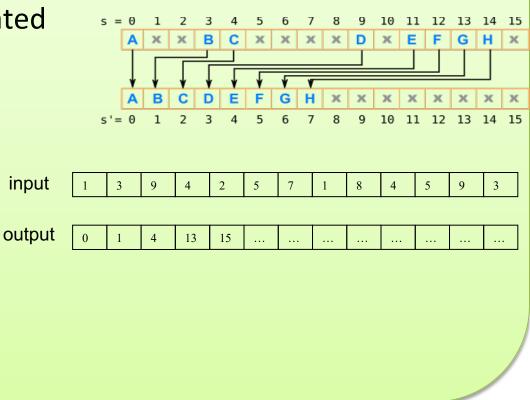
Efficient Programming

- If your program can be constructed this way, you are a winner!
- More often possible than anticipated

9

5

- Stream compaction
- Prefix sums
- Sorting



Fermi: 16 multi-processors à 2x16 SIMD width

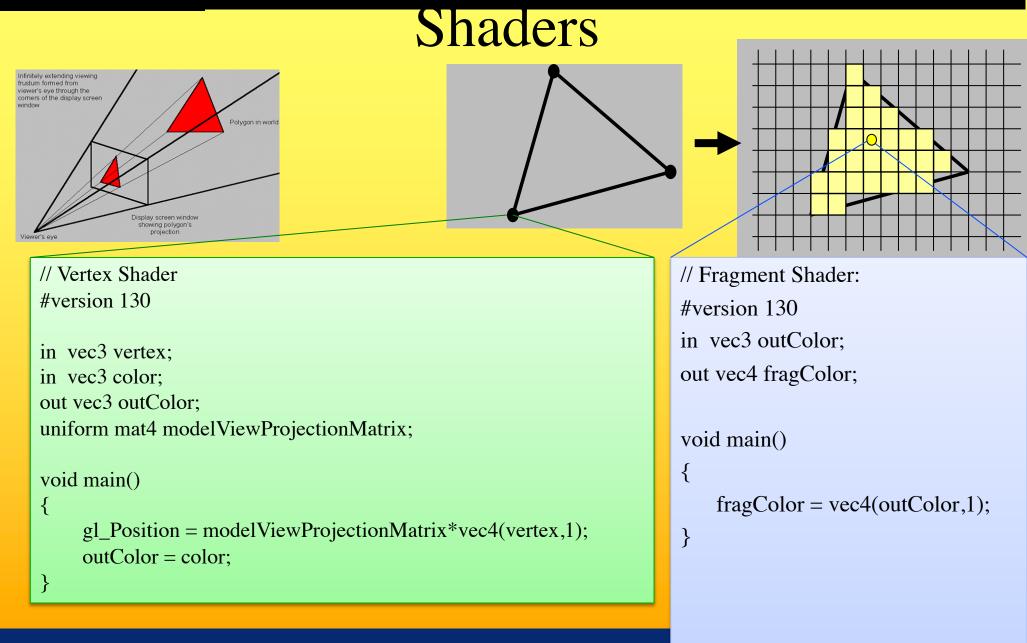
5 100 1 63 79

19 63 79 100

Beyond Programmable Shading

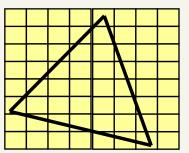
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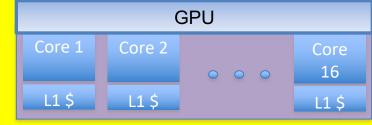
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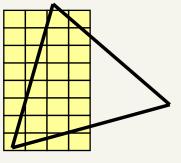


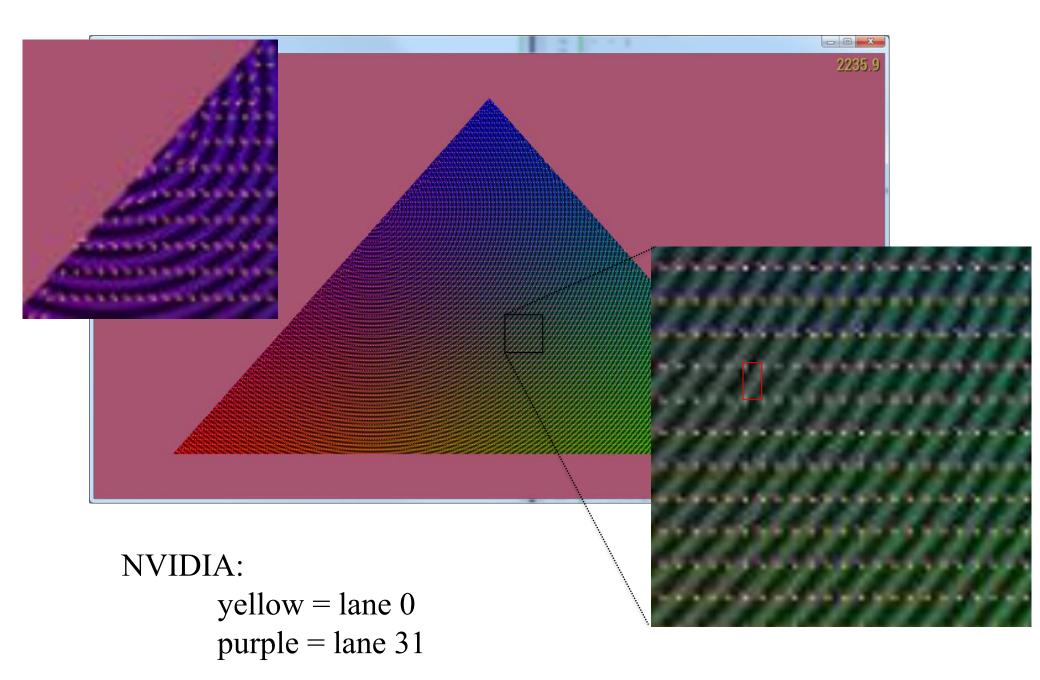
Shaders and coallesced memory accesses

- Each core (e.g. 192-SIMD) executes the same instruction per clock cycle for either a:
 - Vertex shader:
 - E.g. 192 vertices
 - Geometry shader
 - E.g. 192 triangles
 - Fragment shader:
 - E.g. 192 pixels
 in blocks of at least 2x2 pixels
 (to compute texture filter derivatives) .
 Here is an example of blocks
 4x8 = 32 pixels:
- However, many architectures can execute different instructions, of the same shader, for different warps (warp = group of 32 ALUs)







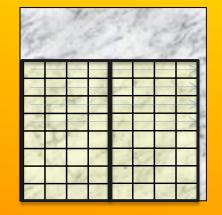


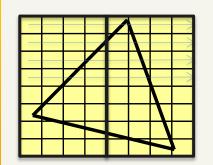
Shaders and coallesced memory accesses

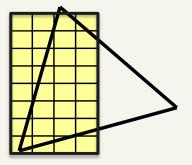
 For mipmap-filtered texture lookups in a fragment shader, this can provide coallesced memory accesses.

GPU							
Core 1	Core 2	0 0 0	Core 16				
L1 \$	L1 \$		L1 \$				









Thread utilization

- Each core executes one program (=shader)
- Each of the 192 ALUs execute one "thread" (a shader for a vertex or fragment)
- Since the core executes the same instruction for at least 32 threads (as far as the programmer is concerned)...
- If (...)

 Then, a = b + c;
 The core must execute both paths if any of the 32 threads need the if and else-path.

 Else

 a = c + d;

 But not if all need the same path.

Summary

Need to know:

 Perspective correct interpolation (e.g. for textures)

Taxonomy:

- Sort first
- sort middle
- sort last fragment
- sort last image
- Bandwidth

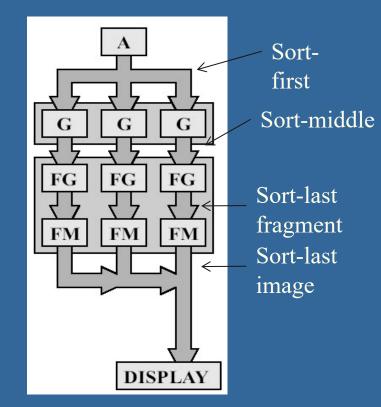
58

- Why it is a problem and how to "solve" it
 - L1 / L2 caches
 - Texture caching with prefetching, (warp switching)
 - Texture compression, Z-compression, Z-occlusion testing (HyperZ)
- Be able to sketch the functional blocks and relation to hardware for a modern graphics card (next slide→)

Linearly interpolate $(u_i/w_i, v_i/w_i, 1/w_i)$ in screenspace from each triangle vertex i. Then at each pixel:

$$u_{ip} = (u/w)_{ip} / (1/w)_{ip}$$
$$v_{ip} = (v/w)_{ip} / (1/w)_{ip}$$

where ip = screen-space interpolated value from the triangle vertices.



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