



Course on Computer Communication and Networks

Lecture 12

Continuously evolving Internet-working

Part B: QoS, traffic engineering, SDN, IoT

EDA344/DIT 423, CTH/GU

Based on the book Computer Networking: A Top Down Approach, Jim Kurose, Keith Ross, Addison-Wesley.

Timing/bandwidth guarantees in networks

aka Quality of Service (QoS): 2-party agreement (NW user – NW provider) on

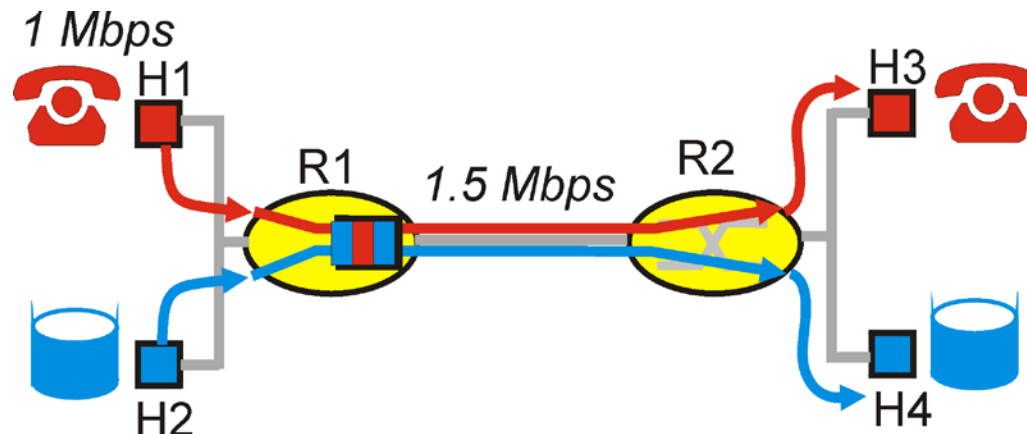
- Traffic characteristics (packet rate, sizes, ...)
- Network service guarantees (delay, jitter, loss rate, ...)

*Model for resource sharing and congestion studies:
questions/principles for QoS in Network Core*

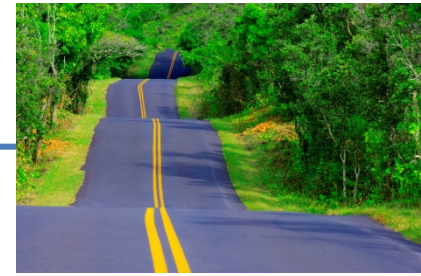
- Distinguish traffic?
- Control offered load? (isolate different "streams"?)
- Allocate: resources? (utilization)
- Control acceptance of new sessions?

Tasks for the NW core:

- Packet classification & scheduling (bandwidth allocation)
- Traffic shaping/policing (enforce contract terms)
- Admission control



Let's hit the road again😊: Roadmap



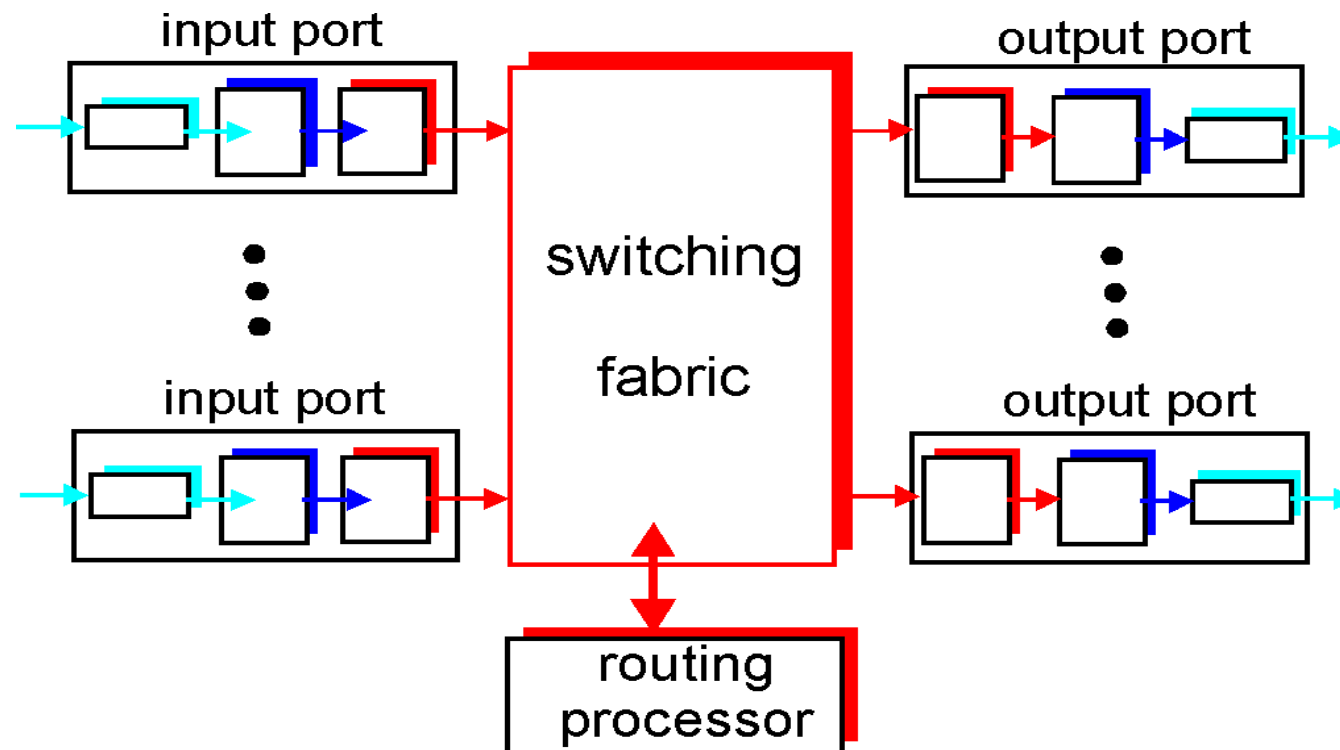
NW support for multimedia / QoS: [Ch. 9.5 (7.5 6/e)]

- **Improving timing/QoS guarantees in Networks** (also related with congestion-control):
Packet scheduling and policing
 - A VC (ATM) approach *[incl. Ch 3.7.2 (6e 3.62-3.6.3)]*
 - Internet approaches
 - Diff-serv, Int-serv + RSVP,
 - Traffic Engineering MPLS *[incl. ch. 6.5 (6/e 5.5)]*
- SDN *[ch 4.4, 5.5 (cf separate notes @pingpong docs, if you do not have access to 7e)]*
- Internet-of-Things in evolution: more types of traffic/devices... *[optional study, just browse example protocols mentioned]*

Where does this go in?

Scheduling = choosing the next packet for transmission on a link (= allocate bandwidth)

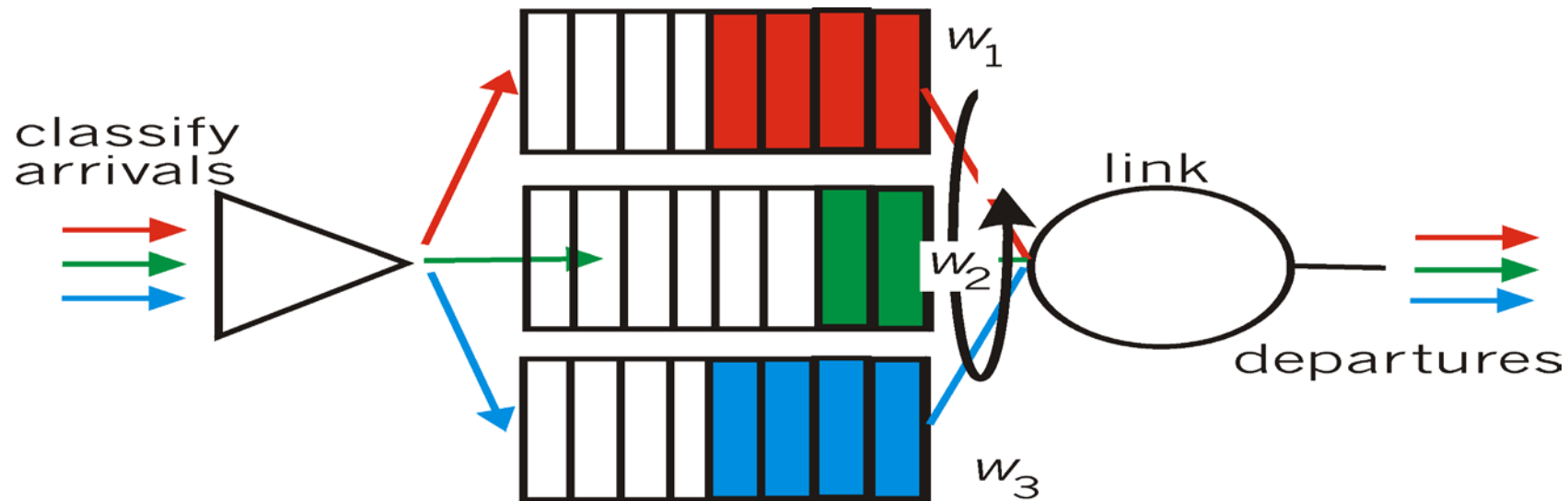
if buffer full: a **discard policy** determines which packet to discard among the arrival and those already queued



Packet Scheduling example: Weighted Fair Queueing

Weighted Fair Queueing: generalized **Round Robin**, including priorities (weights)

- provide each class with a differentiated amount of service
- class i receives a fraction of service $w_i / \sum(w_j)$



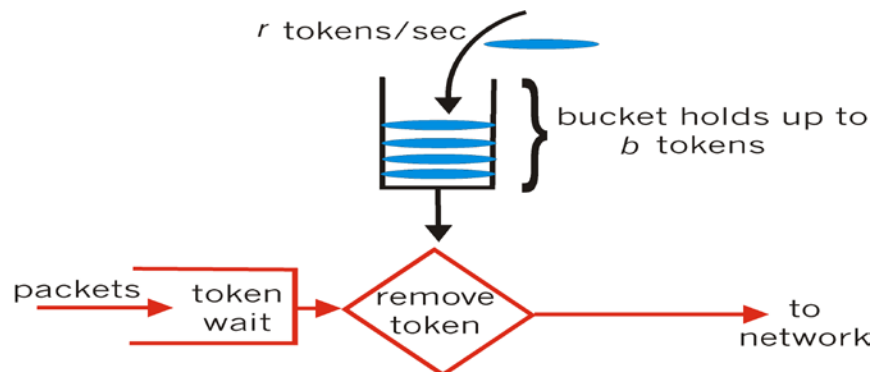
- There are a lot more decision options about packet scheduling: work-conserving policies, delays, ...

Policing Mechanisms

- Idea:** *shape* the packet traffic :network provider does *traffic policing*, ie enforces the "shape" agreed.
- **Traffic shaping**, to limit transmission rates:
 - (Long term) **Average Rate** (e.g.100 pkts/sec or 6000 packets per min)
 - **Peak Rate**: e.g.1500 pkts/sec peak
 - (Max.) **Burst Size**: Max. number of packets sent consecutively, ie over a very short period of time

Policing Mechanisms: LeakyToken Bucket

- Idea:** packets sent by consuming tokens that are produced at constant rate r
- limit input's
 - Burst Size (b = bucket capacity)
 - Average Rate (max admitted #packets over time period t is $b + rt$).



Another way to illustrate token buckets:

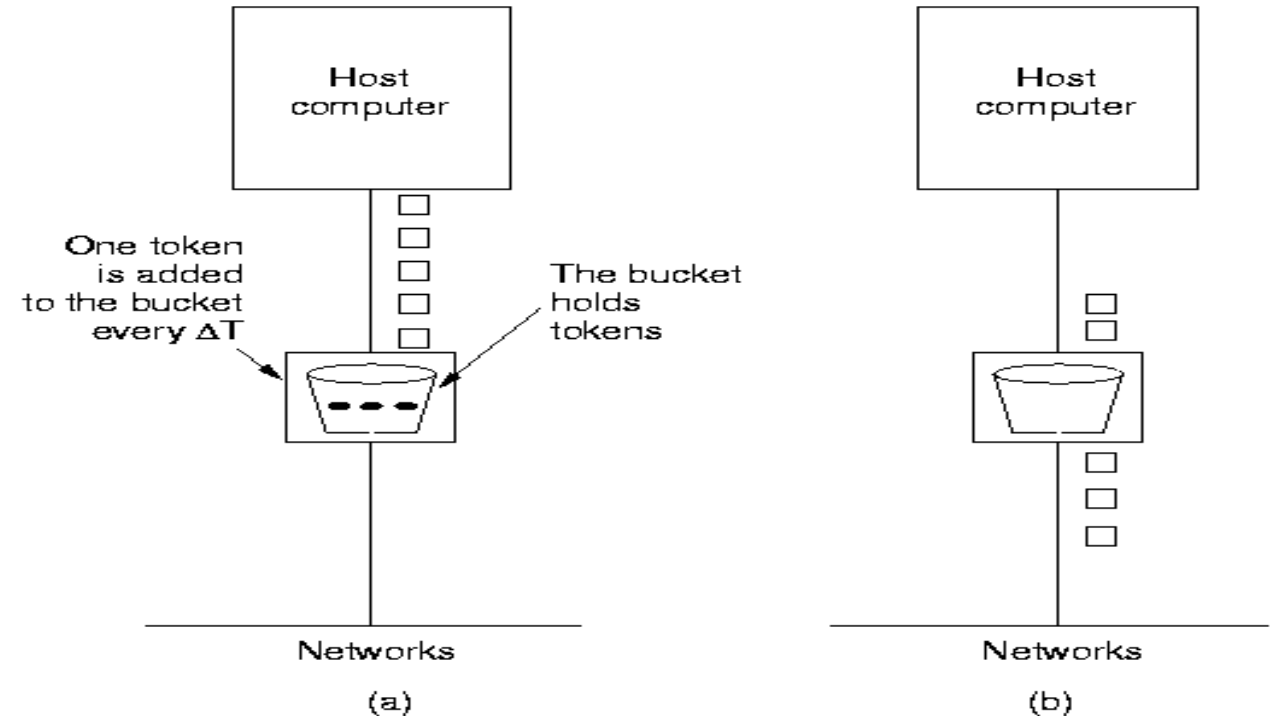


Fig. 5-26. The token bucket algorithm. (a) Before. (b) After.

Policing: the effect of buckets

input

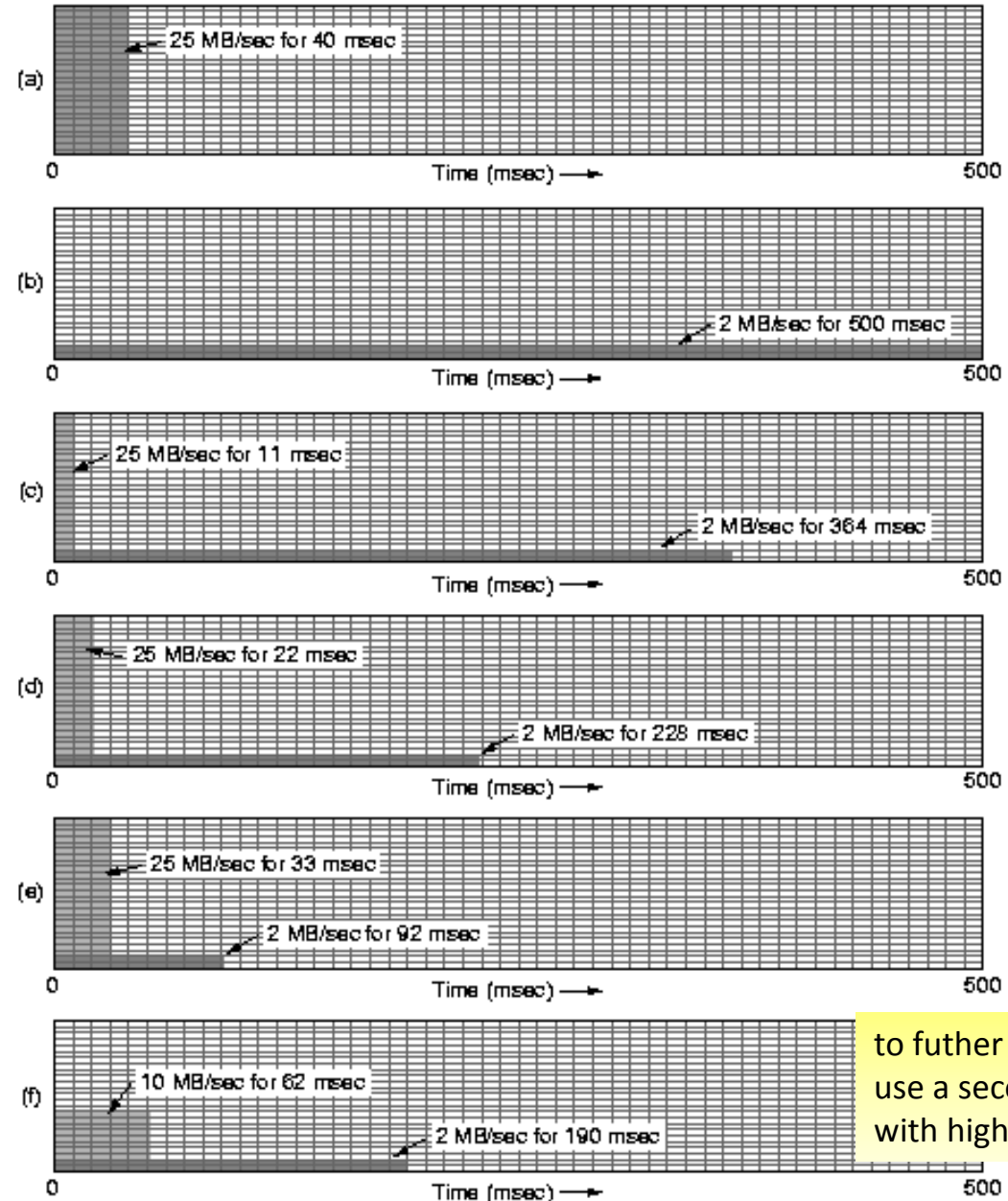
output 0KB token leaky bucket, 2MBps

output 250KB token leaky bucket, 2MBps

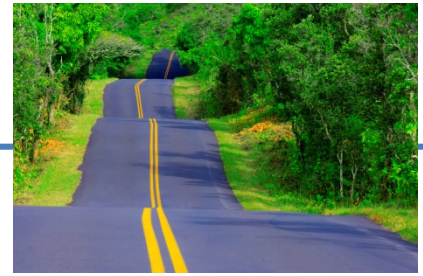
output 500KB token leaky bucket, 2MBps

output 750KB token leaky bucket, 2MBps

output token leaky bucket 500KB, 2MBps,
feeding 0KB, 10MBps token leaky bucket



to further limit burstiness,
use a second leaky bucket
with higher rate



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- **Internet-of-Things in evolution: more types of traffic/devices...** [optional study, just browse example protocols mentioned]

Virtual Circuit example:

ATM: Asynchronous Transfer Mode nets

Internet 's IP:

- today's *de facto* standard for global data networking

1980's:

- telco's develop ATM specifications: competing network standard for carrying high-speed voice/data

ATM principles:

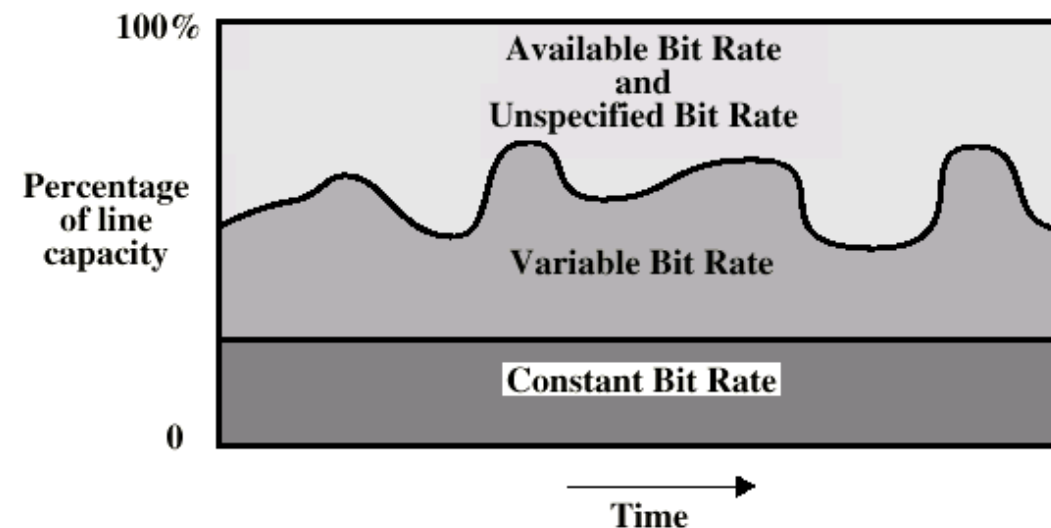
- **virtual-circuit networks**: switches maintain state for each "call"
- small (48 byte payload, 5 byte header) fixed length *cells* (like packets)
 - fast switching
 - small size good for voice
- well-defined interface between "network" and "user" (think of classic telecom)

Example VC technology

ATM Network service models (i.e. transport layer services):

Service Model	Example	Guarantees ?				Congestion control
		Bandwidth	Loss	Order	Timing	
Constant Bit Rate	voice	constant rate	yes	yes	yes	Admission control
VariableBR (RT/nRT)	Video/ "streaming"	guaranteed rate	yes	yes	yes	Admission control
Available BR	www-browsing	guaranteed minimum	no	yes	no	Yes, feedback
UndefinedBR	Background file transfer	none	no	yes	no	discard pkts

With ABR you can get min guaranteed capacity and better, if possible; with UBR you can get better, but you may be thrown out in the middle ☹



ATM (VC) Congestion Control (**hand-in-hand with Bandwidth reservation**)

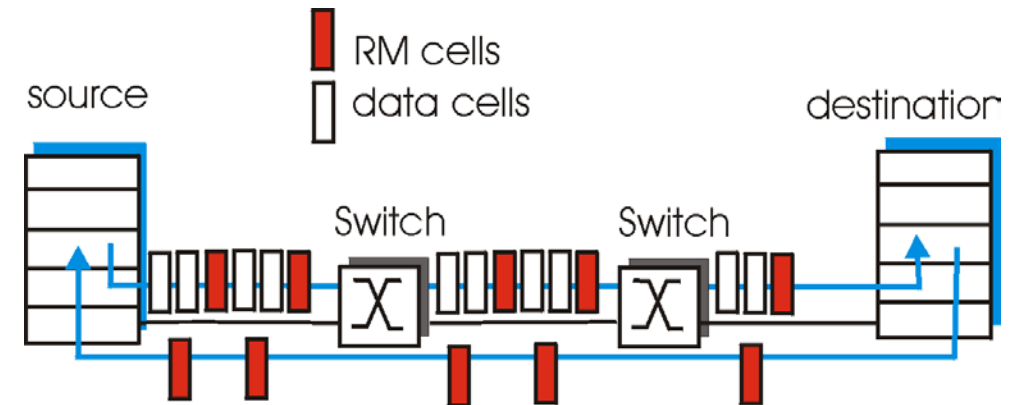
Several different strategies in place :

Admission control and resource reservation (CBR and VBR traffic:

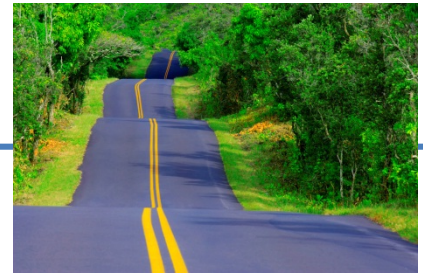
reserve resources when opening a VC; traffic shaping and policing (*use bucket-like methods*)

Rate-based congestion control: (ABR traffic)

- **idea** = feedback to the sender and intermediate stations on the *available* (= *max. acceptable*) rate on the VC.
- similar to “choke packets” (option provided in ICMP, which is not used in implementations...)



Roadmap



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Internet bandwidth-guarantee support possibilities?

Diffserv proposed Architecture

Edge router: marking



□ per-aggr-flow traffic management

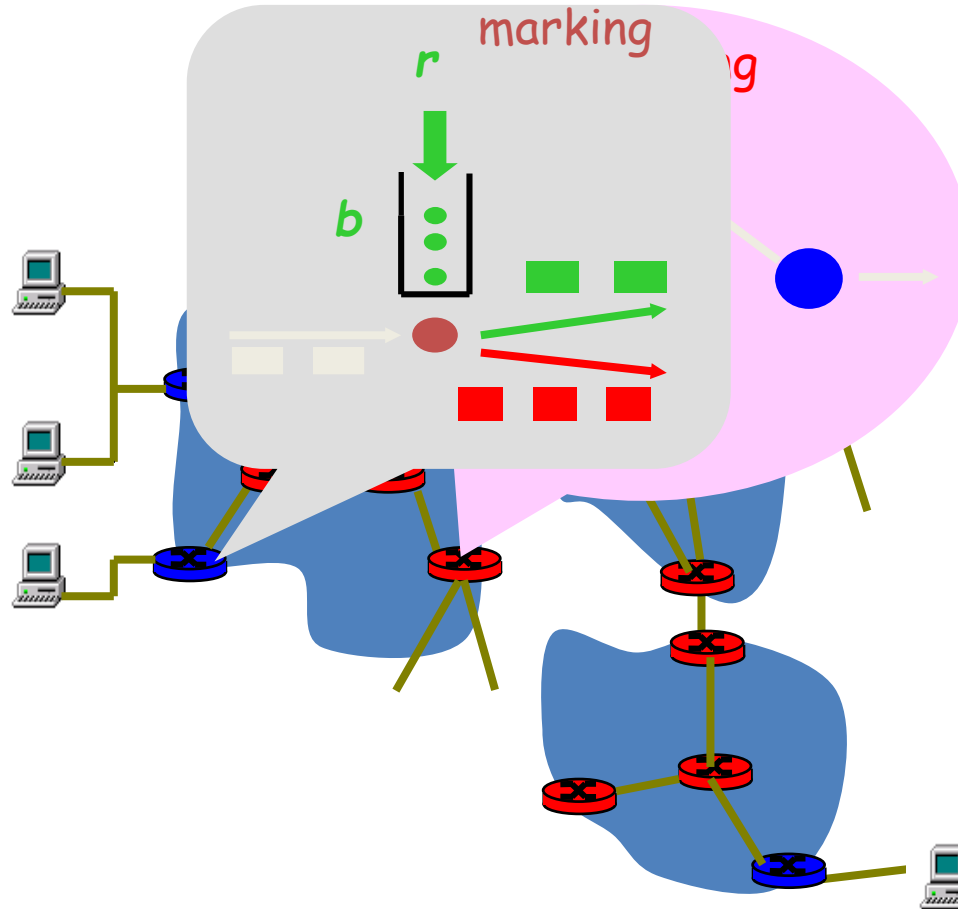
- marks packets as **in-profile** and **out-profile**

Core router: scheduling



□ per class traffic scheduling

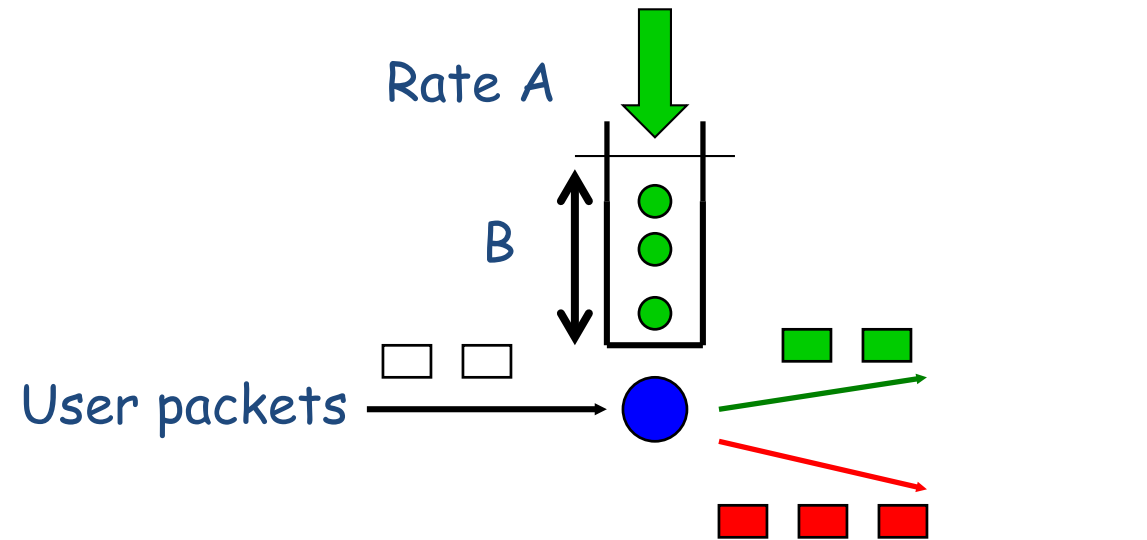
- based on **marking** at edge
- preference given to **in-profile** packets



Diffserv approach: provide functional components to build **service classes**

- **Network core:** stateless, simple
- Combine into **aggregated flows**, classification, shaping, admission: @ network edge

Edge-router Packet Marking



-Class-based marking: packets of different classes marked differently

Profile within class: pre-negotiated rate A, bucket size B

Packet is marked in the Type of Service (TOS) in IPv4, and Traffic Class in IPv6

DiffServ Core Functions

Forwarding: according to “Per-Hop-Behavior” (PHB) strictly based on classification marking

- PHB **does not** specify mechanisms to ensure required PHB performance
- E.g.:
 - Class A gets x% of bandwidth over time intervals of a specified length
 - Class A packets leave before class B packets

- **Advantage:**

No state info to be maintained by routers

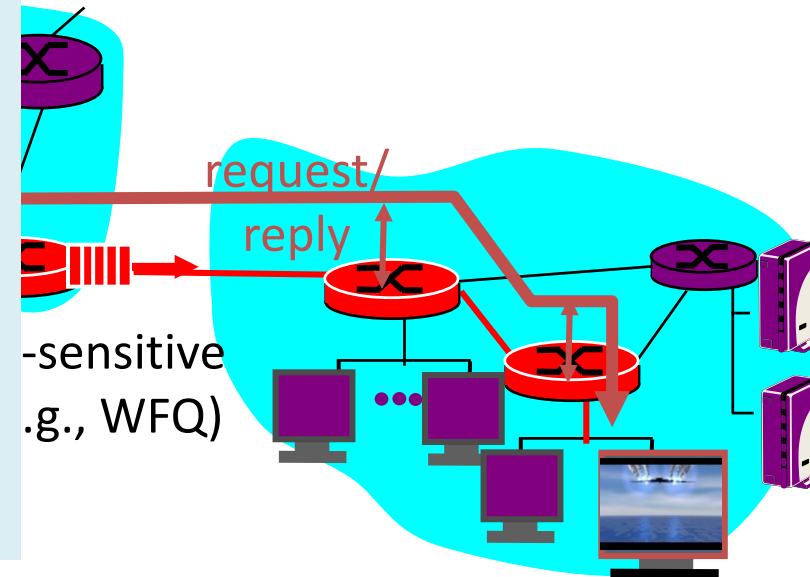
Intserv: QoS guarantee scenario

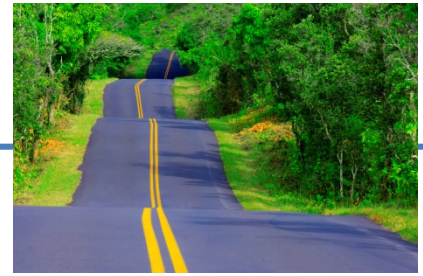
Resource reservation per individual application session
(admission control, continuous)

• "setup, signaling (RSVP)

Maintains state a la VC (but soft state, ie times out)

- responsibility at the client to renew reservations





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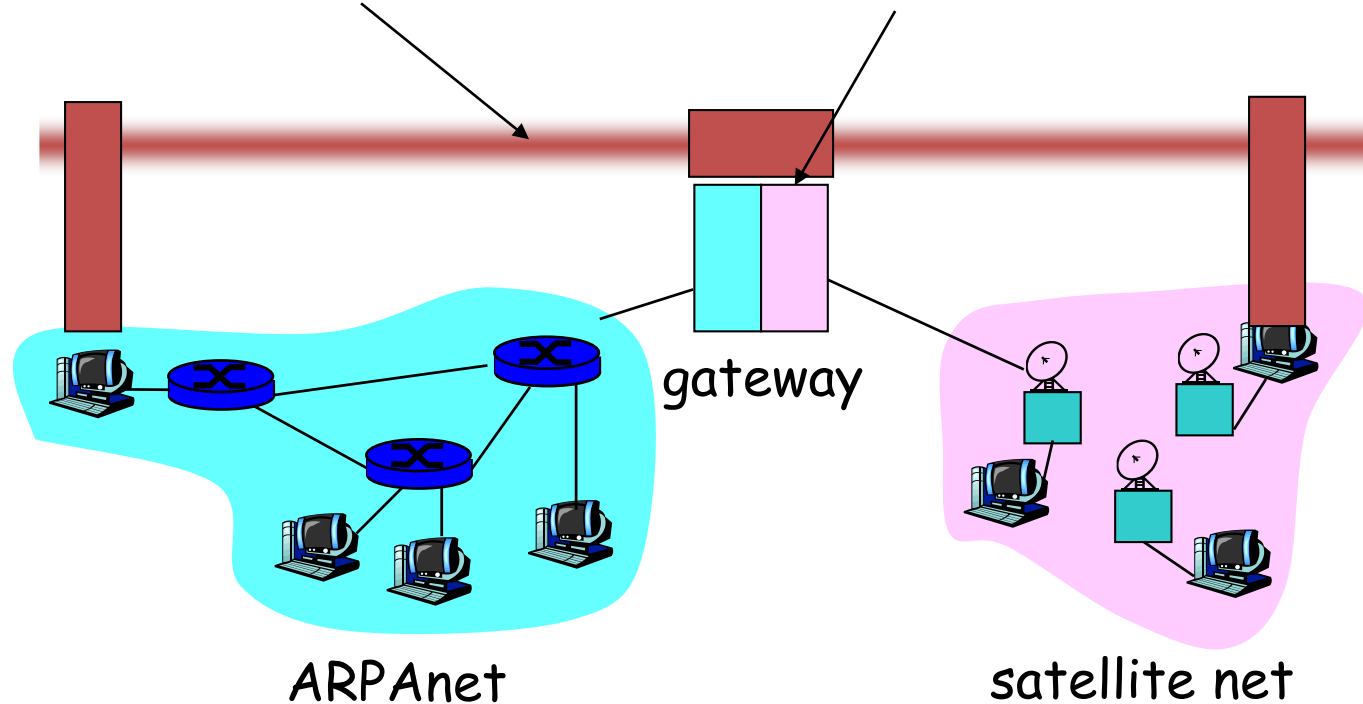
Recall the Internet approach : virtualizing networks

Internetwork layer (IP):

- addressing: internetwork appears as single, uniform entity, despite underlying local network heterogeneity
- network of networks

Gateway:

- “embed internetwork packets in local packet format”
- route (at internetwork level) to next gateway



What happened?

E.g. ATM: network or link layer?

Vision: end-to-end transport: “ATM from desktop to desktop”

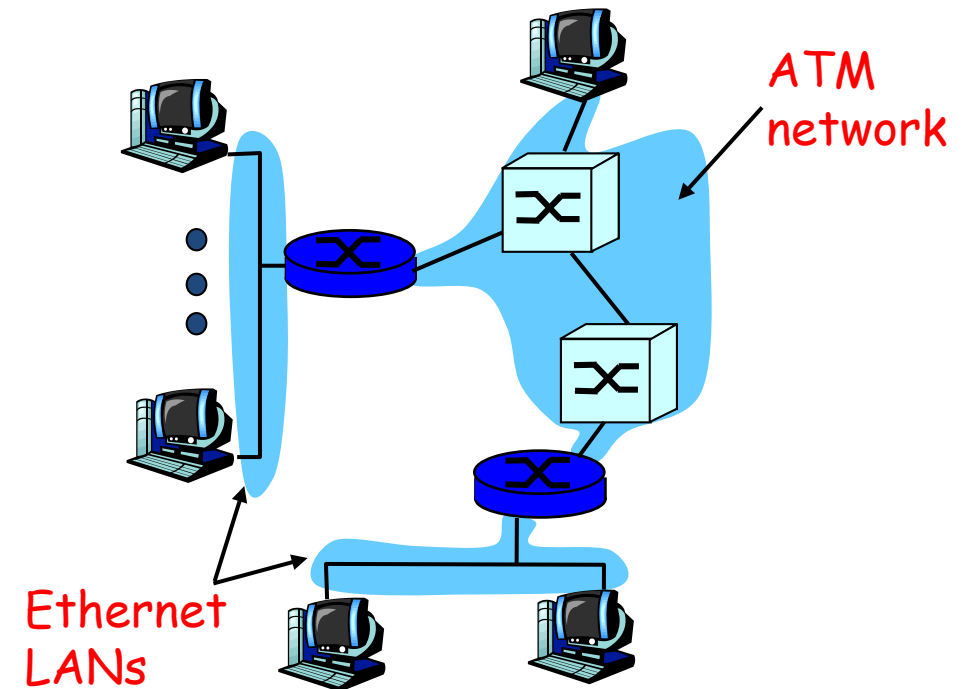
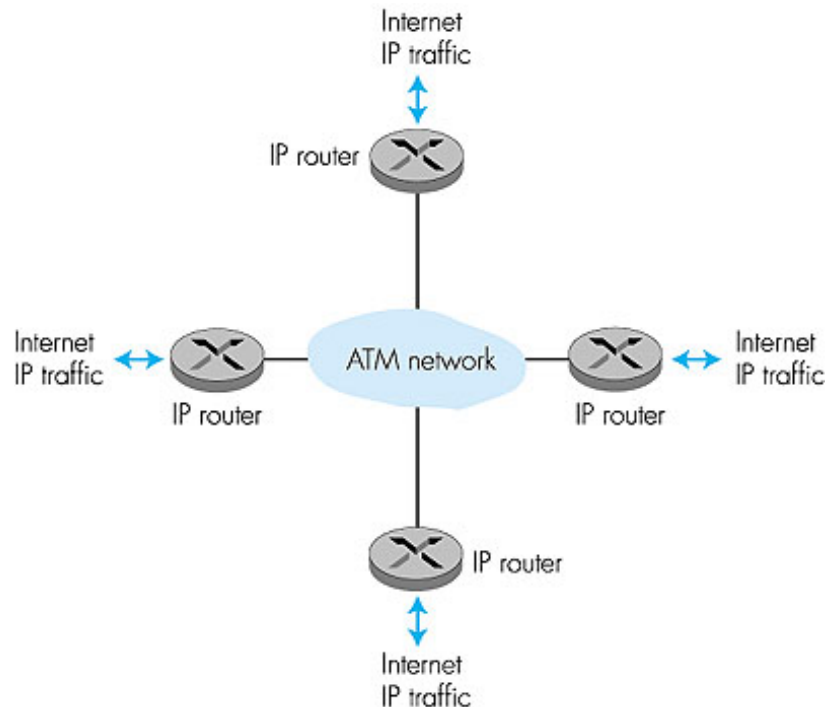
- ATM is a network technology

Reality:

- used to connect IP backbone routers

... or IP over ATM

- replace “network” (e.g., LAN segment) with ATM network, (ATM + IP addresses)
 - Run datagram routing on top of virtual-circuit routing

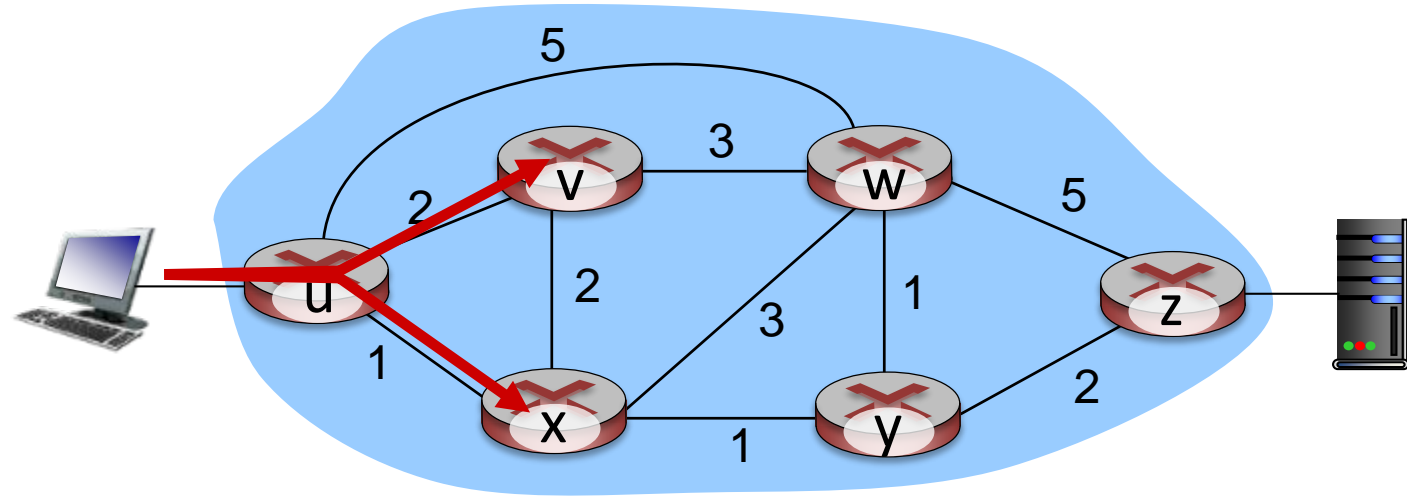


Cerf & Kahn's Internetwork Architecture

What is virtualized?

- two layers of addressing: internetwork and local network
 - new layer (IP) makes everything homogeneous at internetwork layer
 - underlying local network technology
 - Cable, satellite, 56K telephone modem
 - Ethernet, other LAN
 - ATM
 - **More recent: MPLS (Multiprotocol Label Switching Protocol): for traffic engineering**
- ... “invisible” at internetwork layer. Looks like a link layer technology to IP

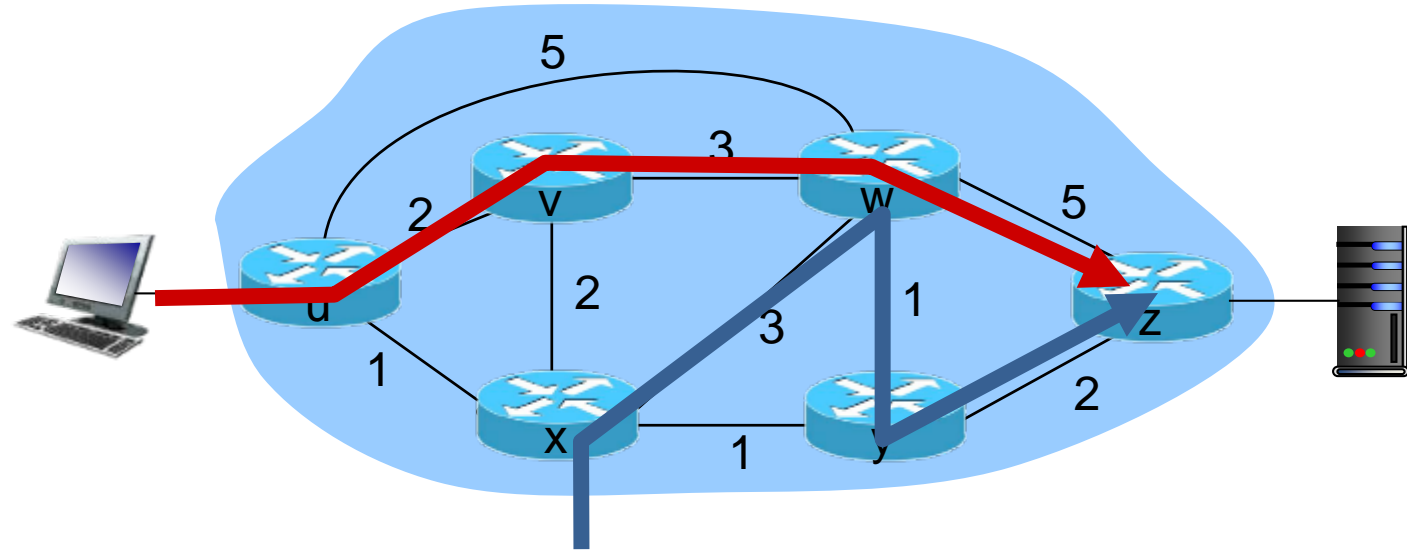
Traffic engineering: difficulties with traditional Internet routing



Q: what if network operator wants to split u-to-z traffic along uvwz *and* uxyz (load balancing)?

A: can't do it (or need a new routing approach...)

Traffic engineering: difficulties with traditional Internet routing

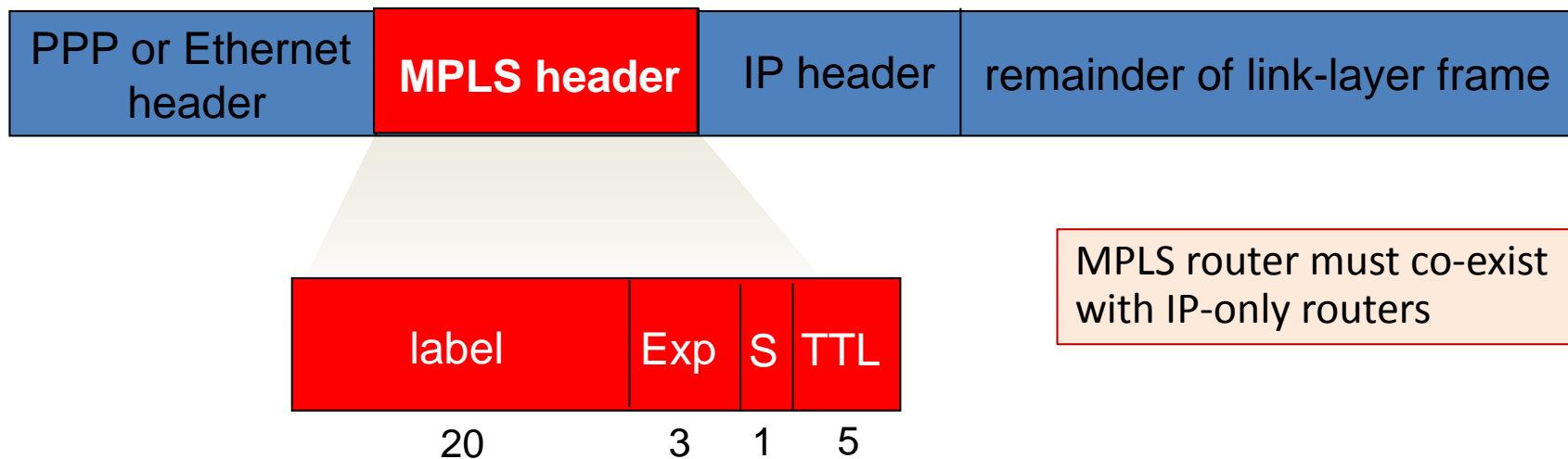


Q: what if w wants to route blue and red traffic differently?

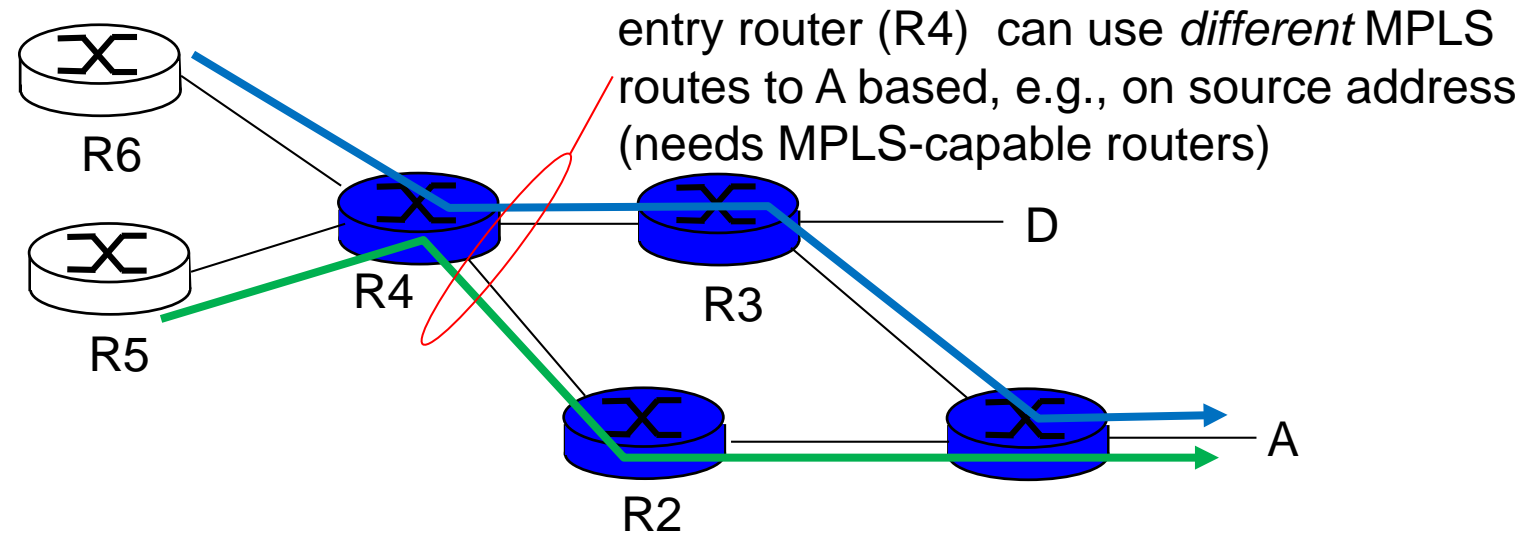
A: can't do it (with destination based forwarding, and LS, DV routing)

Multiprotocol label switching (MPLS) in IP networks: VC-inspired

- goal: utilize multiple S-T paths simultaneously
 - borrow ideas from Virtual Circuit (VC) approach but IP datagram still keeps IP address
- **label-switched router**
 - forwards packets to outgoing interface based only on label value (don't inspect IP address)
 - MPLS protocol's forwarding table distinct from IP forwarding tables



MPLS versus IP paths



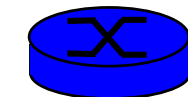
IP routing: path to destination determined by destination address alone

MPLS routing: path can be based on source *and* dest. address

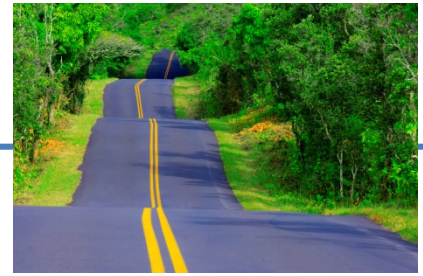
fast reroute: precompute backup routes in case of link failure or congestion (eg for CDN distribution)



IP-only router



MPLS and IP router

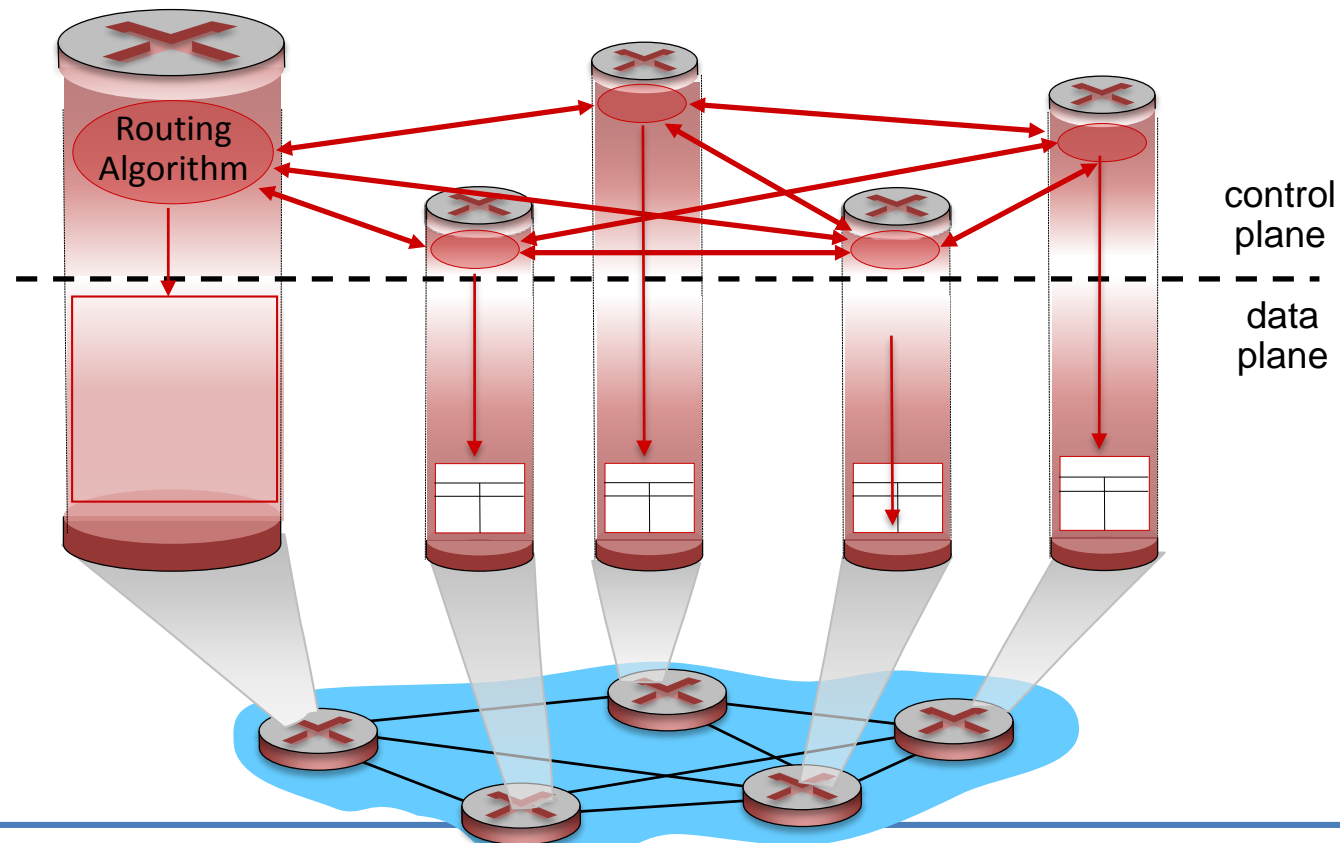


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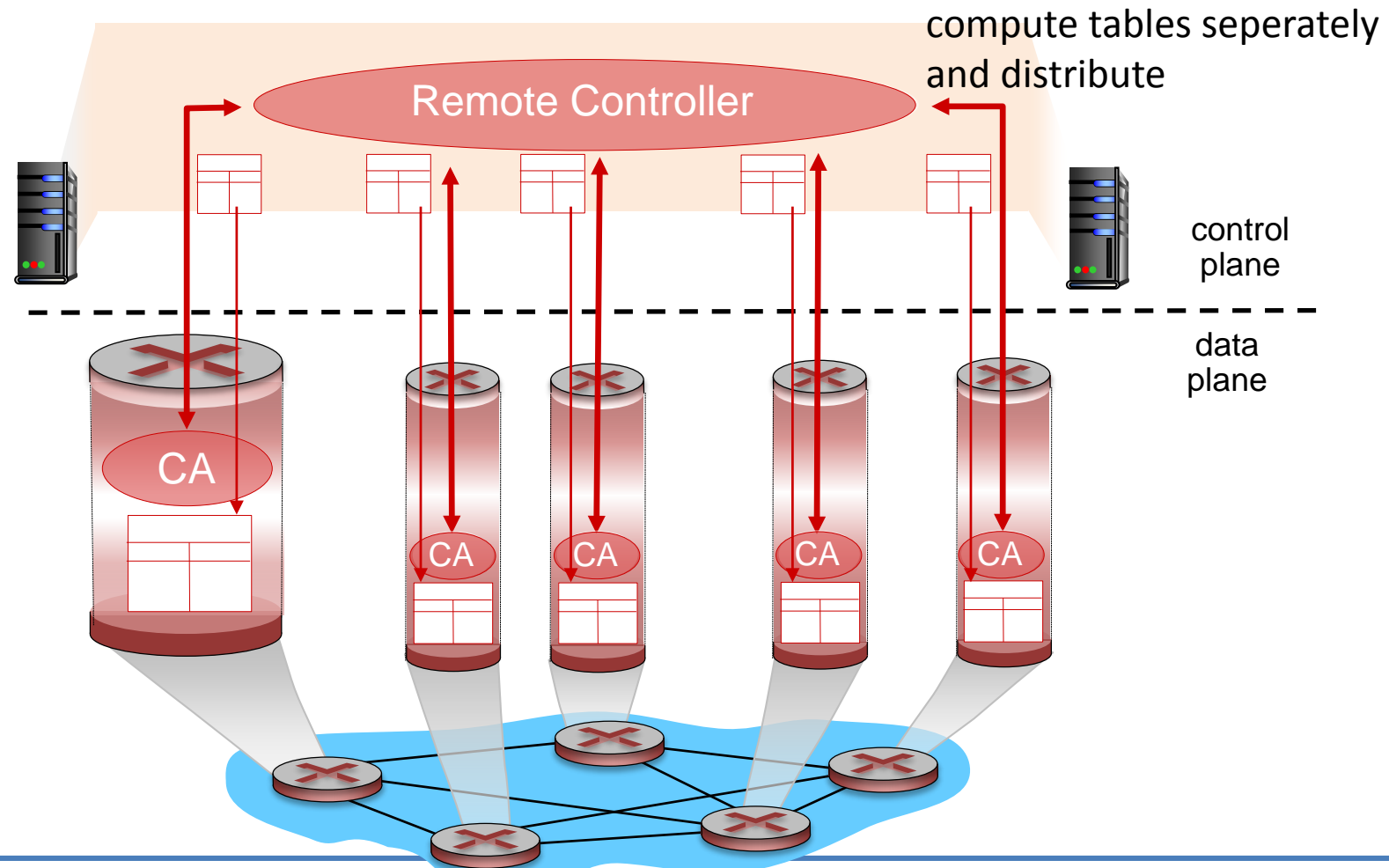
Recall: Traditional Internet, per-router control plane

Individual routing algorithm components *in each and every router* interact with each other in control plane to compute forwarding tables

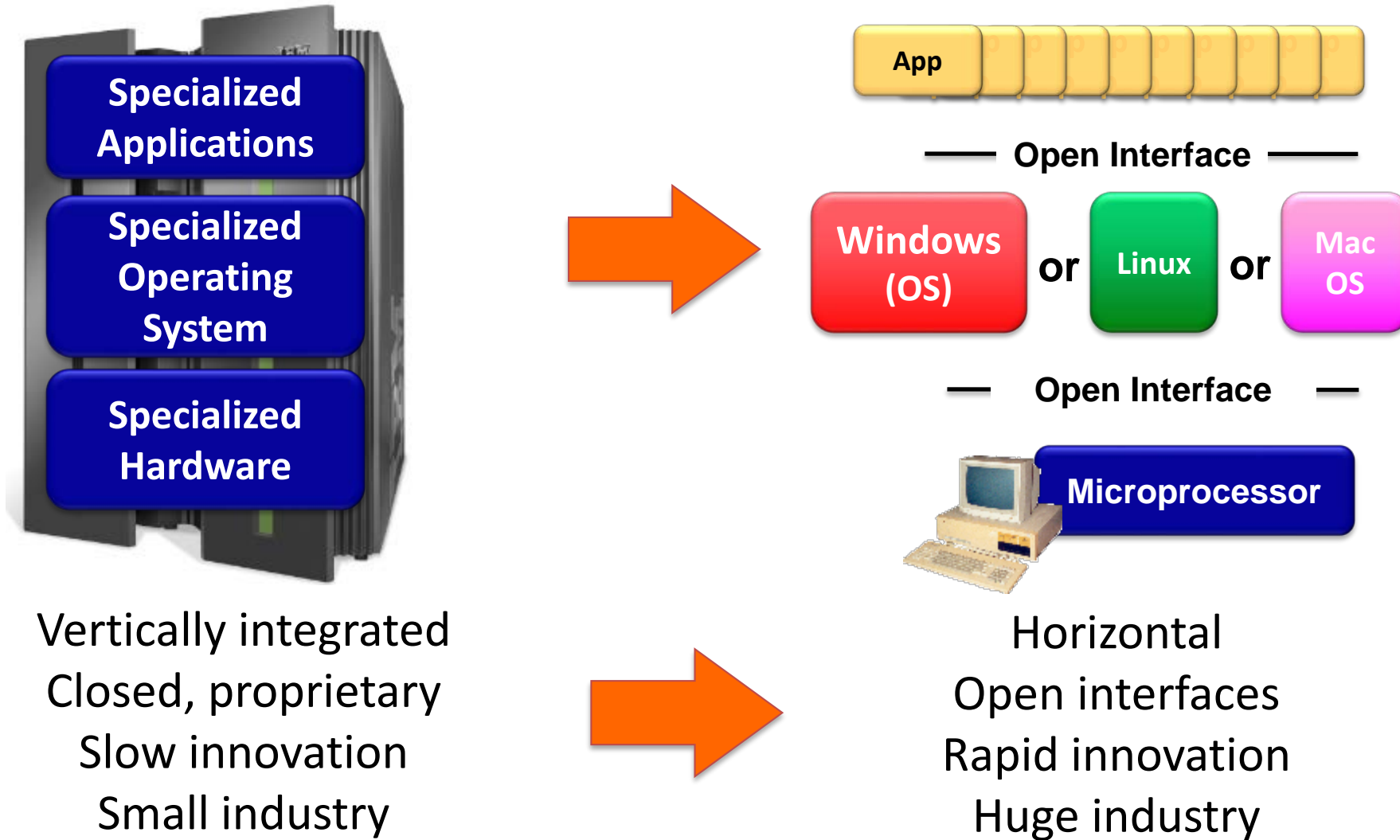


Recall: logically separated control plane

A distinct (typically remote) controller interacts with local **control** agents (CAs) in routers to compute forwarding tables

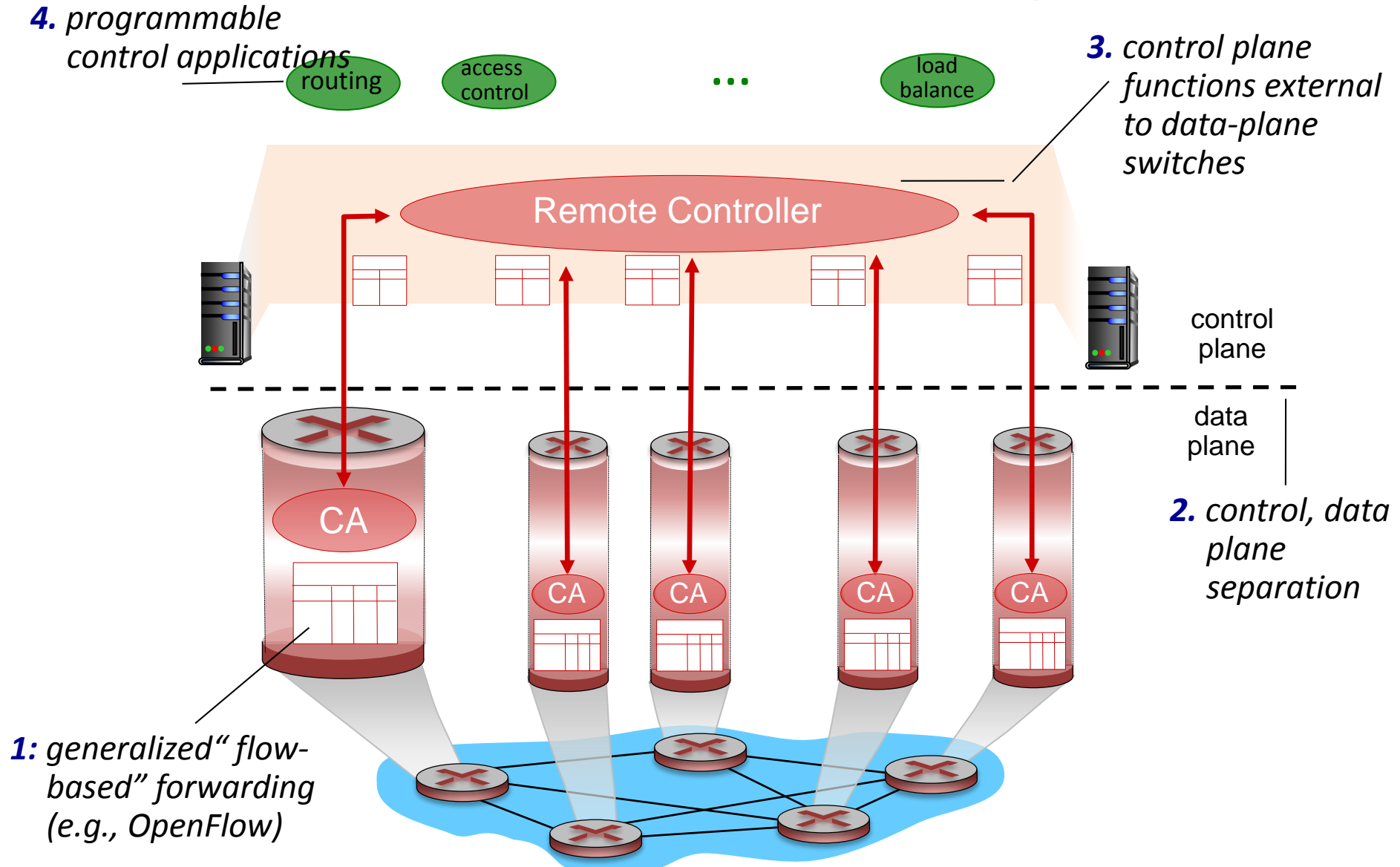


Analogy: mainframe to PC evolution*



* Slide courtesy: N. McKeown

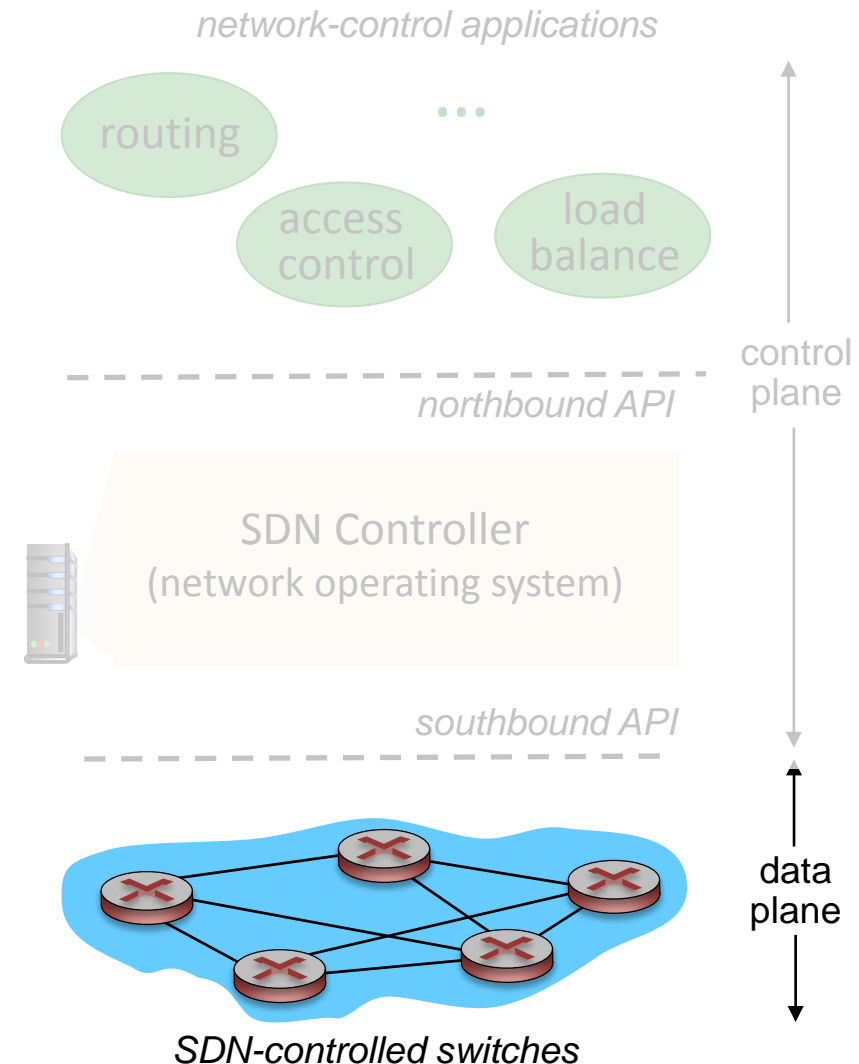
Software defined networking (SDN)



SDN perspective: data plane switches

Data plane switches

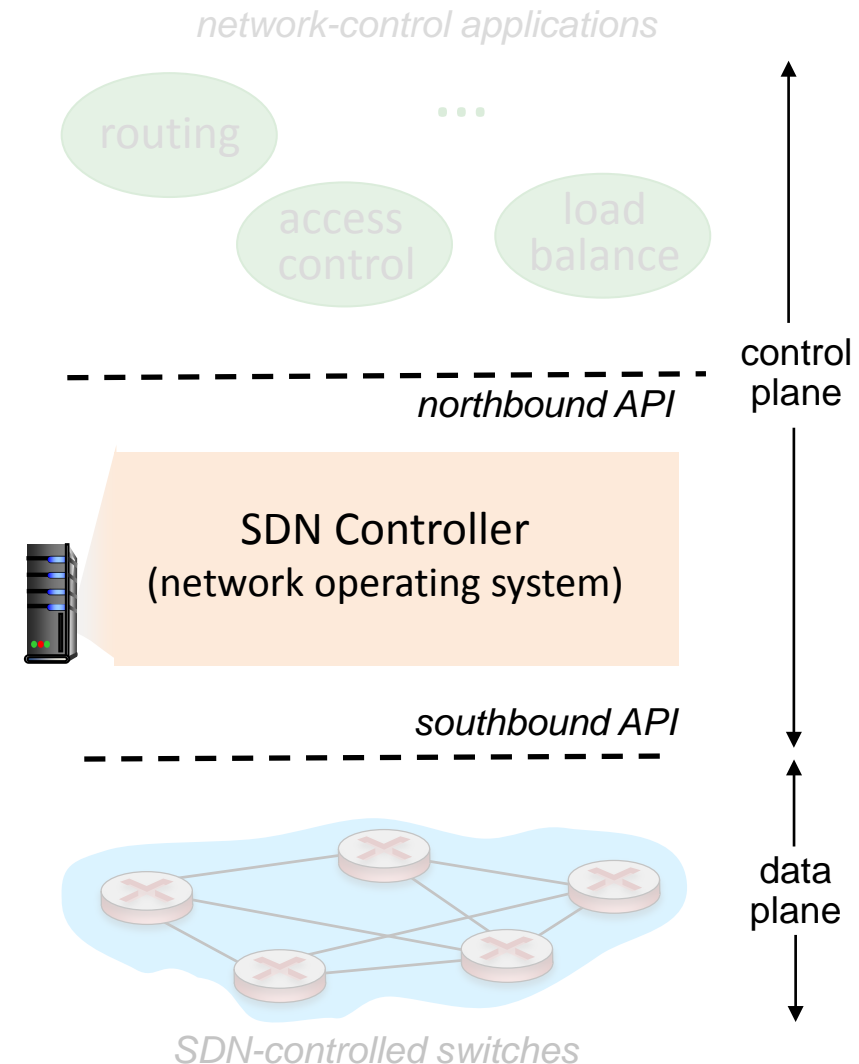
- fast, simple, for data-plane forwarding in H/W
- switch flow table: computed by controller
- API for table-based switch control (e.g., OpenFlow)
- protocol for communicating with controller (e.g., OpenFlow)



SDN perspective: SDN controller

SDN controller (network OS):

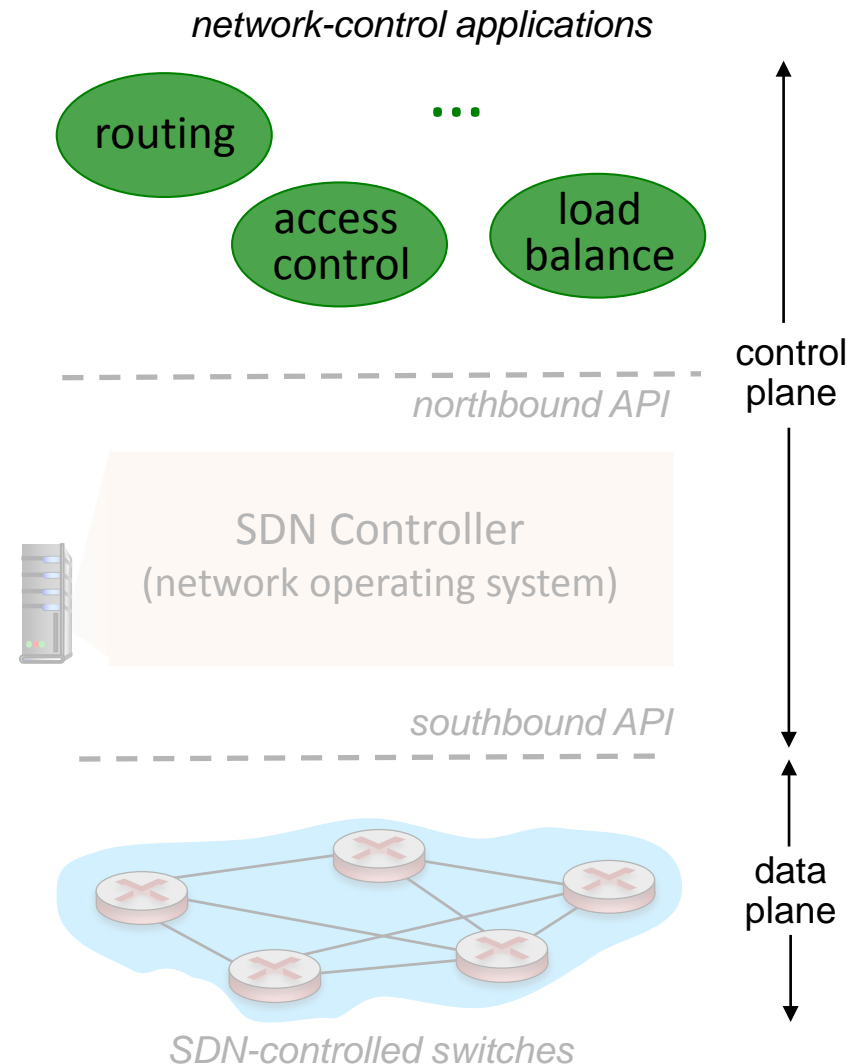
- maintains network state information
- interacts with network control applications “above” (northbound API)
- interacts with network switches “below” (southbound API)
- implemented as ***distributed system*** for performance, scalability, robustness



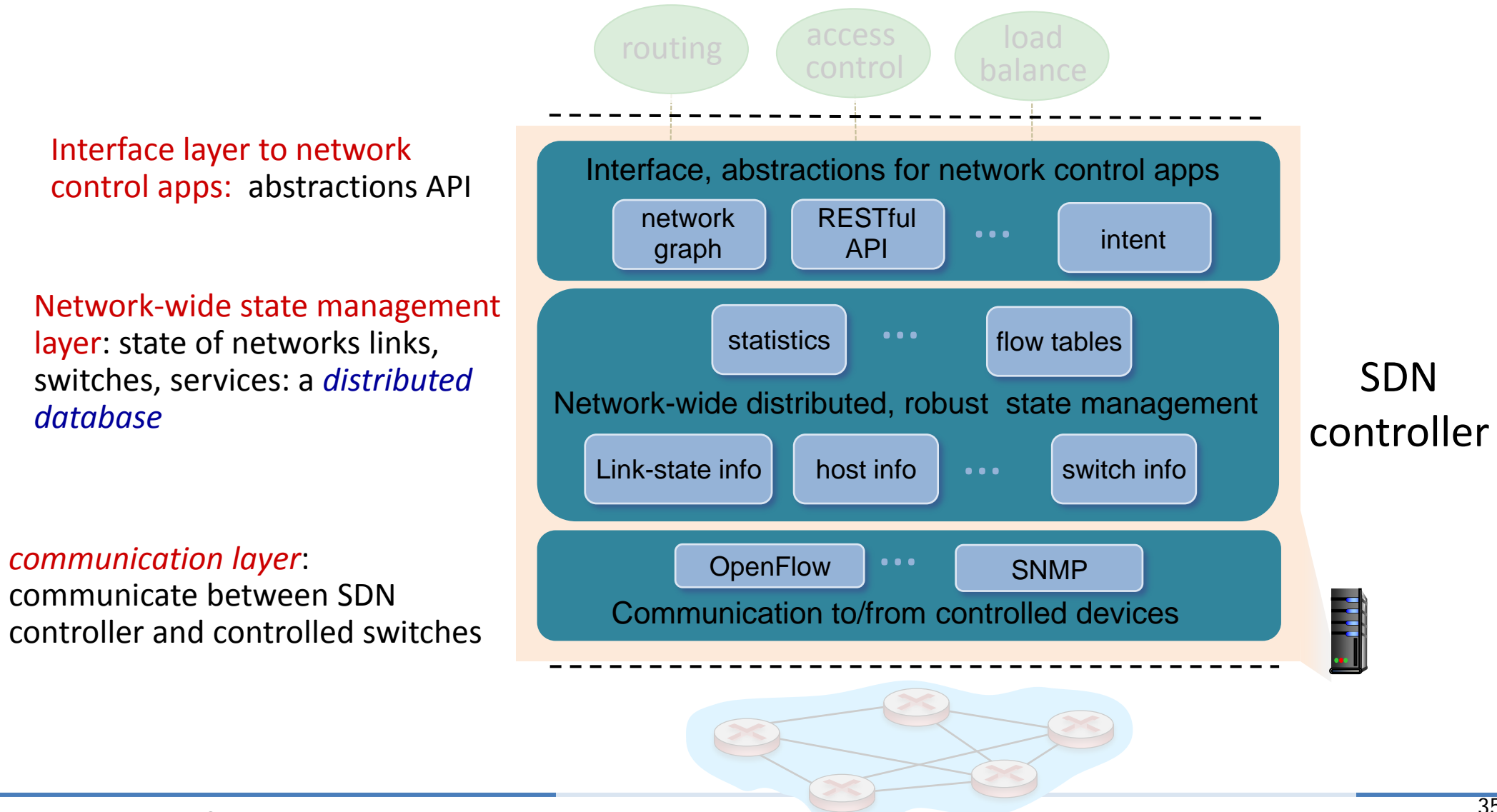
SDN perspective: control applications

network-control apps:

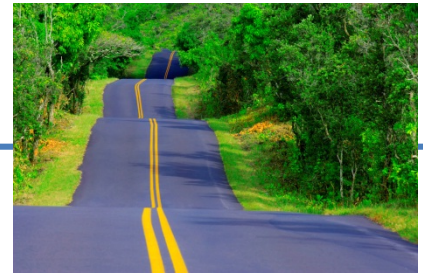
- “brains” of control: for control functions using lower-level services, API provided by SDN controller
- *unbundled*: can be provided by 3rd party: distinct from routing vendor, or SDN controller



Zooming in: components of SDN controller



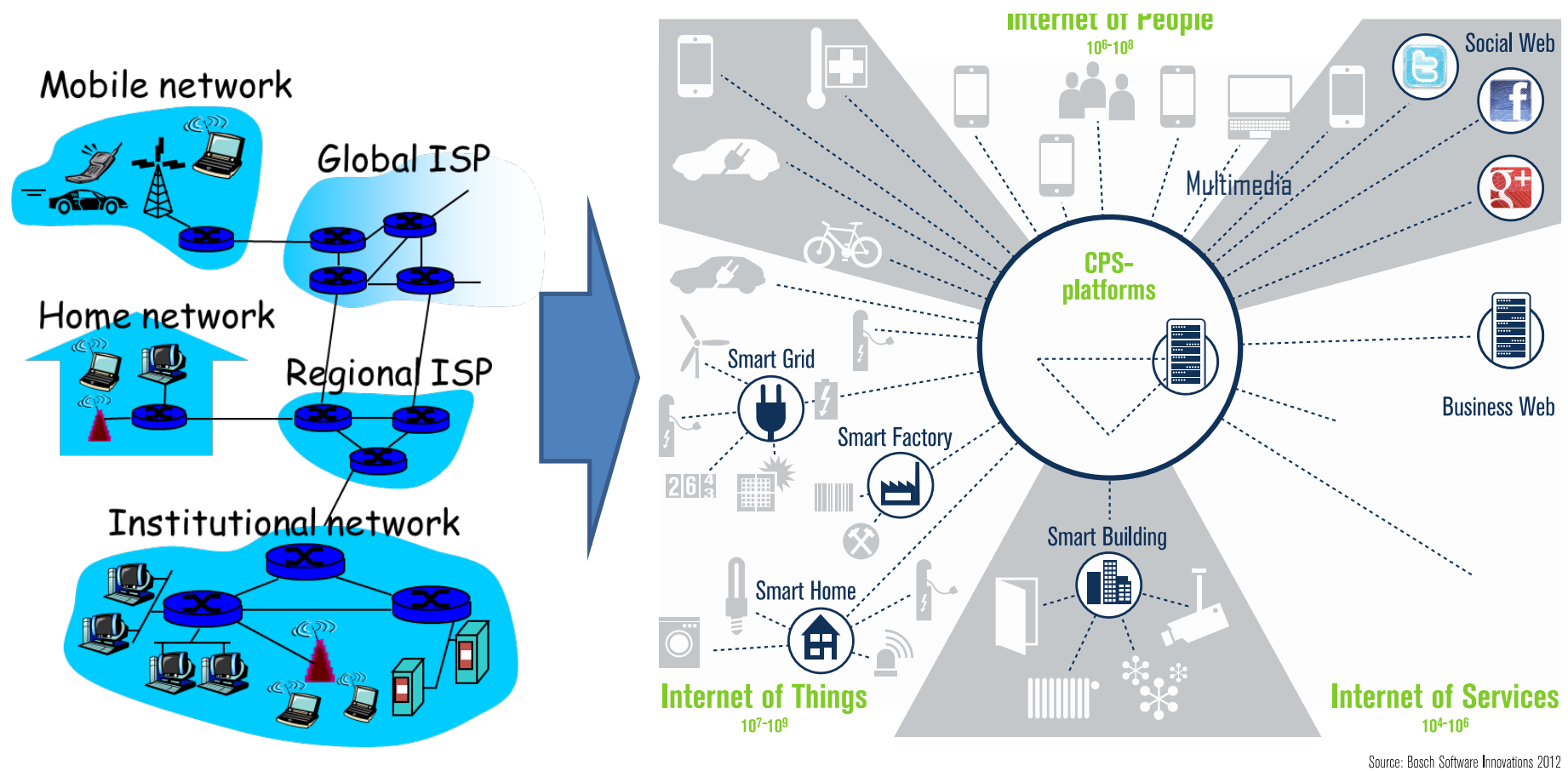
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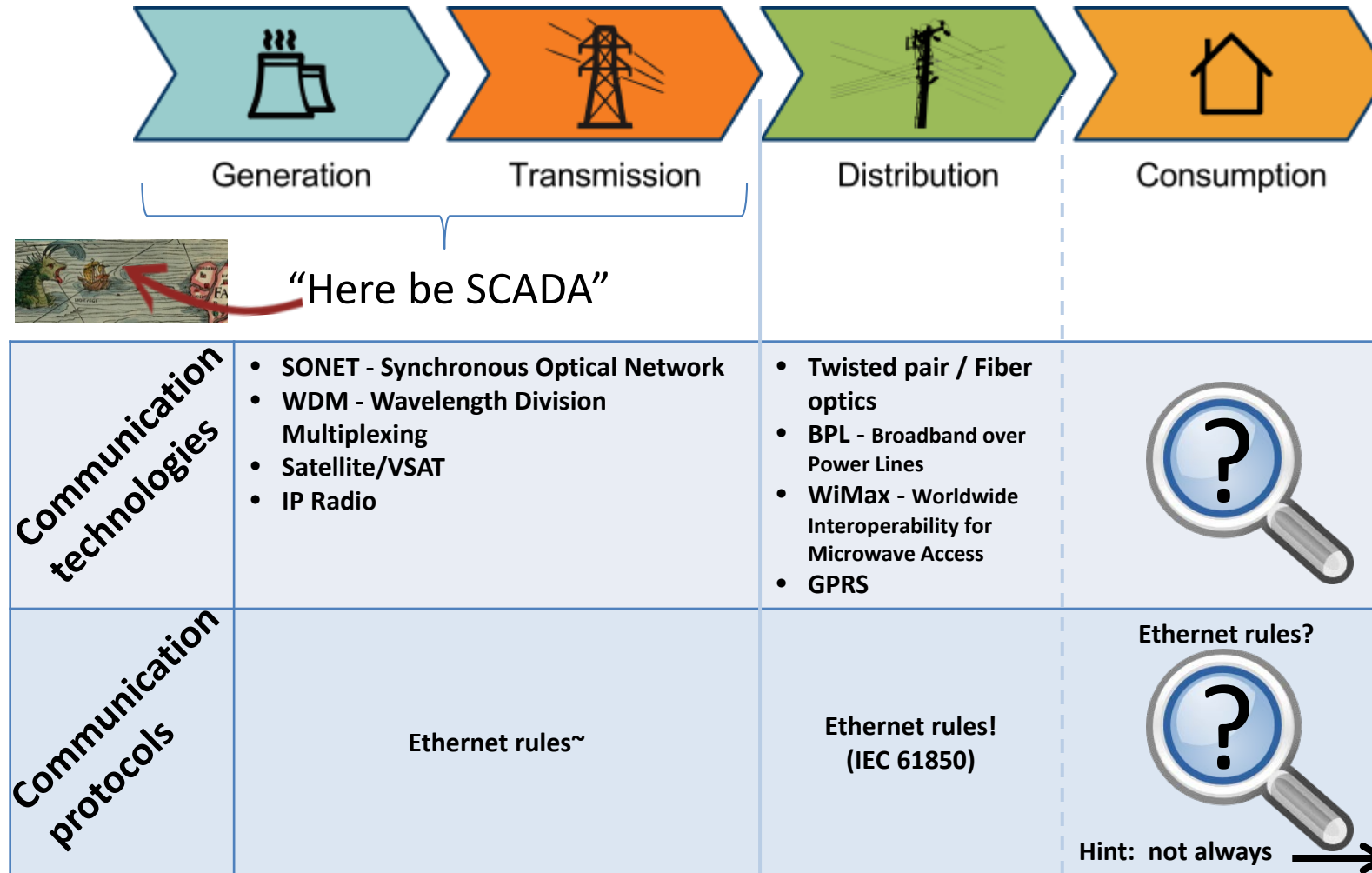
Recall: Internet & its context....



approx 10 yrs ago

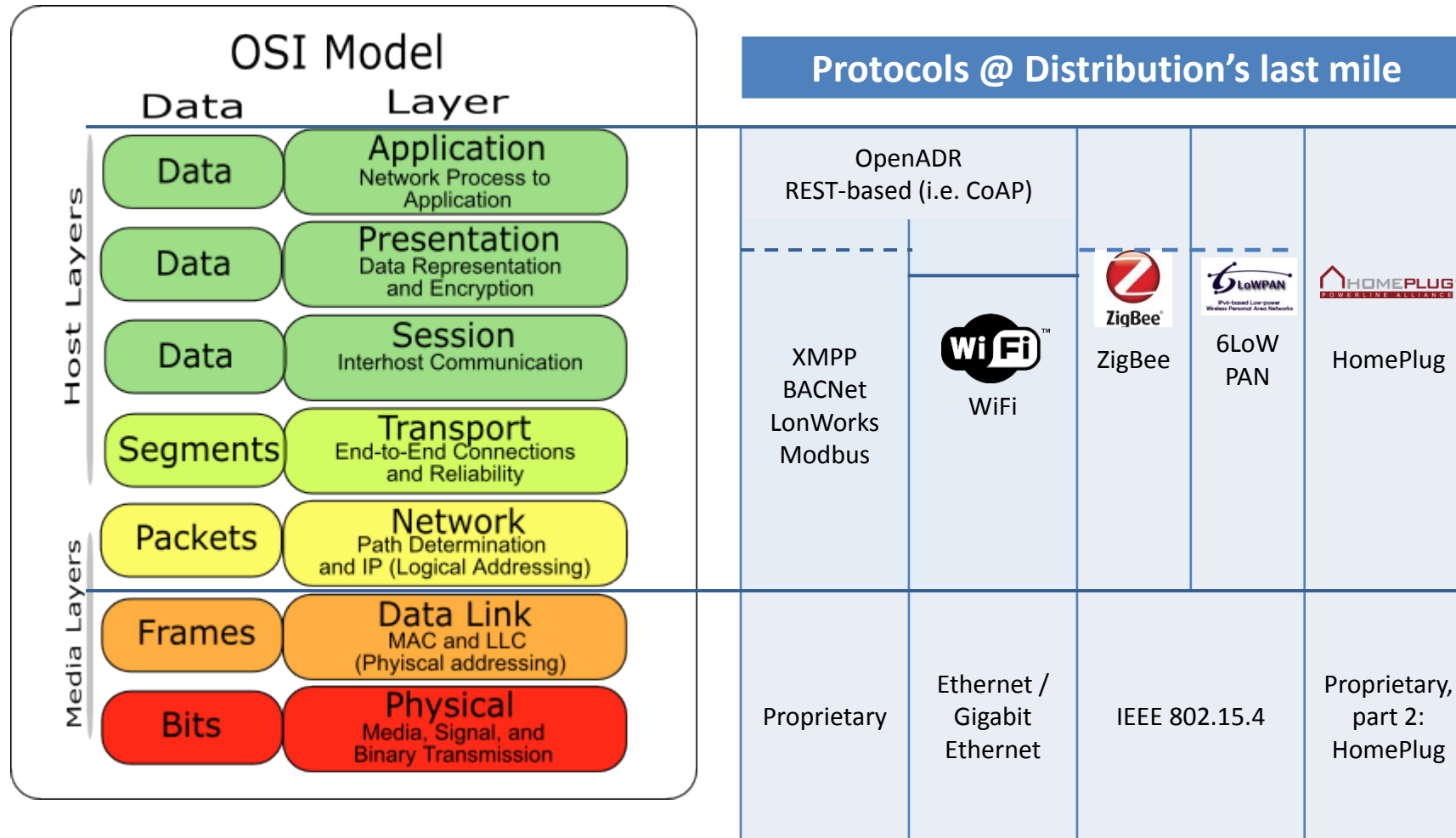
continuous evolution

Example: Data networking technologies in Smart Grids



Slides: Giorgos Georgiadis

Approximate overview of shaping new stacks

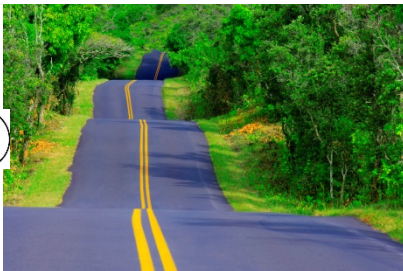


Slides: Giorgos Georgiadis
(see extra slides for more refs¬es)

Summary & Study list

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1. Internet core and transport protocols do not provide guarantees for multimedia streaming traffic
2. Applications/edge take matters into own hands
 - New, evolving methods; new proposals for transport protocols
3. Another type of service @ core (VC-like) would imply a different situation
 - Internet core is re-shaping, for long time ... (Intserv & Diffserv, Traffic engineering, SDN,)
4. **Internet-of-Things in evolution**
 - **even more types of traffic, new needs**

Review questions

bandwidth allocation and congestion

; mechanism and give examples

.S.

ane, the data plane and network

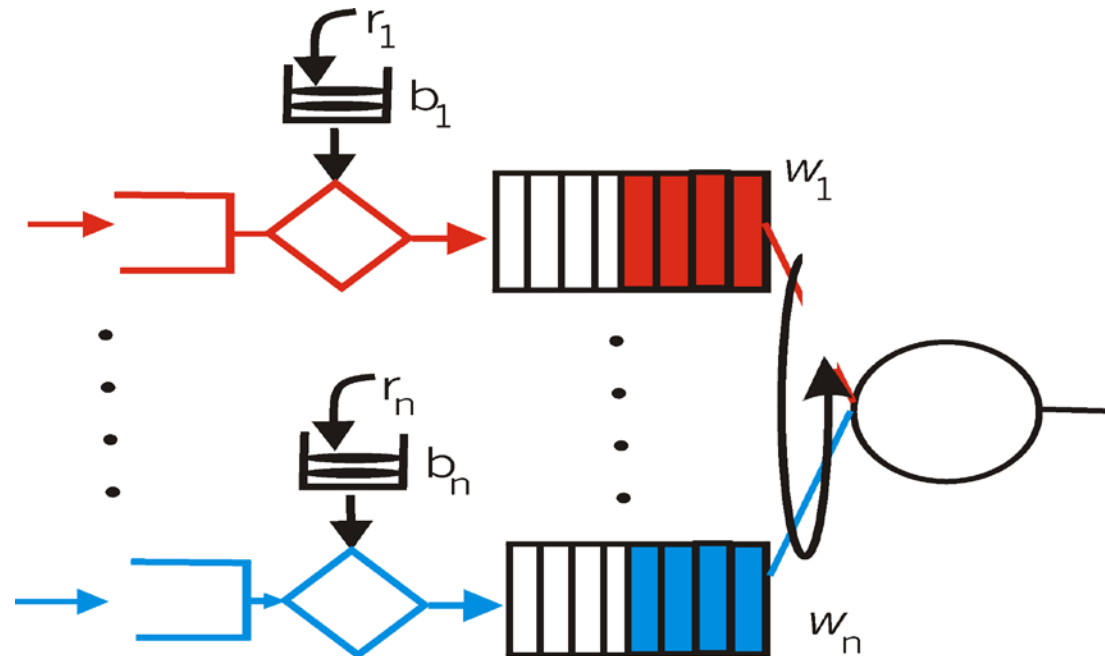
Extra slides/notes for further study

Token bucket + WFQ...

...can be combined to provide upper bound on packet delay in queue:

- b_i packets in queue, packets are serviced at a rate of at least $R \cdot w_i / \sum (w_j)$ packets per second, then the time until the last packet is transmitted is at most

$$b_i / (R \cdot w_i / \sum (w_j))$$

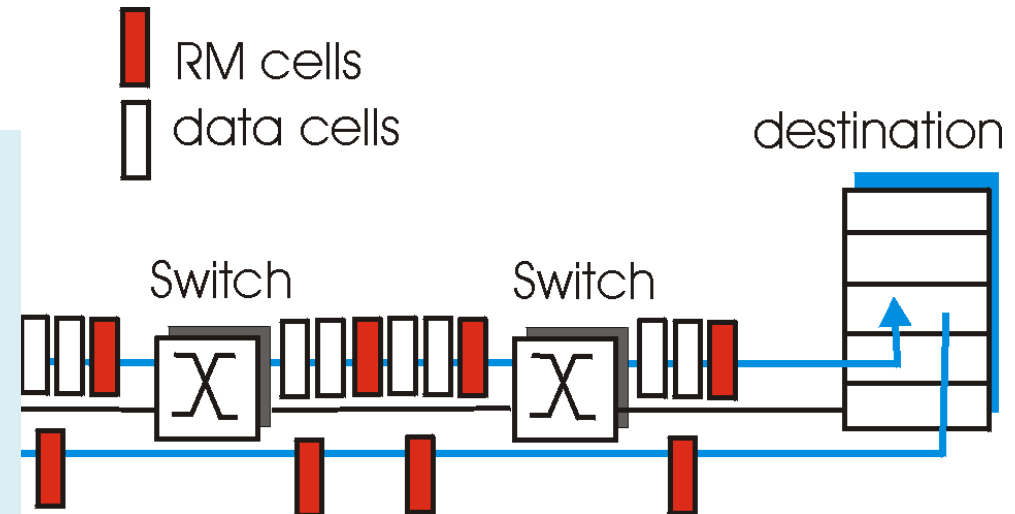


ATM ABR congestion control

ABR: available bit rate:

"elastic service"

source



work-assisted")

ngestion)

yte ER (explicit rate) field in RM cell

in cell

ortable rate on path

Traffic Shaping and Policing in ATM

Enforce the QoS parameters: check if *Peak Cell Rate (PCR)* and *Cell Delay Variation (CDVT)* are within the negotiated limits:

Generic Cell Rate Algo: introduce:
expected next time for a successive cell based on $T = 1/\text{PCR}$

border time $L (= \text{CDVT}) < T$ in which next transmission may start (but never before $T-L$)

A *nonconforming cell* may be discarded or its *Cell Loss Priority* bit be set, so it may be discarded in case of congestion

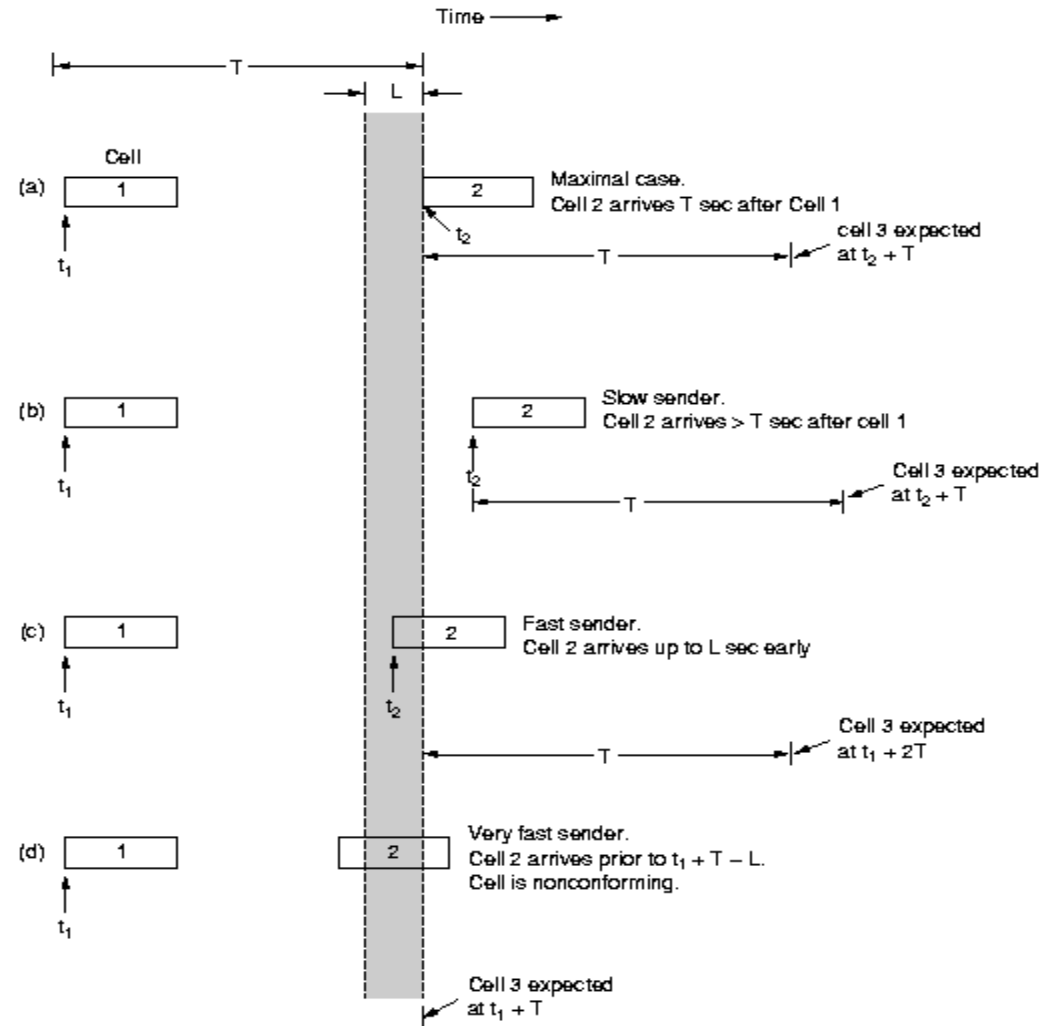


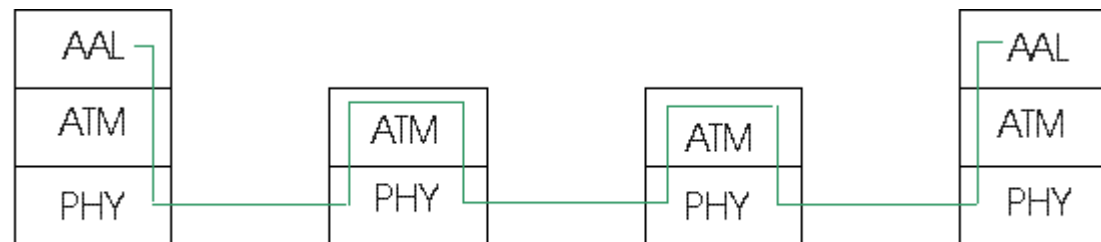
Fig. 5-73. The generic cell rate algorithm.

ATM Adaptation (Transport) Layer: AAL

Basic idea: cell-based VCs need to be "complemented" to be **supportive for applications**.

- r Several ATM Adaptation Layer (AALx) protocols defined, suitable for different classes of applications
 - r AAL1: for CBR (Constant Bit Rate) services, e.g. circuit emulation
 - r AAL2: for VBR (Variable Bit Rate) services, e.g., MPEG video
 - r

- "suitability" has not been very successful
- computer science community introduced AAL5, (simple, elementary protocol), to make the whole ATM stack usable as switching technology for data communication under IP!



Network support for multimedia

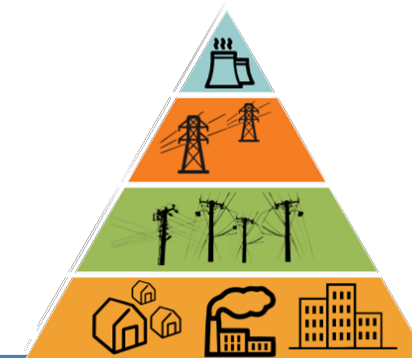
Approach	Granularity	Guarantee	Mechanisms	Complex	Deployed?
Making best of best effort service	All traffic treated equally	None or soft	No network support (all at application)	low	everywhere
Differentiated service	Traffic “class”	None of soft	Packet market, scheduling, policing.	med	some
Per-connection QoS	Per-connection flow	Soft or hard after flow admitted	Packet market, scheduling, policing, call admission	high	little to none

Software defined networking (SDN)

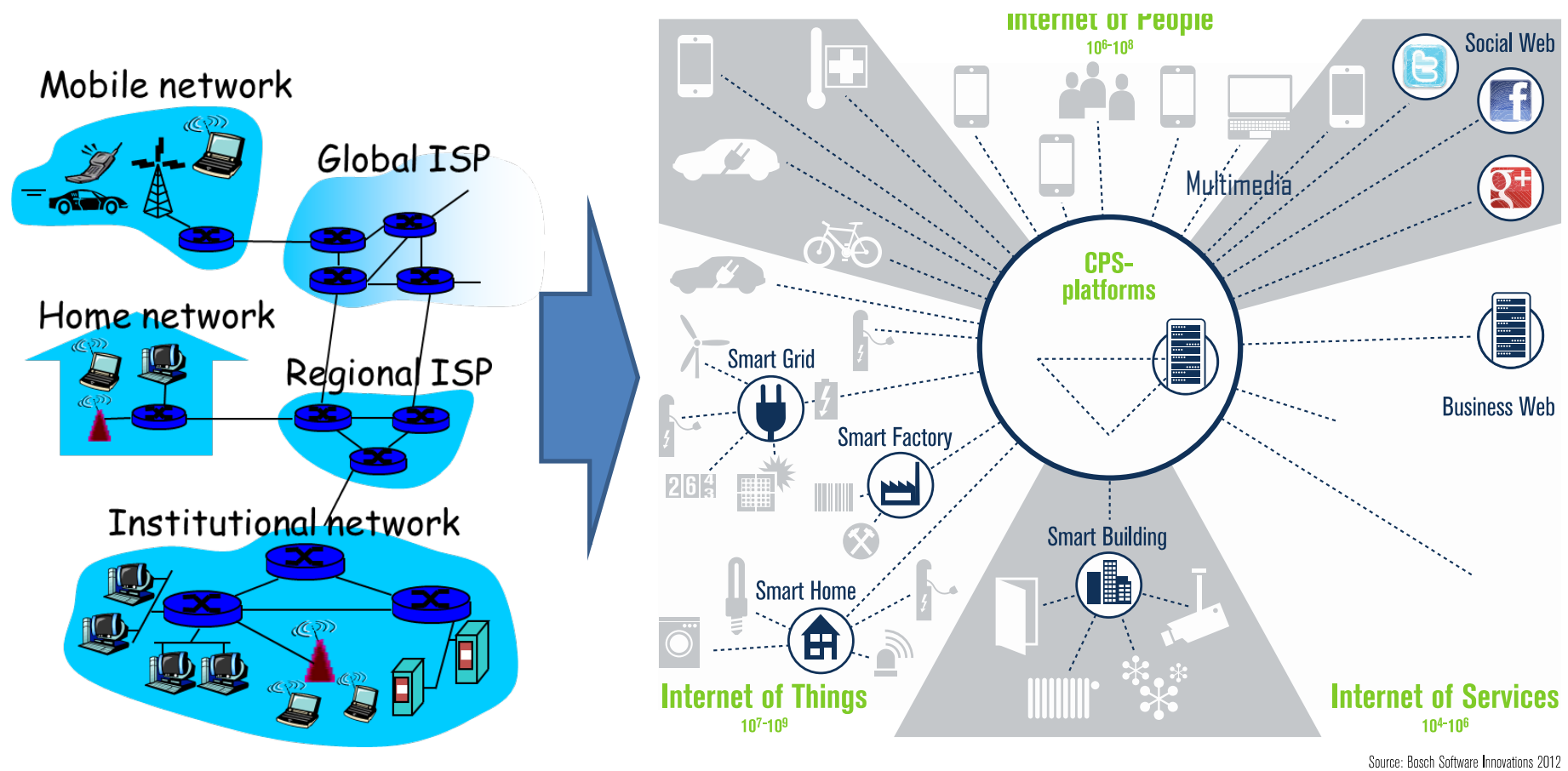
- Internet network layer: historically has been implemented via distributed, per-router approach
 - *monolithic* router contains switching hardware, runs proprietary implementation of Internet standard protocols (IP, RIP, IS-IS, OSPF, BGP) in proprietary router OS (e.g., Cisco IOS)
 - different “middleboxes” for different network layer functions: firewalls, load balancers, NAT boxes, ..
- ~2005: renewed interest in rethinking network control plane

Data networking technologies in Smart Grids

Presentation by
Giorgos Georgiadis
(former CTH / curr. Bosch R&D)



Recall: Internet & its context....

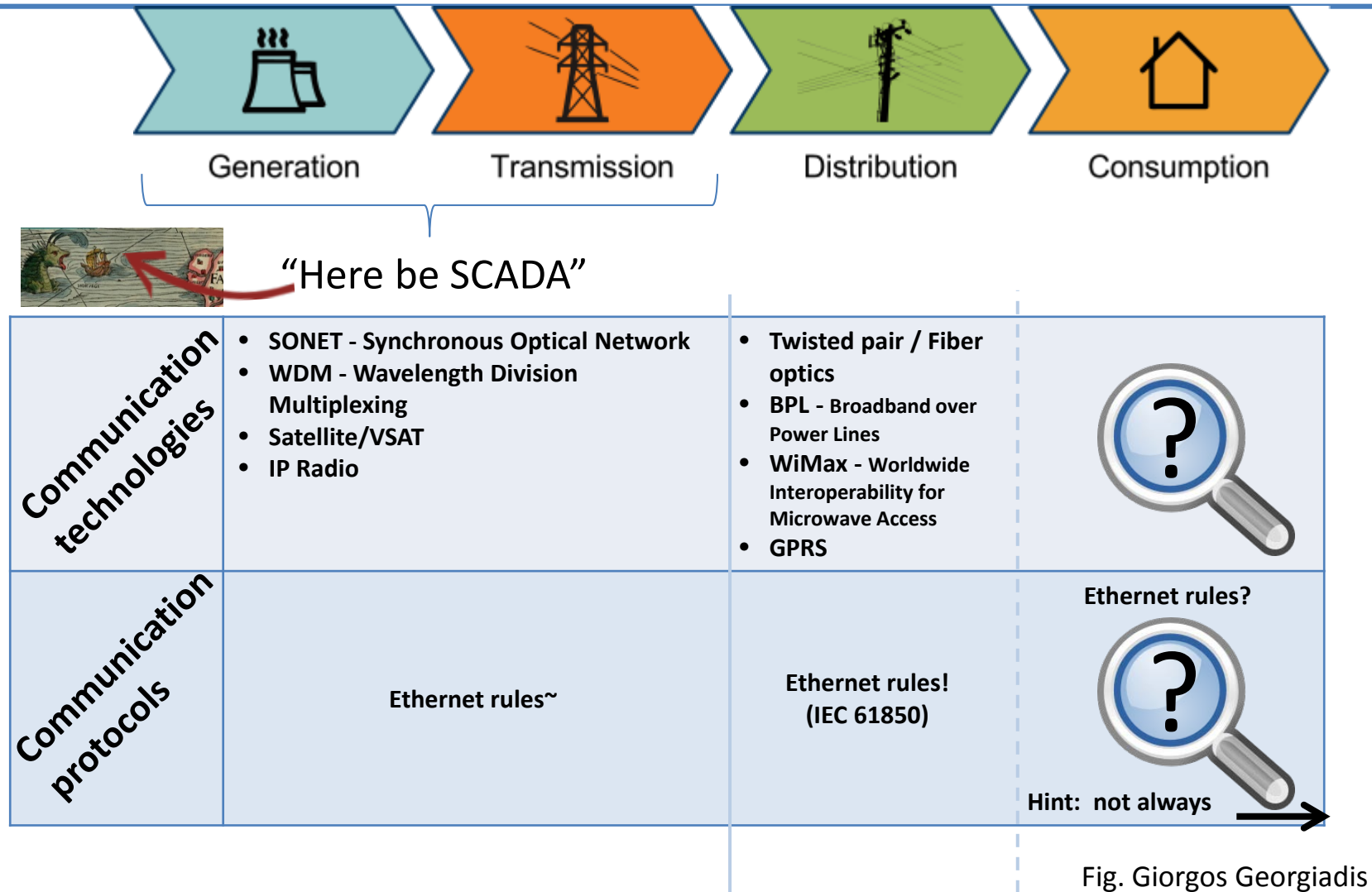


approx 10 yrs ago

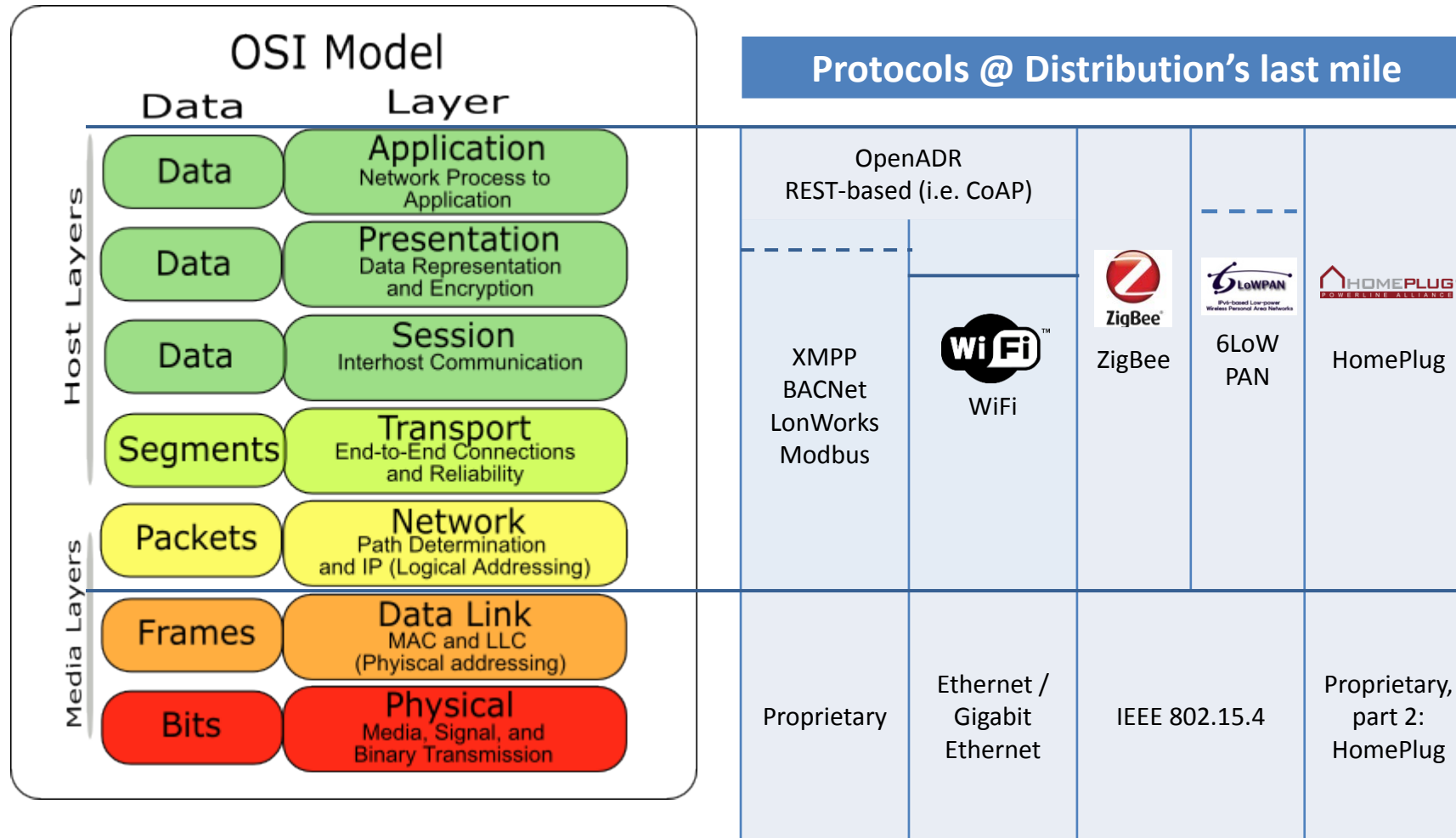
continuous evolution

Introduction

1

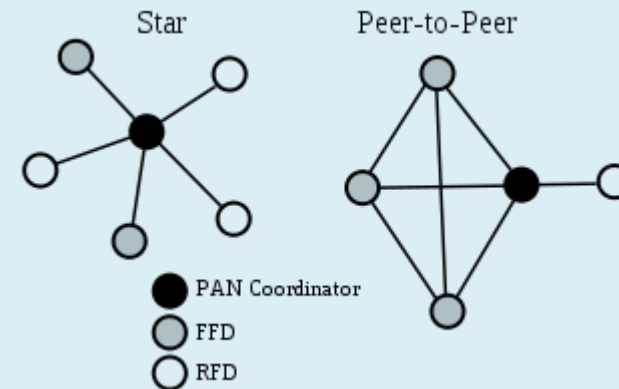


Approximate overview of shaping new stacks

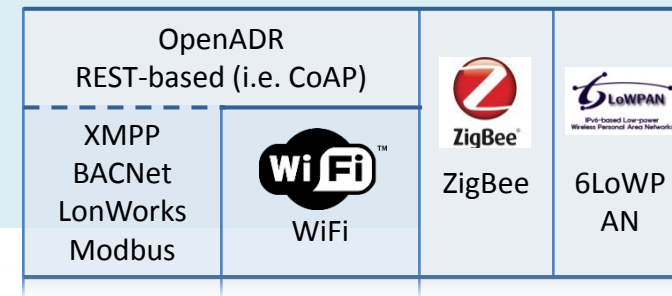


- Ethernet
 - Not much to say
- HomePlug
 - Honorable mention: popular home automation protocol
 - Powerline based
 - Speed: ~200mbps
 - Otherwise, vanilla protocol:
 - i.e. using TDMA,
 - Two kinds of nodes,
 - ...




- IEEE 802.15.4
 - Radio based, usually 2.4GHz
 - Small packets (≤ 127 bytes)
 - Medium speed (~ 250 kbps)
 - Originally DSSS
 - Topologies supported:
 - Star
 - Peer-to-peer
 - Roles supported:
 - Full-function device
 - Reduced-function device



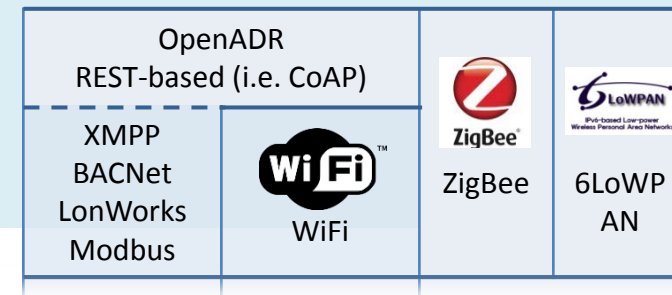
- 6LoWPAN
 - “IPv6 over LoW Power wireless Area Networks”
 - Builds on 802.15.4, IPv6
 - Aimed at low power devices (sensors, controllers)
 - Topologies
 - Star, peer-to-peer + **Mesh**
 - Many Challenges:
 - IP packets ≥ 1280 bytes (!)
 - 128bit IP addresses
 - ...



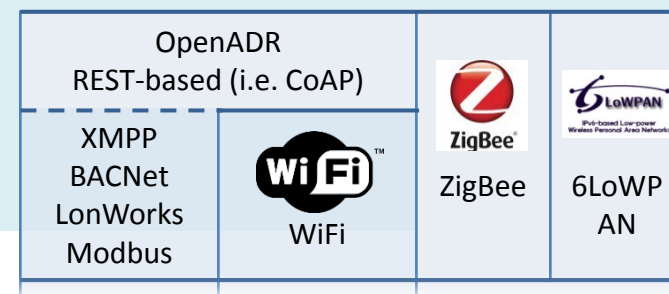
- ZigBee
 - Builds on 802.15.4, **but not IP**
 - Aimed at low power devices too (sensors, controllers)
 - Speed 250kbps
 - Packet 127bytes
 - Battery powered devices (supports sleep)
 - Topologies supported
 - + Mesh (jump to: example)
 - Roles supported
 - Coordinator, router, end node
 - Different profiles exist:
 - ZigBee Home Automation
 - Zigbee Smart Energy
 - **Zigbee IP, ...**

OpenADR REST-based (i.e. CoAP)		 ZigBee	 6LoWPAN AN
XMPP BACNet LonWorks Modbus	 WiFi		

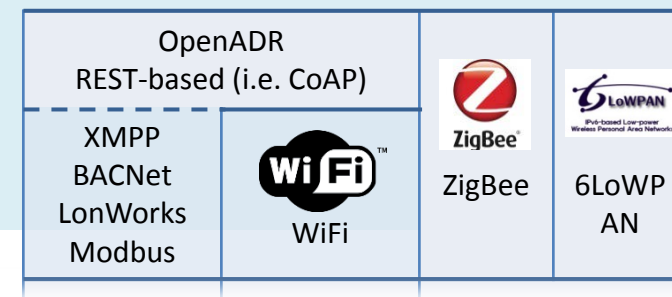
- More protocols, same story:
 - XMPP, BACNet, LonWorks, Modbus, ...
 - Wired
 - Proprietary, build around specific companies (BACNet, LonWorks) or legacy protocols (Modbus)
 - Today gateway devices to “break out” to Ethernet are in use
 - Simple topologies (i.e bus), same roles as before
- But what is the connecting thread over all?
 - Open standards!
 - Internet! (of Things?)



- OpenADR
 - ADR: Advanced Metering Response
 - Trying to ‘unify’ different solutions in a high level protocol
 - Formalizing:
 - Roles
 - Messages
 - Device detection
 - Simple topologies (i.e bus), same roles as before
- REST-based APIs
 - I.e. Costrained Application Protocol
 - Ultimately, HTTP-based
 - Verb oriented: GET, PUT, DELETE, ...



- Ethernet/IP-based integration
 - Remember:
 - Radio band: 2.4GHz (WiFi, ZigBee, 6LoWPAN)
 - Similar topologies, roles
 - Made for low energy devices, but flops/watt/kr increase!
 - Ethernet gateways commonly used
 - Solution: make them (formally) interoperable
 - ZigBee Smart Energy v2.0
 - ZigBee, WiFi, HomePlug on board
 - 6LoWPAN coming soon





- Ethernet + misc communication technologies
- Ethernet vs non-ethernet
 - Why?
 - Design for low energy devices (smaller packets, lower comm speed)
 - Peer to peer, mesh topologies
 - Now + Future?
 - Devices' specs catching up
 - Importance of being connected (to the Internet?)
 - Topologies still important (i.e. reliability)
 - Will probably remain radio-based