Lecture
Computability
(DIT312, DAT415)

Nils Anders Danielsson

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Today

X, a small functional language:

- Concrete and abstract syntax.
- Operational semantics.
- Several variants of the halting problem.
- Representing inductively defined sets.

Concrete and abstract syntax

Rough definitions (in the context of programming languages):

- Concrete syntax: The well-formed strings of a programming language.
- Abstract syntax: The essential information of a program text, ignoring details of notation.

Concrete syntax

Concrete syntax

$$
\begin{aligned}
e ::= {} & x \\
\mid {} & (e_1\ e_2) \\
\mid {} & \lambda x.\, e \\
\mid {} & \mathsf{C}(e_1, ..., e_n) \\
\mid {} & \mathbf{case}\ e\ \mathbf{of}\ \{\mathsf{C}_1(x_1, ..., x_n) \to e_1; ...\} \\
\mid {} & \mathbf{rec}\ x = e
\end{aligned}
$$

Variables (x) and constructors (C) are assumed to come from two disjoint, countably infinite sets.

Sometimes extra parentheses are used, and sometimes parentheses are omitted around applications: $e_1\ e_2\ e_3$ means $((e_1\ e_2)\ e_3)$.

Examples

X	Haskell
$\lambda x.\, e$	`\x -> e`
$\mathsf{True}()$	`True`
$\mathsf{Suc}(n)$	`Suc n`
$\mathsf{Cons}(x, xs)$	`x : xs`
$\mathbf{rec}\ x = e$	`let x = e in x`

Note: Haskell is typed and non-strict, χ is untyped and strict.

Another example

X:

$$\mathbf{case}\ e\ \mathbf{of}\ \{\mathsf{Zero}() \to x; \mathsf{Suc}(n) \to y\}$$

Haskell:

```
case e of
  Zero  -> x
  Suc n -> y
```

And two more

$$\begin{array}{l}\mathbf{rec}\ add = \lambda\, m.\, \lambda\, n.\, \mathbf{case}\ n\ \mathbf{of} \\ \quad \{\, \mathsf{Zero}() \ \to m \\ \quad ;\, \mathsf{Suc}(n) \to \mathsf{Suc}(add\ m\ n) \\ \quad \}\end{array}$$

$$\begin{array}{l}\lambda\, m.\, \mathbf{rec}\ add = \lambda\, n.\, \mathbf{case}\ n\ \mathbf{of} \\ \quad \{\, \mathsf{Zero}() \ \to m \\ \quad ;\, \mathsf{Suc}(n) \to \mathsf{Suc}(add\ n) \\ \quad \}\end{array}$$

What is the value of the following expression?

$$
\begin{array}{l}
(\mathbf{rec}\ foo = \lambda\, m.\, \lambda\, n.\, \mathbf{case}\ n\ \mathbf{of}\ \{ \\
\quad \mathsf{Zero}() \ \rightarrow m; \\
\quad \mathsf{Suc}(n) \rightarrow \mathbf{case}\ m\ \mathbf{of}\ \{ \\
\qquad \mathsf{Zero}() \ \rightarrow \mathsf{Zero}(); \\
\qquad \mathsf{Suc}(m) \rightarrow foo\ m\ n\}\}) \\
\mathsf{Suc}(\mathsf{Suc}(\mathsf{Zero}()))\ \mathsf{Suc}(\mathsf{Zero}())
\end{array}
$$

1. $\mathsf{Zero}()$
2. $\mathsf{Suc}(\mathsf{Zero}())$
3. $\mathsf{Suc}(\mathsf{Suc}(\mathsf{Zero}()))$
4. $\mathsf{Suc}(\mathsf{Suc}(\mathsf{Suc}(\mathsf{Zero}())))$

Abstract syntax

Abstract syntax

$$\frac{x \in Var}{\mathsf{var}\ x \in Exp} \qquad \frac{e_1 \in Exp \qquad e_2 \in Exp}{\mathsf{apply}\ e_1\ e_2 \in Exp}$$

$$\frac{x \in Var \qquad e \in Exp}{\mathsf{lambda}\ x\ e \in Exp} \qquad \frac{x \in Var \qquad e \in Exp}{\mathsf{rec}\ x\ e \in Exp}$$

Var: Assumed to be countably infinite.

Abstract syntax

$$\frac{c \in \mathit{Const} \qquad es \in \mathit{List}\ \mathit{Exp}}{\mathsf{const}\ c\ es \in \mathit{Exp}}$$

$$\frac{e \in \mathit{Exp} \qquad bs \in \mathit{List}\ \mathit{Br}}{\mathsf{case}\ e\ bs \in \mathit{Exp}}$$

$$\frac{c \in \mathit{Const} \qquad xs \in \mathit{List}\ \mathit{Var} \qquad e \in \mathit{Exp}}{\mathsf{branch}\ c\ xs\ e \in \mathit{Br}}$$

Const: Assumed to be countably infinite.

Examples

Concrete syntax	Abstract syntax
$\lambda\, x.\ e$	lambda $\underline{x}$ $\underline{e}$
True()	const $\underline{True}$ nil
Suc(n)	const $\underline{Suc}$ (cons $\underline{n}$ nil)
Cons(x, xs)	const $\underline{Cons}$ (cons $\underline{x}$ (cons $\underline{xs}$ nil))
rec $x = e$	rec $\underline{x}$ $\underline{e}$

Here $\underline{x}$ is the element of Var standing for x,
$\underline{e}$ the element of Exp standing for e, and
$\underline{True}$ the element of $Const$ standing for True.

Another example

Concrete:

$$\mathbf{case}\ e\ \mathbf{of}\ \{\mathsf{Zero}() \to x; \mathsf{Suc}(n) \to y\}$$

Abstract:

case $\underline{e}$
(cons (branch $\underline{Zero}$ nil $\underline{x}$)
(cons (branch $\underline{Suc}$ (cons $\underline{n}$ nil) $\underline{y}$)
nil))

Operational semantics

Operational semantics

- $e \Downarrow v$: e terminates with the value v.
- The expression e terminates (with a value) if $\exists v.\, e \Downarrow v$.
- Note that a "crash" does not count as termination (with a value).

Operational semantics

- The binary relation $\Downarrow$ relates *closed* expressions.
- An expression is closed if it has no free variables.

Free variables

$$
\begin{aligned}
&FV \in Exp \to \wp\, Var \\
&FV\,(\mathsf{var}\ x) &&= \{\, x \,\} \\
&FV\,(\mathsf{apply}\ e_1\ e_2) &&= FV\ e_1 \cup FV\ e_2 \\
&FV\,(\mathsf{lambda}\ x\ e) &&= FV\ e \setminus \{\, x \,\} \\
&FV\,(\mathsf{rec}\ x\ e) &&= FV\ e \setminus \{\, x \,\} \\
&FV\,(\mathsf{const}\ c\ es) &&= \bigcup_{e\ \in\ \underline{es}} FV\ e \\
&FV\,(\mathsf{case}\ e\ bs) &&= FV\ e \cup \bigcup_{b\ \in\ \underline{bs}} FV_{Br}\ b
\end{aligned}
$$

$$
\begin{aligned}
&FV_{Br} \in Br \to \wp\, Var \\
&FV_{Br}\,(\mathsf{branch}\ c\ xs\ e) = FV\ e \setminus \underline{xs}
\end{aligned}
$$

(Here $\underline{xs}$ is a set containing the elements in xs.)

Quiz

Which of the following expressions are closed?

1. y
2. $\lambda\, x.\, \lambda\, y.\, x$
3. $\mathbf{case}\ x\ \mathbf{of}\ \{\mathsf{Cons}(x, xs) \to x\}$
4. $\mathbf{case}\ \mathsf{Suc}(\mathsf{Zero}())\ \mathbf{of}\ \{\mathsf{Suc}(x) \to x\}$
5. $\mathbf{rec}\ f = \lambda\, x.\, f$

Operational semantics (1/3)

$$\frac{}{\mathsf{lambda}\ x\ e \Downarrow \mathsf{lambda}\ x\ e}$$

$$\frac{e_1 \Downarrow \mathsf{lambda}\ x\ e \qquad e_2 \Downarrow v_2 \qquad e\,[x \leftarrow v_2] \Downarrow v}{\mathsf{apply}\ e_1\ e_2 \Downarrow v}$$

$$\frac{e\,[x \leftarrow \mathsf{rec}\ x\ e] \Downarrow v}{\mathsf{rec}\ x\ e \Downarrow v}$$

Substitution

- $e\,[x \leftarrow e']$: Substitute e' for every *free* occurrence of x in e.
- To keep things simple: e' must be closed.
- If e' is not closed, then this definition is prone to *variable capture*.

Substitution

$$\text{var } x\,[x \leftarrow e'] = e'$$
$$\text{var } y\,[x \leftarrow e'] = \text{var } y \quad \text{if } x \neq y$$

$$\text{apply } e_1\ e_2\,[x \leftarrow e'] =$$
$$\quad \text{apply } (e_1\,[x \leftarrow e'])\ (e_2\,[x \leftarrow e'])$$

$$\text{lambda } x\ e\,[x \leftarrow e'] = \text{lambda } x\ e$$
$$\text{lambda } y\ e\,[x \leftarrow e'] =$$
$$\quad \text{lambda } y\ (e\,[x \leftarrow e']) \quad \text{if } x \neq y$$

And so on...

Quiz

What is the result of

$$(\mathbf{rec}\ y = \mathbf{case}\ x\ \mathbf{of}\ \{\mathsf{C}() \to x; \mathsf{D}(x) \to x\})\ [x \leftarrow \lambda\, z.\ z]?$$

$$\begin{array}{lll}
\mathbf{rec}\ y = \mathbf{case}\ x & \mathbf{of}\ \{\mathsf{C}() \to x; & \mathsf{D}(x) \to x\} \\
\mathbf{rec}\ y = \mathbf{case}\ x & \mathbf{of}\ \{\mathsf{C}() \to \lambda\, z.\ z; & \mathsf{D}(x) \to x\} \\
\mathbf{rec}\ y = \mathbf{case}\ \lambda\, z.\ z & \mathbf{of}\ \{\mathsf{C}() \to x; & \mathsf{D}(x) \to x\} \\
\mathbf{rec}\ y = \mathbf{case}\ \lambda\, z.\ z & \mathbf{of}\ \{\mathsf{C}() \to \lambda\, z.\ z; & \mathsf{D}(x) \to x\} \\
\mathbf{rec}\ y = \mathbf{case}\ \lambda\, z.\ z & \mathbf{of}\ \{\mathsf{C}() \to \lambda\, z.\ z; & \mathsf{D}(x) \to \lambda\, z.\ z\} \\
\mathbf{rec}\ y = \mathbf{case}\ \lambda\, z.\ z & \mathbf{of}\ \{\mathsf{C}() \to \lambda\, z.\ z; & \mathsf{D}(\lambda\, z.\ z) \to \lambda\, z.\ z\}
\end{array}$$

Operational semantics (2/3)

$$\frac{es \Downarrow^\star vs}{\mathsf{const}\ c\ es \Downarrow \mathsf{const}\ c\ vs}$$

$$\frac{}{\mathsf{nil} \Downarrow^\star \mathsf{nil}} \qquad \frac{e \Downarrow v \qquad es \Downarrow^\star vs}{\mathsf{cons}\ e\ es \Downarrow^\star \mathsf{cons}\ v\ vs}$$

An example

$$\dfrac{\dfrac{}{\mathsf{lambda}\ x\ (\mathsf{var}\ x) \Downarrow \mathsf{lambda}\ x\ (\mathsf{var}\ x)} \quad \dfrac{\dfrac{}{\mathsf{nil} \Downarrow^{\star} \mathsf{nil}}}{\mathsf{const}\ c\ \mathsf{nil} \Downarrow \mathsf{const}\ c\ \mathsf{nil}} \quad \dfrac{\dfrac{}{\mathsf{nil} \Downarrow^{\star} \mathsf{nil}}}{\mathsf{var}\ x\ [x \leftarrow \mathsf{const}\ c\ \mathsf{nil}] \Downarrow \mathsf{const}\ c\ \mathsf{nil}}}{\mathsf{apply}\ (\mathsf{lambda}\ x\ (\mathsf{var}\ x))\ (\mathsf{const}\ c\ \mathsf{nil}) \Downarrow \mathsf{const}\ c\ \mathsf{nil}}$$

Operational semantics (3/3)

$$\frac{e \Downarrow \mathsf{const}\ c\ vs \qquad Lookup\ c\ bs\ xs\ e' \qquad e'\ [xs \leftarrow vs] \mapsto e'' \qquad e'' \Downarrow v}{\mathsf{case}\ e\ bs \Downarrow v}$$

Operational semantics (3/3)

$$\frac{e \Downarrow \mathsf{const}\ c\ vs \qquad Lookup\ c\ bs\ xs\ e' \qquad e'\ [xs \leftarrow vs] \mapsto e'' \qquad e'' \Downarrow v}{\mathsf{case}\ e\ bs \Downarrow v}$$

The first matching branch, if any:

$$\frac{}{Lookup\ c\ (\mathsf{cons}\ (\mathsf{branch}\ c\ xs\ e)\ bs)\ xs\ e}$$

$$\frac{c \neq c' \qquad Lookup\ c\ bs\ xs\ e}{Lookup\ c\ (\mathsf{cons}\ (\mathsf{branch}\ c'\ xs'\ e')\ bs)\ xs\ e}$$

Operational semantics (3/3)

$$\frac{e \Downarrow \text{const } c\ vs \qquad Lookup\ c\ bs\ xs\ e' \qquad e'\ [xs \leftarrow vs] \mapsto e'' \qquad e'' \Downarrow v}{\text{case } e\ bs \Downarrow v}$$

$e\ [xs \leftarrow vs] \mapsto e'$ holds iff

- there is some n such that
  $xs = \text{cons } x_1\ (...(\text{cons } x_n\ \text{nil}))$ and
  $vs = \text{cons } v_1\ (...(\text{cons } v_n\ \text{nil}))$, and
- $e' = ((e\ [x_n \leftarrow v_n])...)\ [x_1 \leftarrow v_1]$.

Operational semantics (3/3)

$$\frac{e \Downarrow \mathsf{const}\ c\ vs \qquad Lookup\ c\ bs\ xs\ e' \qquad e'\ [xs \leftarrow vs] \mapsto e'' \qquad e'' \Downarrow v}{\mathsf{case}\ e\ bs \Downarrow v}$$

$$\frac{}{e\ [\mathsf{nil} \leftarrow \mathsf{nil}] \mapsto e}$$

$$\frac{e\ [xs \leftarrow vs] \mapsto e'}{e\ [\mathsf{cons}\ x\ xs \leftarrow \mathsf{cons}\ v\ vs] \mapsto e'\ [x \leftarrow v]}$$

Quiz

Which of the following sets are inhabited?

$$\mathbf{case}\ \mathsf{C}()\ \mathbf{of}\ \{\mathsf{C}()\ \to \mathsf{D}();\mathsf{C}() \to \mathsf{C}()\} \Downarrow \mathsf{C}()$$
$$\mathbf{case}\ \mathsf{C}()\ \mathbf{of}\ \{\mathsf{C}()\ \to \mathsf{D}();\mathsf{C}() \to \mathsf{C}()\} \Downarrow \mathsf{D}()$$
$$\mathbf{case}\ \mathsf{C}()\ \mathbf{of}\ \{\mathsf{C}(x) \to \mathsf{D}();\mathsf{C}() \to \mathsf{D}()\} \Downarrow \mathsf{D}()$$
$$\mathbf{case}\ \mathsf{C}(\mathsf{C}(),\mathsf{D}())\ \mathbf{of}\ \{\mathsf{C}(x,x) \to x\} \Downarrow \mathsf{C}()$$
$$\mathbf{case}\ \mathsf{Suc}(\mathsf{False}())\ \mathbf{of}$$
$$\quad \{\mathsf{Zero}() \to \mathsf{True}();\mathsf{Suc}(n) \to n\} \Downarrow \mathsf{False}()$$
$$\mathbf{case}\ \mathsf{Suc}(\mathsf{False}())\ \mathbf{of}$$
$$\quad \{\mathsf{Zero}() \to \mathsf{True}();\mathsf{Suc}() \to \mathsf{False}()\} \Downarrow \mathsf{False}()$$

Some properties

Deterministic

The semantics is deterministic:
$e \Downarrow v_1$ and $e \Downarrow v_2$ imply $v_1 = v_2$.

Values

- An expression e is called a value if $e \Downarrow e$.
- Values can be characterised inductively:

$$\frac{}{Value\ (\mathsf{lambda}\ x\ e)} \qquad \frac{Values\ es}{Value\ (\mathsf{const}\ c\ es)}$$

$$\frac{}{Values\ \mathsf{nil}} \qquad \frac{Value\ e \qquad Values\ es}{Values\ (\mathsf{cons}\ e\ es)}$$

- $Value\ e$ holds iff $e \Downarrow e$.
- If $e \Downarrow v$, then $Value\ v$.

There is a non-terminating expression

- The program $\mathsf{rec}\ x\ (\mathsf{var}\ x)$ does not terminate with a value.
- Recall the rule for rec: $\dfrac{e\,[x \leftarrow \mathsf{rec}\ x\ e] \Downarrow v}{\mathsf{rec}\ x\ e \Downarrow v}$.
- Note that
  $\mathsf{var}\ x\,[x \leftarrow \mathsf{rec}\ x\ (\mathsf{var}\ x)] = \mathsf{rec}\ x\ (\mathsf{var}\ x)$.
- Idea:

$$
\begin{aligned}
&\mathsf{rec}\ x\ (\mathsf{var}\ x) && \rightarrow \\
&\mathsf{var}\ x\,[x \leftarrow \mathsf{rec}\ x\ (\mathsf{var}\ x)] && = \\
&\mathsf{rec}\ x\ (\mathsf{var}\ x) && \rightarrow \\
&\vdots
\end{aligned}
$$

There is a non-terminating expression

- If the program did terminate, then there would be a *finite* derivation of the following form:

$$\dfrac{\dfrac{\dfrac{\vdots}{\mathsf{rec}\ x\ (\mathsf{var}\ x) \Downarrow v}}{\mathsf{rec}\ x\ (\mathsf{var}\ x) \Downarrow v}}{\mathsf{rec}\ x\ (\mathsf{var}\ x) \Downarrow v}$$

- Exercise: Prove more formally that this is impossible, using induction on the structure of the semantics.

The halting problem

The extensional halting problem

There is no closed expression $halts$ such that, for every closed expression p,

- $halts\ (\lambda\, x.\, p) \Downarrow \mathsf{True}()$, if p terminates, and
- $halts\ (\lambda\, x.\, p) \Downarrow \mathsf{False}()$, otherwise.

The extensional halting problem

Note the abuse of notation:

- The variables $halts$ and p are not χ variables.
- *Meta-variables* standing for χ expressions.
- An alternative is to use abstract syntax:

  apply $halts$ (lambda $\underline{x}\ p$) $\Downarrow$ const $\underline{True}$ nil
  apply $halts$ (lambda $\underline{x}\ p$) $\Downarrow$ const $\underline{False}$ nil

  (For *distinct* $\underline{True}$, $\underline{False} \in Const$.)
- More verbose.

The extensional halting problem

- Assume that $halts$ can be defined.
- Define $terminv \in Exp \to Exp$:

$$\begin{aligned} terminv\ p = \mathbf{case}\ halts\ (\lambda x.\ p)\ \mathbf{of} \\ \{\mathsf{True}() \to \mathbf{rec}\ x = x \\ ; \mathsf{False}() \to \mathsf{Zero}() \\ \} \end{aligned}$$

- For any closed expression p:
  $terminv\ p$ terminates iff p does not terminate.

The extensional halting problem

- Now consider the closed expression $strange$ defined by $\mathbf{rec}\ p = terminv\ p$ (where $p \neq x$).
- We get a contradiction:

$$
\begin{array}{lll}
(\exists v.\ strange & \Downarrow v) & \Leftrightarrow \\
(\exists v.\ \mathbf{rec}\ p = terminv\ p & \Downarrow v) & \Leftrightarrow \\
(\exists v.\ terminv\ p\ [p \leftarrow strange] & \Downarrow v) & \Leftrightarrow \\
(\exists v.\ terminv\ strange & \Downarrow v) & \Leftrightarrow \\
\neg\ (\exists v.\ strange & \Downarrow v) &
\end{array}
$$

The extensional halting problem

- Note that we apply $halts$ to a program, not to the source code of a program.
- How can source code be represented?

Representing inductively defined sets

Natural numbers

One method:

- Notation: $\ulcorner n \urcorner \in Exp$ represents $n \in \mathbb{N}$.
- Representation:

$$\begin{aligned} \ulcorner \mathsf{zero} \urcorner &= \mathsf{Zero}() \\ \ulcorner \mathsf{suc}\ n \urcorner &= \mathsf{Suc}(\ulcorner n \urcorner) \end{aligned}$$

Natural numbers

One method:

- Notation: $\ulcorner n \urcorner \in \mathit{Exp}$ represents $n \in \mathbb{N}$.
- Representation:

$$\begin{aligned} \ulcorner \mathsf{zero} \urcorner &= \mathsf{Zero}() \\ \ulcorner \mathsf{suc}\ n \urcorner &= \mathsf{Suc}(\ulcorner n \urcorner) \end{aligned}$$

- Note that the concrete syntax should be interpreted as abstract syntax:

$$\begin{aligned} \ulcorner \mathsf{zero} \urcorner &= \mathsf{const}\ \underline{\mathit{Zero}}\ \mathsf{nil} \\ \ulcorner \mathsf{suc}\ n \urcorner &= \mathsf{const}\ \underline{\mathit{Suc}}\ (\mathsf{cons}\ \ulcorner n \urcorner\ \mathsf{nil}) \end{aligned}$$

(For some distinct $\underline{\mathit{Zero}}, \underline{\mathit{Suc}} \in \mathit{Const}$.)

Lists

If elements in A can be represented, then elements in $List\ A$ can also be represented:

$$\begin{aligned} \ulcorner \mathsf{nil} \urcorner &= \mathsf{Nil}() \\ \ulcorner \mathsf{cons}\ x\ xs \urcorner &= \mathsf{Cons}(\ulcorner x \urcorner, \ulcorner xs \urcorner) \end{aligned}$$

Many inductively defined sets can be treated in the same way.

Variables, constants

- Var: Countably infinite.
- Thus each variable $x \in Var$ can be assigned a unique natural number $number\, x \in \mathbb{N}$.
- Define $\ulcorner x \urcorner = \ulcorner number\, x \urcorner$.
- Similarly for constants.

Variables, constants

- Var: Countably infinite.
- Thus each variable $x \in Var$ can be assigned a unique natural number $number\, x \in \mathbb{N}$.
- Define $\ulcorner x \urcorner^{Var} = \ulcorner number\, x \urcorner^{\mathbb{N}}$.
- Similarly for constants.

Source code

$$
\begin{aligned}
\ulcorner \text{var } x \urcorner &= \text{Var}(\ulcorner x \urcorner) \\
\ulcorner \text{apply } e_1\ e_2 \urcorner &= \text{Apply}(\ulcorner e_1 \urcorner, \ulcorner e_2 \urcorner) \\
\ulcorner \text{lambda } x\ e \urcorner &= \text{Lambda}(\ulcorner x \urcorner, \ulcorner e \urcorner) \\
\ulcorner \text{rec } x\ e \urcorner &= \text{Rec}(\ulcorner x \urcorner, \ulcorner e \urcorner) \\
\ulcorner \text{const } c\ es \urcorner &= \text{Const}(\ulcorner c \urcorner, \ulcorner es \urcorner) \\
\ulcorner \text{case } e\ bs \urcorner &= \text{Case}(\ulcorner e \urcorner, \ulcorner bs \urcorner) \\
\ulcorner \text{branch } c\ xs\ e \urcorner &= \text{Branch}(\ulcorner c \urcorner, \ulcorner xs \urcorner, \ulcorner e \urcorner)
\end{aligned}
$$

Example

- Concrete syntax: $\lambda\, x.\, \mathsf{Suc}(x)$.
- Abstract syntax:

  $$\mathsf{lambda}\ \underline{x}\ (\mathsf{const}\ \underline{Suc}\ (\mathsf{cons}\ (\mathsf{var}\ \underline{x})\ \mathsf{nil}))$$

  (for some $\underline{x} \in Var$ and $\underline{Suc} \in Const$).
- Representation (concrete syntax):

  $$\mathsf{Lambda}(\ulcorner \underline{x} \urcorner, \mathsf{Const}(\ulcorner \underline{Suc} \urcorner, \mathsf{Cons}(\mathsf{Var}(\ulcorner \underline{x} \urcorner), \mathsf{Nil}())))$$

- If $\underline{x}$ and $\underline{Suc}$ both correspond to zero:

  $$\mathsf{Lambda}(\mathsf{Zero}(), \mathsf{Const}(\mathsf{Zero}(), \mathsf{Cons}(\mathsf{Var}(\mathsf{Zero}()), \mathsf{Nil}()))))$$

Example

Representation (abstract syntax):

const $\underline{Lambda}$ (
  cons (const $\underline{Zero}$ nil) (
  cons (const $\underline{Const}$ (
    cons (const $\underline{Zero}$ nil) (
    cons (const $\underline{Cons}$ (
      cons (const $\underline{Var}$ (cons (const $\underline{Zero}$ nil) nil)) (
      cons (const $\underline{Nil}$ nil)
      nil)))
    nil)))
  nil))

Quiz

How is $\mathbf{rec}\ x = x$ represented?
Assume that x corresponds to 1.

1. $\mathsf{Rec}(\mathsf{X}(), \mathsf{X}())$
2. $\mathsf{Rec}(\mathsf{X}(), \mathsf{Var}(\mathsf{X}()))$
3. $\mathsf{Equals}(\mathsf{Rec}(\mathsf{X}()), \mathsf{X}())$
4. $\mathsf{Rec}(\mathsf{Suc}(\mathsf{Zero}()), \mathsf{Suc}(\mathsf{Zero}()))$
5. $\mathsf{Rec}(\mathsf{Suc}(\mathsf{Zero}()), \mathsf{Var}(\mathsf{Suc}(\mathsf{Zero}())))$
6. $\mathsf{Equals}(\mathsf{Rec}(\mathsf{Suc}(\mathsf{Zero}())), \mathsf{Suc}(\mathsf{Zero}()))$

The halting problem, take two

The intensional halting problem (with self-application)

There is no closed expression $halts$ such that, for every closed expression p,

- $halts \ulcorner p \urcorner \Downarrow \mathsf{True}()$, if $p \ulcorner p \urcorner$ terminates, and
- $halts \ulcorner p \urcorner \Downarrow \mathsf{False}()$, otherwise.

With self-application

- Assume that $halts$ can be defined.
- Define the closed expression $terminv$:

$$\begin{aligned} terminv = \lambda\, p.\, \mathbf{case}\ halts\ p\ \mathbf{of} \\ \{\, \mathsf{True}() \rightarrow \mathsf{rec}\ x = x \\ ;\, \mathsf{False}() \rightarrow \mathsf{Zero}() \\ \} \end{aligned}$$

- For any closed expression p:
  $terminv\ulcorner p\urcorner$ terminates iff
  $p\ulcorner p\urcorner$ does not terminate.
- Thus $terminv\ulcorner terminv\urcorner$ terminates iff
  $terminv\ulcorner terminv\urcorner$ does not terminate.

The intensional halting problem

There is no closed expression $halts$ such that, for every closed expression p,

- $halts \ulcorner p \urcorner \Downarrow \mathsf{True}()$, if p terminates, and
- $halts \ulcorner p \urcorner \Downarrow \mathsf{False}()$, otherwise.

The intensional halting problem

- Assume that $halts$ can be defined.
- If we can use $halts$ to solve the previous variant of the halting problem, then we have found a contradiction.

The intensional halting problem

Exercise: Define a closed expression $code$ satisfying

$$code \ulcorner p \urcorner \Downarrow \ulcorner \ulcorner p \urcorner \urcorner$$

for any closed expression p.

The intensional halting problem

Exercise: Define a closed expression $code$ satisfying

$$code \ulcorner p \urcorner \Downarrow \ulcorner \ulcorner p \urcorner \urcorner$$

for any closed expression p.

Example:

$$\ulcorner \ulcorner \lambda x.\, x \urcorner \urcorner$$

The intensional halting problem

Exercise: Define a closed expression $code$ satisfying

$$code \ulcorner p \urcorner \Downarrow \ulcorner \ulcorner p \urcorner \urcorner$$

for any closed expression p.

Example:

$$\ulcorner \ulcorner \mathsf{lambda}\ \underline{x}\ (\mathsf{var}\ \underline{x}) \urcorner \urcorner$$

The intensional halting problem

Exercise: Define a closed expression $code$ satisfying

$$code \ulcorner p \urcorner \Downarrow \ulcorner \ulcorner p \urcorner \urcorner$$

for any closed expression p.

Example:

$$\ulcorner \mathsf{Lambda}(\ulcorner \underline{x} \urcorner, \mathsf{Var}(\ulcorner \underline{x} \urcorner)) \urcorner$$

The intensional halting problem

Exercise: Define a closed expression $code$ satisfying

$$code \ulcorner p \urcorner \Downarrow \ulcorner \ulcorner p \urcorner \urcorner$$

for any closed expression p.

Example:

$$\ulcorner \mathsf{Lambda}(\mathsf{Zero}(), \mathsf{Var}(\mathsf{Zero}())) \urcorner$$

The intensional halting problem

Exercise: Define a closed expression $code$ satisfying

$$code \ulcorner p \urcorner \Downarrow \ulcorner \ulcorner p \urcorner \urcorner$$

for any closed expression p.

Example:

$\ulcorner$ const $\underline{Lambda}$ (
    cons $\ulcorner$ Zero() $\urcorner$ (
    cons $\ulcorner$ Var(Zero()) $\urcorner$
    nil)) $\urcorner$

The intensional halting problem

Exercise: Define a closed expression $code$ satisfying

$$code \ulcorner p \urcorner \Downarrow \ulcorner \ulcorner p \urcorner \urcorner$$

for any closed expression p.

Example:

$\ulcorner$ const $\underline{Lambda}$ (
  cons (const $\underline{Zero}$ nil) (
  cons $\ulcorner$ Var(Zero()) $\urcorner$
  nil)) $\urcorner$

The intensional halting problem

Exercise: Define a closed expression $code$ satisfying

$$code \ulcorner p \urcorner \Downarrow \ulcorner \ulcorner p \urcorner \urcorner$$

for any closed expression p.

Example:

$\ulcorner$ const $\underline{Lambda}$ (
cons (const $\underline{Zero}$ nil) (
cons (const $\underline{Var}$ (cons (const $\underline{Zero}$ nil) nil))
nil)) $\urcorner$

The intensional halting problem

Exercise: Define a closed expression $code$ satisfying

$$code \ulcorner p \urcorner \Downarrow \ulcorner \ulcorner p \urcorner \urcorner$$

for any closed expression p.

Example:

$$\begin{array}{l}
\mathsf{Const}(\ulcorner \underline{Lambda} \urcorner, \\
\quad \mathsf{Cons}(\mathsf{Const}(\ulcorner \underline{Zero} \urcorner, \mathsf{Nil}()), \\
\quad \mathsf{Cons}(\mathsf{Const}(\ulcorner \underline{Var} \urcorner, \\
\qquad \mathsf{Cons}(\mathsf{Const}(\ulcorner \underline{Zero} \urcorner, \mathsf{Nil}()), \\
\qquad \mathsf{Nil}())), \\
\quad \mathsf{Nil}())))
\end{array}$$

The intensional halting problem

Exercise: Define a closed expression $code$ satisfying

$$code\ulcorner p\urcorner \Downarrow \ulcorner\ulcorner p\urcorner\urcorner$$

for any closed expression p.

Example:

Const(Suc(Zero()),
  Cons(Const(Suc(Suc(Zero())), Nil()),
  Cons(Const(Suc(Suc(Suc(Zero())))),
    Cons(Const(Suc(Suc(Zero())), Nil()),
    Nil()))),
  Nil())))

The intensional halting problem

Define the closed expression $halts'$ by

$$\lambda\, p.\, halts\ \mathsf{Apply}(p, code\ p).$$

For any closed expression p:

$$\begin{array}{lll}
p\ \ulcorner p \urcorner \text{ terminates} & & \Rightarrow \\
halts\ \ulcorner p\ \ulcorner p \urcorner \urcorner & \Downarrow \mathsf{True}() & \Rightarrow \\
halts\ \mathsf{Apply}(\ulcorner p \urcorner, \ulcorner \ulcorner p \urcorner \urcorner) & \Downarrow \mathsf{True}() & \Rightarrow \\
halts\ \mathsf{Apply}(\ulcorner p \urcorner, code\ \ulcorner p \urcorner) & \Downarrow \mathsf{True}() & \Rightarrow \\
halts'\ \ulcorner p \urcorner & \Downarrow \mathsf{True}() &
\end{array}$$

The intensional halting problem

Define the closed expression $halts'$ by

$$\lambda\, p.\, halts\ \mathsf{Apply}(p, code\ p).$$

For any closed expression p:

$$\begin{array}{lll}
p\ \ulcorner p \urcorner \text{ does not terminate} & & \Rightarrow \\
halts\ \ulcorner p\ \ulcorner p \urcorner \urcorner & \Downarrow \mathsf{False}() & \Rightarrow \\
halts\ \mathsf{Apply}(\ulcorner p \urcorner, \ulcorner \ulcorner p \urcorner \urcorner) & \Downarrow \mathsf{False}() & \Rightarrow \\
halts\ \mathsf{Apply}(\ulcorner p \urcorner, code\ \ulcorner p \urcorner) & \Downarrow \mathsf{False}() & \Rightarrow \\
halts'\ \ulcorner p \urcorner & \Downarrow \mathsf{False}() &
\end{array}$$

Thus $halts'$ solves the previous variant of the halting problem, and we have found a contradiction.

Summary

- Concrete and abstract syntax.
- Operational semantics.
- Several variants of the halting problem.
- Representing inductively defined sets.