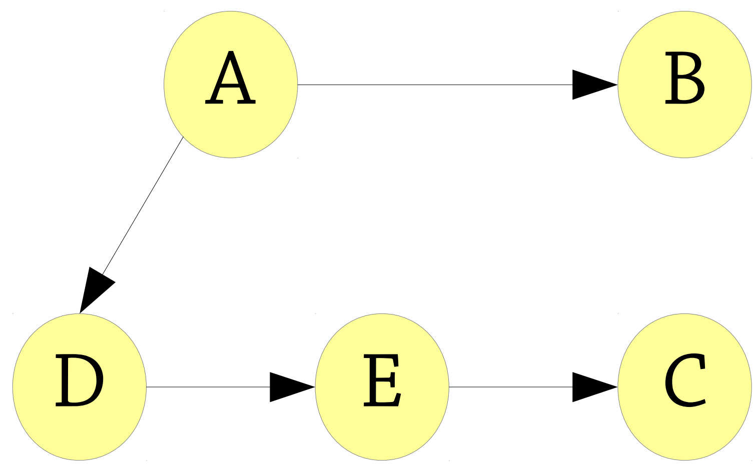
Graphs

Graphs

A graph is a data structure consisting of *nodes* (or vertices) and *edges*

- An edge is a connection between two nodes

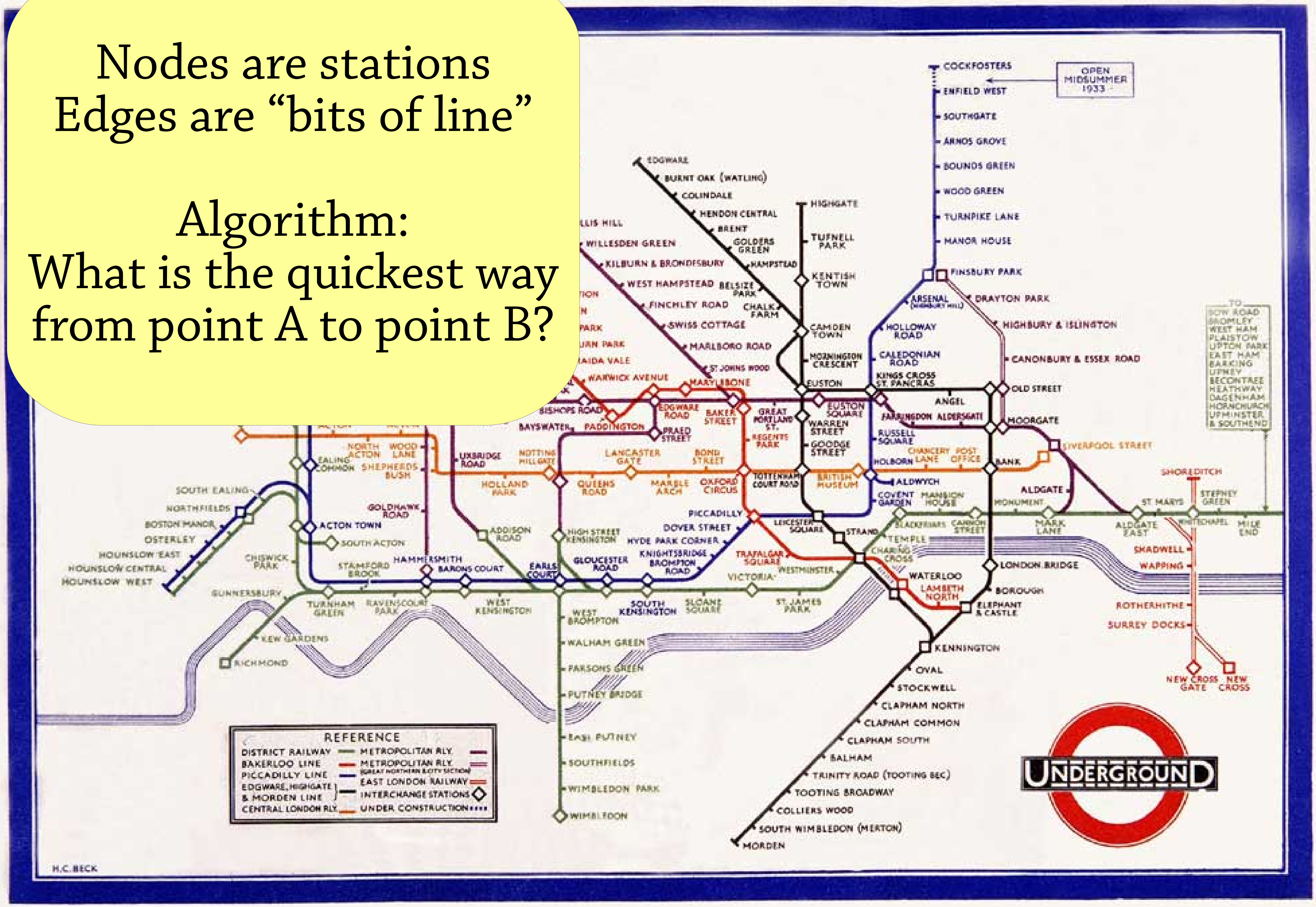


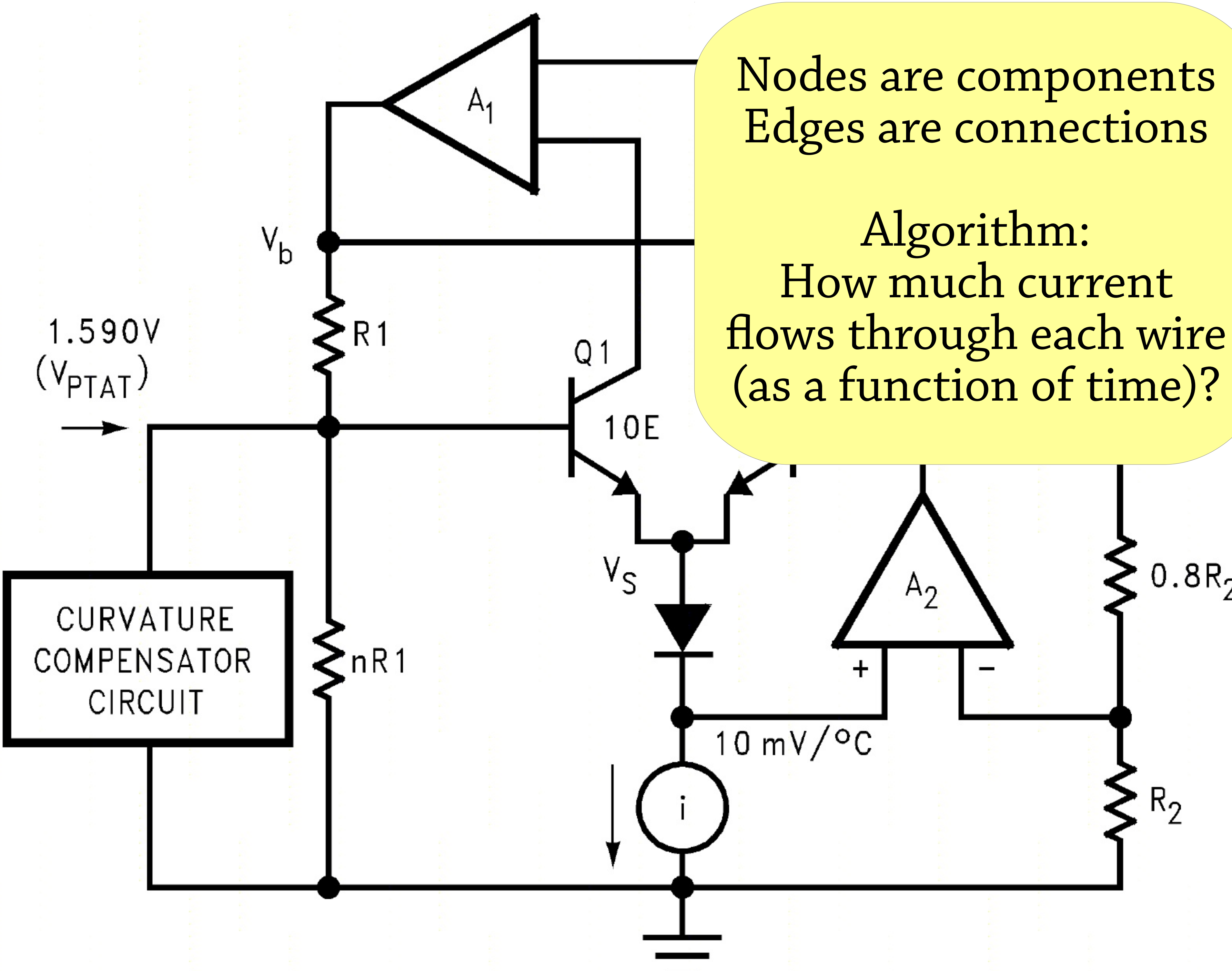
Nodes: A, B, C, D, E

Edges: (A, B), (A, D), (D, E), (E, C)

Nodes are stations
Edges are “bits of line”

Algorithm:
What is the quickest way from point A to point B?





Nodes are components
Edges are connections

Algorithm:
How much current flows through each wire (as a function of time)?

Graphs

Graphs are used all over the place:

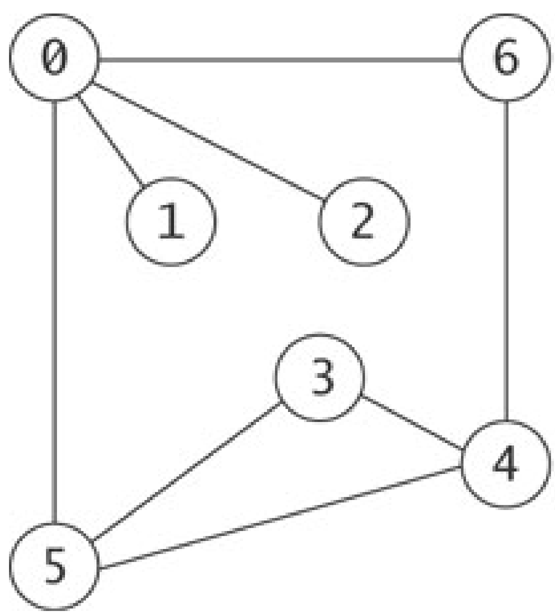
- communications networks
- many of the algorithms behind the internet
- maps, transport networks, route finding
- etc.

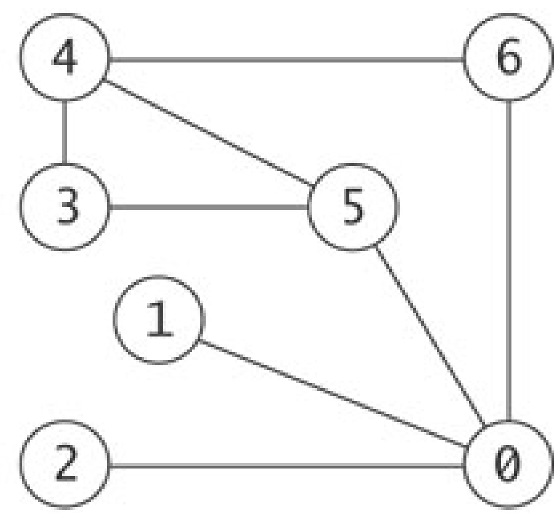
Anywhere where you have connections or relationships!

Normally the vertices and labels are *labelled* with relevant information!

Graphs

We only care what nodes and edges the graph has, not how it's drawn – these two are the *same graph*





V = {0, 1, 2, 3, 4, 5, 6}

E = {(0, 1), (0, 2), (0, 5), (0, 6), (3, 5), (3, 4), (4, 5), (4, 6)}

Graphs

Graphs can be *directed* or *undirected*

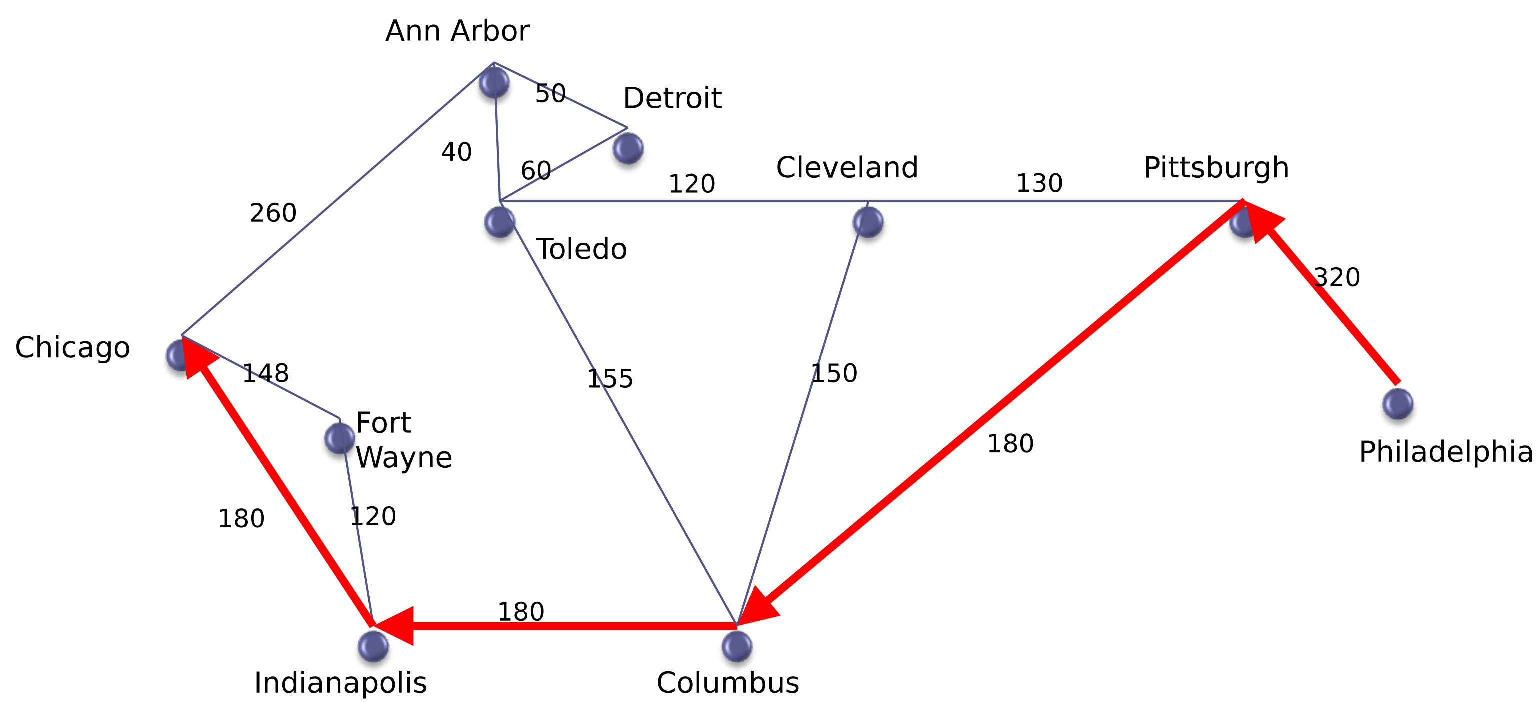
- In an undirected graph, an edge connects two nodes symmetrically (we draw a line between the two nodes)
- In a directed graph, the edge goes from the *source node* to the *target node* (we draw an arrow from the source to the target)

A tree is a special case of a directed graph

- Edge from parent to child

Paths

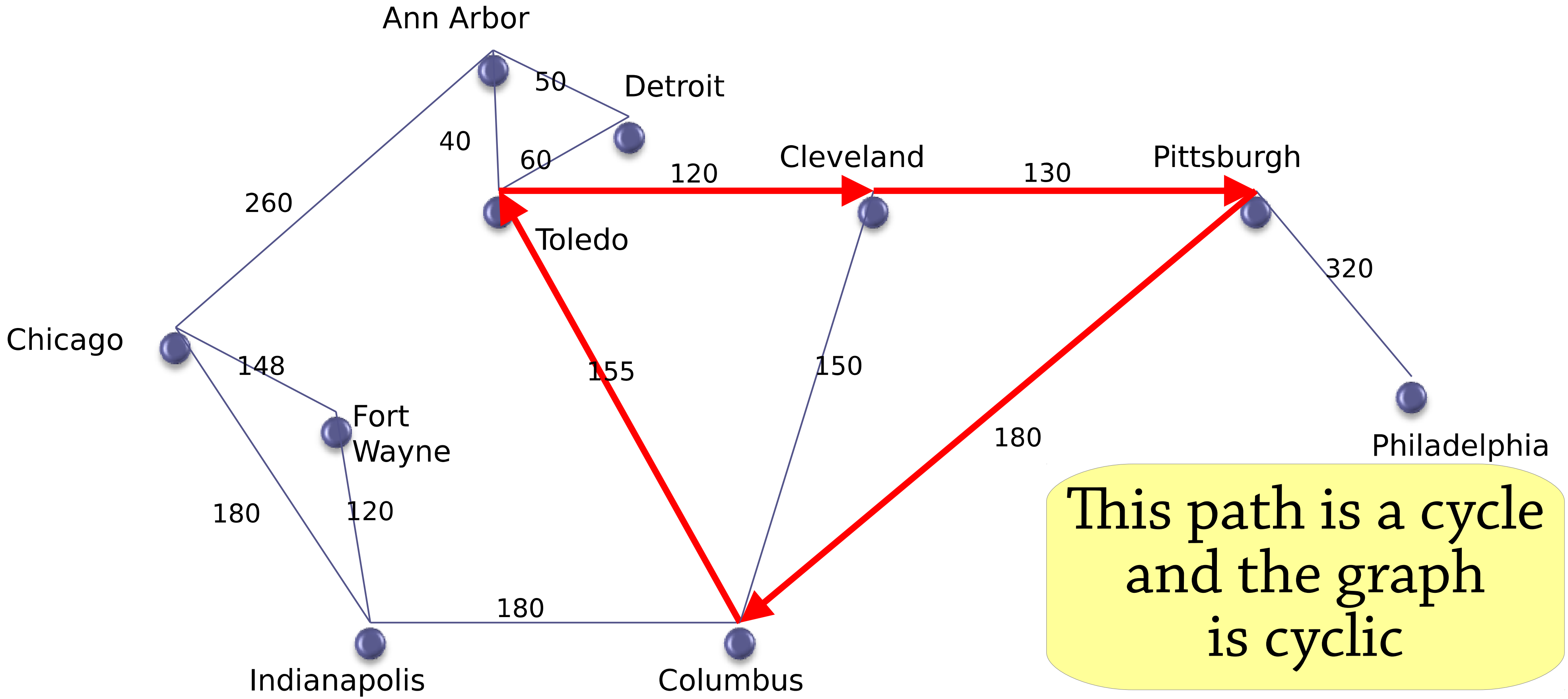
A *path* is a sequence of edges that take you from one node to another



If there is a path from node A to node B, we say that B is *reachable* from A

Cyclic graphs

A graph is *cyclic* if there is a path from a node to itself; we call the path a *cycle*. Otherwise the graph is *acyclic*.



Cyclic graphs

A path is only a cycle if:

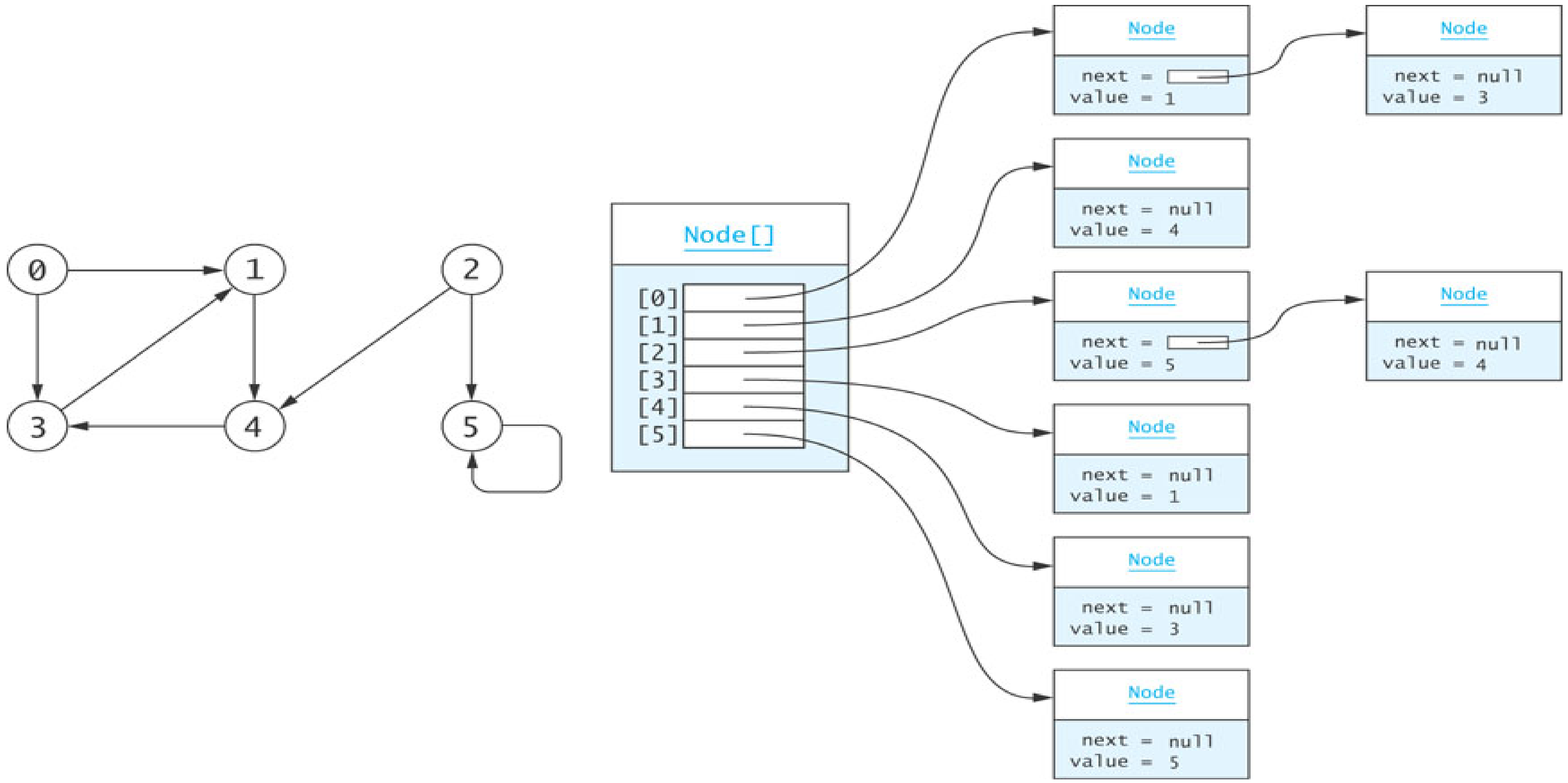
- it starts and ends at the same node
  (otherwise it's definitely not a cycle!)
- it's non-empty
  (otherwise all graphs would be cyclic)
- it is a *simple path:* it doesn't pass through the same node or edge twice, except for the first and last node
  (otherwise the following graph would be cyclic, by going from 4 to 5 and back again:

)

How to implement a graph

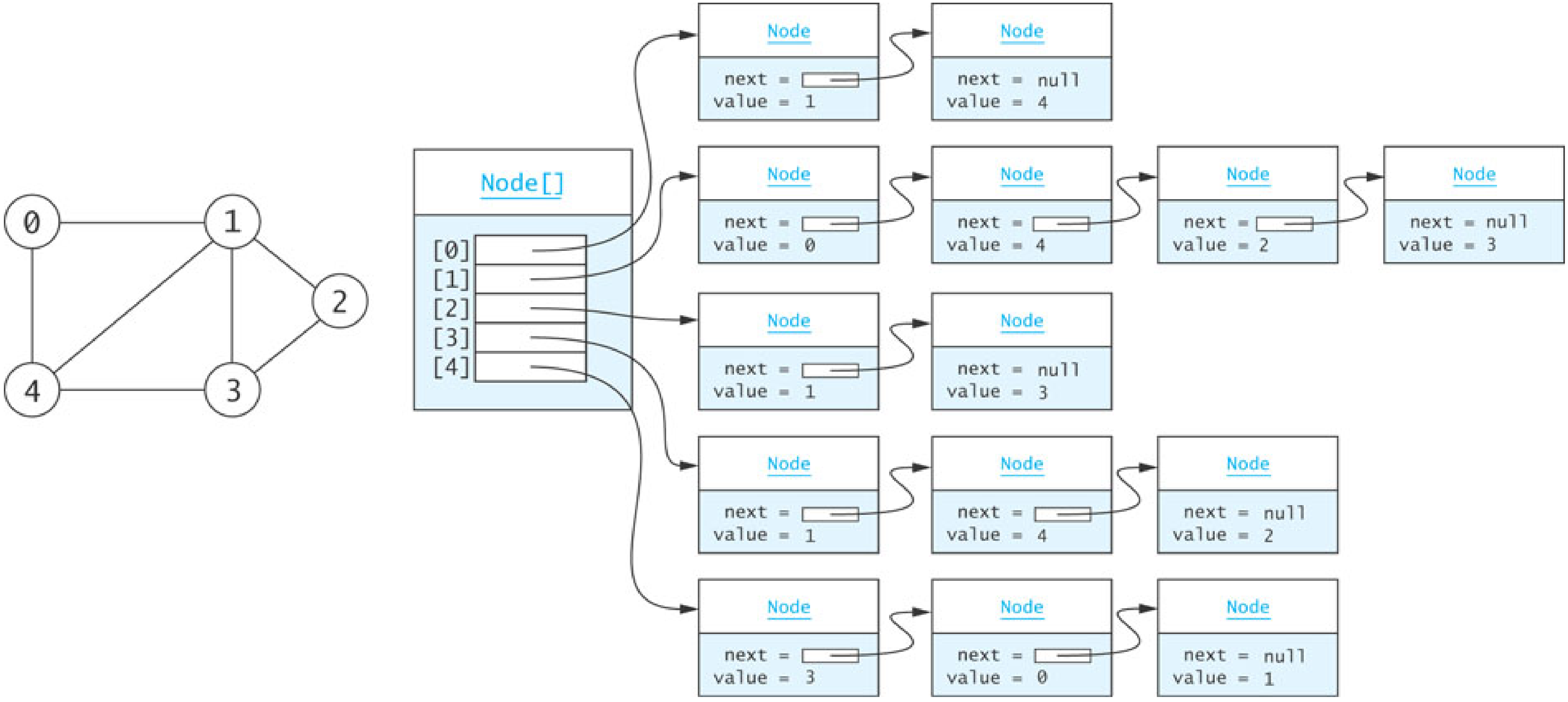
Typically: *adjacency list*

- List of all nodes in the graph, and with each node store all the edges having that node as source



Adjacency list – undirected graph

Each edge appears twice, once for the source and once for the target node



Graph algorithms:
depth-first search,
reachability,
connected components

Reachability

How can we tell what nodes are reachable from a given node?

We can start exploring the graph from that node, but we have to be careful not to (e.g.) get caught in cycles

Depth-first search is one way to explore the part of the graph reachable from a given node

Depth-first search

Depth-first search is a *traversal* algorithm

- This means it takes a node as input, and enumerates all nodes reachable from that node

It comes in two variants, *preorder* and *postorder* – we'll start with preorder

To do a *preorder* DFS starting from a node:

- visit the node
- for each outgoing edge from the node, recursively DFS the target of that edge, unless it has already been visited

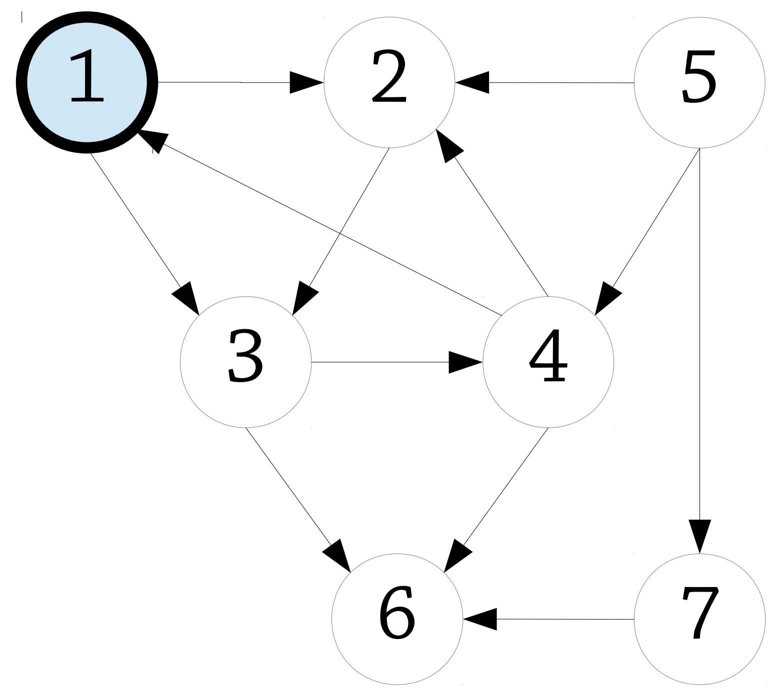
It's called preorder because we visit each node *before* its outgoing edges

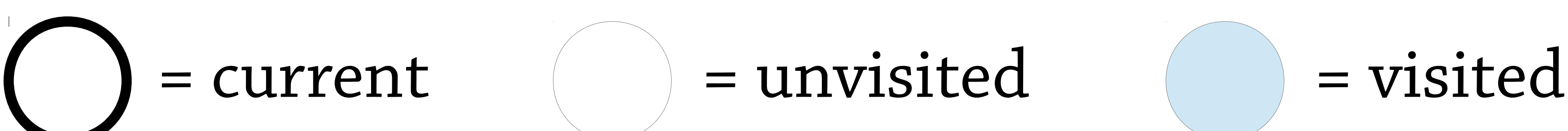
Example of a depth-first search

Visit order: 1

DFS node 1

(By the way, is 5 reachable from 1?)

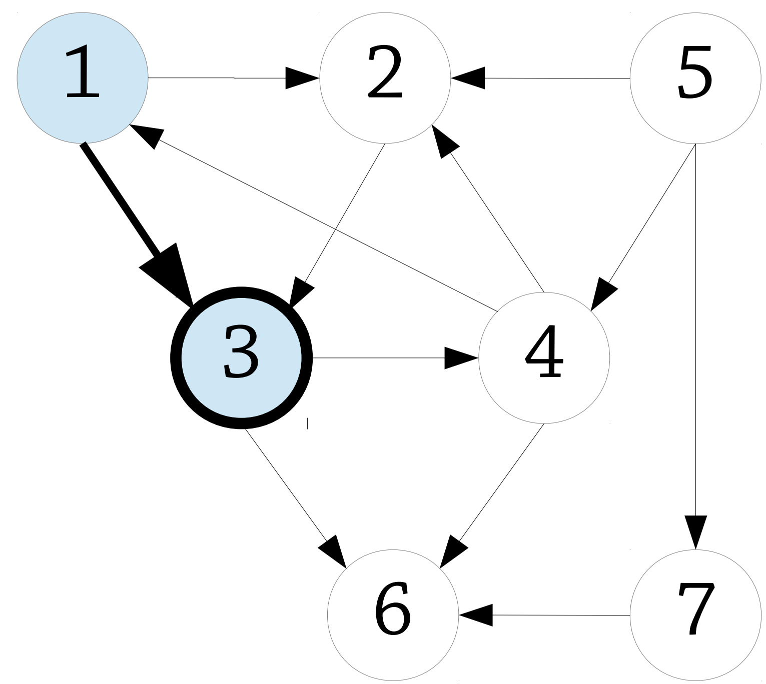




Example of a depth-first search

Visit order: 1 3

Follow edge 1 → 3,
recursively DFS node 3

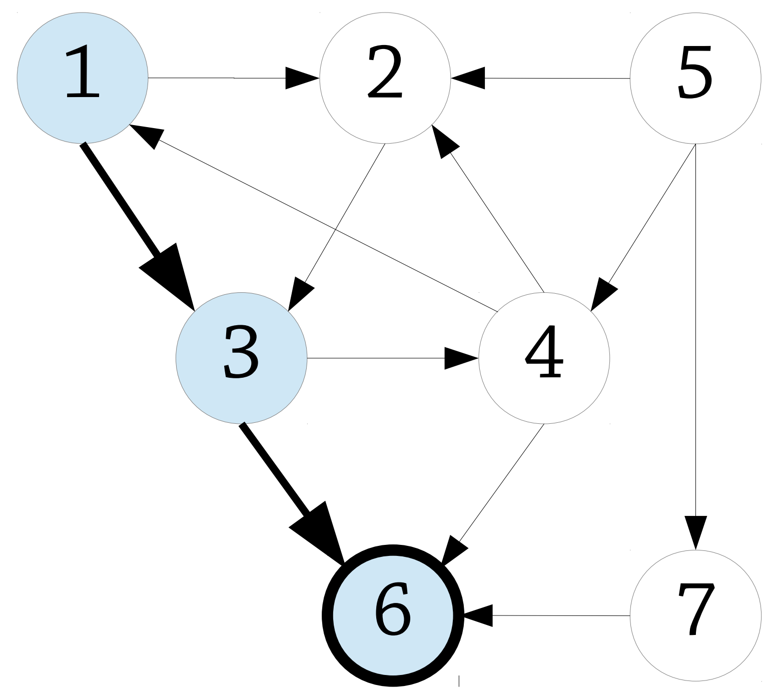


= current   = unvisited   = visited

Example of a depth-first search

Visit order: 1 3 6

Follow edge 3 → 6,
recursively DFS node 6

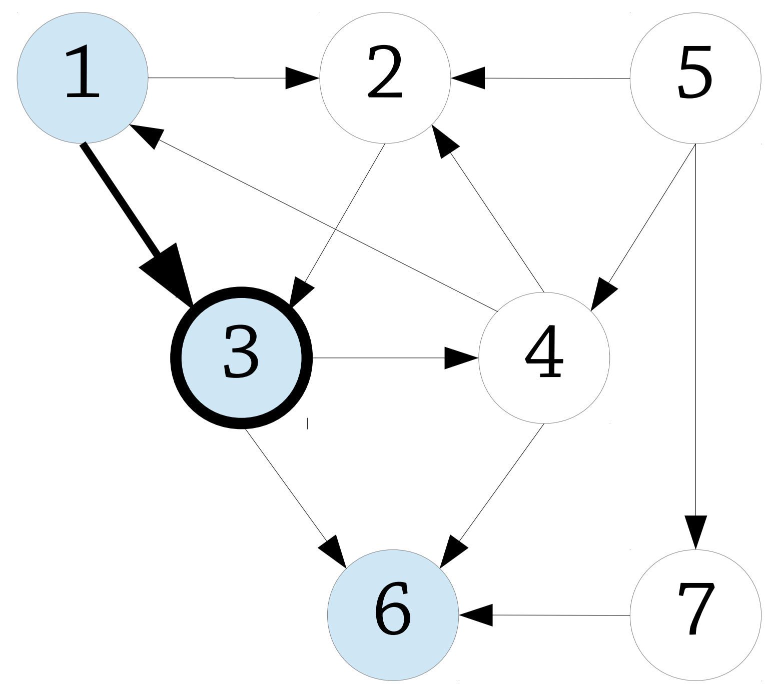


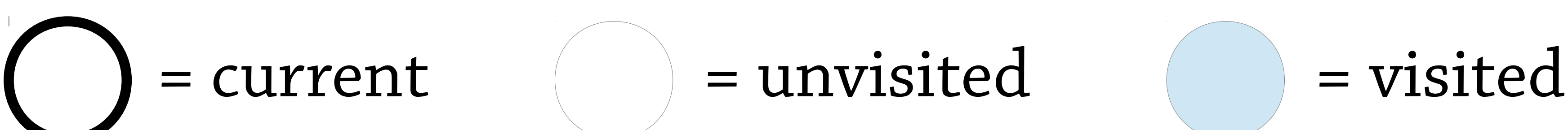
= current = unvisited = visited

Example of a depth-first search

Visit order: 1 3 6

Recursion backtracks to 3

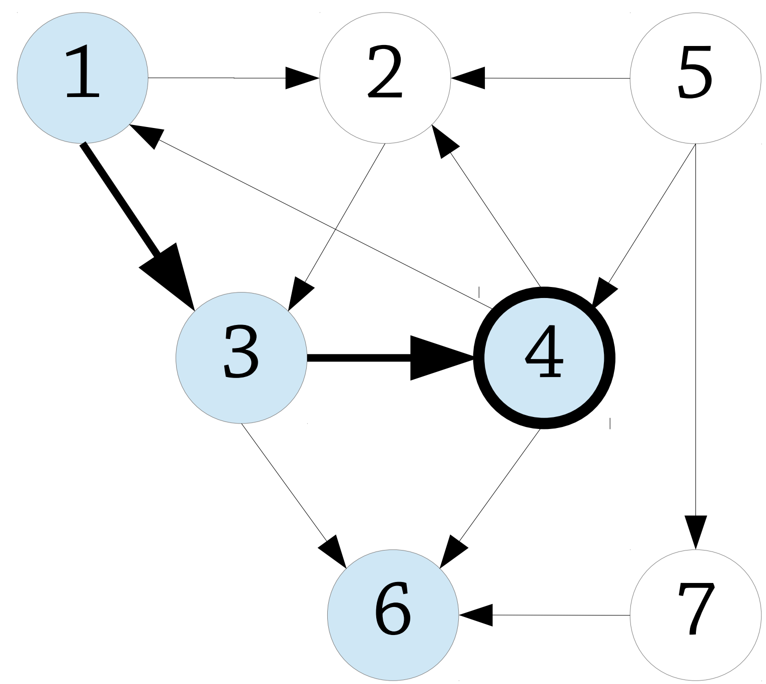




Example of a depth-first search

Visit order: 1 3 6 4

Follow edge 3 → 4,
recursively DFS node 4



= current = unvisited = visited

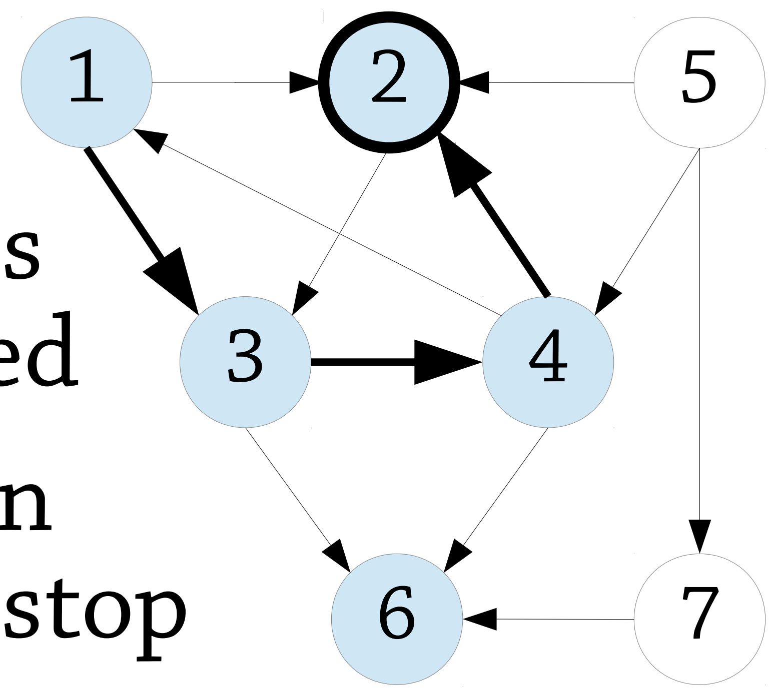
Example of a depth-first search

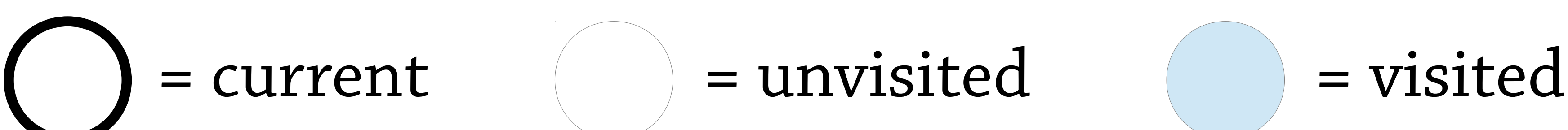
Visit order: 1 3 6 4 2

Follow edge 4 → 2, recursively DFS node 2

We don't follow 4 → 6 or 2 → 3, as those nodes have already been visited

Eventually the recursion backtracks to 1 and we stop





Reachability revisited

How can we tell what nodes are reachable from a given node?

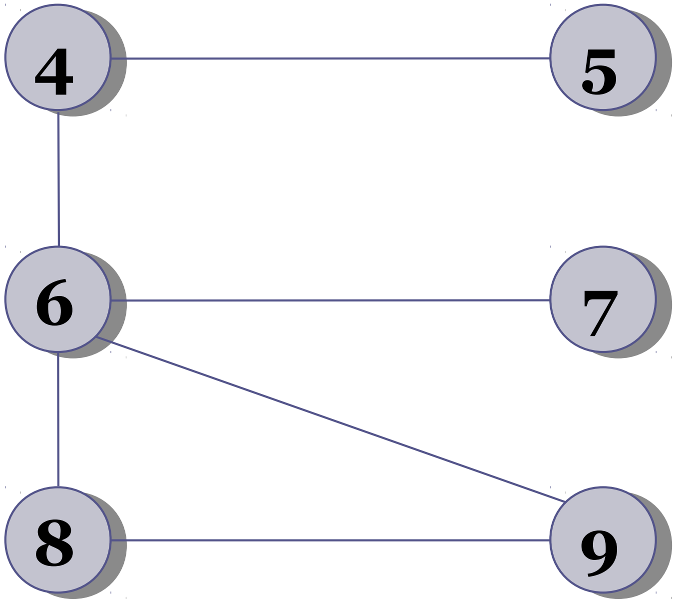
Answer:

Perform a depth-first search starting from node A, and the nodes visited by the DFS are exactly the reachable nodes

Connectedness

An *undirected* graph is called *connected* if there is a path from every node to every other node

This graph is connected

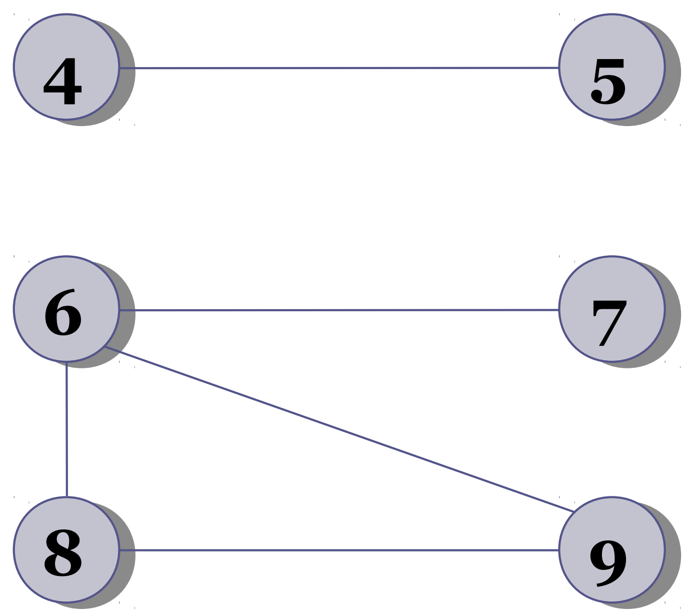


How can we tell if a graph is connected?

Connectedness

An *undirected* graph is called *connected* if there is a path from every node to every other node

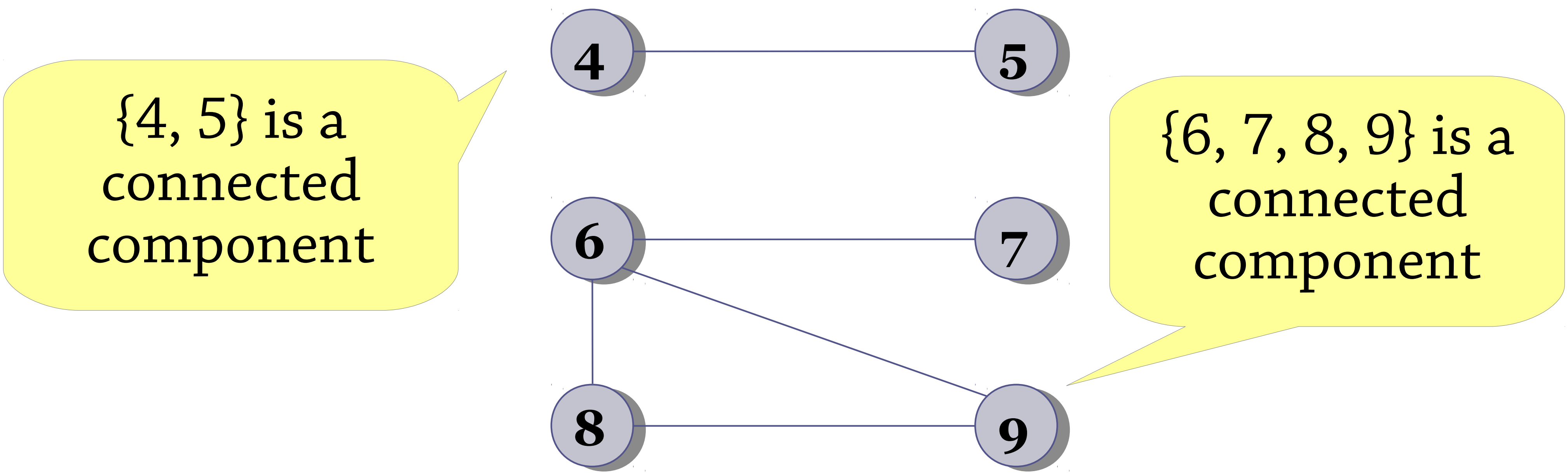
This graph is **not** connected



How can we tell if a graph is connected?

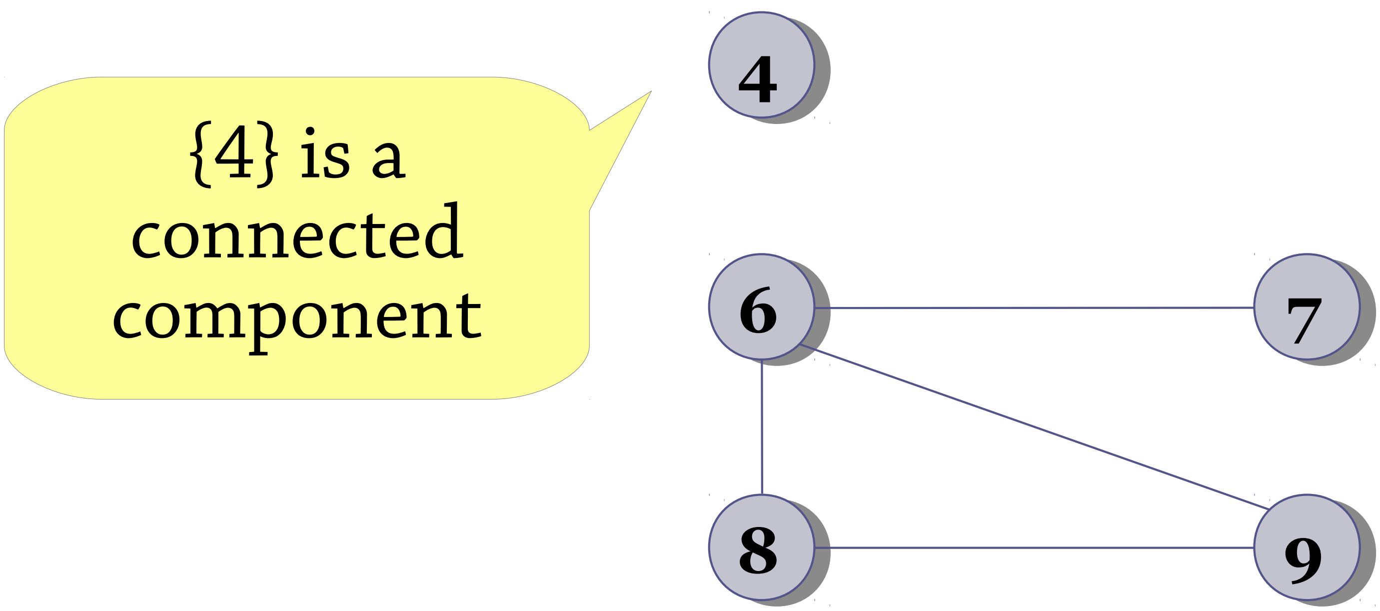
Connectedness

If an undirected graph is unconnected, it still consists of *connected components*



Connectedness

A single unconnected node is a connected component in itself



Connected components

How can we find:

- the connected component containing a given node?
- all connected components in the graph?

Connected components

To find the connected component containing a given node:

- Perform a DFS starting from that node
- The set of visited nodes is the connected component

To find all connected components:

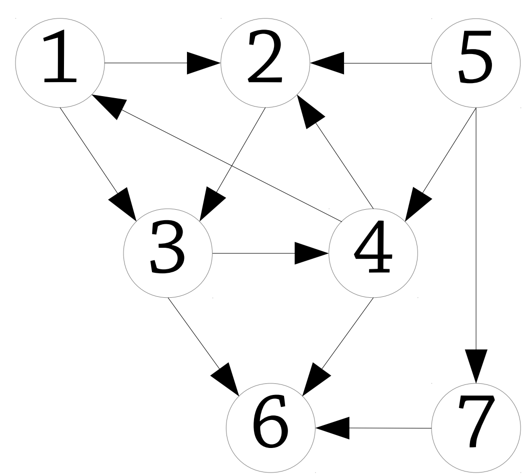
- Pick a node that doesn't have a connected component yet
- Use the algorithm above to find its connected component
- Repeat until all nodes are in a connected component

Strongly-connected components

In a directed graph, there are two notions of connectedness:

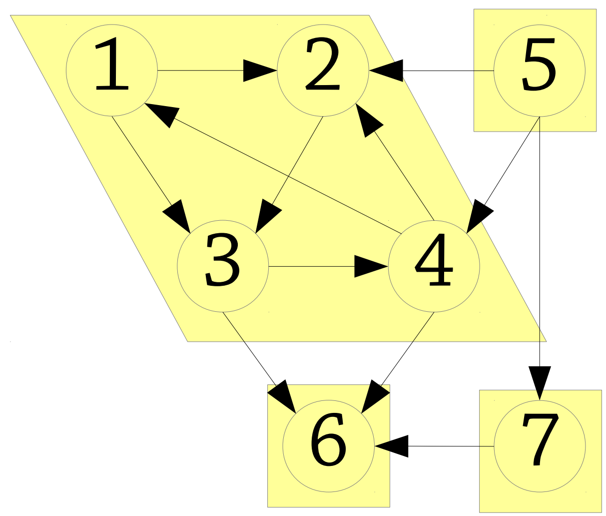
- *strongly connected* means there is a path from every node to every other node
- *weakly connected* means the graph is connected if you ignore the direction of the edges (the equivalent undirected graph is connected)

This graph is weakly connected, but not strongly connected (why?)



Strongly-connected components

You can always divide a directed graph into its *strongly-connected components (SCCs)*:



In each strongly-connected component, every node is reachable from every other node

- The relation "nodes A and B are both reachable from each other" is an *equivalence relation* on nodes
- The SCCs are the equivalence classes of this relation

Strongly-connected components

To find the SCC of a node A, we take the intersection of:

- the set of nodes reachable from A
- the set of nodes which A can be reached from (the set of nodes "backwards-reachable" from A)

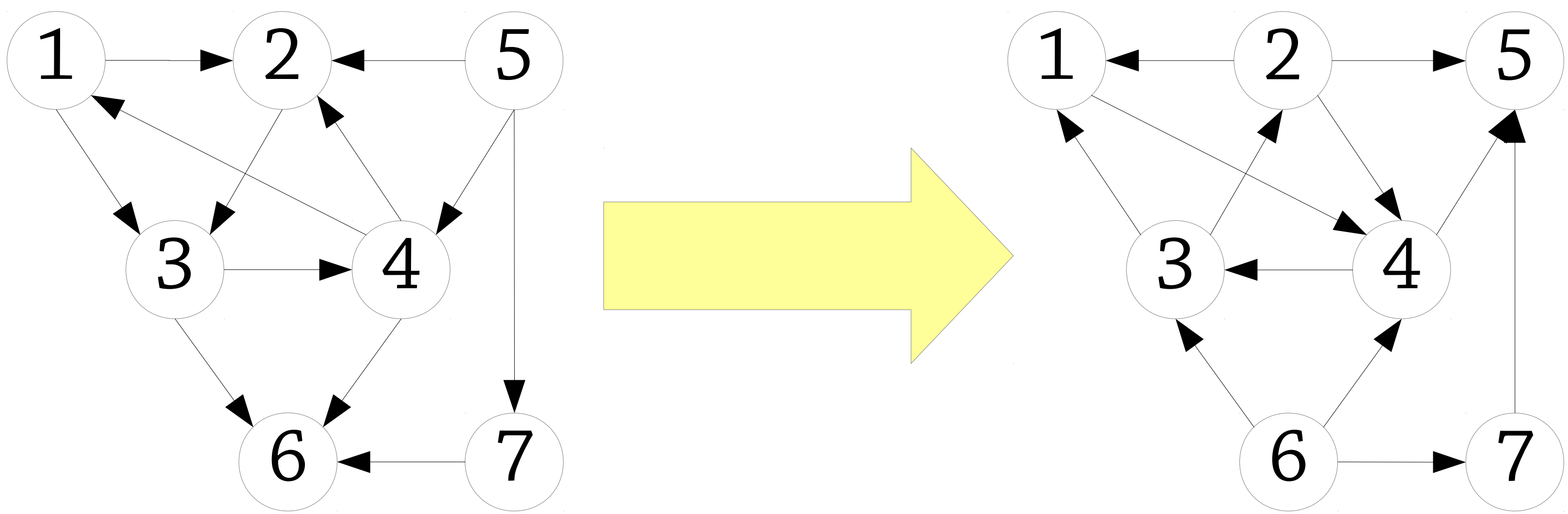
This gives us all the nodes B such that:

- there is a path from A to B, and
- there is a path from B to A

To find the set of nodes backwards-reachable from A, we will use the idea of the *transpose* of a graph

Transpose of a graph

To find the transpose of a directed graph, flip the direction of all the graph's edges:

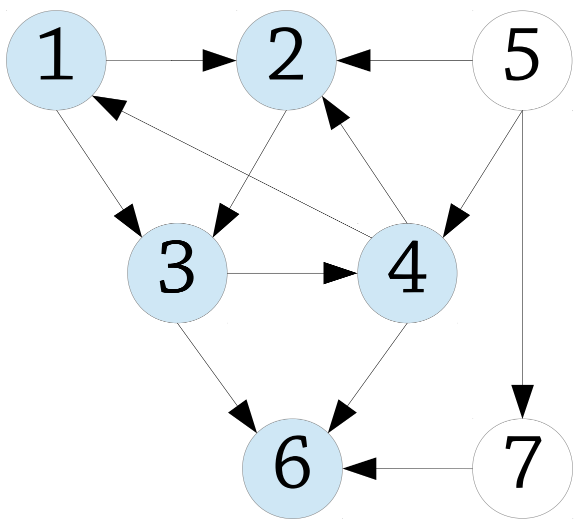


Graph   Transpose

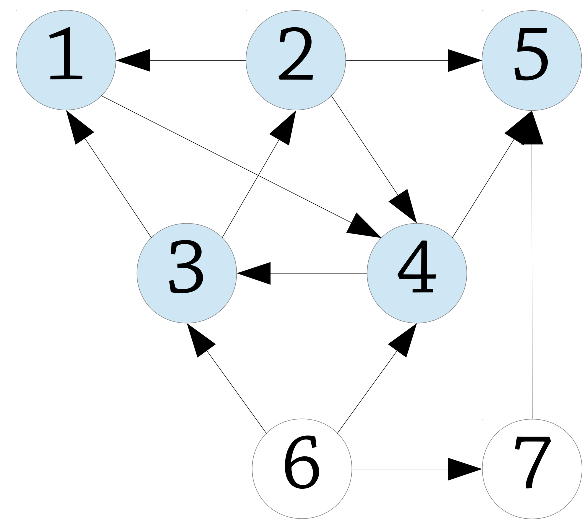
Note that: there is a path from A to B in the original graph iff there is a path from B to A in the transpose graph!

Strongly-connected components

To find the SCC of a node (such as 2), perform a DFS in the graph and the transpose graph:



Graph



Transpose

The nodes visited in both DFSs are the SCC – in this case {1, 2, 3, 4}

Strongly-connected components

To find the SCC of a node A:

- Find the set of nodes reachable from A, using DFS
- Find the set of nodes which have a path to A, by doing a DFS in the *transpose* graph
- Take the intersection of these two sets

Implementation in practice:

- When doing the DFS in the transpose graph, we restrict the search to the nodes that were reachable from A in the original graph

What do SCCs mean?

The SCCs in a graph tell you about the *cycles* in that graph!

- If a graph has a cycle, all the nodes in the cycle will be in the same SCC
- If an SCC contains two nodes A and B, there is a path from A to B and back again, so there is a cycle

A directed graph is acyclic iff:

- All the SCCs have size 1, and
- no node has an edge to itself (SCCs do not take any notice of self-loops)

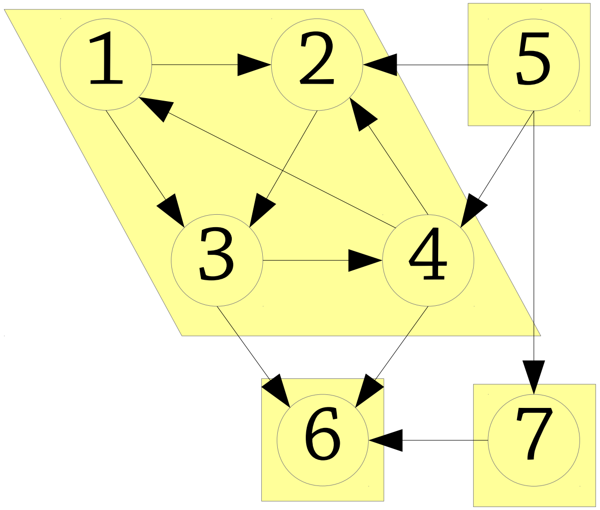
Cycles and SCCs

Here is the directed graph from before.

Notice that:

- The big SCC is where all the cycles are
- The acyclic "parts" of the graph have SCCs of size 1

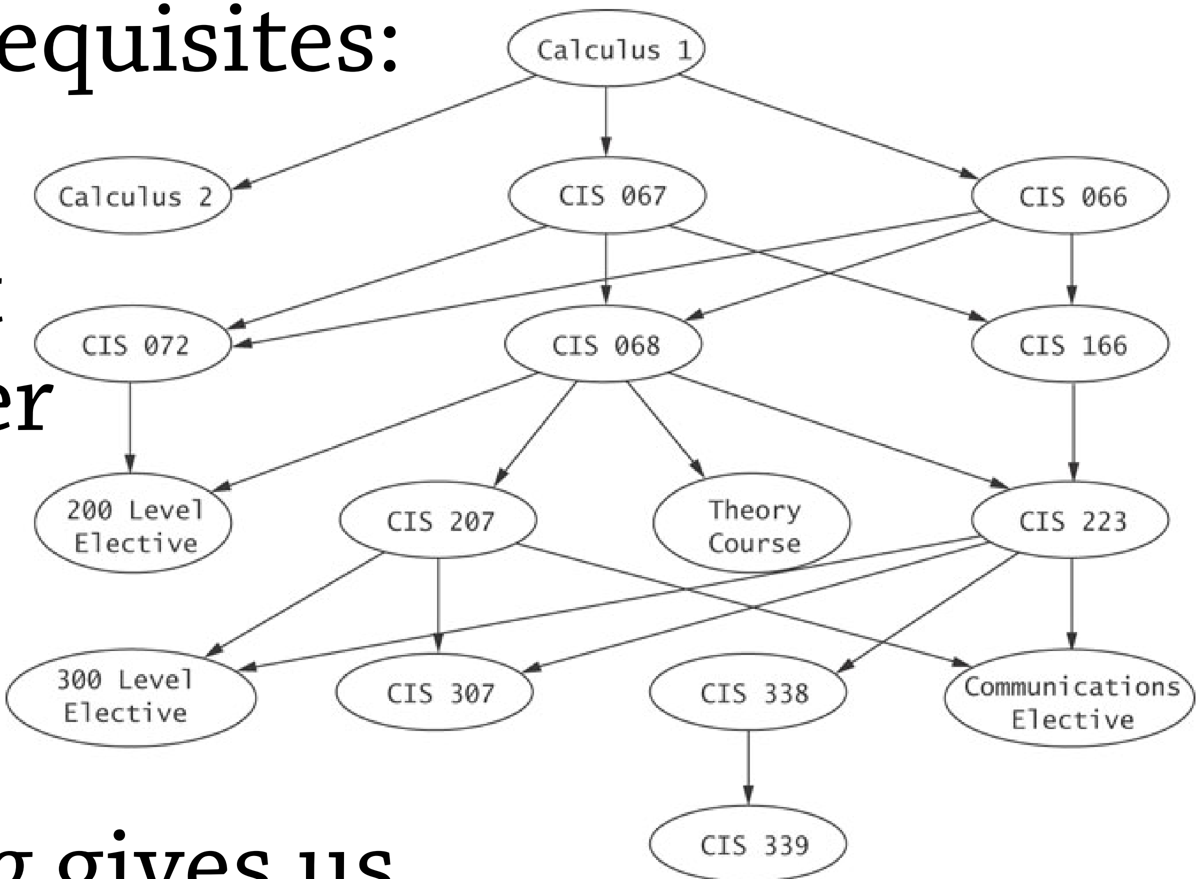
The SCCs characterise the cycles in the graph!



Graph algorithms: postorder DFS, detecting cycles, topological sorting

Topological sorting

Here is a *directed acyclic graph (DAG)* with courses and prerequisites:



We might want to find out: what is a possible order to take these courses in?

This is what *topological sorting* gives us.

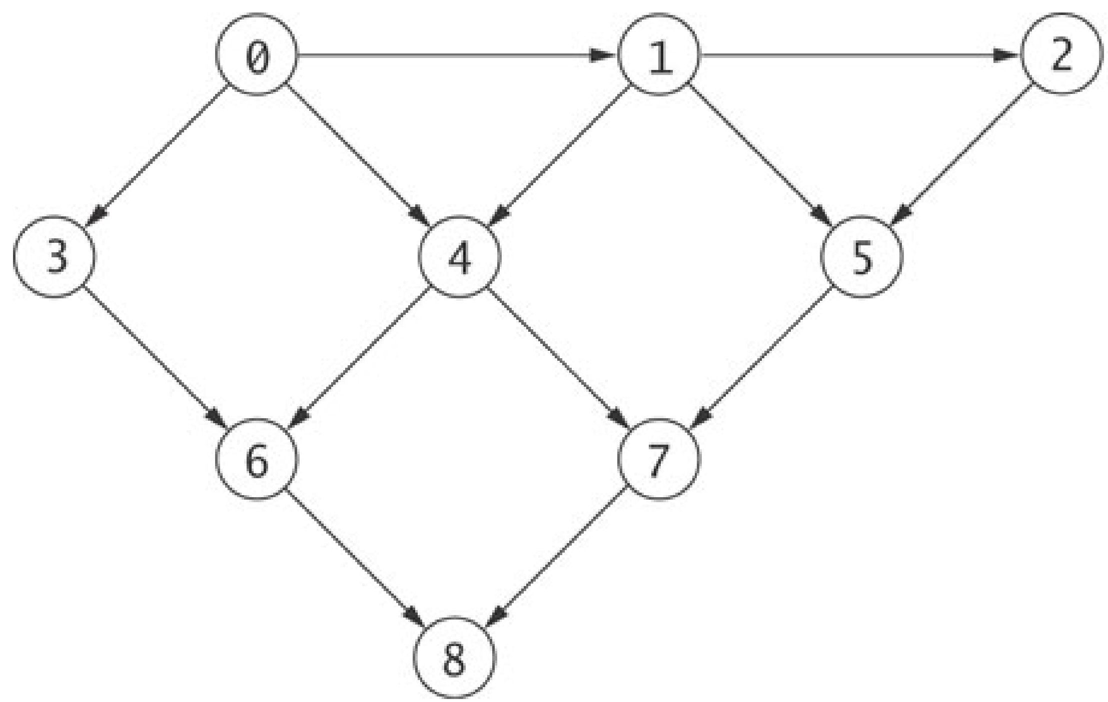
Note that the graph must be acyclic!

Example: topological sort

A topological sort of the nodes in a DAG is a list of all the nodes, so that *if there is a path from u to v, then u comes before v in the list*

Every DAG has a topological sort, often several

012345678 is a topological sort of this DAG, but 015342678 isn't.



Postorder depth-first search

To implement topological sorting we'll need a variant of DFS called *postorder* depth-first search

To do a postorder DFS starting from a node:

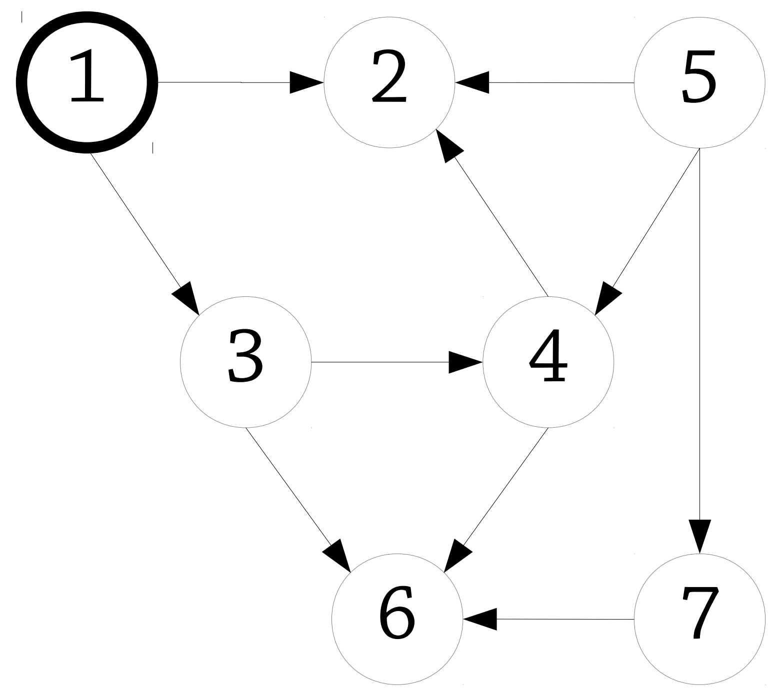
- mark the node as reached
- for each outgoing edge from the node, recursively DFS the target of that edge, unless it has already been reached
- visit the node

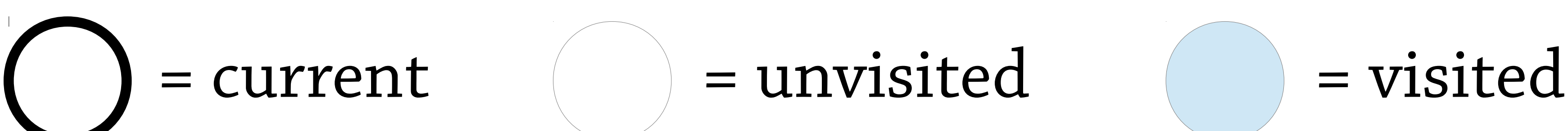
In postorder DFS, we visit each node *after* we visit its outgoing edges!

Postorder depth-first search

Visit order:

DFS node 1 (don't visit it yet, but remember that we have reached it)

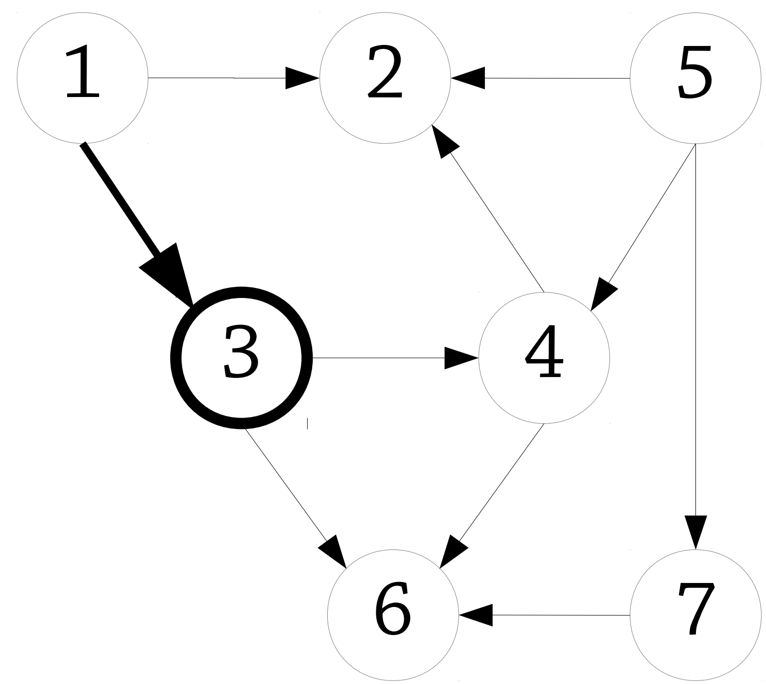




Postorder depth-first search

Visit order:

Follow edge 1 → 3,
recursively DFS node 3



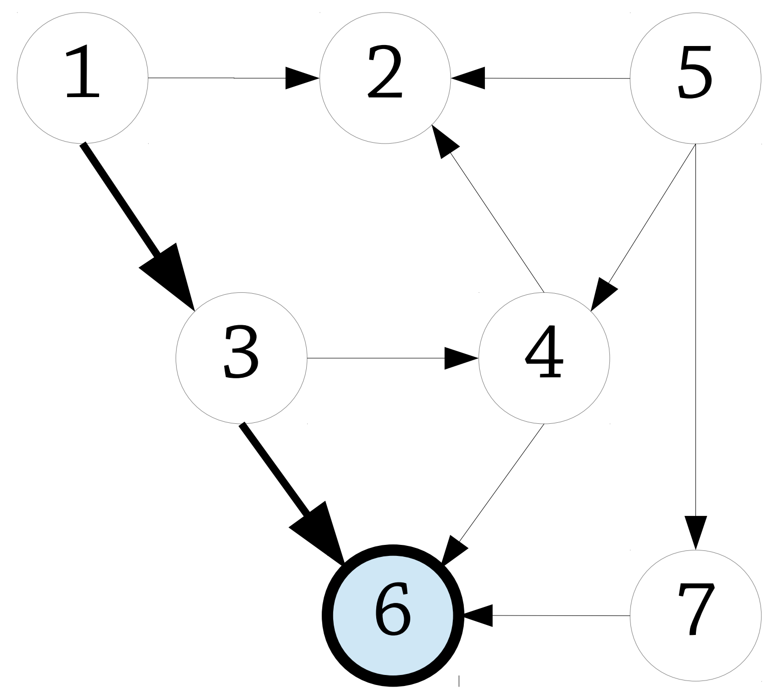
= current = unvisited = visited

Postorder depth-first search

Visit order: 6

Follow edge 3 → 6, recursively DFS node 6

The recursion bottoms out, visit 6!

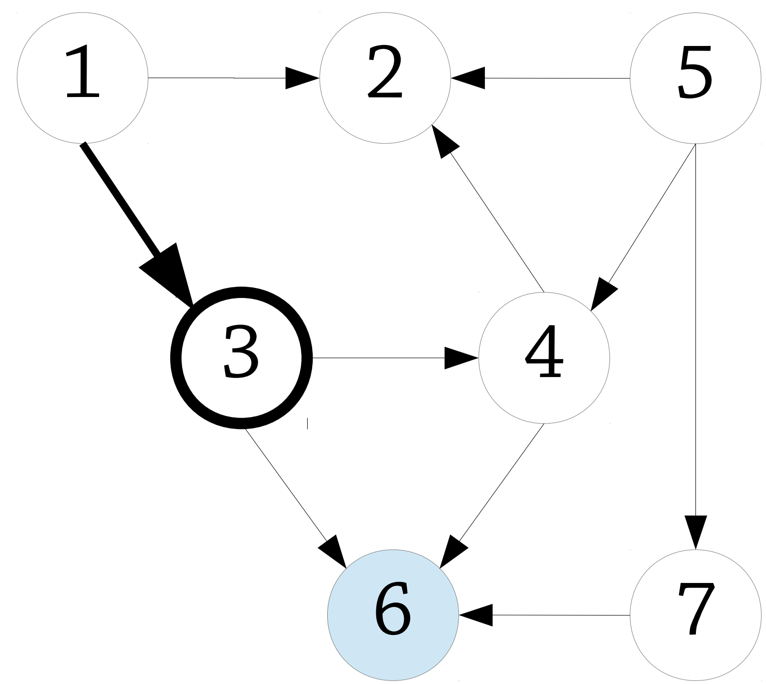


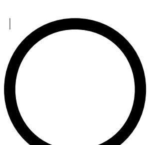
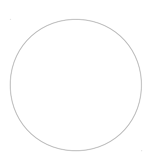
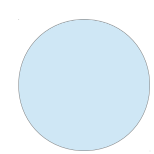
= current = unvisited = visited

Postorder depth-first search

Visit order: 6

Recursion backtracks to 3

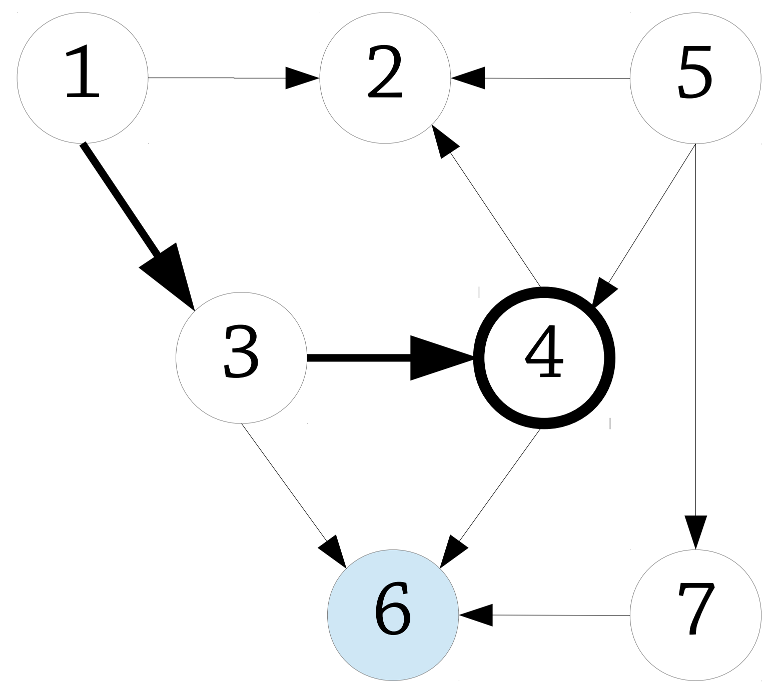


 = current  = unvisited  = visited

Postorder depth-first search

Visit order: 6

Follow edge 3 → 4,
recursively DFS node 4



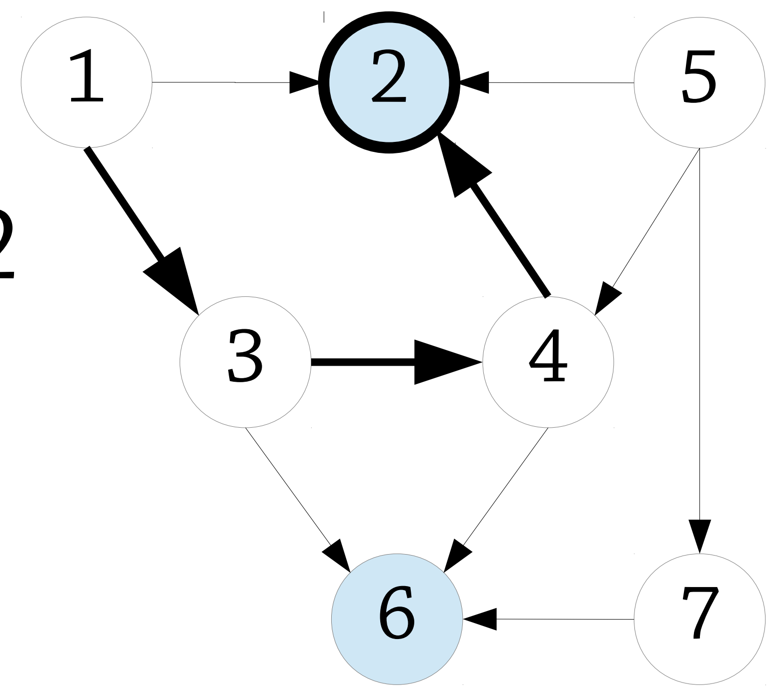
= current = unvisited = visited

Postorder depth-first search

Visit order: 6 2

Follow edge 4 → 2, recursively DFS node 2

The recursion bottoms out again and we visit 2



= current

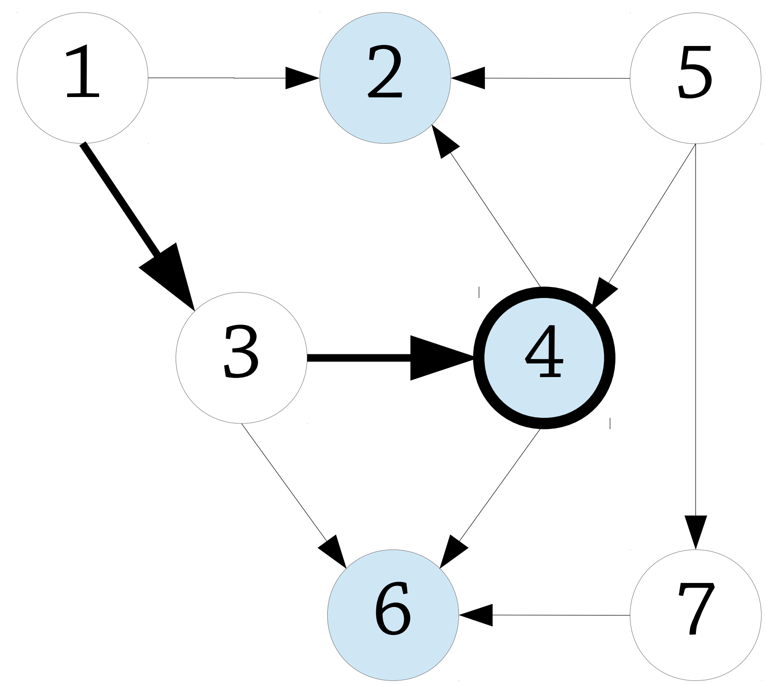
= unvisited

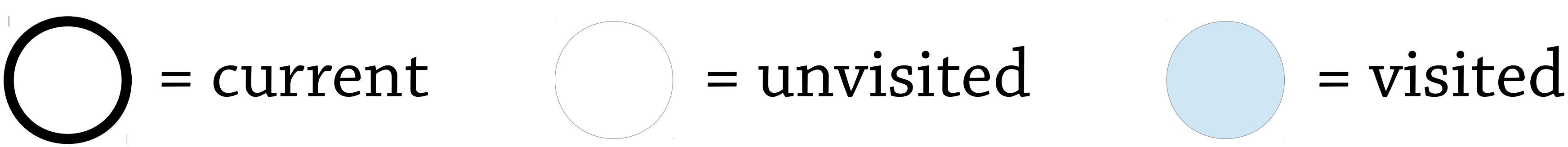
= visited

Postorder depth-first search

Visit order: 6 2 4

The recursion backtracks and now we visit 4

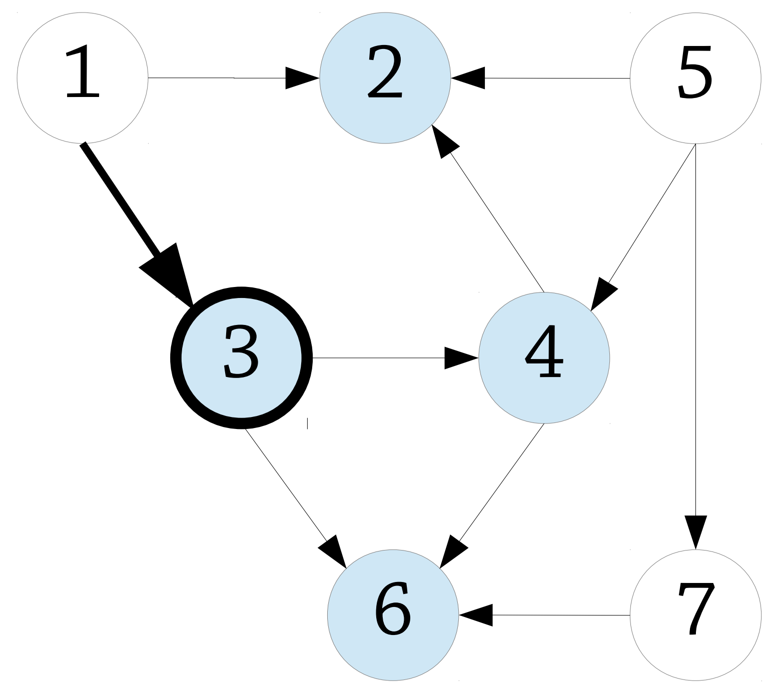


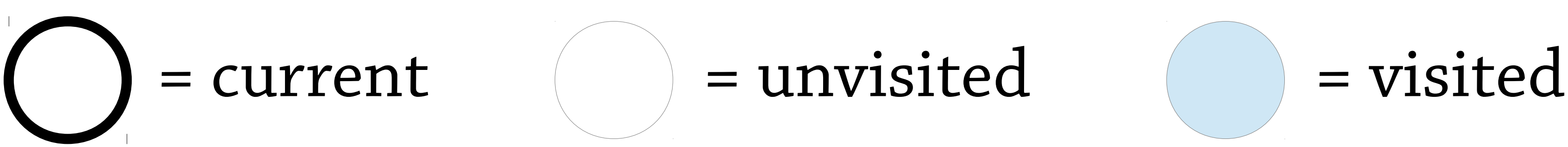


Postorder depth-first search

Visit order: 6 2 4 3

The recursion backtracks and now we visit 3

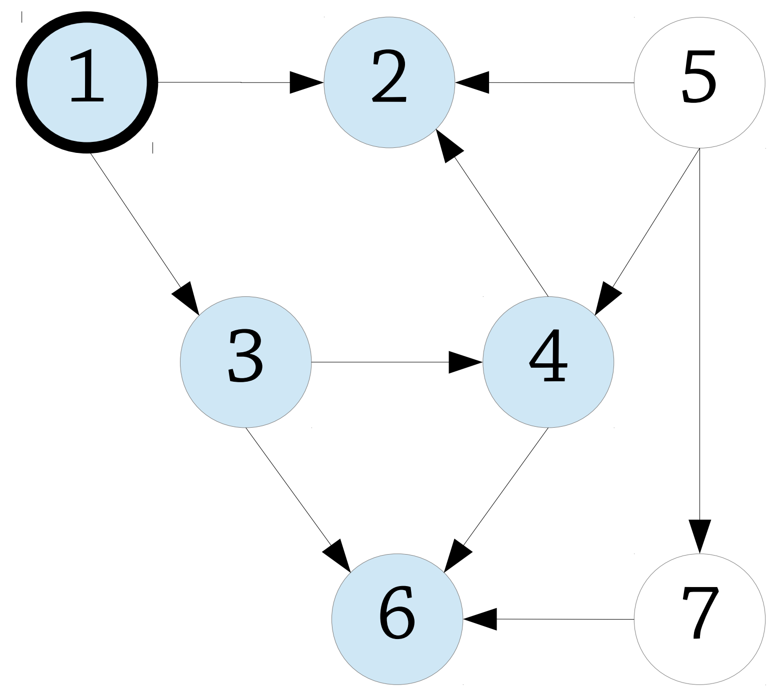


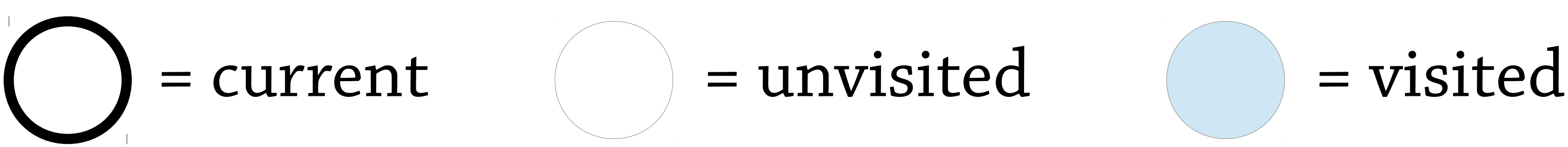


Postorder depth-first search

Visit order: 6 2 4 3 1

The recursion backtracks and now we visit 1





Why postorder DFS?

In postorder DFS:

- We only visit a node *after* we recursively DFS its successors (the nodes it has an edge to)

If we look at the order the nodes are visited (rather than the calls to DFS):

- If the graph is acyclic, we visit a node only after we have visited all its successors

If we look at the list of nodes in the order they are visited, each node comes after all its successors (look at the previous slide)

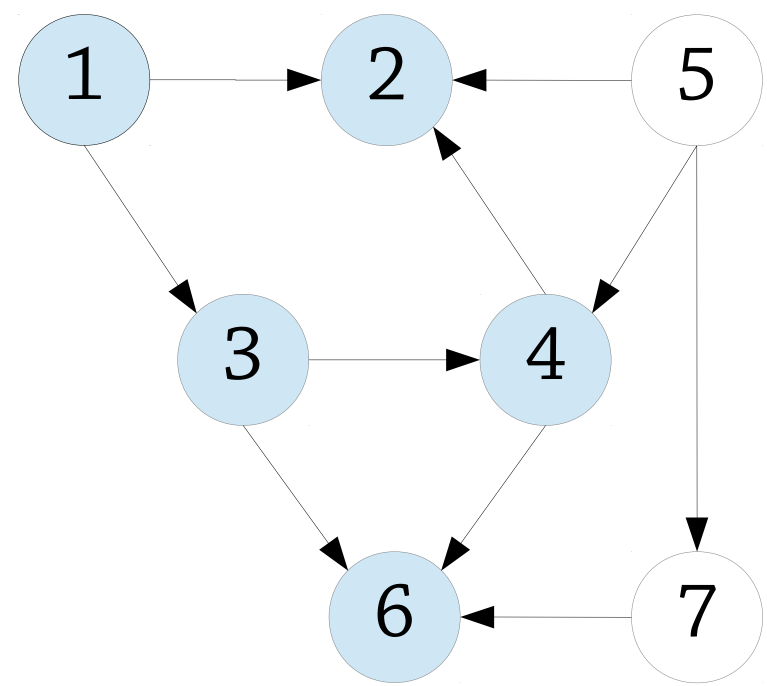
Topological sorting

Visit order: 6 2 4 3 1

In topological sorting, we want each node to come *before* its successors...

With postorder DFS, each node is visited *after* its successors!

Idea: to topologically sort, do a postorder DFS, look at the order the nodes are visited in and *reverse* it



Small problem: not all nodes are visited!
Solution: pick a node we haven't visited and DFS it

Topological sorting

To topologically sort a DAG:

- Pick a node that we haven't visited yet
- Do a postorder DFS on it
- Repeat until all nodes have been visited

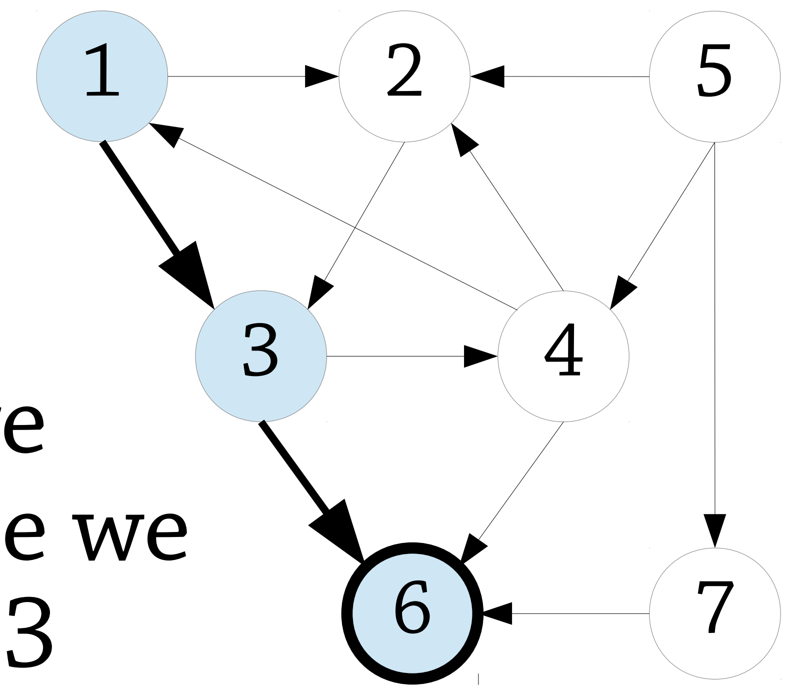
Then take the list of nodes in the order they were visited, and reverse it

If the graph is acyclic, the list is topologically sorted:

- If there is a path from node A to B, then A comes before B in the list

Preorder vs postorder

You might think that in preorder DFS, we visit each node *before* we visit its successsors



But this is not the case, in this example from earlier we visited 6 before its predecessor 4, because we happened to go through 3

Preorder DFS visits the nodes in "any old order" – postorder is more well-behaved

Detecting cycles in graphs

We can only topologically sort *acyclic* graphs – how can we detect if a graph is cyclic?

Easiest answer: topologically sort the graph and check if the result is actually topologically sorted

- Does any node in the result list have an edge to a node *earlier* in the list? If so, the topological sorting failed, and the graph must be cyclic
- Otherwise, the graph is acyclic

Kosaraju's algorithm **(not on exam)**

Kosaraju's algorithm finds *all* the SCCs in a directed graph in linear time

Recall our algorithm to find the SCC of a node A:

- Do a DFS starting from node A
- Do a DFS starting from node A in the *transpose* graph
- Take the intersection of the two visited sets

In Kosaraju's algorithm, we first do a DFS starting from node A, giving a set S of visited nodes

Then we find the SCCs of *all* nodes in S, by doing *several* DFSes in the transpose graph!

Kosaraju's algorithm **(not on exam)**

Start with a node A, do a *topological sort* starting from A

Now take the visited nodes in topological order, and for each node:

- If we have already assigned the node an SCC, skip it
- Otherwise, do a DFS starting from that node in the transpose graph
- The SCC of that node is the intersection of the two visited sets

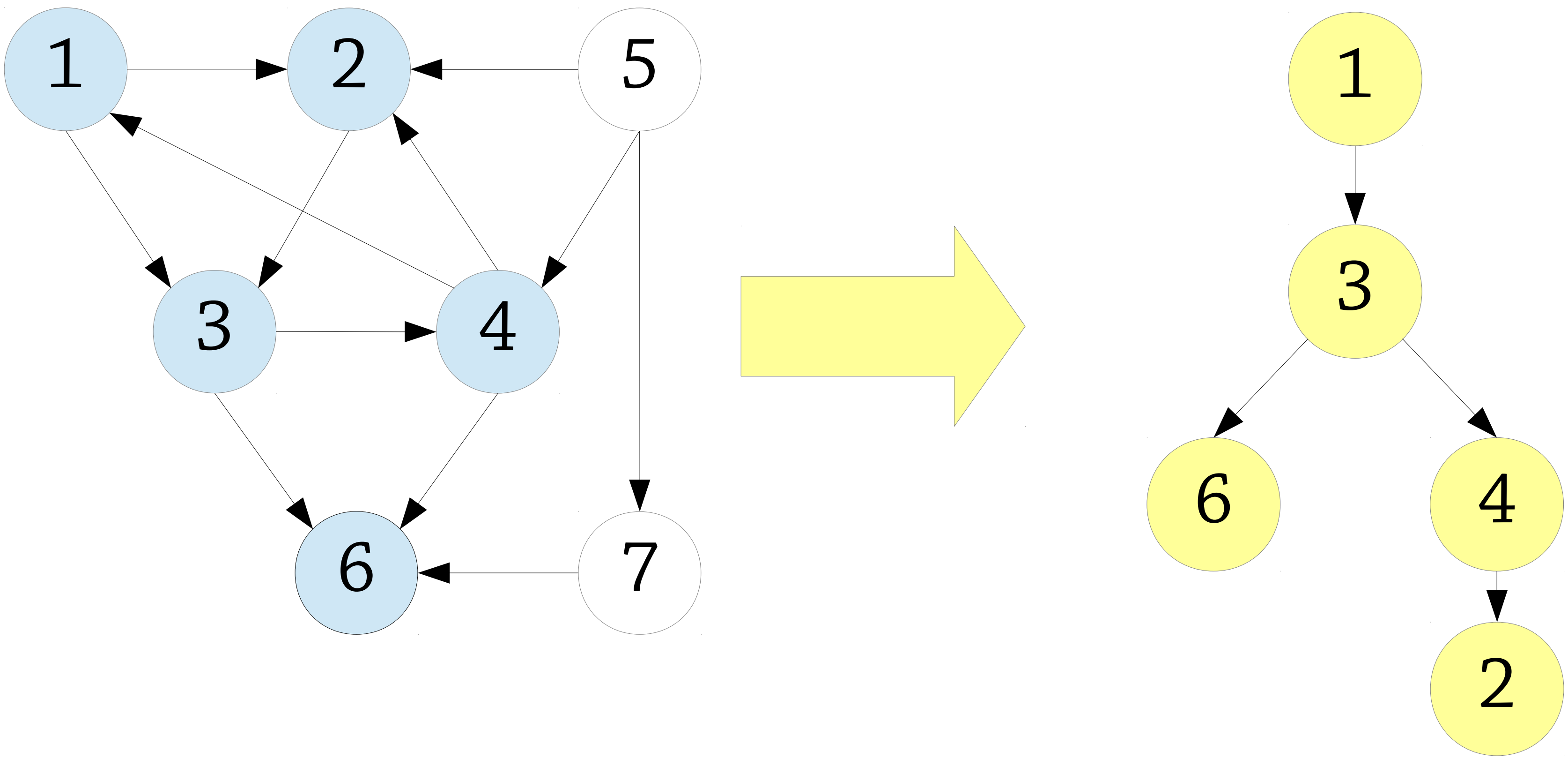
Read up on it if you're interested!

- http://scienceblogs.com/goodmath/2007/10/30/computing-strongly-connected-c/

An alternative:
depth-first forests
(not on exam)

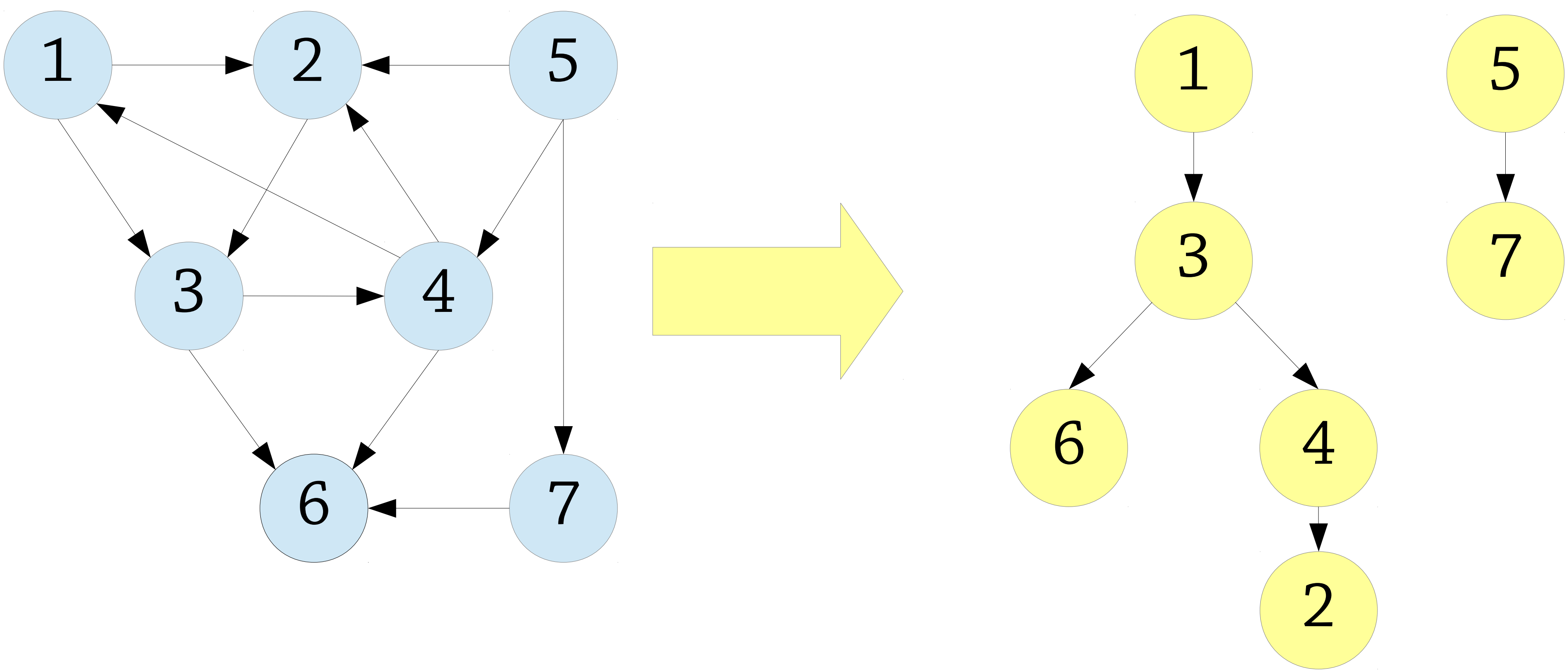
Depth-first forests

Instead of producing a *list of nodes*, DFS can return a *tree* that shows how the nodes were explored (the recursion structure):



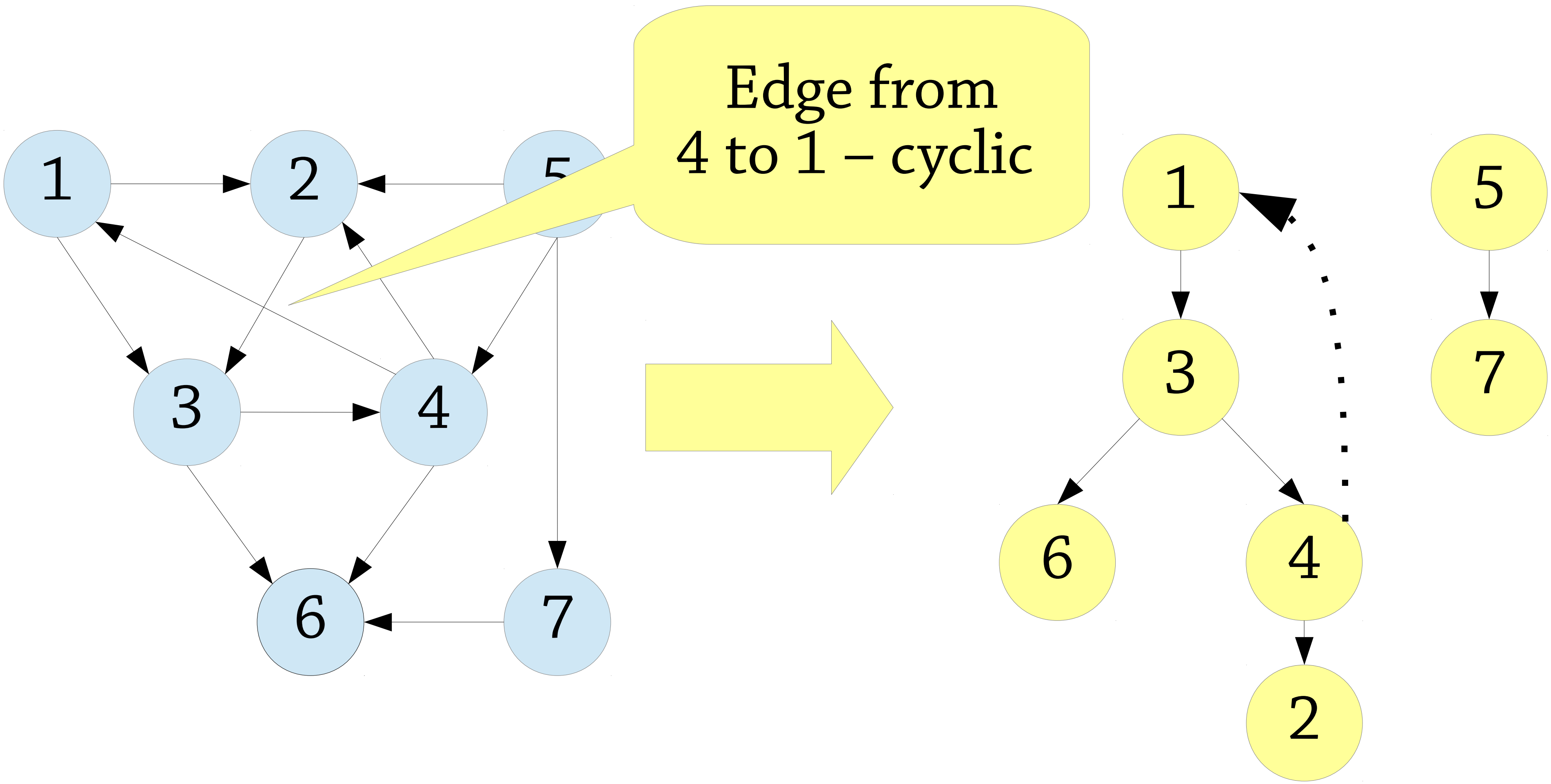
Depth-first forests

Repeating until all nodes have been visited, we get a *forest* (set of trees):



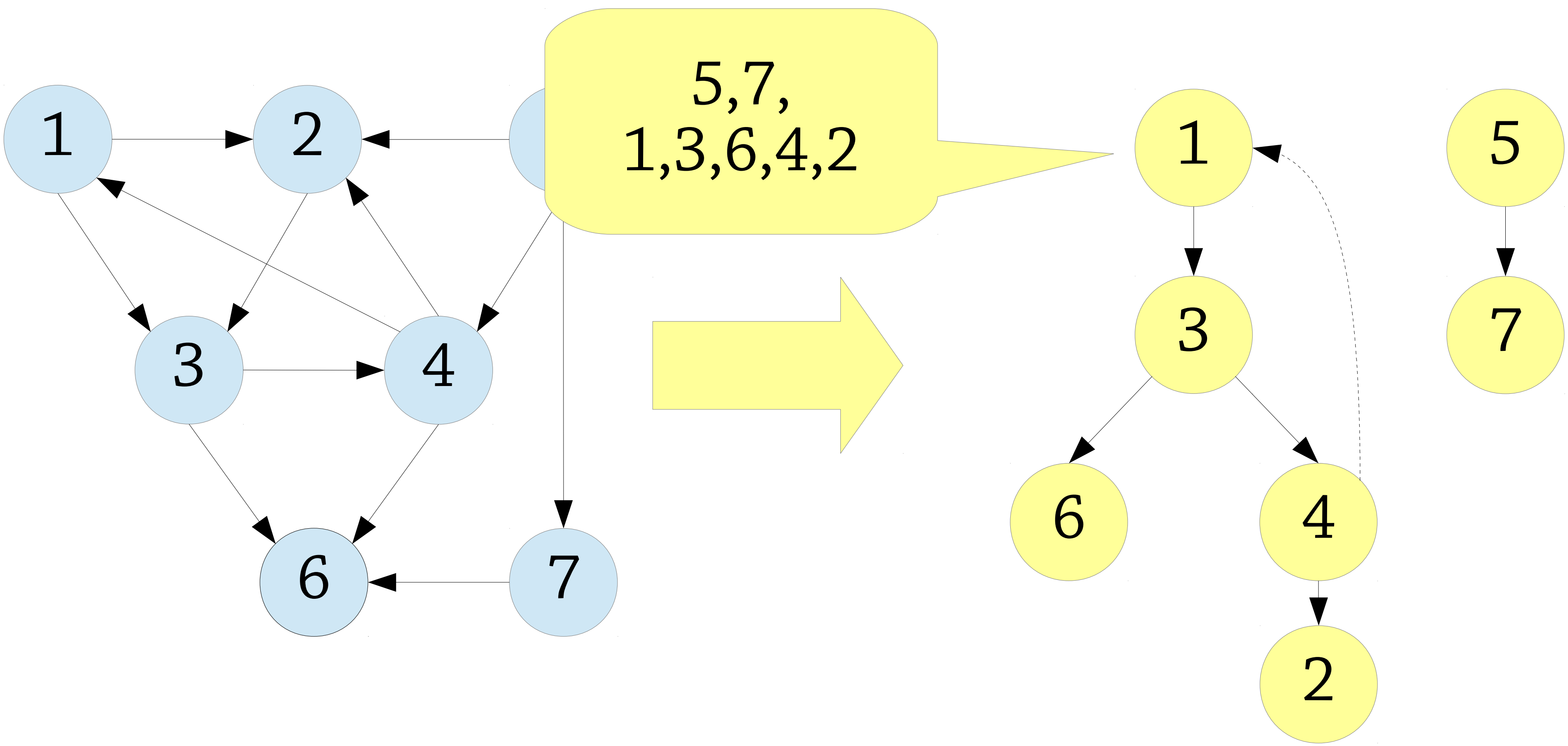
Depth-first forests

A graph is cyclic iff the graph has an edge from a node in the tree to its ancestor:



Depth-first forests

You can also topologically sort a graph by flattening the forest into a list!



Depth-first forests

The idea: make DFS return a forest of nodes, instead of a list

- Pre/post-order? Those are just different ways to flatten the forest

Many algorithms based on DFS come out pretty elegant that way

- Especially in a functional setting, where trees are very easy to deal with
- You can view the graph as a forest, plus some extra edges that go upwards, downwards or sideways in the tree

If you're interested, you can read the paper

"Graph algorithms with a functional flavour"

which is on the course webpage

Summary

Graphs are extremely useful!

- Common representation: adjacency lists (or just implicitly as references between the objects in your program)

Several important graph algorithms:

- Reachability – can I get from node A to B?
- Does the graph have a cycle?
- Strongly-connected components – where are the cycles in the graph?
- Topological sorting – how can I order the nodes in an acyclic graph?

All based on depth-first search!

- Enumerate the nodes reachable from a starting node
- Preorder: visit each node before its successors
- Postorder: visit each node after its successors, gives nicer order
- Common pattern in these algorithms: repeat DFS from different nodes until all nodes have been visited