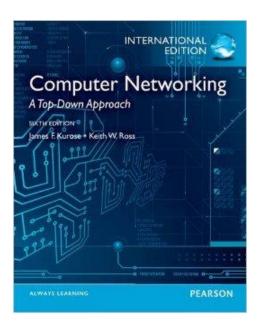
# Chapter 8 Network Security



# Roadmap



- 8.1 What is network security?
- 8.2 Principles of cryptography
- 8.3 Message integrity

Security protocols and measures:

- □ Securing TCP connections: SSL
- □ Network layer security: IPsec
- ☐ Firewalls

Network Layer 4-2

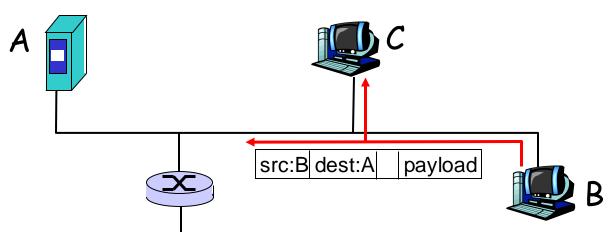
# What is security? CIA!

- Confidentiality: only sender, intended receiver should "understand" message contents
  - sender encrypts message
  - o receiver decrypts message
- Integrity: sender, receiver want to ensure message not altered (in transit, or afterwards) without detection
- Availability: services must be accessible and available to users

The book also includes Authentication: it is normally seen as a mechanism to implement the services above

#### Packet sniffing:

- broadcast media
- promiscuous NIC reads all packets passing by
- o can read all unencrypted data (e.g. passwords)
- e.g.: C sniffs B's packets

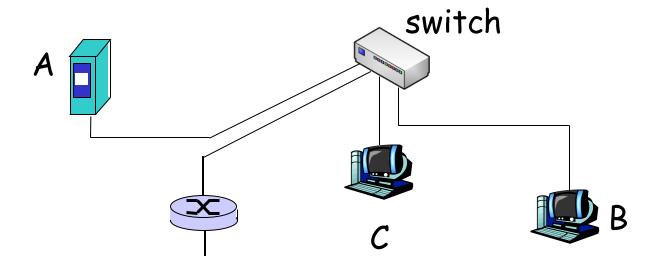


Countermeasures?



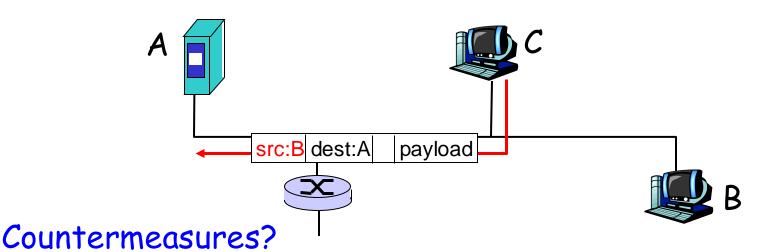
#### Packet sniffing: countermeasures

- One host per segment of broadcast media
   Use switches (not hubs)
- Segment network
  - Use routers
- Encryption



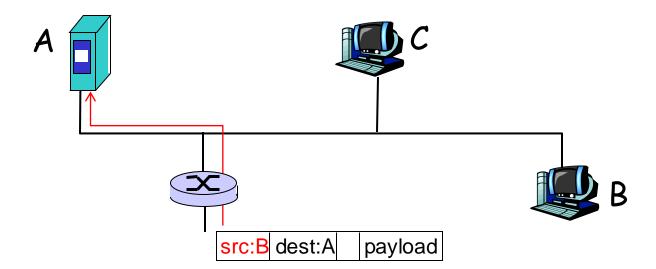
#### IP Spoofing:

- can generate "raw" IP packets directly from application, putting any value into IP source address field
- receiver can't tell if source is spoofed,
   e.g.: C pretends to be B

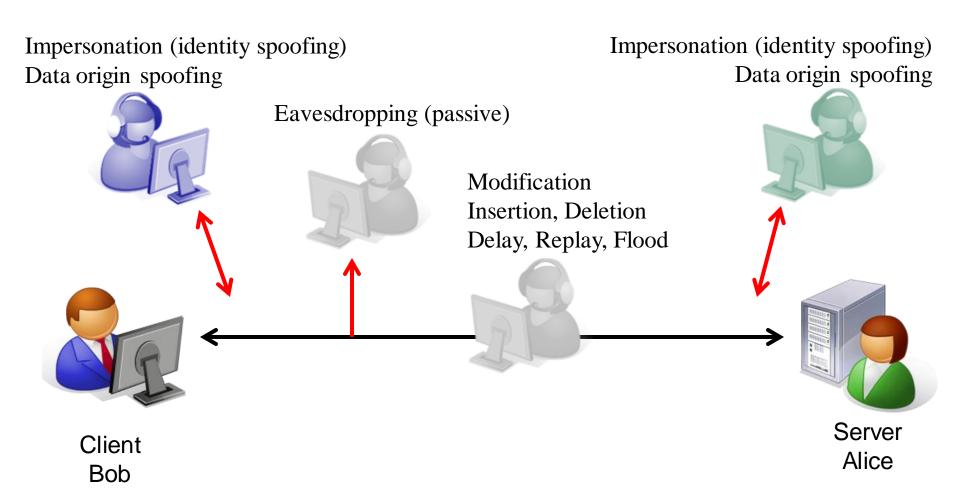


#### IP Spoofing: ingress filtering

- routers should not forward incoming and outgoing packets with invalid addresses
  - Outgoing datagram source address not in router's network (egress filtering)
  - Incoming datagram has internal address as source address (ingress filtering)



# Communication threats - Summary



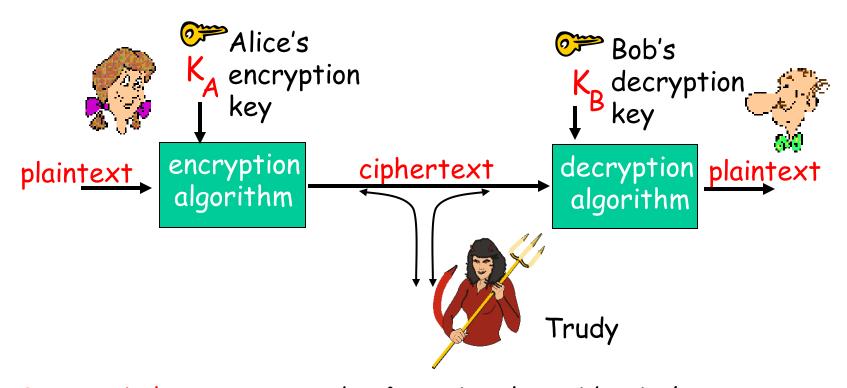
# Roadmap

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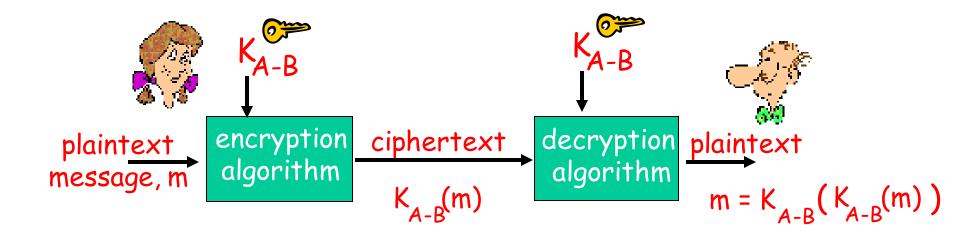
### The language of cryptography



Symmetric key crypto: sender & receiver keys identical

Asymmetric key crypto (or Public-key crypto):
One key for encryption, another for decryption.
One of the keys can be public, the other private.

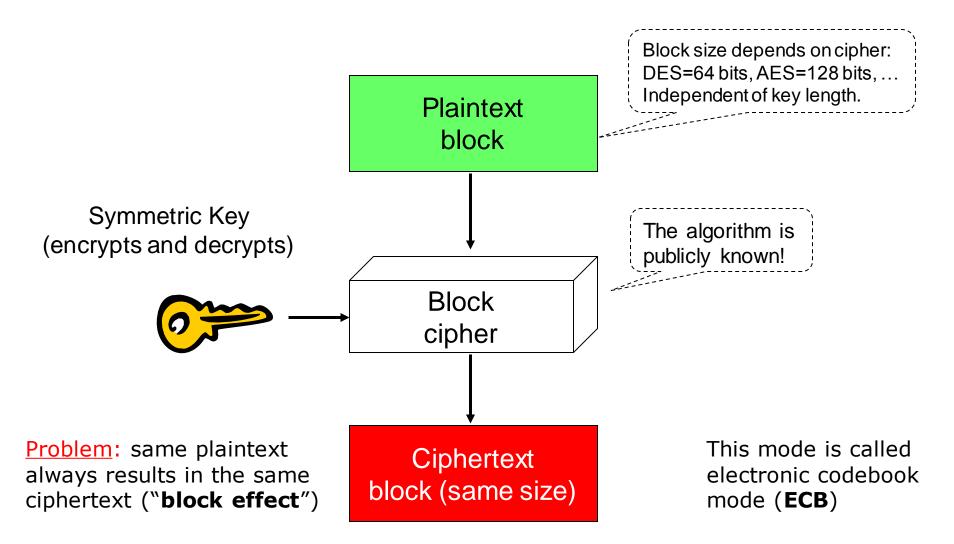
#### Symmetric key cryptography



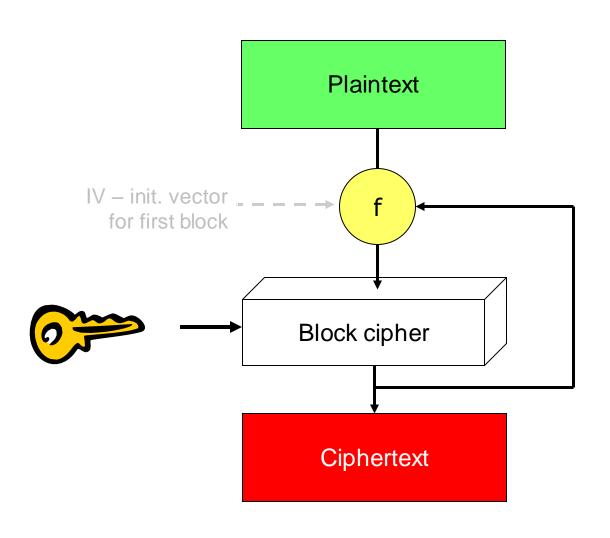
symmetric key crypto: Bob and Alice share the same (symmetric) key:  $K_{A-B}$ 

Q: how do Bob and Alice agree on key value?

# Block Encryption (ECB mode)



### CBC - Cipher block chaining mode

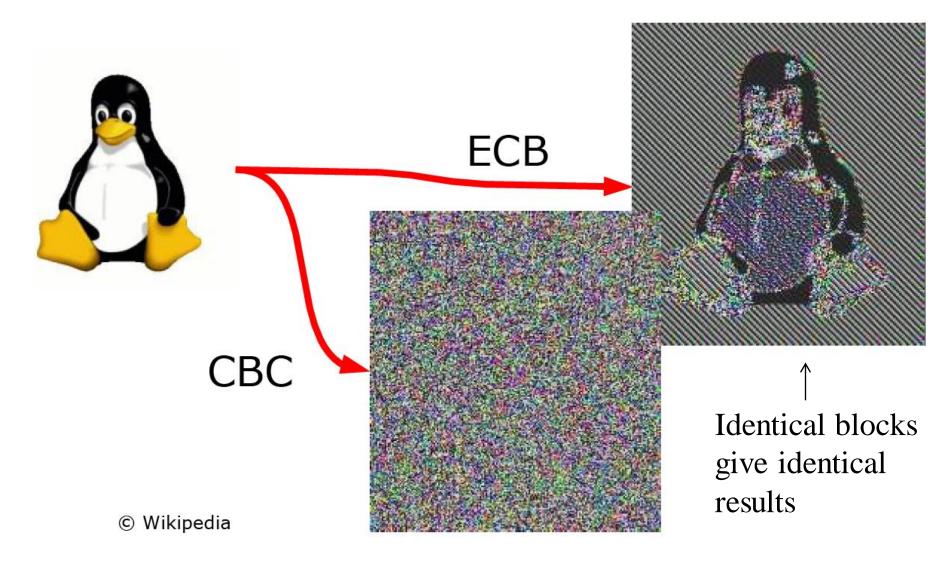


Identical blocks now encrypted differently.

May not always be practical, for example for hard disk encryption.

Note that there is no protection against replays and alteration!

# ECB vs. CBC



# Symmetric Key Ciphers

- DES (Data Encryption Standard)
  - Designed by IBM 1975, Adopted by NIST\* 1977
  - Criticized for key length  $(64 \rightarrow 56)$  and mysterious "S-boxes"
  - Turned out to have protection against differential cryptanalysis (found 1990)
  - o Probably more effort is spent on cracking DES than on all other ciphers together
  - Today key length is a major problem: 56-bit keys can be cracked

EFF DES cracker. Jan 19, 1999: 22h15m

- 3-DES (repeating DES three times with different keys)
  - 3-DES probably secure today but too computational intensive
- □ AES (Advanced Encryption Standard)
  - Replaces DES as of 2001
  - Result of an official competition
  - Key lengths: 128, 192 or 256 bits
  - Brute force decryption: if DES takes 1 second, AES-128 takes 149 trillion years, AES-256 would take 10<sup>52</sup> years
- RC4, RC5, RC6
  - RC4 is considered weak but it is fast

\*NIST = National Institute of Standards and Technology, US, formerly NBS

J ...

#### Key Length and Number of Possible Keys

Key Length in Bits	Number of Possible Keys
1	2
2	4
40	1,099,511,627,776
56	72,057,594,037,927,900
112	5,192,296,858,534,830,000,000,000,000,000,000
168	3.74144E+50
<b>950</b> 256	1.15792E+77
512	1.3408E+154

# Asymmetric key encryption

- One key is used to encrypt, the other to decrypt
- One key can be public the other kept secret
- Based on mathematically hard problems
  - Factorization of very large primes (RSA)
- Slow because of the large numbers involved
  - 1024 bits and up (RSA), 384 bits (ECC)
  - $\circ$  2<sup>1024</sup> = 10<sup>308</sup> which means >300 digit numbers
- Ciphers:
  - RSA Rivest, Shamir, Adleman (Patent expired 2000)
  - ECC Elliptic Curve Cryptosystem
- 768-bit RSA was reported cracked Jan 2010:
  - They generated a five-terabyte decryption table. It would have taken around 1,500 years using a single AMD Opteron-based PC (they used a cluster)
- 1024-bit RSA is too short to protect against extremely large organizations
  - Use 2048-bit RSA keys in sensitive applications



"the overall effort [as] sufficiently low that even for short-term protection of data of little value, 768-bit RSA moduli can no longer be recommended."

# Asymmetric key encryption

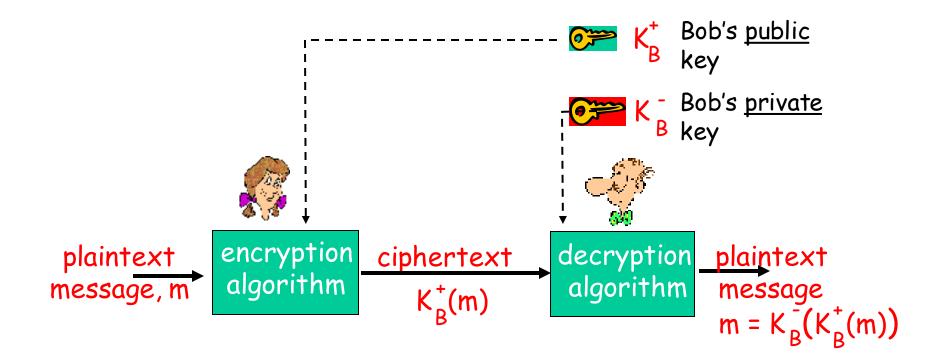
 One key is normally made public ("Public key encryption")



- You decide whether it is the encryption or decryption key that is public:
- 1. Encryption key public: everyone can send encrypted messages to owner of the private key
- 2. Decryption key public: only one can encrypt, everyone can verify that the secret key has been used.

  Useful?
  - Can be used to sign documents and data.

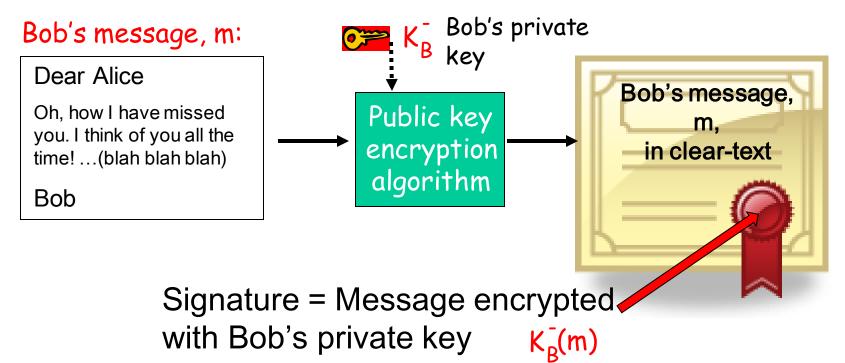
### Example 1: Public Key Encryption



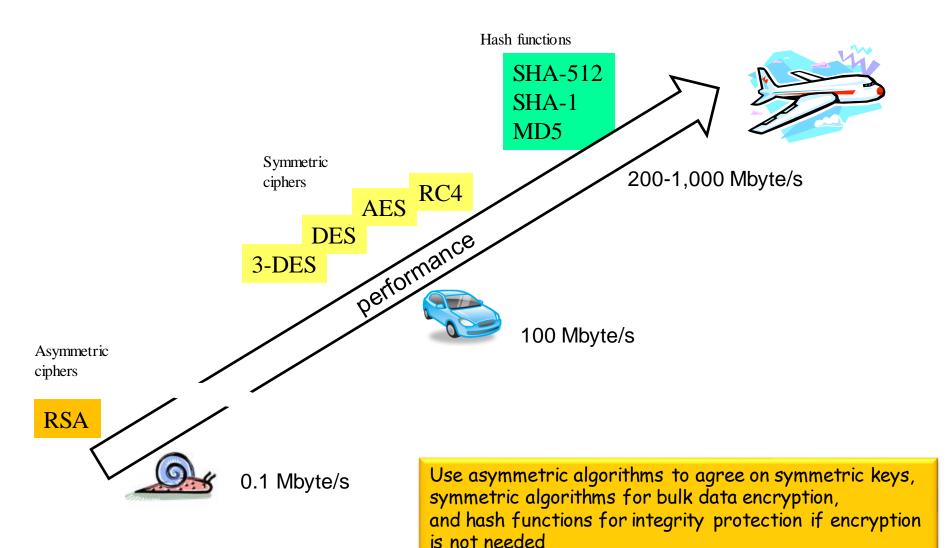
# Example 2: Digital Signatures

#### Simple digital signature for message m:

 $\square$  Bob signs m by encrypting with his private key  $K_B$ , creating "signed" message,  $K_B$ (m)



# Relative performance



# Roadmap

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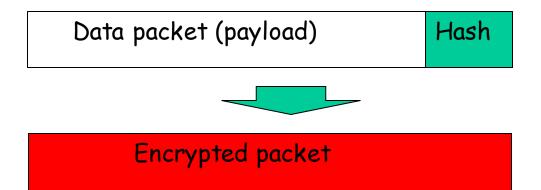
### Message Integrity

Bob receives msg from Alice, wants to ensure:

- □ message originally came from Alice
- message not changed since sent by Alice

#### Just encryption is not enough!

- Contents can be changed even if it is encrypted
- □ Solution: add some kind of checksum (hash) to the message before it is encrypted:



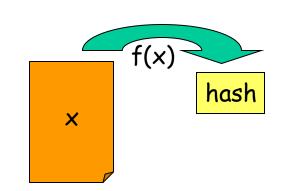
# (Cryptographic) hash functions

- Input: arbitrary length bit-stringOutput: fixed length bit-string
  - Not a one-to-one mapping, output space typically 128 bits



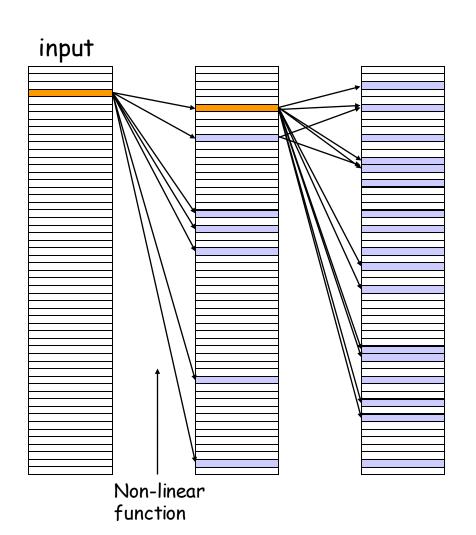
- Computationally efficient: Typically > 10 times faster than symmetric ciphers
- Must be repeatable (same input → same output)
- Impossible to reverse the computation (preimage resistant)
- Infeasible to find an input X with a given hash
- Infeasible to find two inputs resulting in the same hash (pseudorandomness)
- Today's hash functions are not based on mathematical foundations - may lead to problems

"SSL broken! Hackers create rogue CA certificate using MD5 collisions" [www.zdnet.com]



### Hash functions

Even a single bit change should give a completely different result  $\rightarrow$  avalanche effect



SHA-512 has 80 rounds

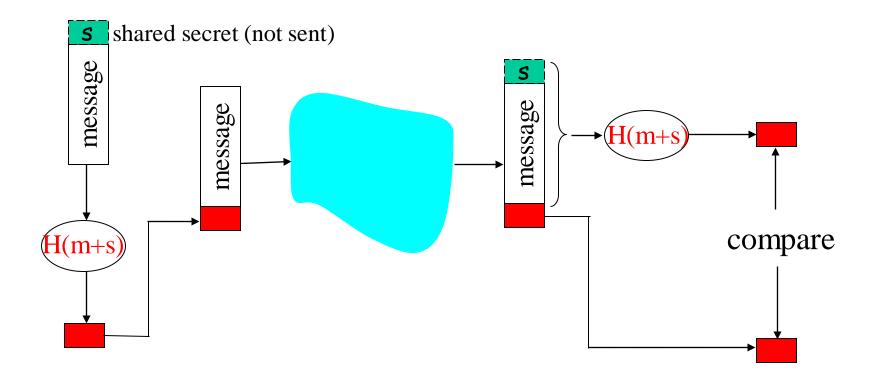
### Hash functions

- Even just one changed bit gives a completely different result:
  - $\circ$  md5("hello") = 5d41402abc4b2a76b9719d911017c592
  - o md5("Hello") = 8b1a9953c4611296a827abf8c47804d7
- MD5 Message Digest 5 (RFC 1321, 1992)
  - 128-bit message digest  $\rightarrow$  10<sup>38</sup> different hashes
  - Avoid in new implementations weak
- SHA-1 Secure Hash Algorithm
  - Designed by NSA, became NIST standard 1995: FIPS-180-2
  - 160-bit message digest  $\rightarrow$  10<sup>48</sup> different hashes
  - Avoid if collisions may cause problems in application, otherwise ok
- □ **SHA-2** (family name for SHA-224, SHA-256, SHA-384 and SHA-512)
  - Similar design as SHA-1, but at least today SHA-1 attacks not applicable
- □ SHA-3 next generation hash functions
  - Keccak winner of open competition (NIST draft 2014)
  - Arbitrary digest size (standard proposes 224, 256, 384 and 512 bit digests)

"As of 2012, an estimated cost of \$2.77M to break a single hash value by renting CPU power from cloud servers."

- SHA-1, Wikipedia

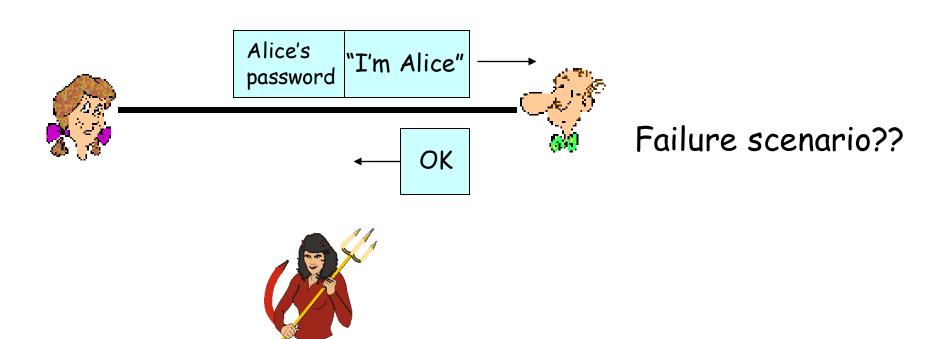
#### Keyed Hash - No need to encrypt message



- Authenticates sender
- □ Verifies message integrity
- No encryption!
- Example: HMAC (Key-Hashing for Message Authentiction)

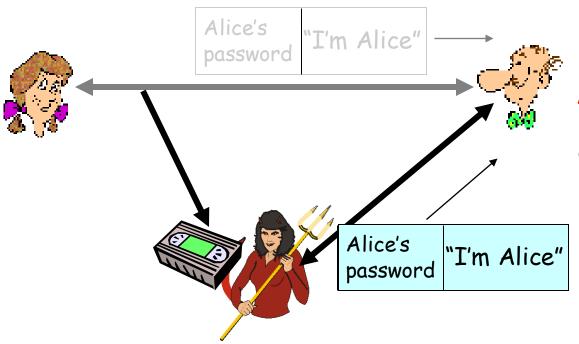
### End point (User) Authentication

Alice says "I am Alice" and sends her secret password to "prove" it.
(Just like the FTP protocol)



### End point (User) Authentication

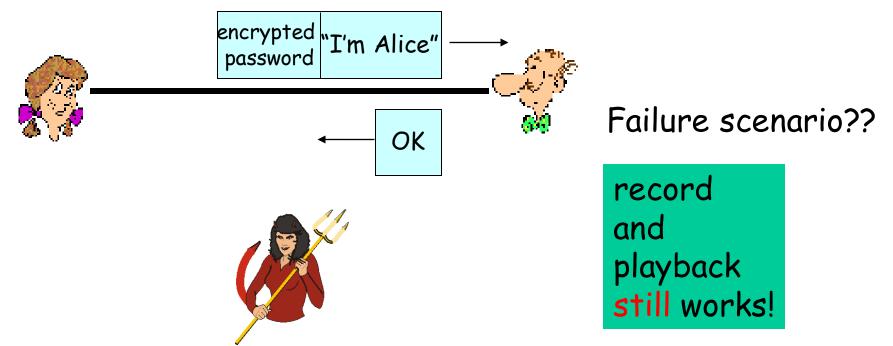
Alice says "I am Alice" and sends her secret password to "prove" it.



playback attack: Trudy records Alice's packet and later plays it back to Bob

#### Authentication: another try

<u>Another attempt:</u> Alice says "I am Alice" and sends her <u>encrypted</u> secret password to "prove" it.

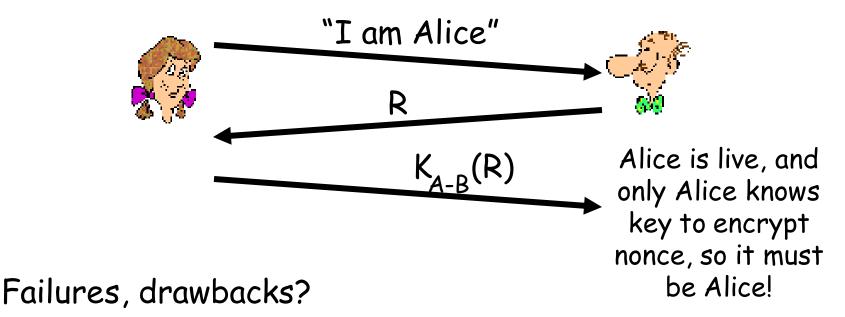


#### Authentication: Challenge response

Goal: avoid playback attack

Nonce: number (R) used only once-in-a-lifetime

To prove Alice is "live", Bob sends Alice nonce, R. Alice must return R, encrypted with shared secret key



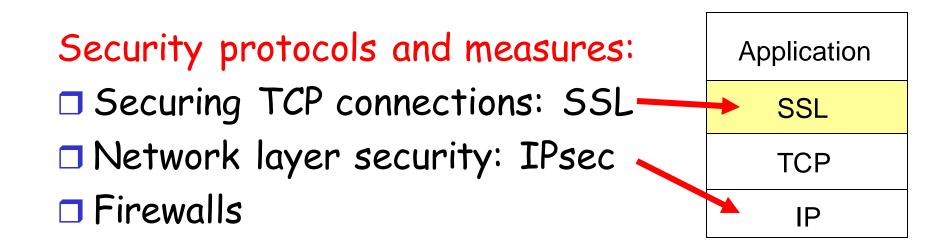
# Summary

- Encryption for confidentiality
- Hashes for data integrity
- Sequence numbers for replay protection
- Authentication (mutual) for identity protection

- □ Symmetric encryption for bulk data
- Asymmetric encryption for key negotiation

# Roadmap

- 8.1 What is network security?
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# SSL: Secure Sockets Layer

- widely deployed security protocol
  - supported by almost all browsers, web servers
  - https
  - billions \$/year over SSL
- mechanisms: [Woo 1994], implementation: Netscape
- variation -TLS: transport layer security, RFC 2246
- provides
  - confidentiality
  - integrity
  - authentication

- original goals:
  - Web e-commerce transactions
  - encryption (especially credit-card numbers)
  - Web-server authentication
  - optional client authentication
  - minimum hassle in doing business with new merchant
- available to all TCP applications
  - secure socket interface

### SSL and TCP/IP

**Application** 

**TCP** 

IP

normal application

Application

SSL

**TCP** 

IΡ

application with SSL

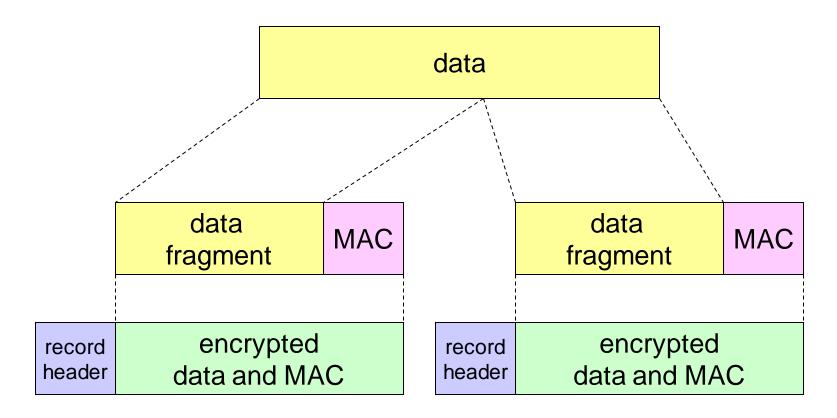
- SSL provides application programming interface (API) to applications
- C and Java SSL libraries/classes readily available

#### handshake: ClientHello Real SSL handshake: ServerHello connection handshake: Certificate handshake: ServerHelloDone handshake: ClientKeyExchange ChangeCipherSpec everything handshake: Finished henceforth ChangeCipherSpec is encrypted handshake: Finished application data application\_data

**TCP FIN follows** 

Alert: warning, close\_notify

# SSL record protocol



record header: content type; version; length

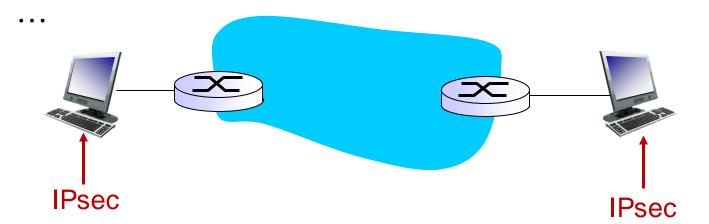
MAC: includes sequence number, MAC key  $M_x$ 

fragment: each SSL fragment 2<sup>14</sup> bytes (~16 Kbytes)

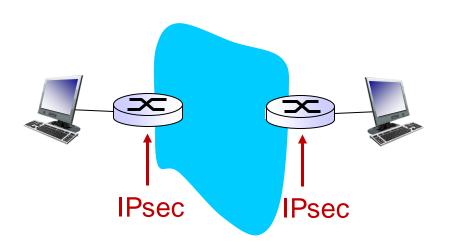
## What is network-layer confidentiality?

### between two network entities:

- sending entity encrypts datagram payload, payload could be:
  - TCP or UDP segment, ICMP message, OSPF message ....
- all data sent from one entity to other would be hidden:
  - web pages, e-mail, P2P file transfers, TCP SYN packets

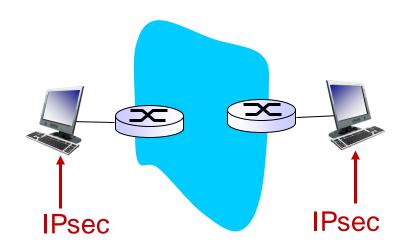


## The two modes of IPSec





- edge routers IPsec-aware
- protects communication gw-to-gw (over Internet)
- Virtual Private Network (VPN)



## Transport mode

- hosts IPsec-aware
- protects communication all the way from end-toend

# **IPsec services**

- data integrity
- confidentiality

- origin authentication
- replay attack prevention

## two protocols providing different service models:

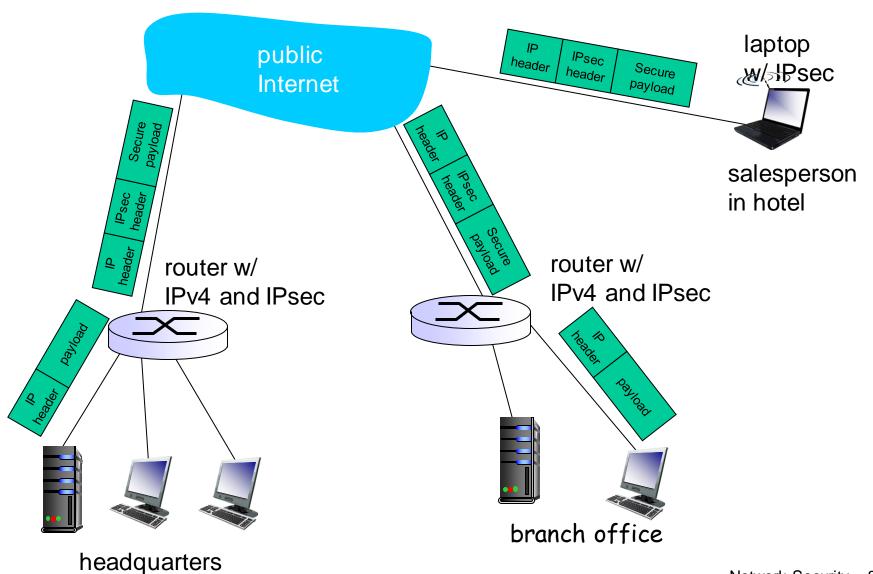
- Authentication Header (AH) protocol
  - provides source authentication & data integrity but *not* confidentiality
- Encapsulation Security Protocol (ESP)
  - provides source authentication, data integrity, and confidentiality
  - more widely used than AH

# Virtual Private Networks (VPNs)

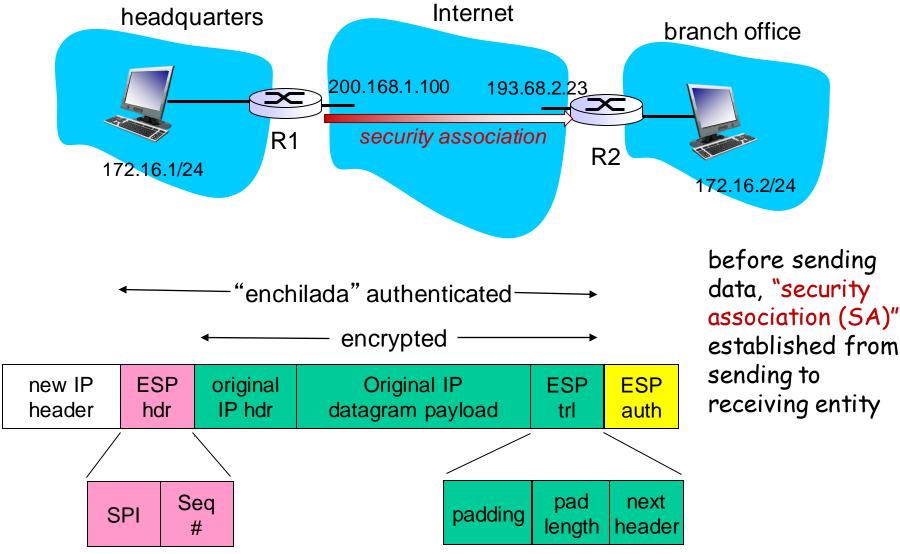
### motivation:

- institutions often want private networks for security.
  - costly: separate routers, links, DNS infrastructure.
- VPN: institution's inter-office traffic is sent over public Internet instead
  - encrypted before entering public Internet
  - logically separate from other traffic

## Virtual Private Networks (VPNs)



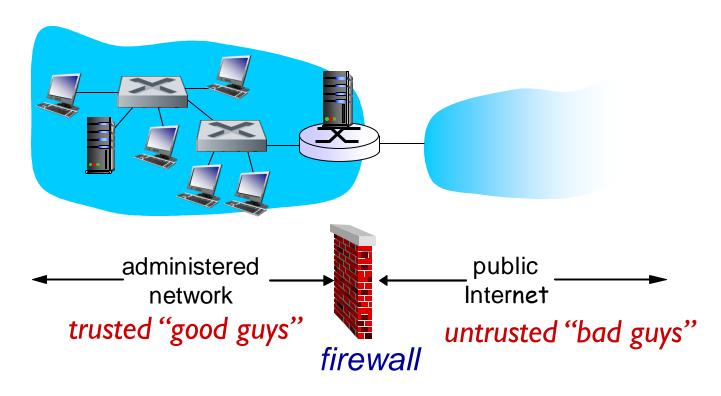
# What happens?



# **Firewalls**

## firewall

isolates organization's internal net from larger Internet, allowing some packets to pass, blocking others



# Firewalls: why

### prevent denial of service attacks:

SYN flooding: attacker establishes many bogus TCP connections, no resources left for "real" connections

### prevent illegal modification/access of internal data

- \* e.g., attacker replaces CIA's homepage with something else allow only authorized access to inside network
  - set of authenticated users/hosts

## three types of firewalls:

- stateless packet filters
- stateful packet filters
- application gateways

# Security Courses at Chalmers

http://www.cse.chalmers.se/edu/master/secspec/



at Chalmers and University of Gothenburg





We are proud to possess multifaceted security expertise at Chalmers University of Technology and University of Gothenburg, home to a world-leading research environment on computer and network security.

Based on this expertise, we offer a security specialization that consists of the following course package\*

#### Computer Security

The course provides basic knowledge in the security area, i.e. how to protect systems against attacks. Attacks may change or delete resources (data, programs, hardware, etc), get unauthorized access to confidential information or make unauthorized use of the system's services. The course covers threats and vulnerabilities, as well as rules, methods and mechanisms for protection. Modeling and assessment of security and dependability as well as metrication methods are covered. A holistic security approach is presented and organizational, business-related, social, human, legal and ethical aspects are treated.

#### Cryptography

The course covers cryptographic primitives such as private-key and public-key ciphers, hash functions, MAC's and signatures and how to embed these in cryptographic protocols to achieve basic goals such as confidentiality, authentication and non-repudiation, but also more elaborate services, such as key management, digital cash and electronic voting. Many examples of broken protocols are also discussed to enhance understanding of the engineering difficulties in building secure systems.

Runs in study period 2

#### Language-based Security

The course covers the principles of programming language-based techniques for computer security. The goal is understanding such application-level attacks as races, buffer overruns, covert channels, and code injection as well as mastering the principles behind such language-based protection techniques as static analysis, program transformation, and reference monitoring. The dual perspective of attack vs. protection is threaded through the lectures, laboratory assignments, and projects.

Runs in study period 4.

#### Network security

Why is it possible to break into networked applications and computer systems? What weaknesses are used? And what makes one protocol more secure than another? This course answers these questions and many more. We look at weaknesses that have plagued wired and wireless networked systems for years and investigate the security of countermeasures like firewalls and security protocols such as SSL, SSH and IPsec. Knowledge about possible threats and countermeasures is important for understanding what level of security a system and an application can offer.

Runs in study period 4

Runs in study period 3

Security is becoming increasingly important for system design and development. System architects and designers must have security expertise, so that the systems they design do not fall victims to attacks. Software developers and engineers must have security expertise, so that the code they produce cannot be exploited. Security and network specialists must have critical knowledge of security principles and practice, in order to ensure the security of the systems they are responsible for.

# Some review questions

- What is the difference between symmetric and asymmetric cryptography? Give examples of applications where each can be applied.
- Can you "decrypt" a hash of a message to get the original message? Explain your answer.
- What is the purpose of a nonce in an end-point authentication protocol? What is the main property of a nonce?

# Reading instructions

- Kurose Ross (Course book)
  - 8.1 8.4, 8.6, 8.7, 8.9.1
- Other resources (extra material):
  - http://shattered.it/ information about practical SHA-I collision
  - https://thehackernews.com/ security news