# Finite Automata Theory and Formal Languages TMV027/DIT321 - LP4 2018 

## Regular Languages

## Assignment 4 - Deadline: Sunday 29th of April 23:59 <br> Assignments should be done and submitted individually!

For obtaining full points the answers should contain enough explanation/description so that they are easy to understand.

1. (2.5pts) Show as formal and clear as possible that $a^{*}\left(b+a b^{*}\right)=b+a a^{*} b^{*}$.
2. Give concrete examples of languages $\mathcal{L}_{1}$ and $\mathcal{L}_{2}$ over the alphabet $\{0,1\}$ such that
(a) (1pt) Both $\mathcal{L}_{1}$ and $\mathcal{L}_{2}$ are non-regular, but $\mathcal{L}_{1} \cup \mathcal{L}_{2}$ is regular;
(b) (1pt) Both $\mathcal{L}_{1}$ and $\mathcal{L}_{2}$ are non-regular, and $\mathcal{L}_{1} \cap \mathcal{L}_{2}$ is infinite and non-regular;
(c) (1pt) $\mathcal{L}_{1}$ is regular, $\mathcal{L}_{2}$ is non-regular, and $\mathcal{L}_{1} \cup \mathcal{L}_{2}$ is regular

Justify your answer by explicitly giving the resulting language in each case. If it is not easily clear why the languages are regular/non-regular, you need to even justify this.
3. (2.25pts) Use the Pumping lemma for regular languages to show that the language $\left\{w \in\{0,1,2\}^{*} \mid \#_{0}(w)+\#_{1}(w)=\#_{2}(w)\right\}$ is not a regular language, where $\#_{0}, \#_{1}$ and $\#_{2}$ are functions counting the number of 0 's, 1 's and of 2's, respectively, in a word.
4. Minimise the following automaton:

|  | 0 | 1 |
| ---: | :---: | :---: |
| $\rightarrow^{*} q_{0}$ | $q_{1}$ | $q_{3}$ |
| $q_{1}$ | $q_{4}$ | $q_{2}$ |
| ${ }^{*} q_{2}$ | $q_{1}$ | $q_{5}$ |
| $q_{3}$ | $q_{0}$ | $q_{4}$ |
| $q_{4}$ | $q_{4}$ | $q_{4}$ |
| $q_{5}$ | $q_{2}$ | $q_{4}$ |

(a) (1.5pts) Show the table that identifies the distinguishable states;
(b) (0.3pts) Indicate the equivalent clases of states resulting from the information in the table;
(c) (0.45pts) Give the minimised automaton.

