

Formal Methods for Software Development

Verification with SPIN

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SPIN: Previous Lecture vs. This Lecture

Previous lecture

SPIN appeared as a PROMELA **simulator**

This lecture

Intro to SPIN as a **model checker**

What Does A Model Checker Do?

Model Checker (MC) is designed to prove the user wrong.

MC does *not* mainly try to prove correctness properties.
It mainly tries the opposite.

MC tuned to **find counter example** to correctness property.

Why can MC **prove** correctness properties?

MC's **search** for counter examples is **exhaustive**.

⇒ **Finding no counter example proves stated correctness properties.**

What does 'exhaustive search' mean here?

exhaustive search

=

resolving non-determinism in all possible ways

For model checking PROMELA code,
two kinds of non-determinism to be resolved:

- ▶ explicit, local:

if/do statements

:: guardX -> ...

:: guardY -> ...

- ▶ implicit, global:

scheduling of concurrent processes
(see next lecture)

Model Checker for This Course: SPIN

SPIN: “Simple Promela Interpreter”

The name is a serious understatement!

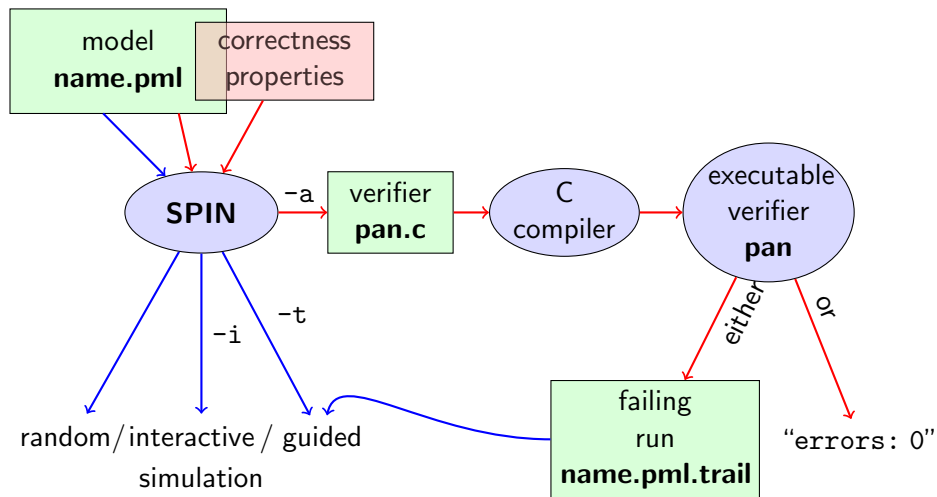
Main functionality of SPIN:

- ▶ simulating a model (randomly/interactively/**guided**)
- ▶ generating a **verifier**

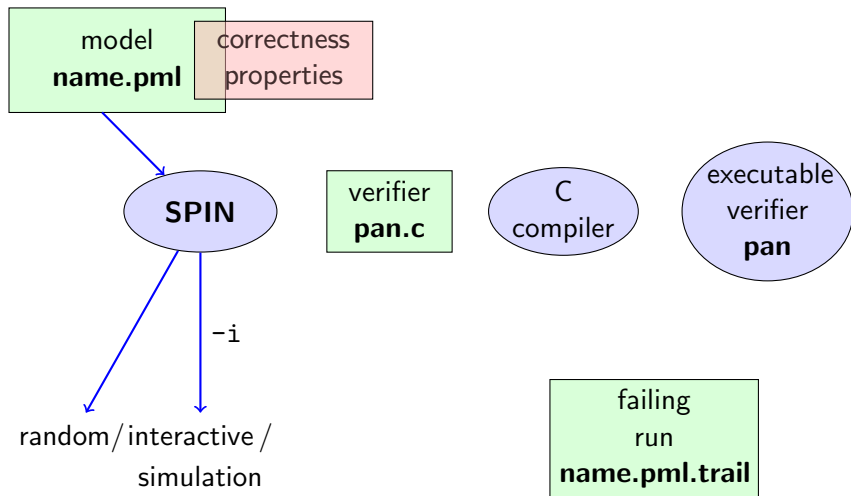
Verifier generated by SPIN is a C program performing **model checking**:

- ▶ exhaustively **checks** PROMELA **model** against correctness properties
- ▶ in case the check is negative:
generates a **failing run** of the model, **to be simulated by SPIN**

SPIN Workflow: Overview



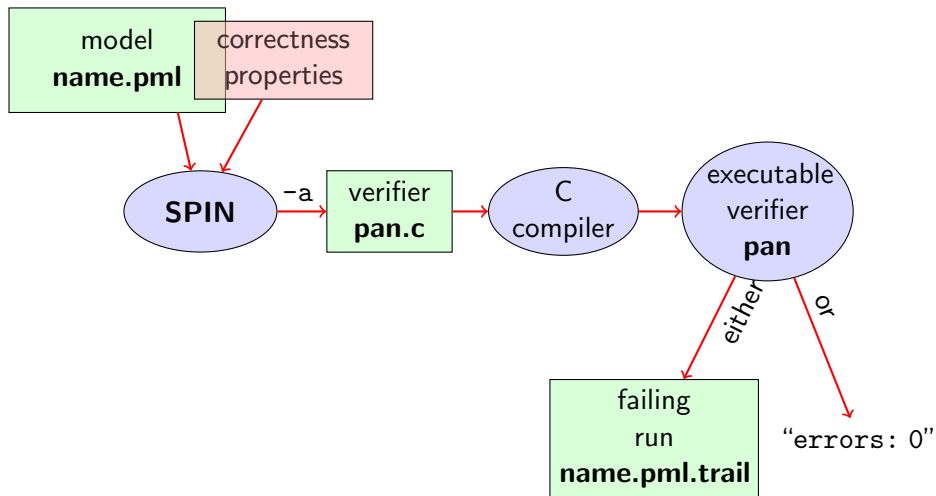
Plain Simulation with SPIN



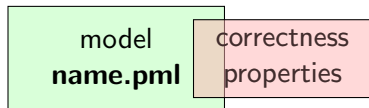
Rehearsal: Simulation Demo

- ▶ run example, random and interactive
`zero.pml`

Model Checking with SPIN



Stating Correctness Properties



Correctness properties can be stated **within**, or **outside**, the model.

stating properties within model, using

- ▶ **assertion statements** (today)
- ▶ meta labels
 - ▶ **end labels** (today)
 - ▶ accept labels
 - ▶ progress labels

stating properties outside model, using

- ▶ never claims
- ▶ temporal logic formulas

Assertion Statements

Definition (Assertion Statements)

Assertion statements in PROMELA are statements of the form

`assert(expr)`

where *expr* is any PROMELA expression.

Typically, *expr* is of type `bool`.

`assert(expr)` can appear wherever a statement is expected.

```
...                               ...
stmt1;                           if
assert(max == a);                :: b1 -> stmt3;
stmt2;                           assert(x < y)
...                               :: b2 -> stmt4
                                ...
```

Meaning of **General** Assertion Statements

`assert(expr)`

- ▶ has **no effect** if *expr* evaluates to **true** **non-zero value**
- ▶ triggers an **error message** if *expr* evaluates to **false** **0**

This holds in both, simulation and model checking mode.

Recall:

`bool true false` are syntactic sugar for

`bit 1 0`

⇒ general case covers Boolean case

Instead of using 'printf's for Debugging ...

```
/* after choosing a,b from {1,2,3} */  
if  
  :: a >= b -> max = a  
  :: a <= b -> max = b  
fi;  
printf("the maximum of %d and %d is %d\n",  
      a, b, max)
```

Command Line Execution

(simulate, inject fault, simulate again)

```
> spin [-i] max.pml
```

... we can employ **Assertions**

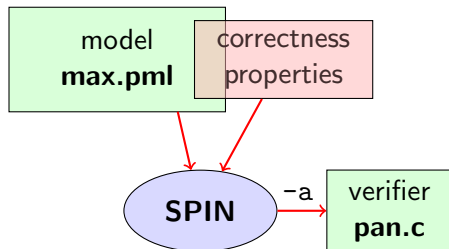
```
/* after choosing a,b from {1,2,3} */  
if  
  :: a >= b -> max = a  
  :: a <= b -> max = b  
fi;  
assert( max == (a>b -> a : b) )
```

Now, we have a first example with a formulated **correctness property**.

We can do **model checking**, for the first time!

(Historic moment in the course.)

Generate Verifier in C



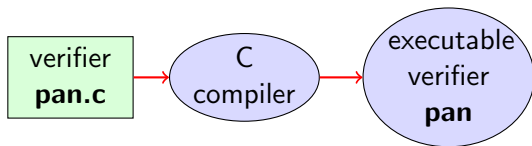
Command Line Execution

Generate Verifier in C

```
> spin -a max2.pml
```

SPIN generates **Verifier** in C, called **pan.c**
(plus helper files)

Compile To Executable Verifier



Command Line Execution

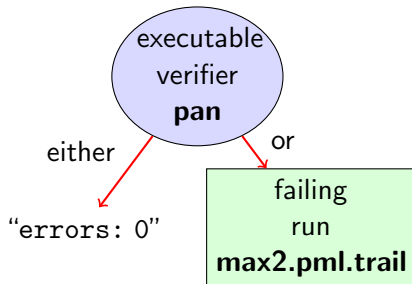
compile to executable verifier

```
> gcc -o pan pan.c
```

C compiler generates **executable verifier pan**

pan: historically “**p**rotocol **a**nalyzer”, now “**p**rocess **a**nalyzer”

Run Verifier (= Model Check)



Command Line Execution

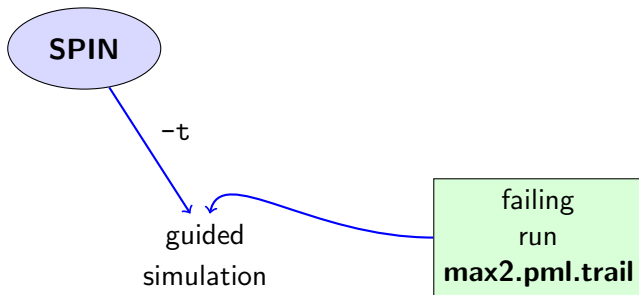
run verifier pan

> ./pan or > pan

- ▶ prints "errors: 0" \Rightarrow Correctness Property verified!
- ▶ prints "errors: n " ($n > 0$) \Rightarrow counter example found!
records failing run in **max2.pml.trail**

Guided Simulation

To **examine failing run**: employ **simulation mode**, “guided” by trail file.



Command Line Execution

inject a fault, re-run verification, and then:

```
> spin -t -p -l max2.pml
```

Output of Guided Simulation

can look like:

Starting P with pid 0

```
1: proc 0 (P) line 8 "max2.pml" (state 1) [a = 1 ]
      P(0):a = 1
2: proc 0 (P) line 14 "max2.pml" (state 7) [b = 2 ]
      P(0):b = 2
3: proc 0 (P) line 23 "max2.pml" (state 13) [((a<=b))]
3: proc 0 (P) line 23 "max2.pml" (state 14) [max = a ]
      P(0):max = 1
spin: max2.pml:22, Error: assertion violated
spin: text of failed assertion:
      assert((max==( ((a>b)) -> (a) : (b) )))
```

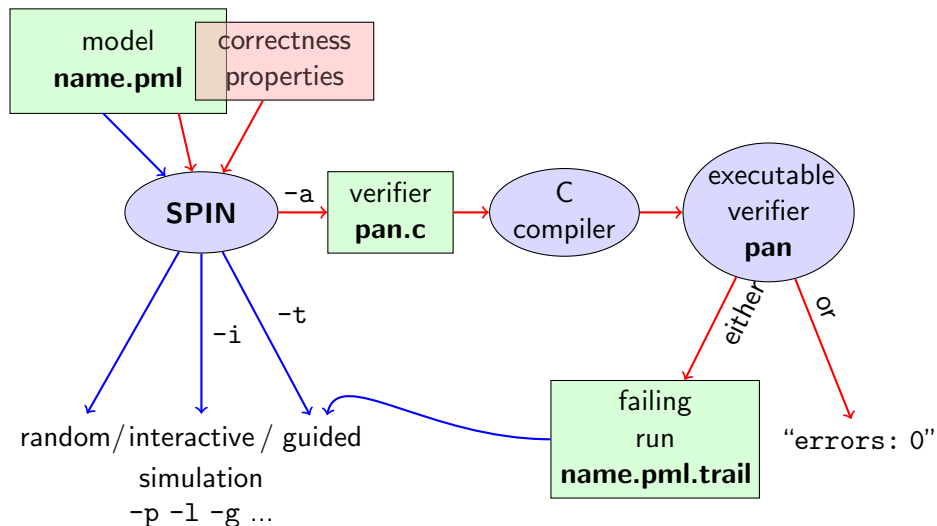
assignments in the run

values of variables whenever updated

(If output doesn't mention max variable, re-verify with ./pan -E)

What did we do so far?

following whole cycle (most primitive example, assertions only)



Further Examples: Integer Division

```
int dividend = 19;
int divisor  = 6;
int quotient, remainder;

quotient = 0;
remainder = dividend;
do
    :: remainder > divisor ->
        quotient++;
        remainder = remainder - divisor
    :: else ->
        break
od;
printf("%d divided by %d = %d, remainder = %d\n",
       dividend, divisor, quotient, remainder)
```

simulate, add assertion, ...

Further Examples: Greatest Common Divisor

greatest common divisor of x and y

```
int a, b;  
a = x; b = y;  
do  
  :: a > b -> a = a - b  
  :: b > a -> b = b - a  
  :: a == b -> break  
od;  
printf("The GCD of %d and %d = %d\n", x, y, a)
```

full functional specification w. assertion not possible (why?)

still, assertions can perform **sanity check**

⇒ **typical for model checking**

Typical Command Lines

typical command line sequences:

random simulation

```
spin name.pml
```

interactive simulation

```
spin -i name.pml
```

model checking

```
spin -a name.pml
```

```
gcc -o pan pan.c
```

```
./pan
```

and in case of error

```
spin -t -p -l -g name.pml
```

SPIN Reference Card

Ben-Ari produced **Spin Reference Card**, summarizing

- ▶ **typical command line sequences**
- ▶ **options for**
 - ▶ **SPIN**
 - ▶ gcc
 - ▶ pan
- ▶ PROMELA
 - ▶ datatypes
 - ▶ operators
 - ▶ statements
 - ▶ guarded commands
 - ▶ processes
 - ▶ channels
- ▶ temporal logic syntax

⇒ **available from course page** (see 'Links, Papers, and Software')

Why SPIN?

- ▶ SPIN targets software, instead of hardware verification (“Formal Methods for *Software* Development”)
- ▶ 2001 ACM Software Systems Award (other winning systems include: Unix, TCP/IP, WWW, Tcl/Tk, Java, GCC, T_EX, Coq)
- ▶ used for safety critical applications
- ▶ distributed freely as research tool, well-documented, actively maintained, large user-base in academia and in industry
- ▶ annual SPIN user workshops series held since 1995
- ▶ based on standard theory of (ω -)automata and linear temporal logic

Why SPIN? (Cont'd)

- ▶ PROMELA and SPIN are rather simple to use
- ▶ availability of good course book (Ben-Ari)
- ▶ availability of front end JSPIN (also Ben-Ari)
- ▶ and: availability of our own web interface

What is JSPIN?

- ▶ graphical user interface for SPIN
- ▶ developed for pedagogical purposes
- ▶ written in JAVA
- ▶ simple user interface
- ▶ SPIN options automatically supplied
- ▶ fully configurable
- ▶ supports graphics output of transition system
- ▶ makes back-end calls transparent

Command Line Execution

calling JSPIN

```
> java -jar /usr/local/jSpin/jSpin.jar
```

(with path adjusted to your setting)

or use shell script:

```
> jspin
```

play around with similar examples ...

Meaning of Correctness w.r.t. Properties

Given PROMELA model M , and correctness properties C_1, \dots, C_n .

- ▶ Be R_M the set of **all possible runs** of M .
- ▶ For each correctness property C_i ,
 R_{M,C_i} is the set of all **runs** of M **satisfying** C_i .
($R_{M,C_i} \subseteq R_M$)
- ▶ M is **correct wrt.** C_1, \dots, C_n iff $R_M \subseteq (R_{M,C_1} \cap \dots \cap R_{M,C_n})$.
- ▶ If M is not correct wrt. C_1, \dots, C_n , then
each $r \in (R_M \setminus (R_{M,C_1} \cap \dots \cap R_{M,C_n}))$ is a **counter example**.

Catching A Different Type of Error

quoting from file **max3.pml**:

```
/* after choosing a,b from {1,2,3} */  
if  
  :: a >= b -> max = a  
  :: b <= a -> max = b  
fi;  
printf("the maximum of %d and %d is %d\n",  
       a, b, max)
```

simulate a few times

⇒ crazy “timeout” message sometimes

generate and execute **pan**

⇒ reports “errors: 1”

????

Catching A Different Type of Error

Further inspection of **pan** output:

```
...  
pan: invalid end state (at depth 1)  
pan: wrote max3.pml.trail  
...
```

Legal and Illegal Blocking

A process may *legally* block, **as long as some other process can proceed.**

Blocking for letting others proceed is useful for concurrent and distributed models (i.p. protocols).

But

It is *illegal* if a process blocks while no other process can proceed.

⇒ “Deadlock”

In **max3.pml**, there exists a blocking run where no process can take over.

(Fix error)

Valid End States

Definition (Valid End State)

An **end state** of a **run** is valid iff the location counter of **each processes** is at an **end location**.

Definition (End Location)

End locations of a process P are:

- ▶ P's textual end
- ▶ each location marked with an **end label**: "endxxx:"

End labels not useful in **max3.pml**, but elsewhere, they are.

Example: `end.pml`

Can get SPIN to ignore 'invalid end state' error: `./pan -E`

Literature for this Lecture

Ben-Ari Chapter 2, Sections 4.7.1, 4.7.2