

CHALMERS UNIVERSITY OF TECHNOLOGY | UNIVERSITY OF GOTHENBURG

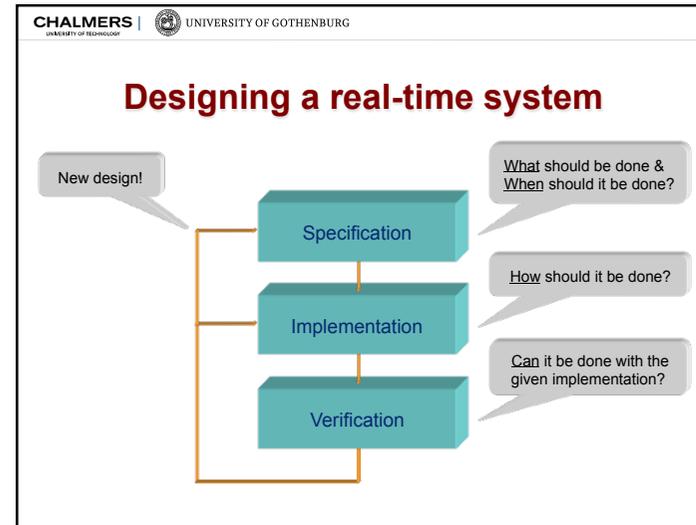


Parallel & Distributed Real-Time Systems

Lecture #2

Professor Jan Jonsson

Department of Computer Science and Engineering
Chalmers University of Technology



CHALMERS UNIVERSITY OF TECHNOLOGY | UNIVERSITY OF GOTHENBURG

Verification

Since timeliness is such an important characteristic of a real-time system: how do we verify that the timing constraints are met for a given system implementation?



... so we don't miss that hard deadline ...



... so we don't miss too many soft deadlines ...



... while we at the same time avoid analyzing all possible software execution scenarios

CHALMERS UNIVERSITY OF TECHNOLOGY | UNIVERSITY OF GOTHENBURG

Verification

What is needed for formal verification?

- A good timing model
 - Enables expressing the timing properties of the application in a syntactically unambiguous way
 - Enables timing constraints to be reflected at all design levels: from specification level (**end-to-end constraints**) to implementation level
- A good schedulability analysis
 - Enables prediction of required processing capacity, e.g. # and speed of processors, of the hardware (**when software is known**)
 - Enables prediction of required resource usage from the software (**when hardware implementation is known**)

CHALMERS UNIVERSITY OF TECHNOLOGY | UNIVERSITY OF GOTHENBURG

Verification

What sources of uncertainty exist in formal verification?

- Non-determinism in tasks' WCET (undisturbed execution)
 - Input data and internal state controls execution paths
 - Memory access patterns control delays in processor architecture (pipelines and cache memories)
- Non-determinism in tasks' execution interference (pseudo-parallel execution)
 - Run-time execution model controls interference pattern
- Conflicts in tasks' demands for shared resources
 - (Pseudo-)parallel task execution may give rise to uncontrolled blocking of shared hardware and software resources

CHALMERS UNIVERSITY OF TECHNOLOGY | UNIVERSITY OF GOTHENBURG

Verification

How do we simplify formal verification?

- Concurrent and reactive programming paradigm
 - Suitable schedulable unit of concurrency (task, thread, ...)
 - Language constructs for expressing application constraints for schedulable unit (priorities, delays, ...)
 - WCET must be derivable for schedulable unit (special caution with usage of dynamic language constructs)
- Deterministic task execution
 - Time tables or static/dynamic task priorities
 - Preemptive task execution
 - Run-time protocols for access to shared resources (dynamic priority adjustment and non-preemptable code sections)

CHALMERS UNIVERSITY OF TECHNOLOGY | UNIVERSITY OF GOTHENBURG

Verification

How do we perform schedulability analysis?

- Introduce abstract models of system components:
 - Task model (computation requirements, timing constraints)
 - Processor model (resource capacities)
 - Run-time model (task states, dispatching)
- Predict whether task executions will meet constraints
 - Use abstract system models
 - Make sure that computation requirements never exceed resource capacities
 - Generate (partly or completely) run-time schedule resulting from task executions and detect worst-case scenarios

CHALMERS UNIVERSITY OF TECHNOLOGY | UNIVERSITY OF GOTHENBURG

Task model

Implementation

→

Abstract model

```

void task1(Object *self, int p) {
    Action1();
    SEND(Period1, Deadline1, self, task1, p);
}

void task2(Object *self, int p) {
    Action2();
    SEND(Period2, Deadline2, self, task2, p);
}

void kickoff(Object *self, int p) {
    AFTER(Offset1, app1, p);
    AFTER(Offset2, app2, p);
}

main() {
    TINYTIMBER(app_main, kickoff, 0);
}
    
```

τ_1

$\tau_1 = \{C_1, T_1, D_1, O_1\}$

→

τ_2

$\tau_2 = \{C_2, T_2, D_2, O_2\}$

CHALMERS UNIVERSITY OF TECHNOLOGY | UNIVERSITY OF GOTHENBURG

Task model

The task model expresses the timing behavior of a task:

- The static parameters describe characteristics of a task that apply independent of other tasks.
 - Derived from the specification or implementation of the system
 - For example: period, deadline, WCET
- The dynamic parameters describe effects that occur during the execution of the task.
 - Is a function of the run-time system and the characteristics of other tasks
 - For example: start time, completion time, response time

CHALMERS UNIVERSITY OF TECHNOLOGY | UNIVERSITY OF GOTHENBURG

Task model

Static task parameters:

C_i : (undisturbed) WCET
 T_i : period
 D_i : (relative) deadline
 O_i : (absolute) time offset

$\tau_i = \{C_i, T_i, D_i, O_i\}$

CHALMERS UNIVERSITY OF TECHNOLOGY | UNIVERSITY OF GOTHENBURG

Task model

Dynamic task parameters:

$a_{i,k}$: arrival time of k^{th} instance
 $s_{i,k}$: start time of k^{th} instance
 $f_{i,k}$: completion time of k^{th} instance
 $R_{i,k}$: response time of k^{th} instance

$\tau_i = \{C_i, T_i, D_i, O_i\}$

$\tau_{i,k}$: the k^{th} instance of τ_i

$a_{i,k} = O_i + (k-1) \cdot T_i$ $R_{i,k} = f_{i,k} - a_{i,k}$

$R_i = \max_{\tau_i \in \mathbf{T}, k \geq 1} \{R_{i,k}\}$ (worst-case response time)

CHALMERS UNIVERSITY OF TECHNOLOGY | UNIVERSITY OF GOTHENBURG

Task model

Different types of tasks:

- Periodic tasks
 - A periodic task arrives with a time interval T_i
- Sporadic tasks
 - A sporadic task arrives with a time interval $\geq T_i$
- Aperiodic tasks
 - An aperiodic task has no guaranteed minimum time between two subsequent arrivals

\Rightarrow A *priori* schedulable real-time systems can only contain periodic and sporadic tasks.

CHALMERS UNIVERSITY OF TECHNOLOGY | UNIVERSITY OF GOTHENBURG

Processor model

Homogeneous processors:

- Identical processors
 - WCET is a constant

Heterogeneous processors:

- Uniform processors
 - WCET is the product of a basic execution time and a scaling factor
- Unrelated processors
 - WCET is not related for different processors

CHALMERS UNIVERSITY OF TECHNOLOGY | UNIVERSITY OF GOTHENBURG

Run-time model

Task states:

- Waiting
 - Task has not yet arrived for the first time, or has finished executing but not re-arrived
- Ready
 - Task has arrived and can potentially execute on the processor (kept waiting in a **ready queue**)
- Running
 - Task is currently executing on the processor

Dispatcher:

- A run-time mechanism that takes the first element (task) in the ready queue and executes it on the processor.

CHALMERS UNIVERSITY OF TECHNOLOGY | UNIVERSITY OF GOTHENBURG

Evaluating a real-time system

An important part of real-time system design is to have techniques that generate good schedules.

Is this a good schedule? What do we need to decide the quality?

CHALMERS UNIVERSITY OF TECHNOLOGY | UNIVERSITY OF GOTHENBURG

Evaluating a real-time system

How do we measure and compare performance?

- Quantify system performance
 - Choose useful performance measures (metrics)
- Perform objective performance analysis
 - Choose suitable evaluation methodology
 - Examples: theoretical and/or experimental analysis
- Compare performance of different designs
 - Make trade-off analysis using chosen performance measures
- Identify fundamental performance limitations
 - Find “bottleneck” mechanisms that affect performance

CHALMERS UNIVERSITY OF TECHNOLOGY | UNIVERSITY OF GOTHENBURG

Performance measures

“Yardsticks” by which the performance of a system is expressed.

Why do we need it?

- To objective evaluate different design solutions and choose the “best” one
- To rubberstamp a system with performance potential or quality guarantees (cf. “Intel inside”, “ISO 9000”)

CHALMERS UNIVERSITY OF TECHNOLOGY | UNIVERSITY OF GOTHENBURG

Performance measures

What is required by a performance measure?

- Must be concise to avoid ambiguity
 - preferably a single number
 - use a weighted sum of constituent local performance measures
 - should reflect user-perceived utility
 - no artificial measures should be used
 - some measures are contradictory
 - processing speed vs. power consumption in a handheld computer
 - some measures are misleading
 - MIPS (million instructions executed per second)

CHALMERS UNIVERSITY OF TECHNOLOGY | UNIVERSITY OF GOTHENBURG

Performance measures

What is required by a performance measure?

- Must provide efficient coding of information
 - determine relevance of individual pieces
- Must provide objective basis for ranking
 - use same set of applications for evaluations
- Must provide objective optimization criteria for design
 - identify application-sensitive criteria
- Must provide verifiable facts
 - use measures that can be derived for a real system

CHALMERS UNIVERSITY OF TECHNOLOGY | UNIVERSITY OF GOTHENBURG

Performance measures

Traditional performance measures:

Throughput
Average # of operations/data processed by system per time unit

Reliability
Probability that system will not fail in a given time interval

Availability
Fraction of time for which system is up (providing service)

These measures do not take deadlines into account!

CHALMERS UNIVERSITY OF TECHNOLOGY | UNIVERSITY OF GOTHENBURG

Performance measures

Suitable real-time performance measures:

Laxity $X = \min_{\tau_i \in \mathbf{T}} \{D_i - C_i\}$
 Amount of time that the start of a task can be delayed without it missing its deadline (calculated before scheduling)

Lateness $L = \max_{\tau_i \in \mathbf{T}} \{R_i - D_i\}$
 Amount of time by which a task completes after its deadline (calculated after scheduling)

Successful tasks $N_{\text{success}} = |\{\tau_i \in \mathbf{T} : R_i - D_i \leq 0\}|$
 Number of tasks that complete on or before their deadline (calculated after scheduling)

Jitter $J_{\text{output}} = \max_{\tau_i \in \mathbf{T}, k \geq 1} \left\{ |(f_{i,k+1} - f_{i,k}) - T_i| \right\}$
 Amount of deviation from expected periodicity of a task's completion (calculated after scheduling)

CHALMERS UNIVERSITY OF TECHNOLOGY | UNIVERSITY OF GOTHENBURG

Performance measures

Cost function – a general real-time performance measure

Cumulative value: $C = \sum_{\tau_i \in \mathbf{T}} v(f_i)$
 Value associated with a task as a function of its completion time

Non real-time

CHALMERS UNIVERSITY OF TECHNOLOGY | UNIVERSITY OF GOTHENBURG

Performance measures

Cost function – a general real-time performance measure

Cumulative value: $C = \sum_{\tau_i \in \mathbf{T}} v(f_i)$
 Value associated with a task as a function of its completion time

Soft real-time

CHALMERS UNIVERSITY OF TECHNOLOGY | UNIVERSITY OF GOTHENBURG

Performance measures

Cost function – a general real-time performance measure

Cumulative value: $C = \sum_{\tau_i \in \mathbf{T}} v(f_i)$
 Value associated with a task as a function of its completion time

Hard real-time