Course on Computer Communication and Networks

Lecture 14
Summary - flashback

EDA344/DIT 420, CTH/GU

Based on the book Computer Networking: A Top Down Approach, Jim Kurose, Keith Ross, Addison-Wesley.

Important for the exam

When/where: wednesday March 18, 14.00-18.00, M

You may have with you:

- English-X dictionary
- no calculators, PDAs, etc (if/where numbers matter, do rounding)

Grading

- 30-40, 41-50, 51-60 (out of 60)= 3, 4, 5 (CTH)
- 30-44, 45-60 (out of 60) = G, VG (GU)

To think during summary-study

Overview; critical eye; explain, ask yourselves: why is this so? / How does it work?

Flashback

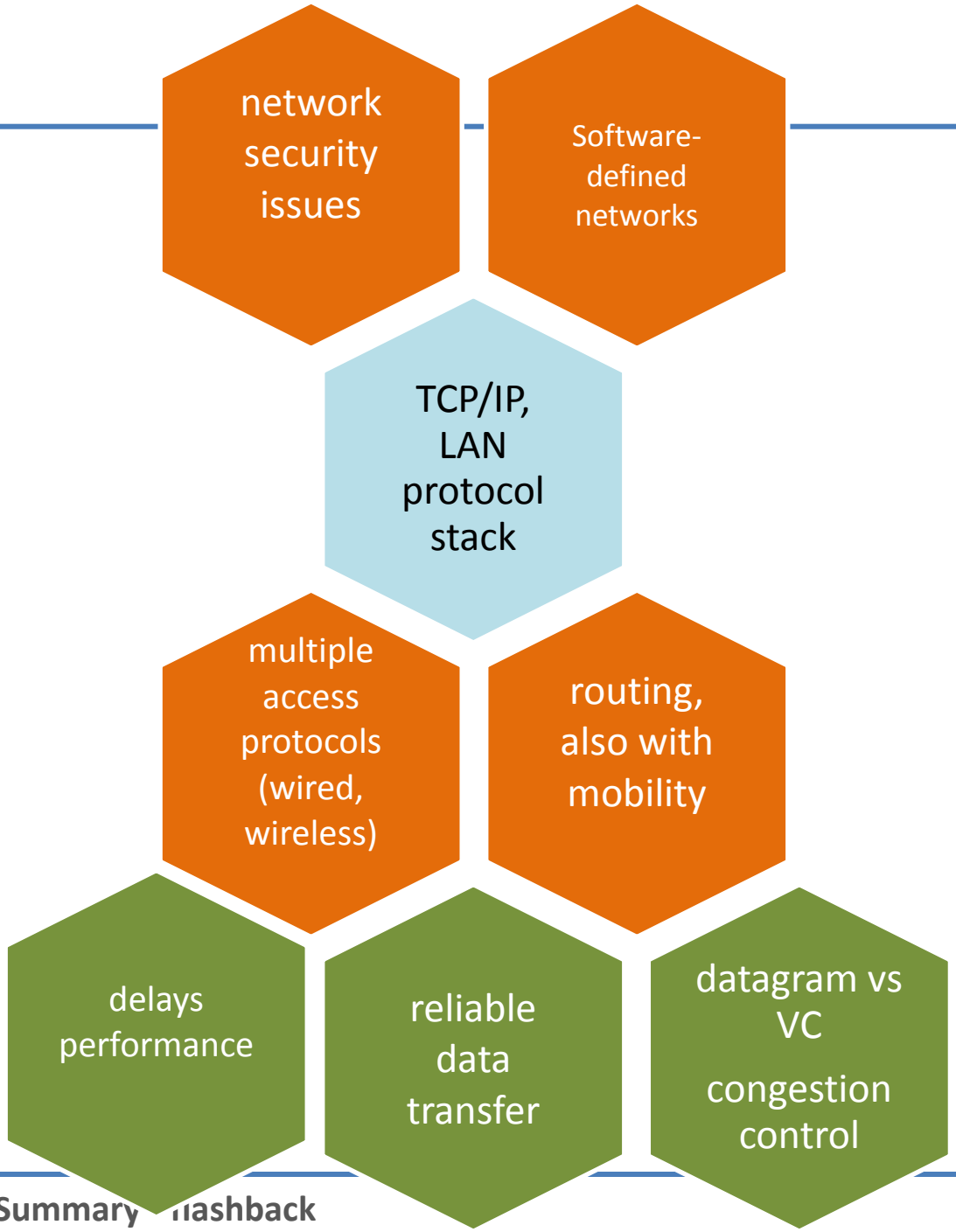
Principles, Organisation

Network Problems (in the order faced in the 1st intro):

- producer-consumer problems, flow and error control,
- manage access to shared (broadcast) transmission media ,
- routing,
- congestion,
- connecting transparently different networks,
- serving different types of traffic,
- performance,
- mobility
- security

Layering : principle, why

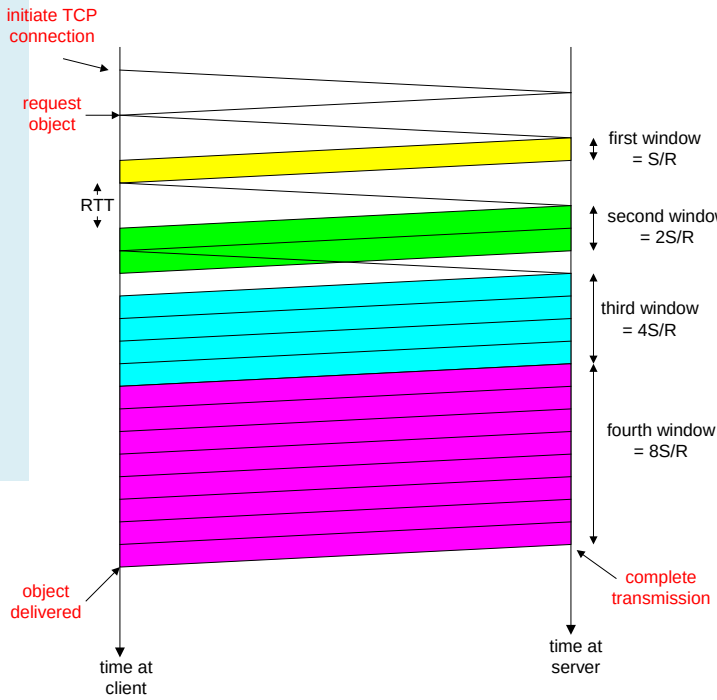
Highlights

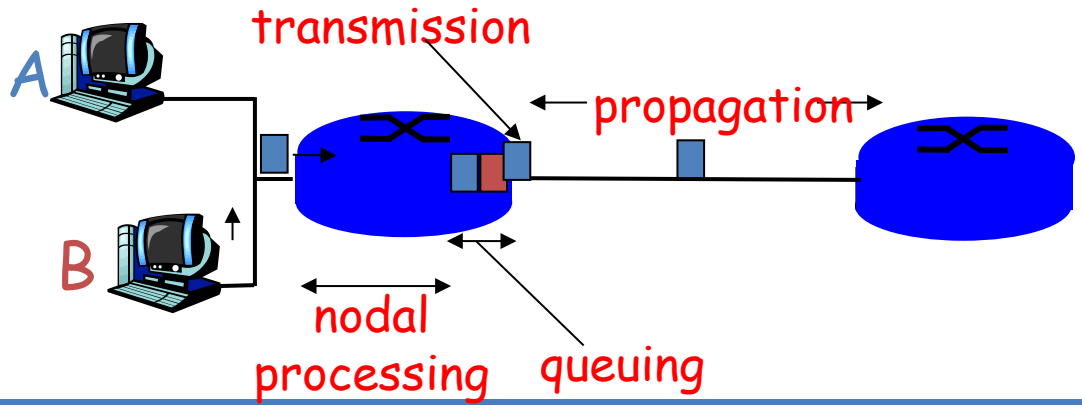

network security issues
Software-defined networks
TCP/IP, LAN protocol stack
multiple access protocols (wired, wireless)
routing, also with mobility
delays performance
reliable data transfer
datagram vs VC
congestion control

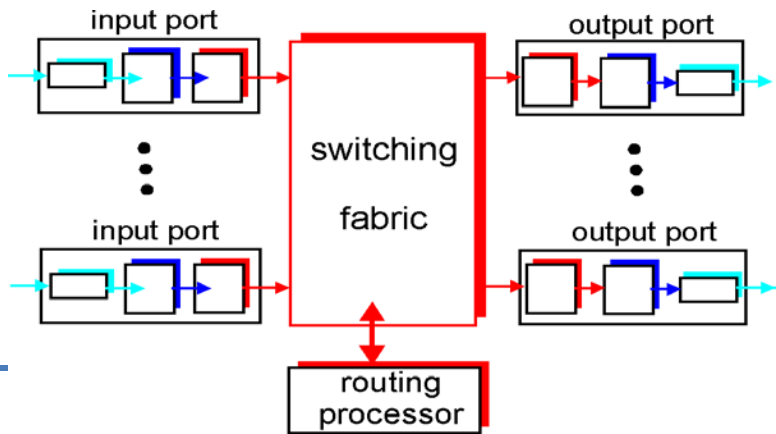
Types of delay; performance

delays performance

- Propagation, transmission, queueing, processing
- Throughput -- effective bandwidth
- Utilization -- efficiency
- Packet-switching: impact of store&forward
- TCP's slow start
- Sliding windows performance





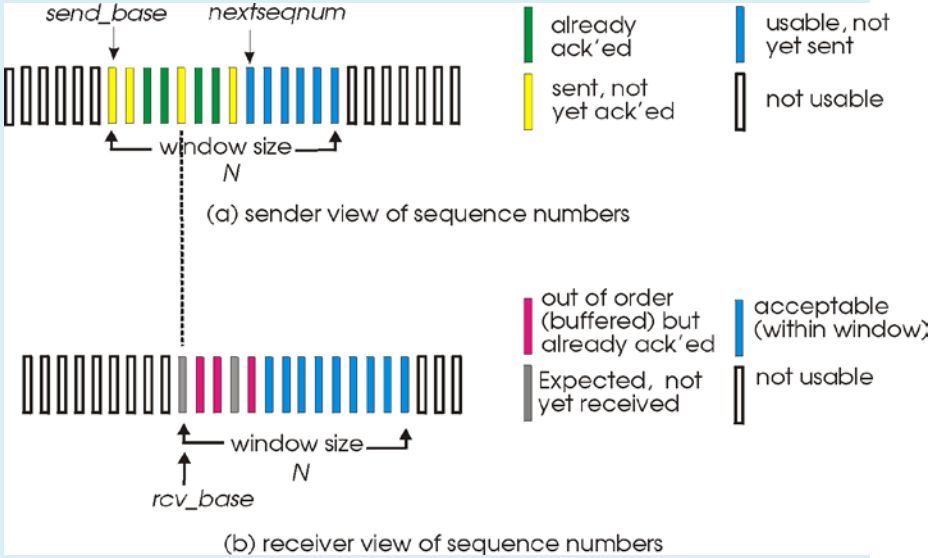


Reliable data transfer

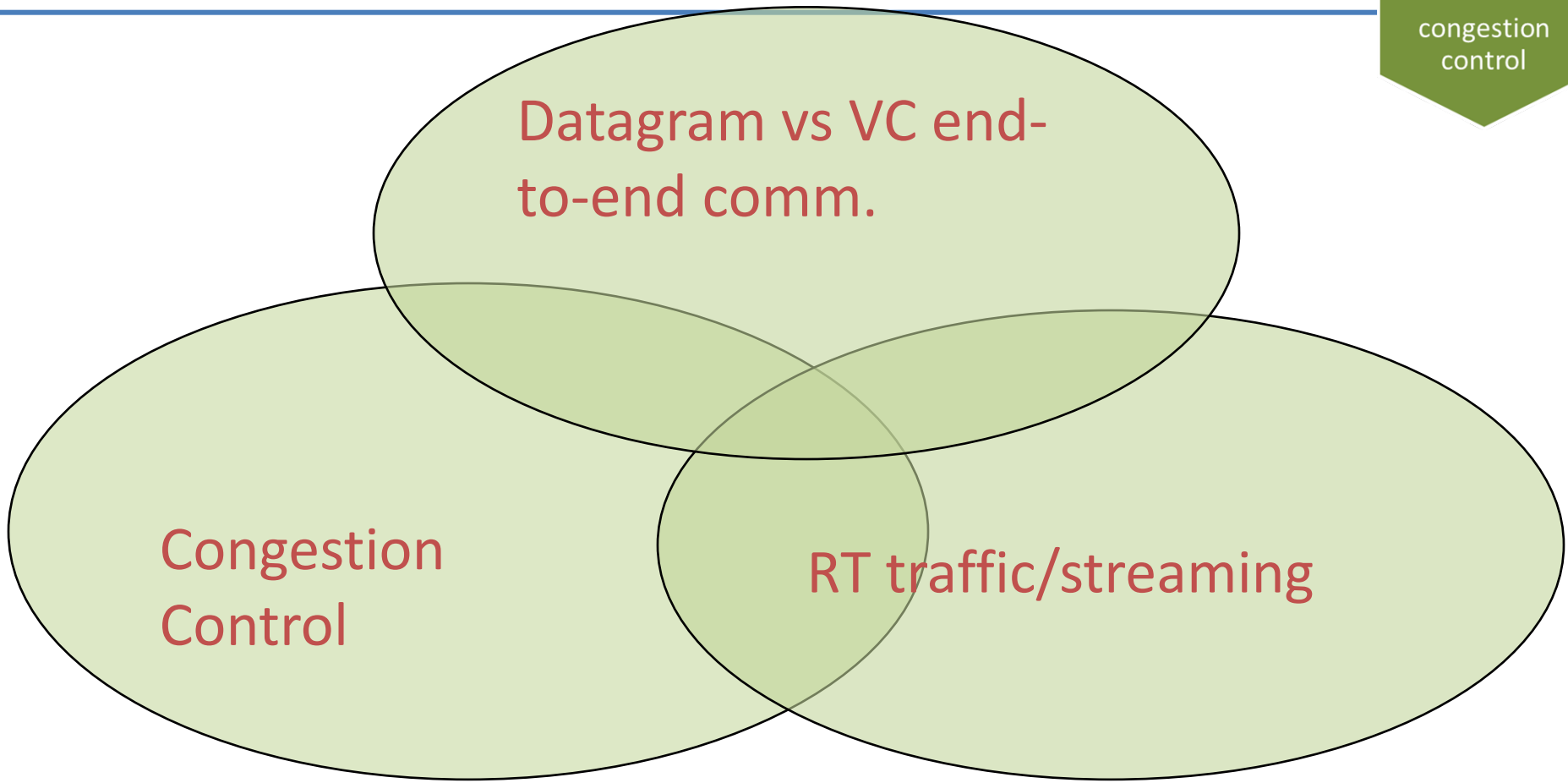
reliable data transfer

Guaranteed, in-order, correct delivery:

- stop&wait
- sliding windows
- sequence numbers
- window sizes
- dynamic windows (TCP)
- performance
- **Flow control**

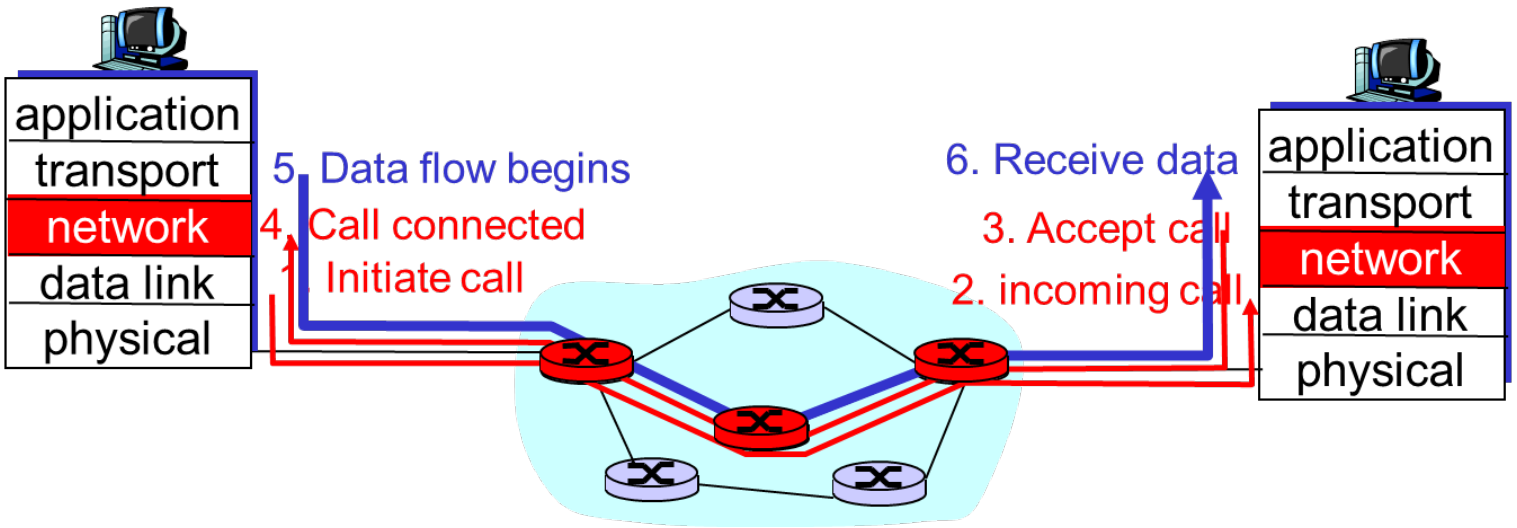


- **Error detection**: checksums, CRC
- **Error control**: go-back-n, selective repeat, FEC methods


Datagram vs VC end-to-end comm.
Congestion Control
RT traffic/streaming

Datagram vs VC end-to-end communication

- Conceptual differences
- Decisions, comparison

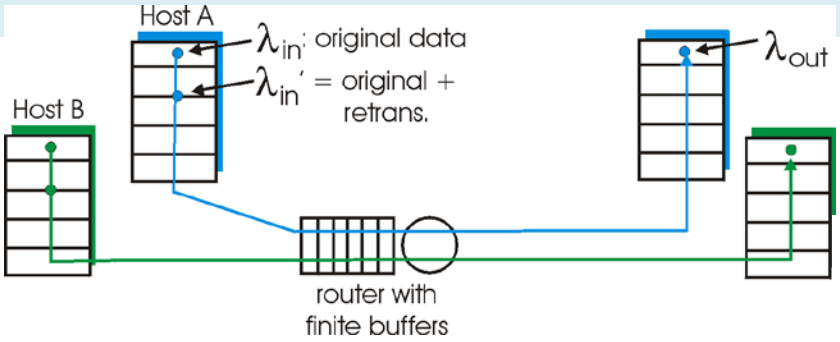


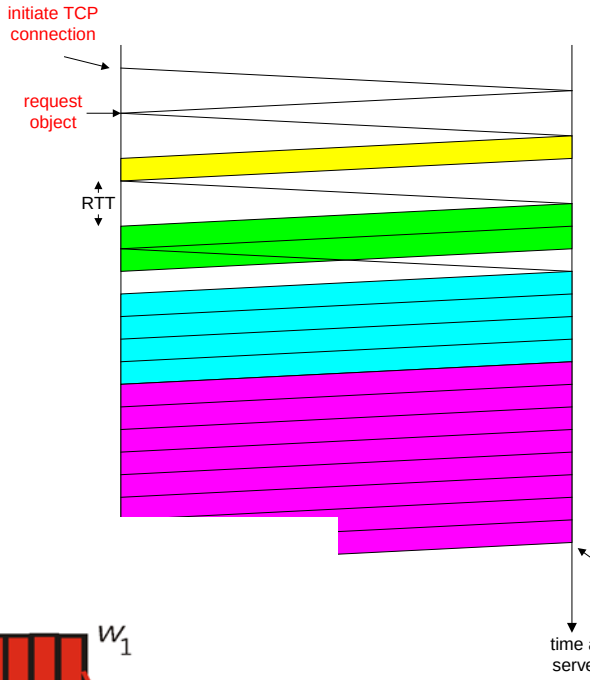
Congestion control (CC)

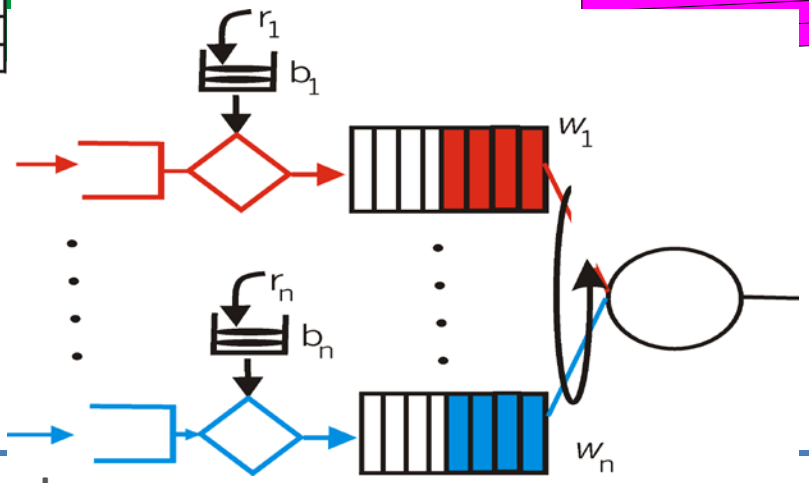
datagram vs VC

congestion control

- why, how congestion occurs
- CC in TCP and performance; implied weaknesses
- CC in other ways, e.g. VC-based networks
  - Real-time (RT)-traffic resource reservation: traffic shaping and policing
  - rate-based







RT/streaming traffic

datagram vs VC

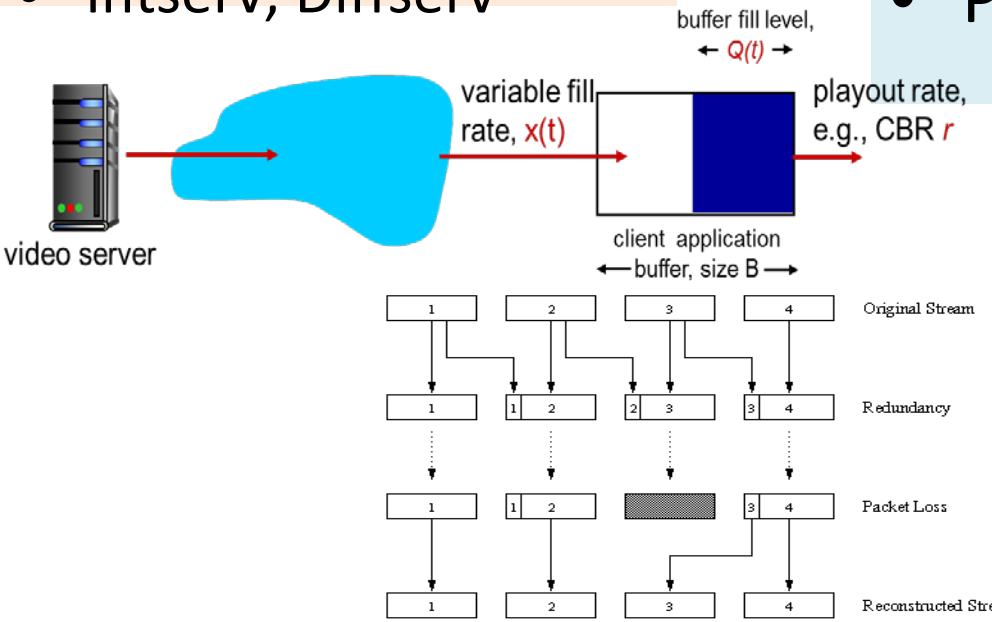
congestion control

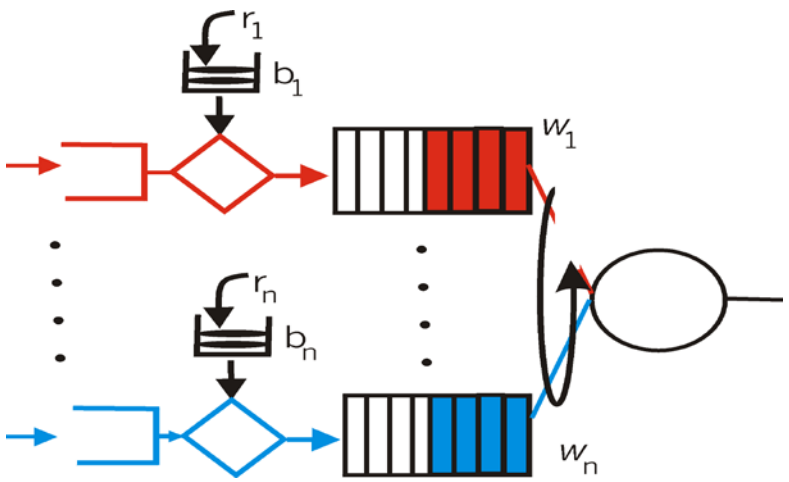
Internet context

- Application-level solutions (FEC, playout delay, caching-CDN)
- Intserv, Diffserv

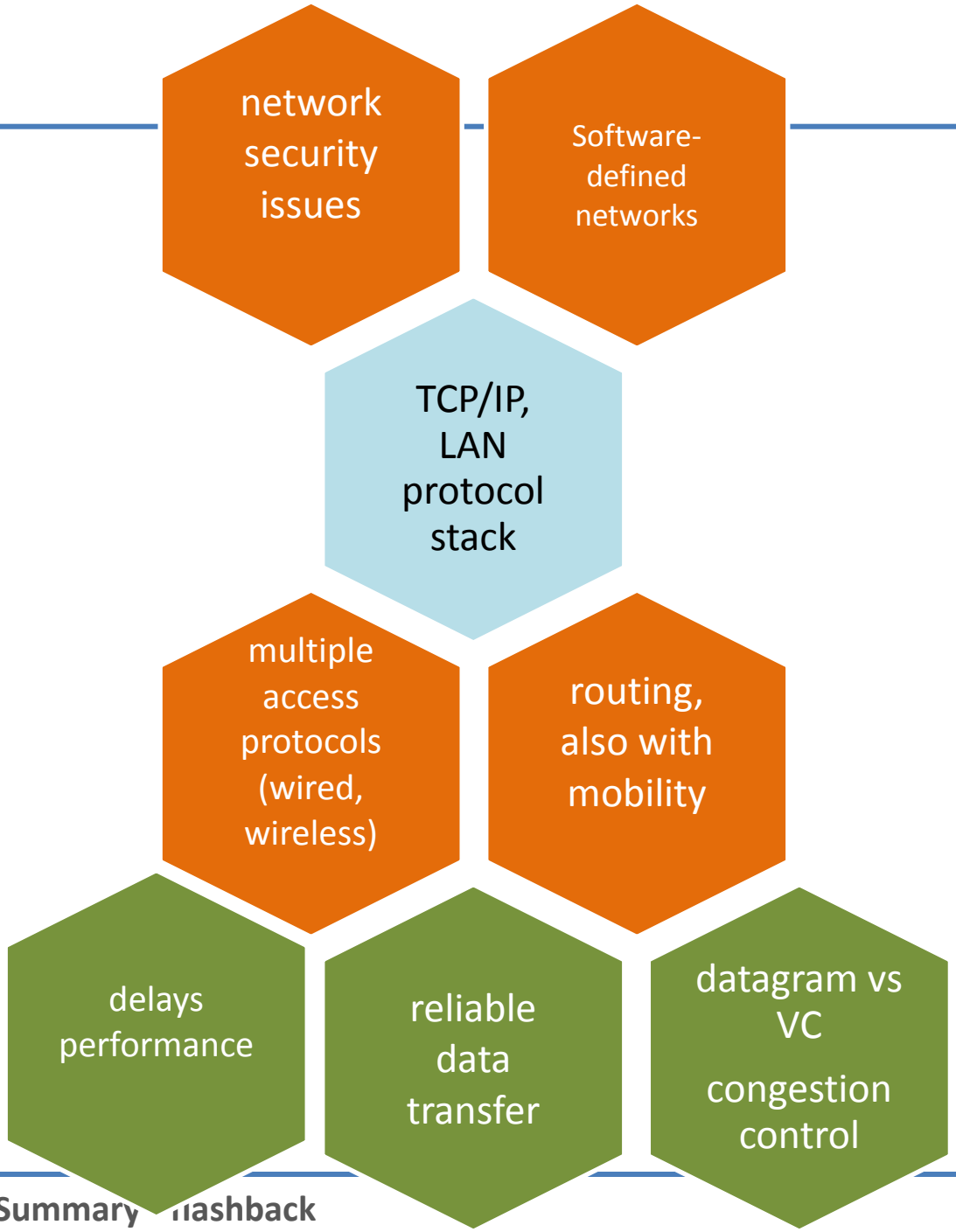
Conceptual needs:

- packet/flow marking
- Admission control
- Traffic shaping & policing
- Packet scheduling





Highlights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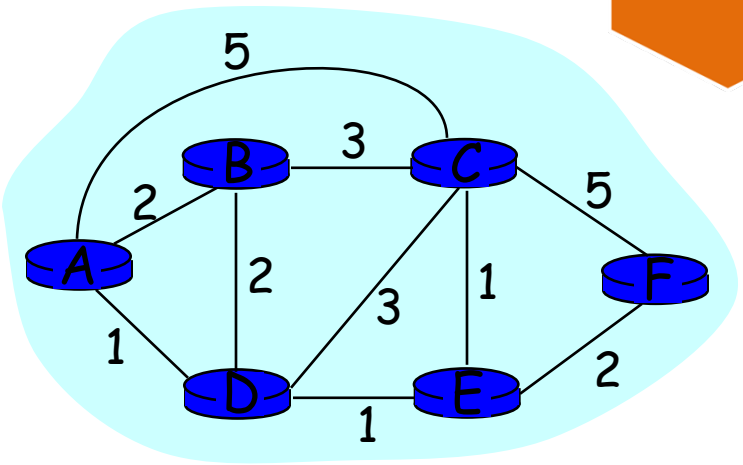
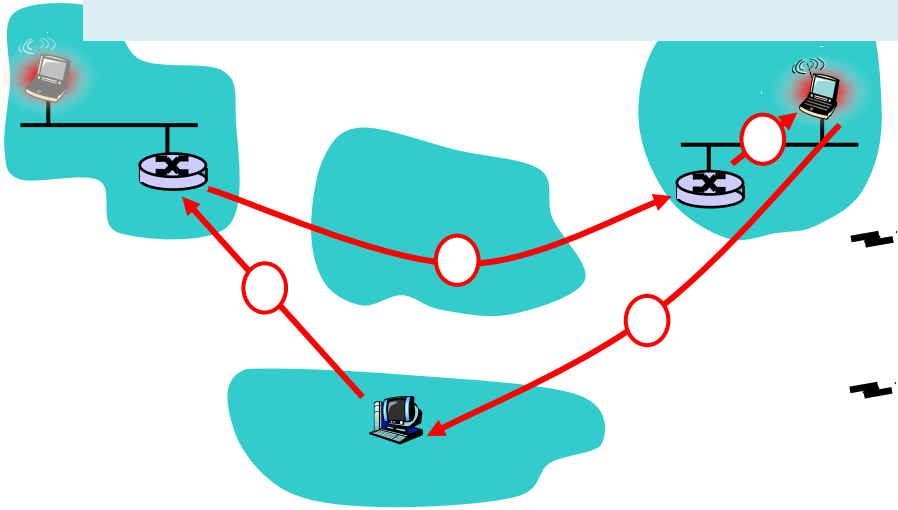
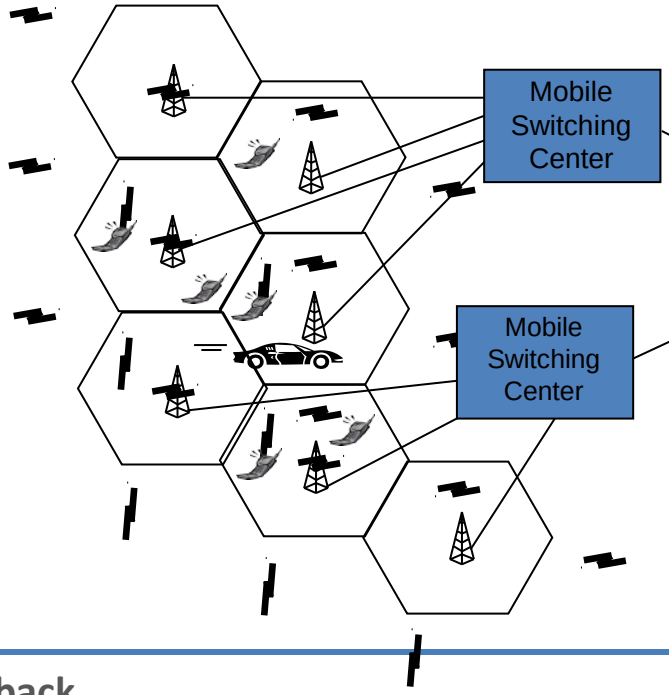
Routing, also with mobility

routing, also with mobility

- Routing algorithms, protocols
- Forwarding in routers
- Resource, policy issues
- Addressing mobility, tunneling







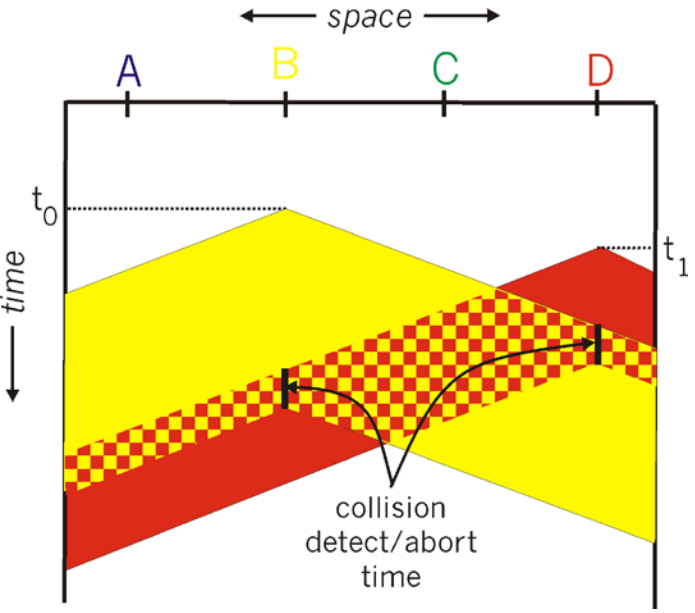
Complementary video links
- IP addresses and subnets
http://www.youtube.com/watch?v=ZTJIkjgyuZE&list=PLE9F3F05C381ED8E8&feature=plcp
- How does BGP choose its routes
http://www.youtube.com/watch?v=RGe0qt9Wz4U&feature=plcp

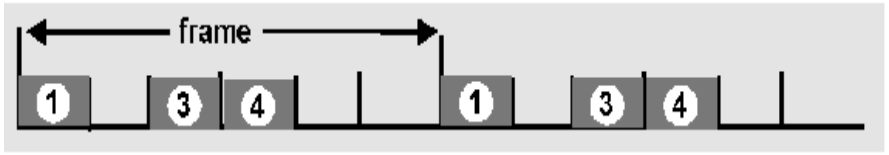
Medium access: multiple access methods

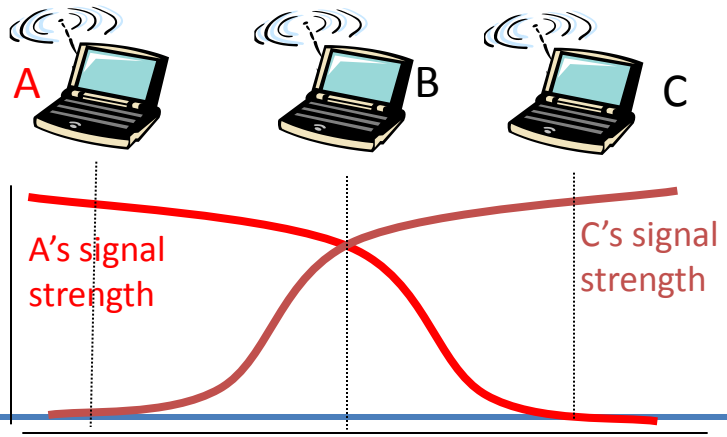
multiple access protocols (wired, wireless)

Strategies: (functionality, appropriateness)

- **Contention-based (random access), wired/wireless:**
  - Aloha, CSMA(CD/CA)
- **Collision-free:**
  - **Channel partitioning:** TDMA, FDMA, CDMA
  - **Taking turns:** e.g. tokens, reservation-based







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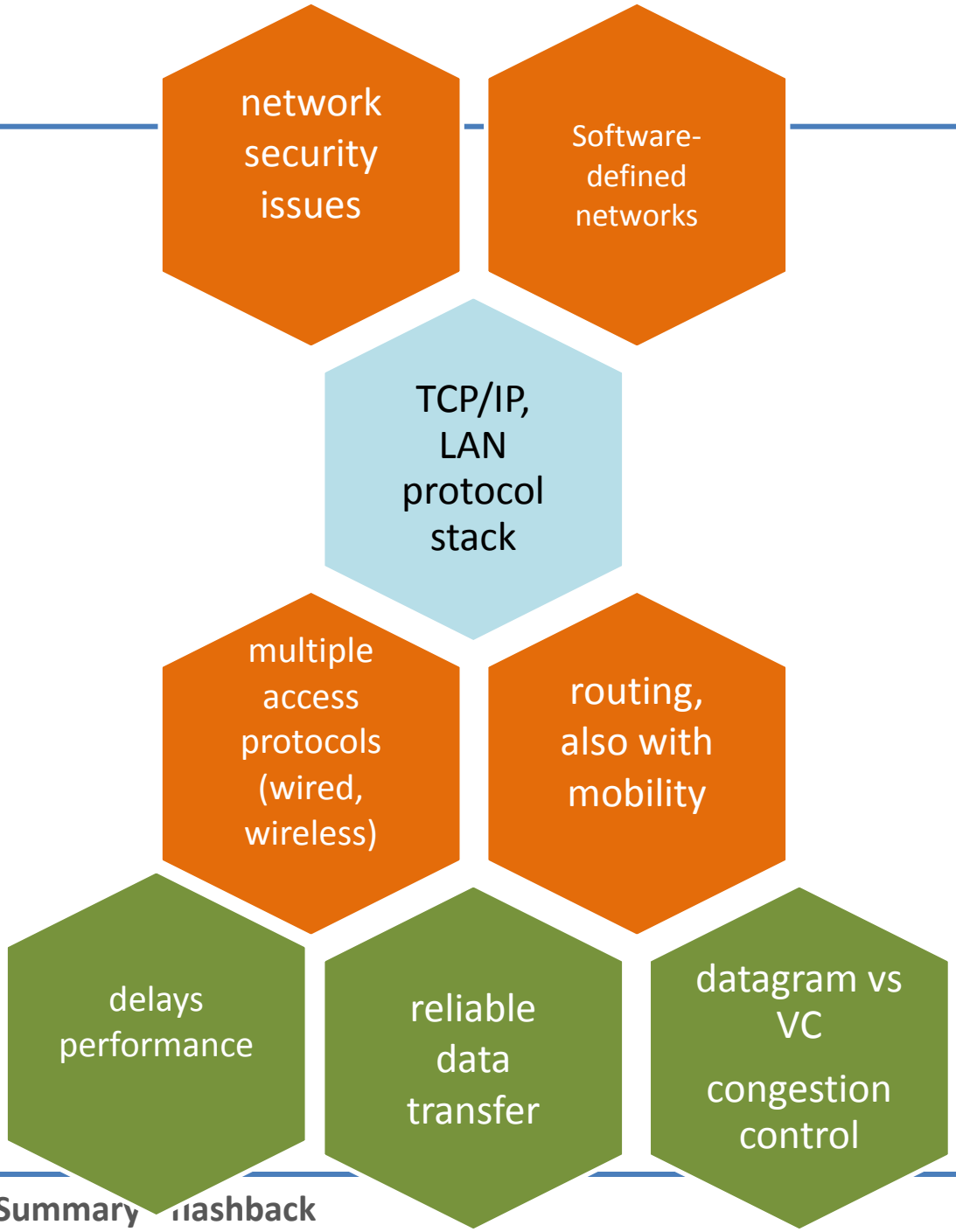
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LANs & related link technologies

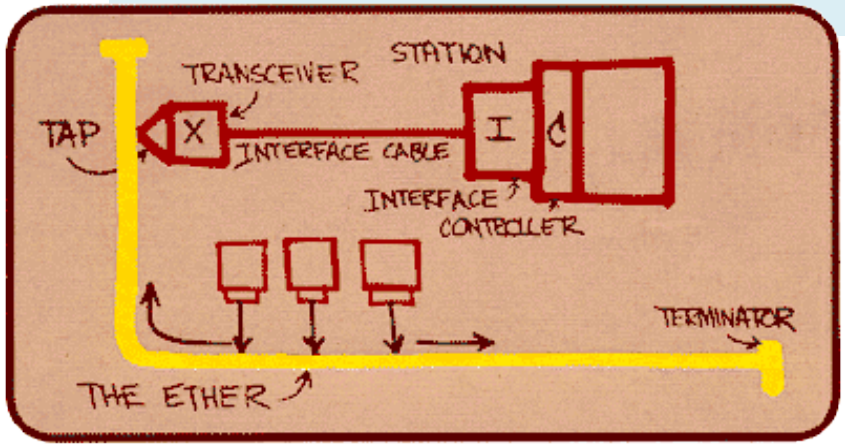
TCP/IP, LAN protocol stack

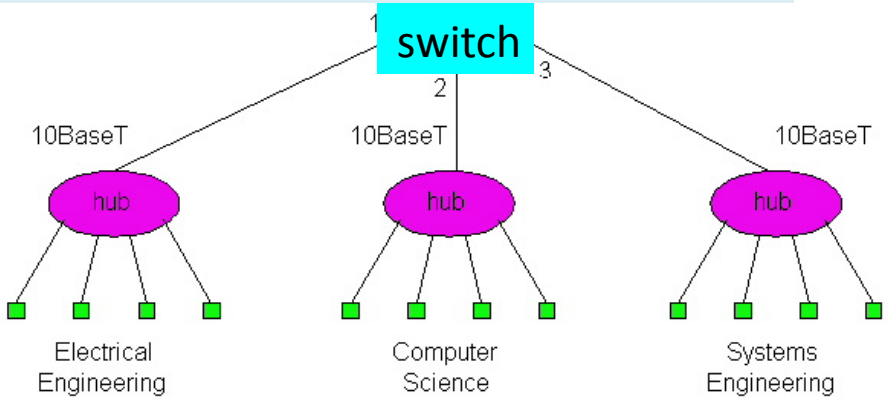
- **Protocol Examples: wired, wireless**

  Ethernet, 802.11 (+ 802.16 wimax), GSM:

  Functionality, performance under low/high load

- **Connecting devices;**
  - functionalities and differences (Hubs, switches)
  - Algorithms for switch-"routing": learning& forwarding of packets
- **ARP**





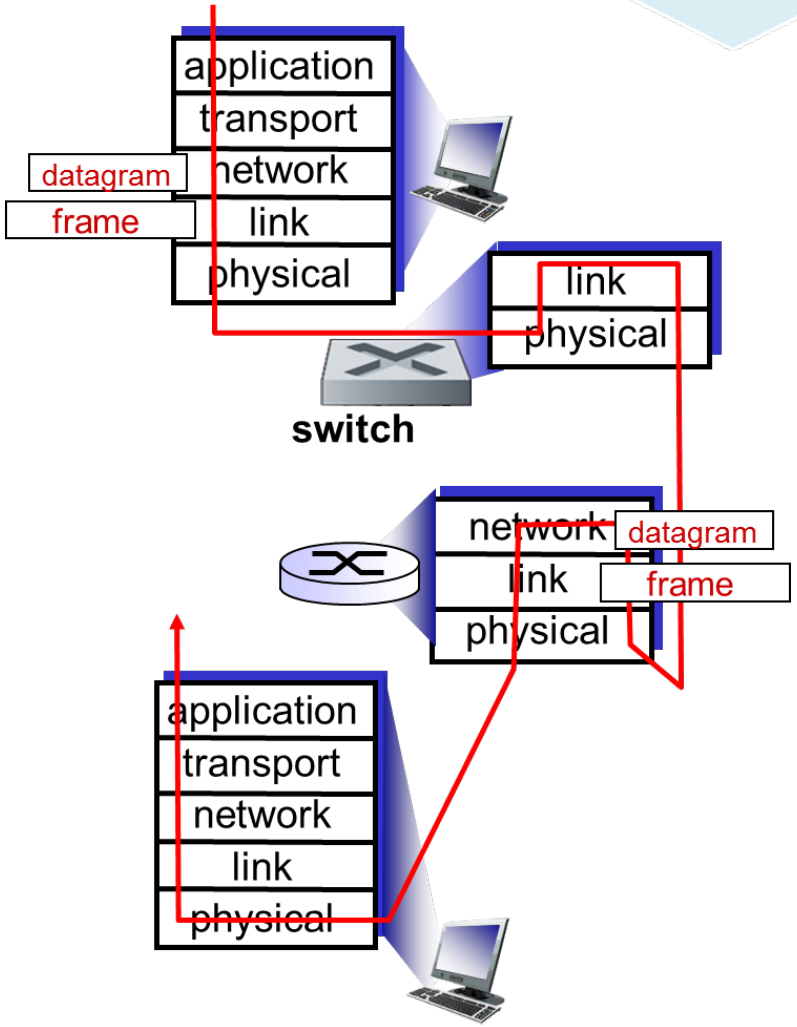
TCP/IP protocol stack, applications, evolution

TCP/IP, LAN protocol stack

- Instantiation of network- solutions

(Routing, Congestion Control, Flow & error control, applications, link layer technologies)

- Advantages, limitations, updates
- New types of applications and how they function given the existing state of Internet: multimedia/streaming applications, CDNs, P2P applications, overlays



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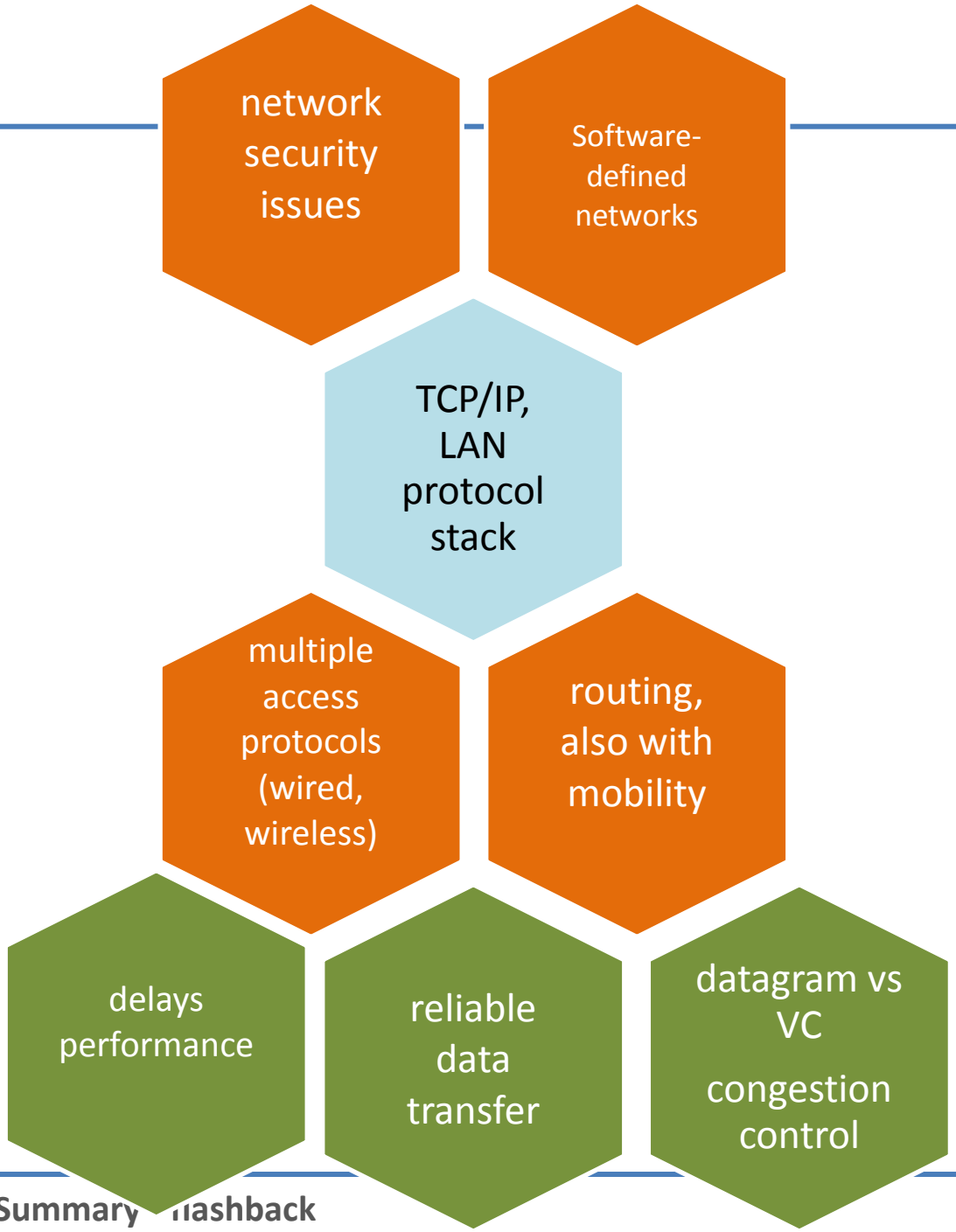
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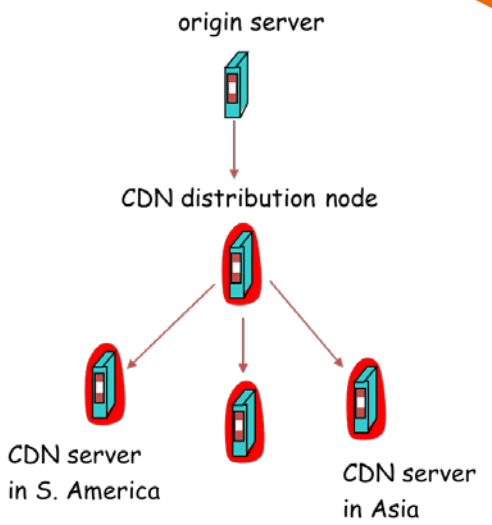
Overlays, software-defined networks

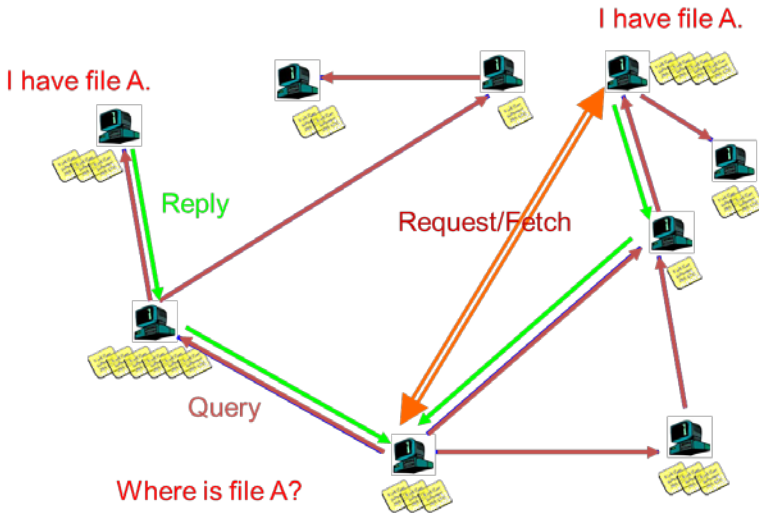
Software-defined networks

- P2P applications
- multimedia/streaming application-infrastructure

… complement the networking infrastructure …

…taking advantage of the network resources at the edge of the network…

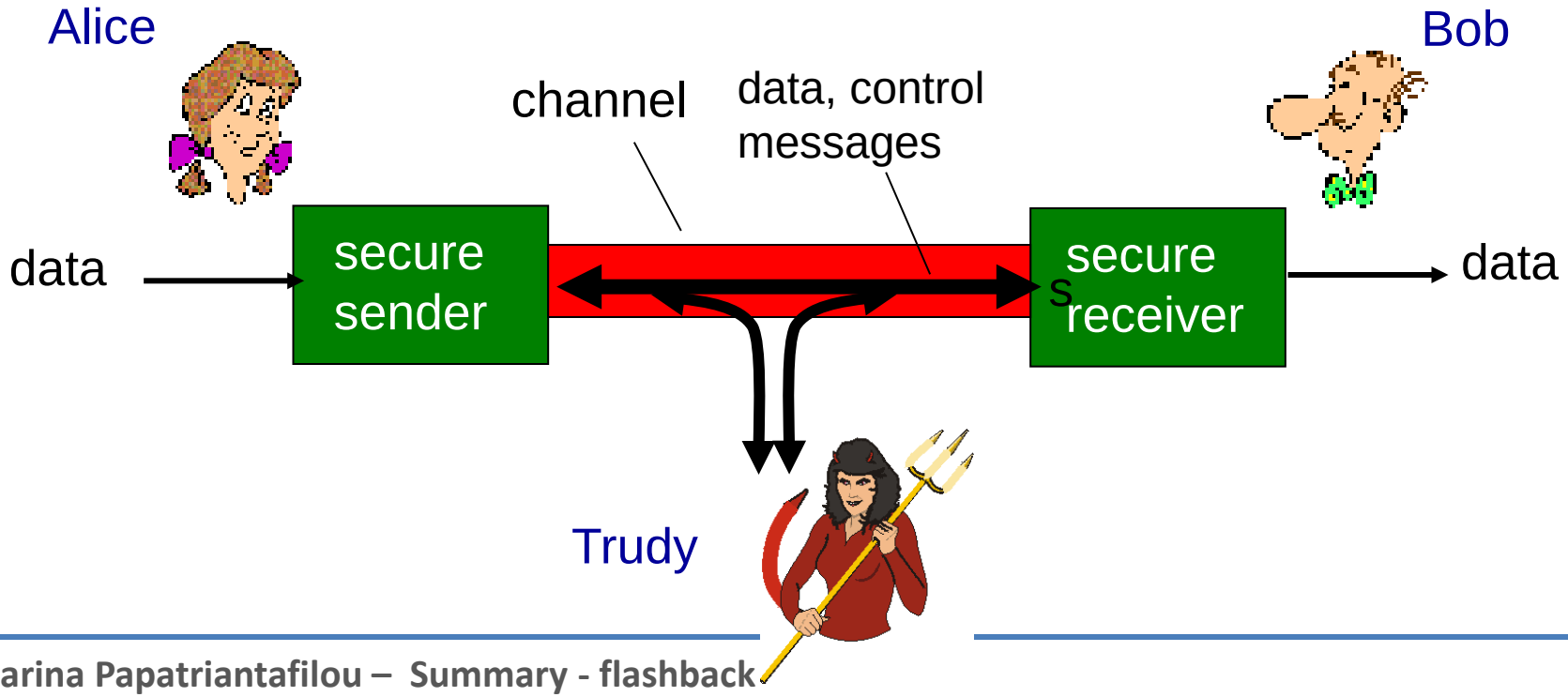




Security issues

network security issues

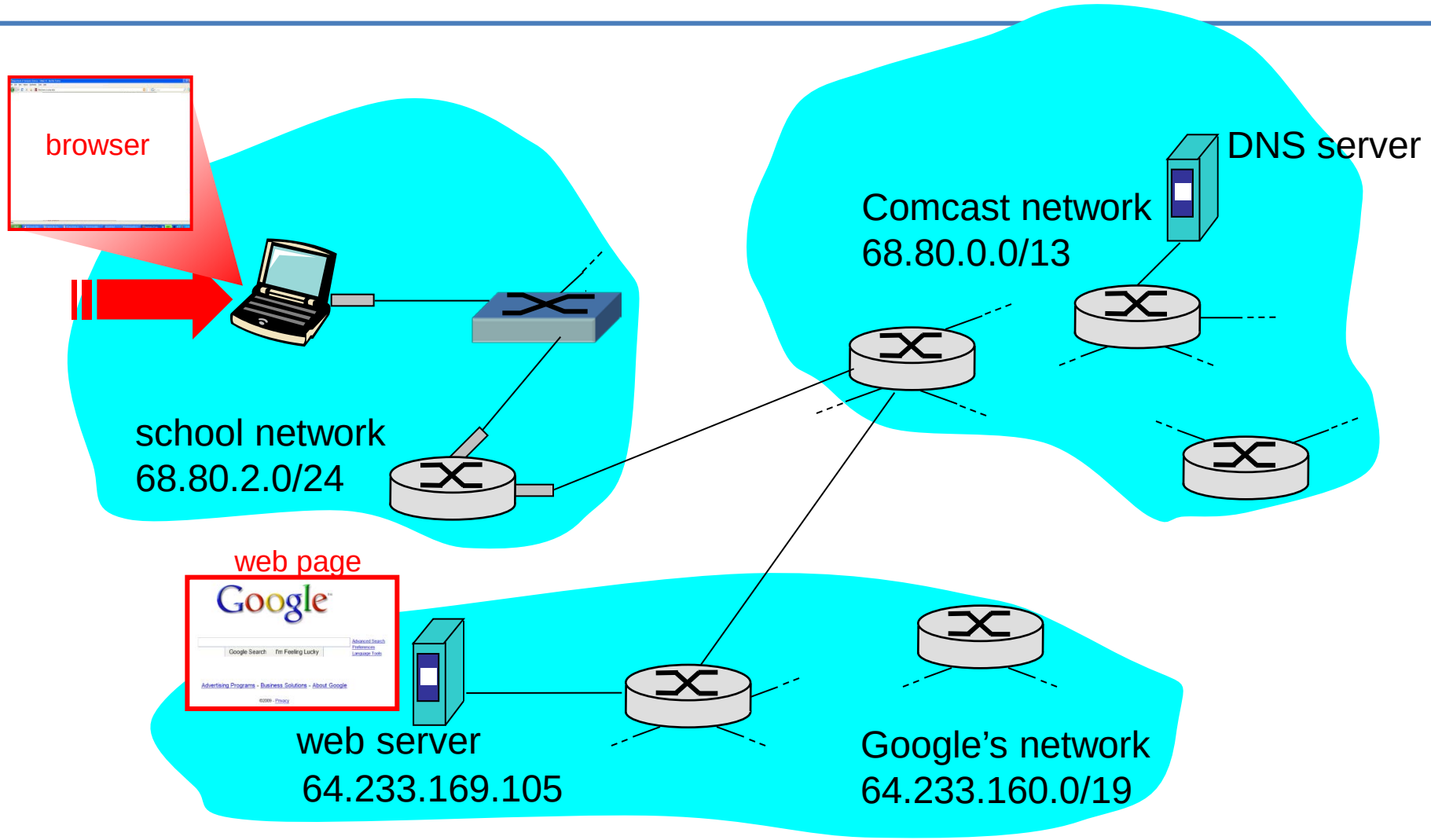
- **C, I, A** and methods to achieve them
  - Threats
  - The language of cryptography
  - Message integrity, signatures
- Instantiation in Internet: SSL, IPsec



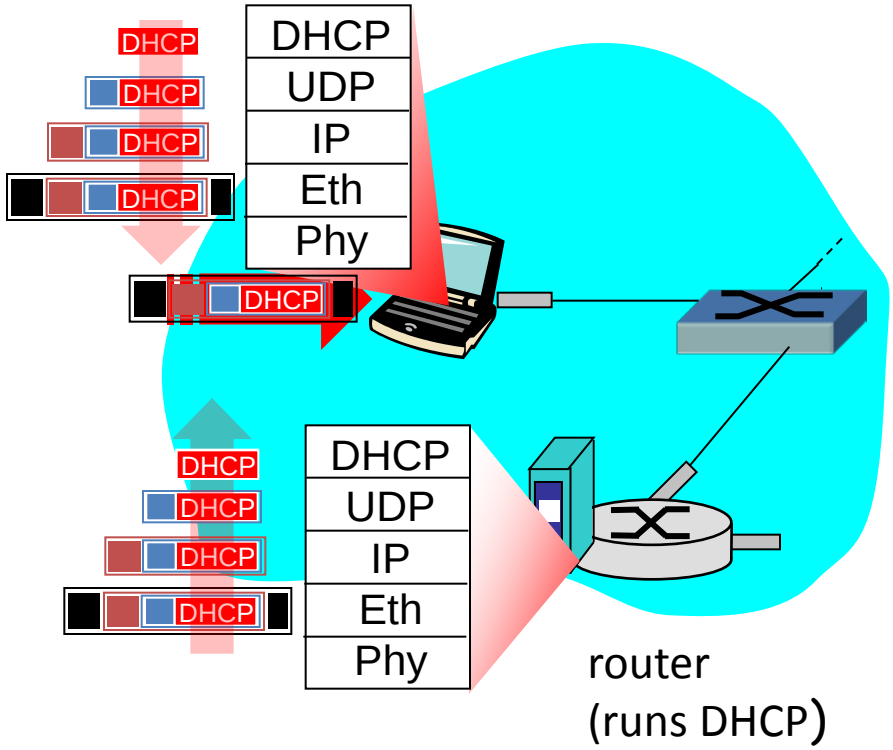
Synthesis: a day in the life of a web request

- putting-it-all-together: synthesis!
  - *goal:* identify, review protocols (at all layers) involved in seemingly simple scenario: requesting www page
  - *scenario:* student attaches laptop to campus network, requests/receives www.google.com

A day in the life …. : scenario



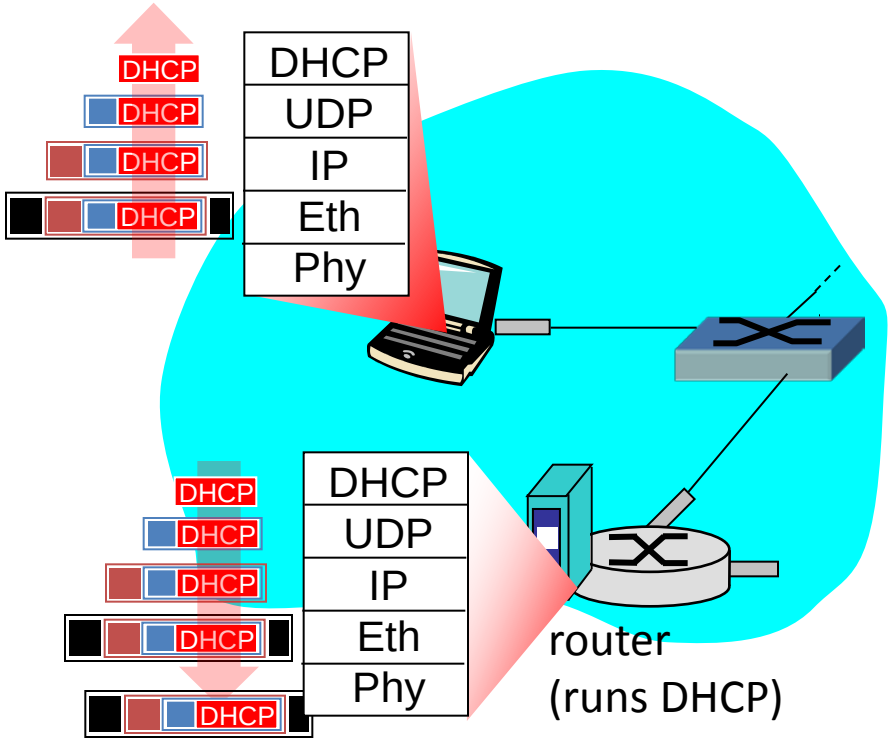
A day in the life… connecting to the Internet



connecting laptop needs to get its own IP address: use ***DHCP***

- DHCP request encapsulated in UDP, encapsulated in IP, encapsulated in Ethernet
- Ethernet frame broadcast (dest: FFFFFFFFFFFF) on LAN, received at router running DHCP server
- Ethernet demux'ed to IP demux'ed to UDP demux'ed to DHCP

A day in the life… connecting to the Internet

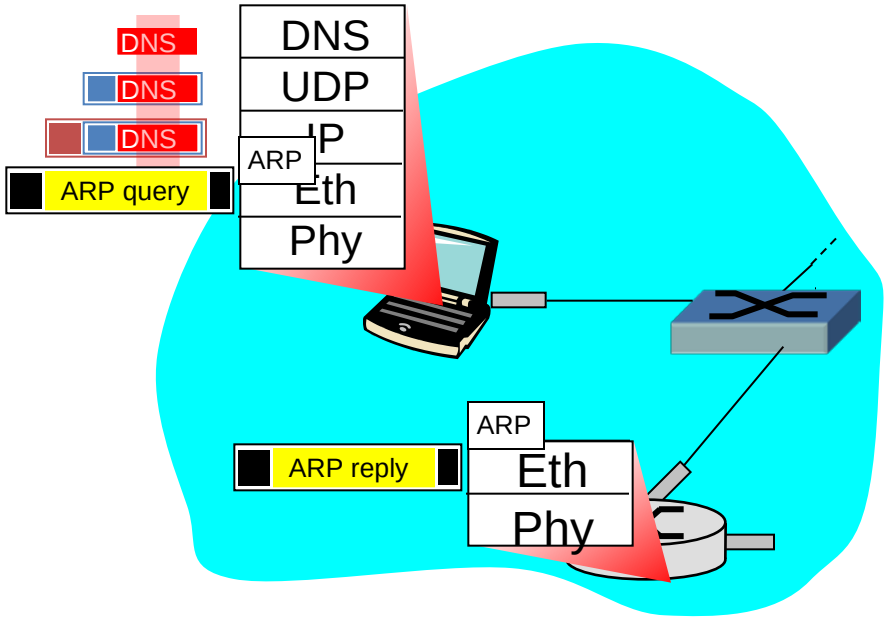


DHCP server formulates ***DHCP ACK*** containing client's IP address (**and also IP address of first-hop router for client, name & IP address of DNS server**)

- frame forwarded (switch learning) through LAN, demultiplexing at client
- DHCP client receives DHCP ACK reply

Client now has IP address, knows name & addr of DNS server, IP address of its first-hop router

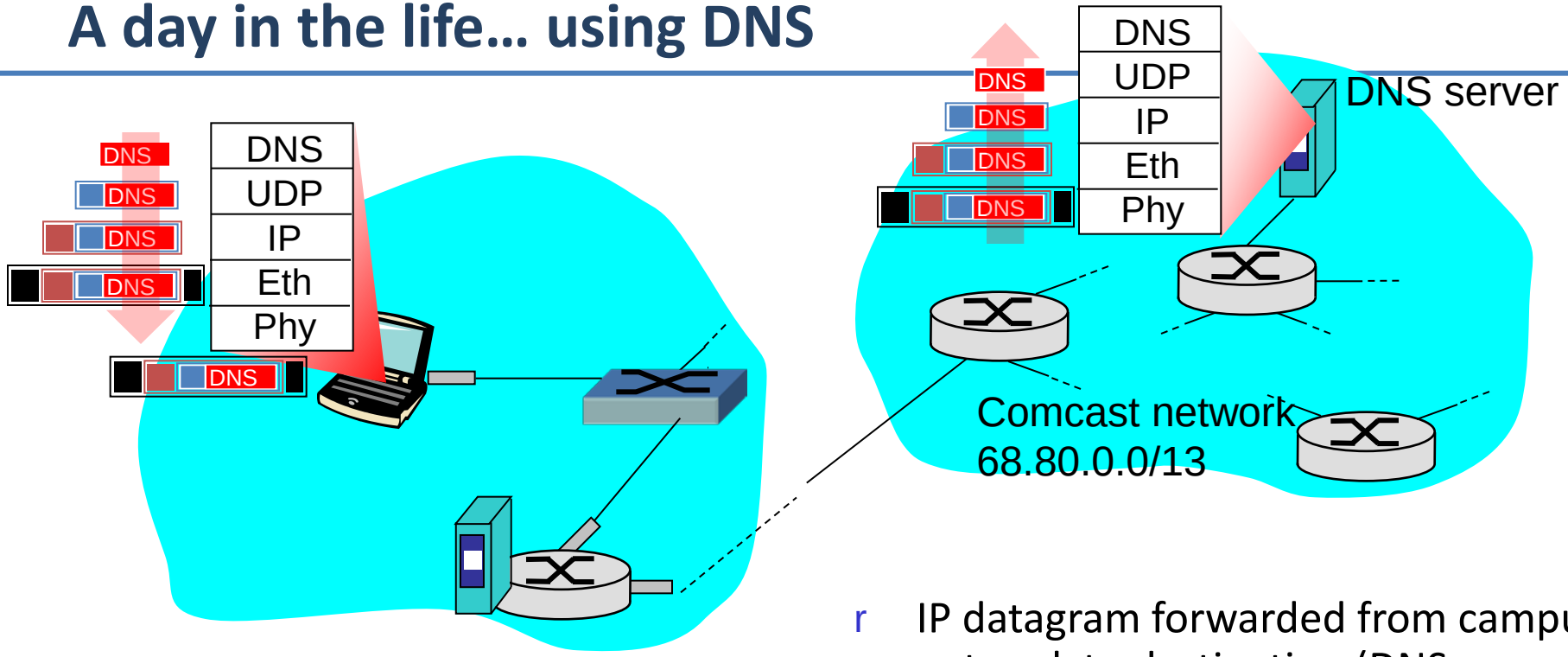
A day in the life… ARP (before DNS, before HTTP)



before sending ***HTTP*** request, need IP address of www.google.com: ***DNS***

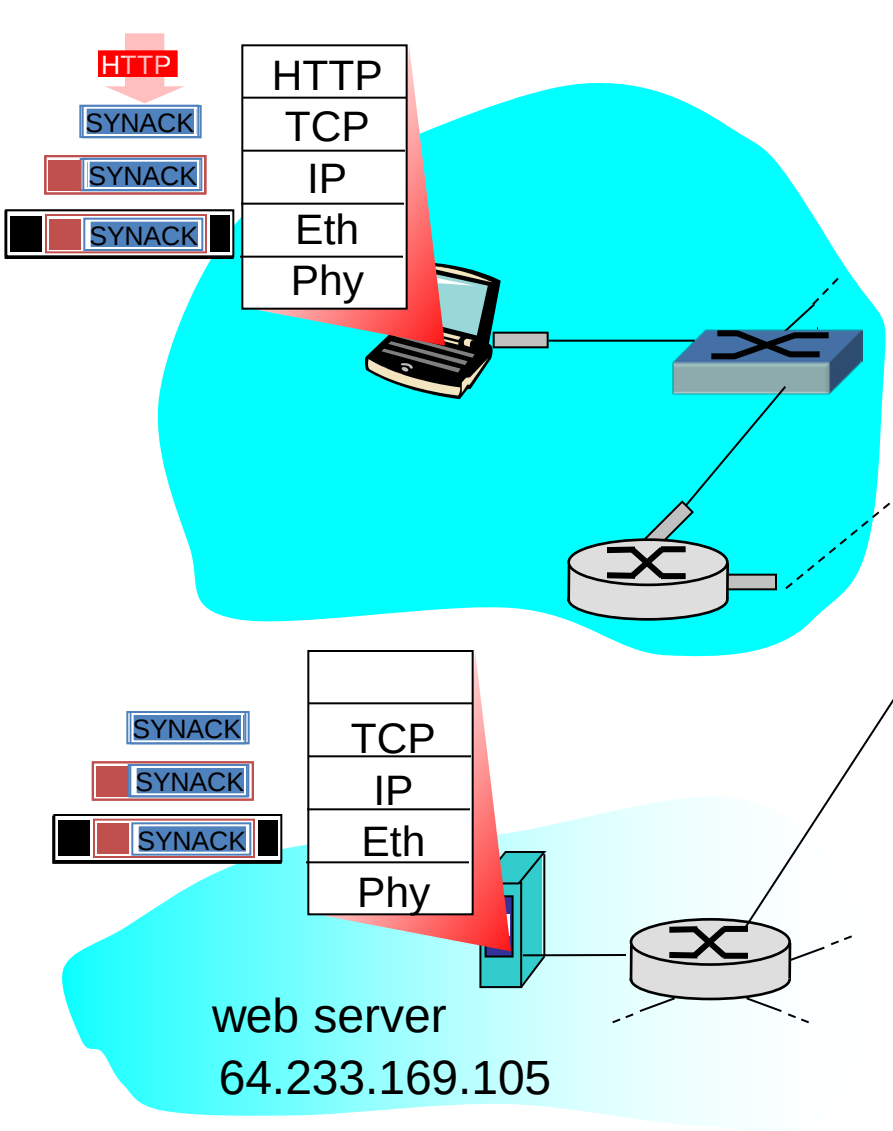
- DNS query created, encapsulated in UDP, encapsulated in IP, encasulated in Eth. In order to send frame to router, need MAC address of router interface: ARP
- ARP query broadcast, received by router, which replies with ARP reply giving MAC address of router interface
- client now knows MAC address of first hop router, so can now send frame containing DNS query

A day in the life… using DNS



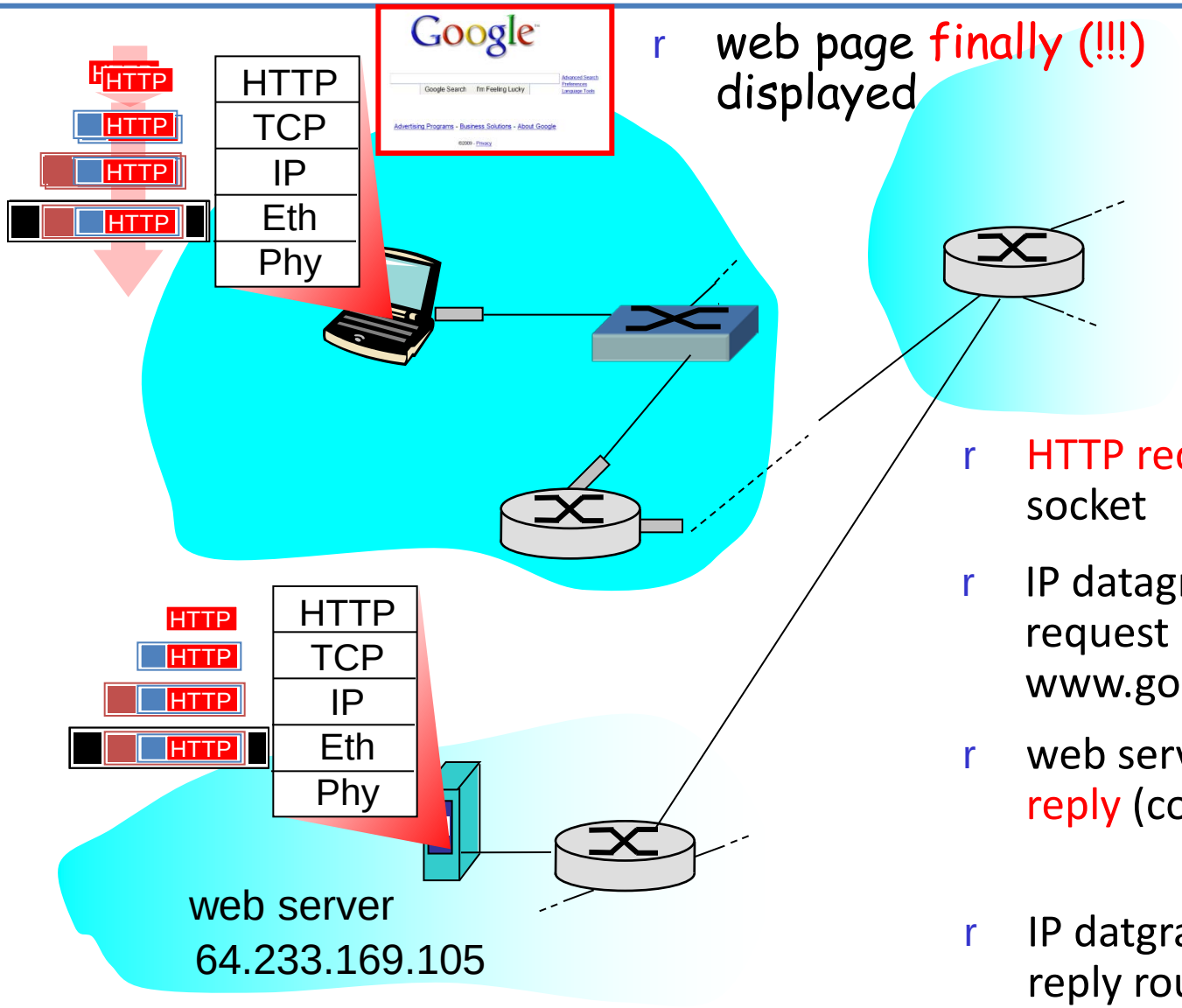
- IP datagram containing DNS query forwarded via LAN switch from client to 1st hop router
- IP datagram forwarded from campus network to destination (DNS-server) network, routed (tables created by RIP, OSPF and BGP routing protocols) to DNS server
- demux'ed to DNS server
- DNS server replies to client with IP address of www.google.com

A day in the life… TCP connection carrying HTTP



- to send HTTP request, client first opens TCP socket to web server
- TCP SYN segment (step 1 in 3-way handshake) inter-domain routed to web server
- web server responds with TCP SYNACK
- TCP connection established!

A day in the life… HTTP request/reply



- web page finally (!!!) displayed

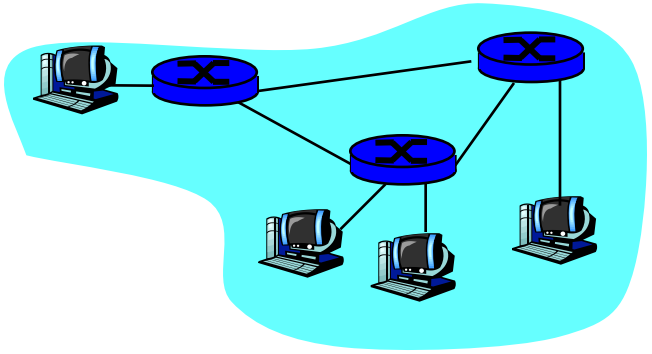
- HTTP request sent into TCP socket
- IP datagram containing HTTP request routed to www.google.com
- web server responds with HTTP reply (containing web page)
- IP datgram containing HTTP reply routed back to client

Synthesis cont.

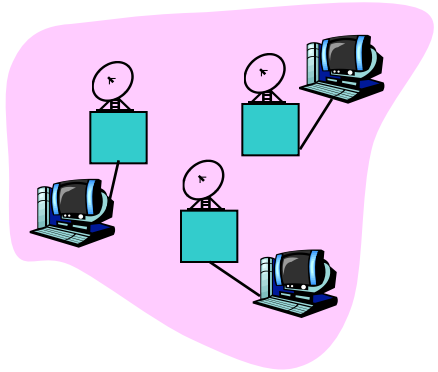
The Internet: virtualizing networks

1974: multiple unconnected nets
- ARPAnet
- data-over-cable networks
- packet satellite network (Aloha)
- packet radio network

… differing in:
- addressing conventions
- packet formats
- error recovery
- routing



ARPAnet



satellite net

"A Protocol for Packet Network Intercommunication", V. Cerf, R. Kahn, IEEE Transactions on Communications, May, 1974, pp. 637-648.

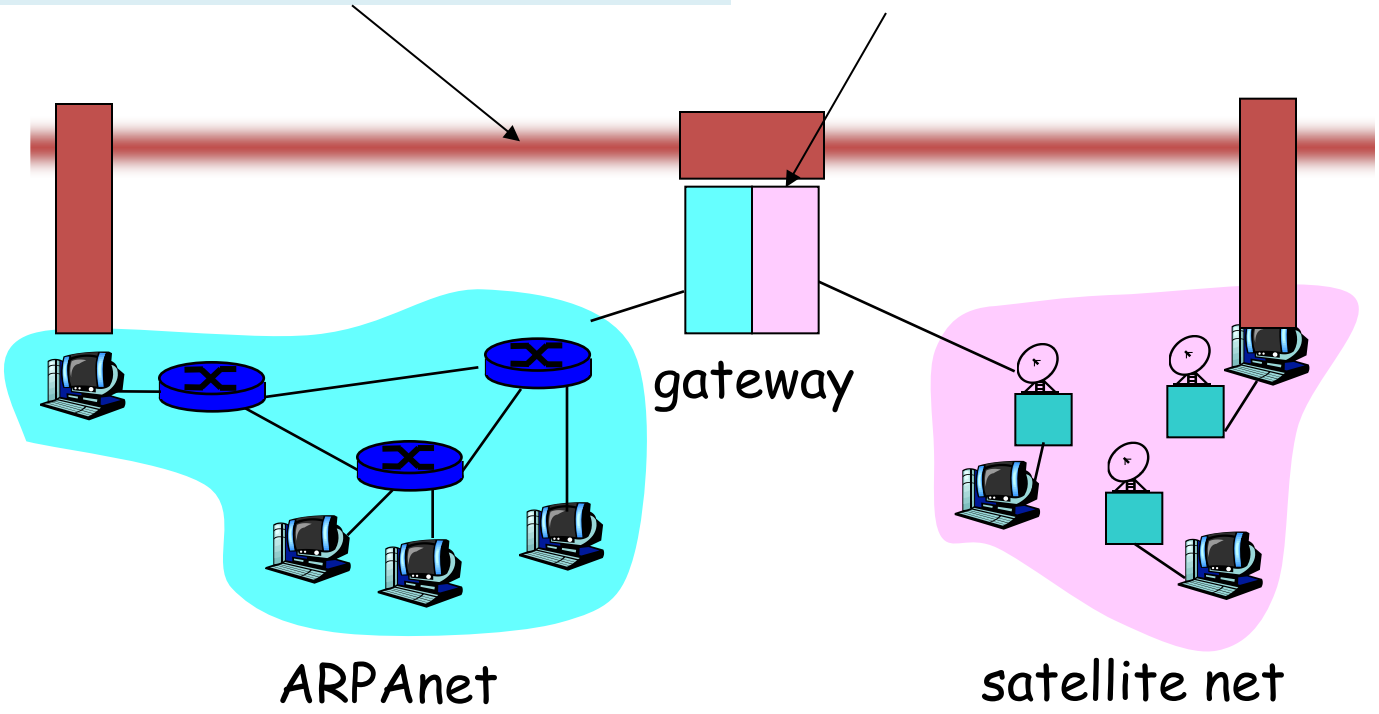
The Internet: virtualizing networks

Internetwork layer (IP):

- addressing: internetwork appears as single, uniform entity, despite underlying local network heterogeneity
- network of networks

Gateway:

- "embed internetwork packets in local packet format"
- route (at internetwork level) to next gateway



Cerf & Kahn's Internetwork Architecture

What is virtualized?

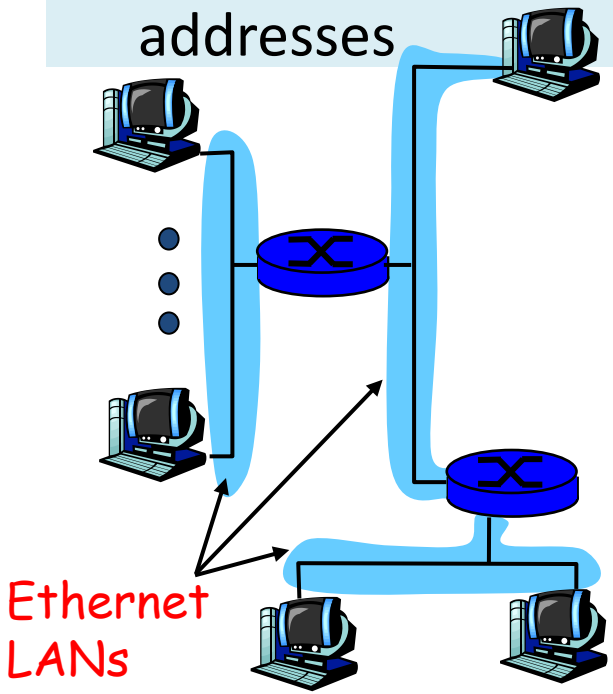
- two layers of addressing: internetwork and local network
- new layer (IP) makes everything homogeneous at internetwork layer
- underlying local network technology
  - Cable, satellite, 56K telephone modem
  - Ethernet, other LAN
  - ATM/ MPLS (Multiprotocol Label Switching Protocol)

… "invisible" at internetwork layer. Looks like a link layer technology to IP

e.g. IP-Over-ATM

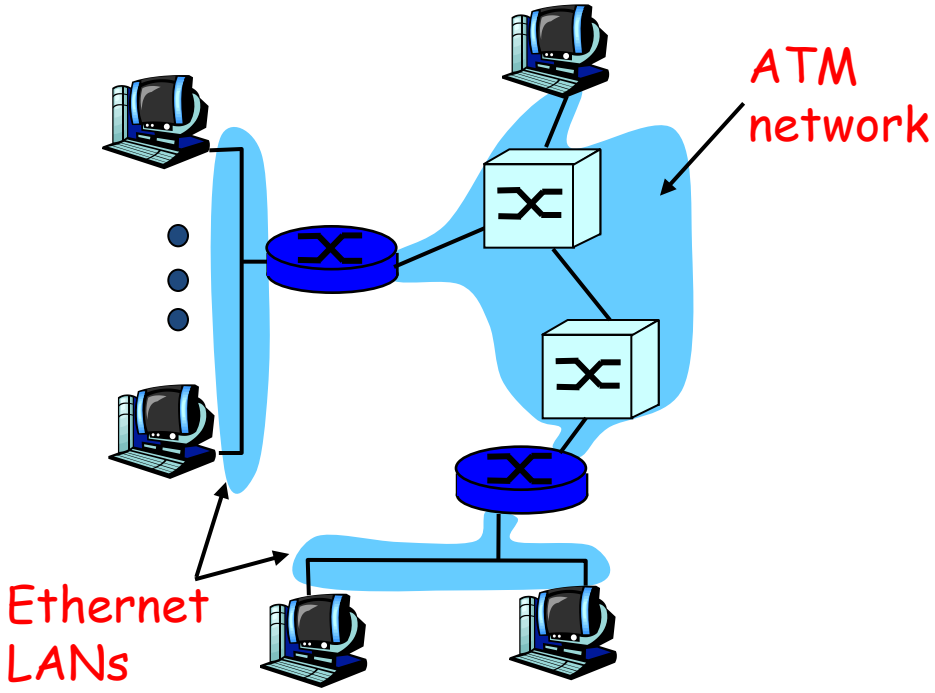
"Classic" IP over eg Ethernet

- 3 "networks" (e.g., LAN segments)
- MAC (eg802.3) and IP addresses



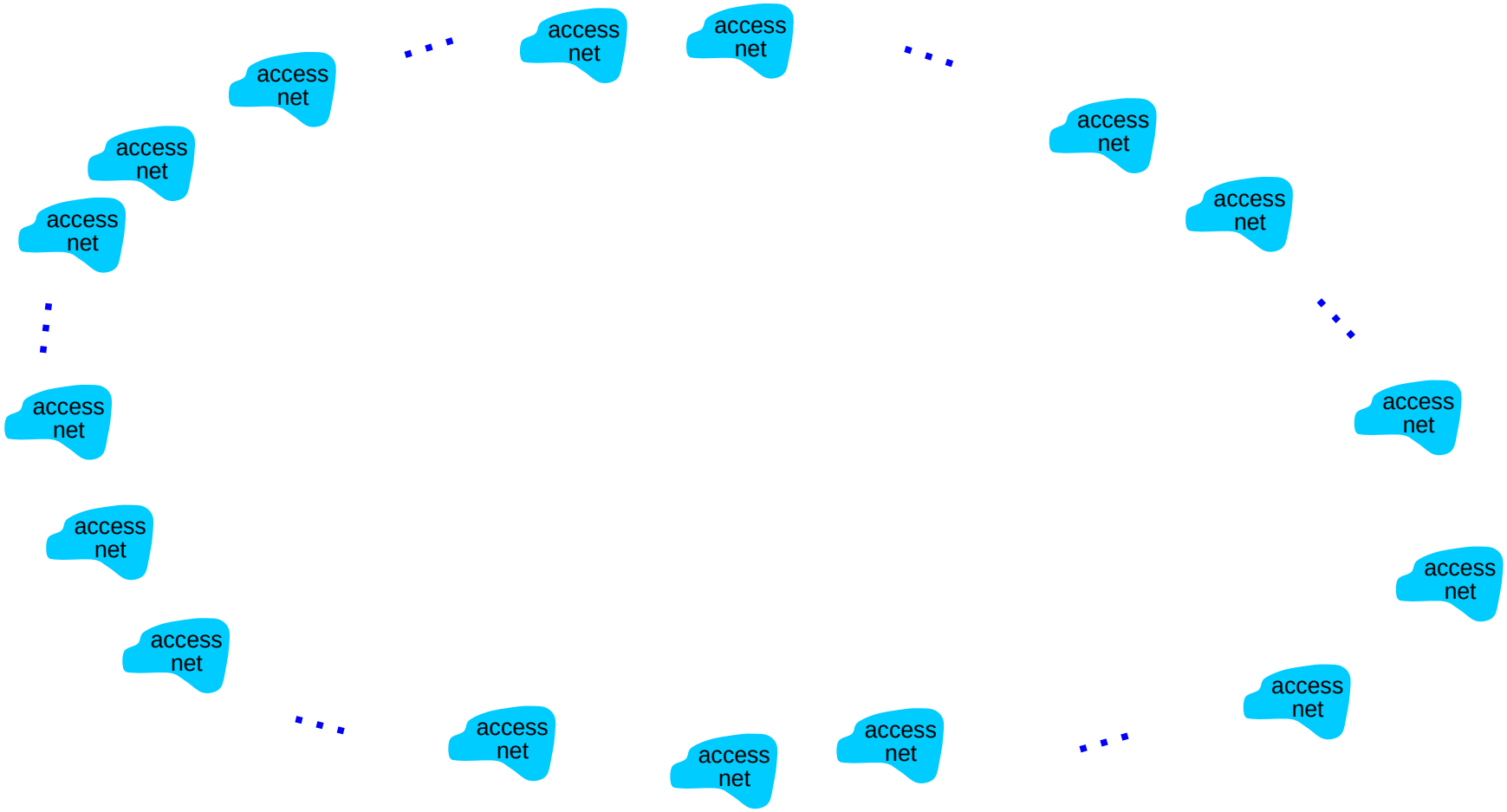
IP over ATM

- replace "network" (e.g., LAN segment) with ATM network
- ATM addresses (as MAC addresses), IP addresses



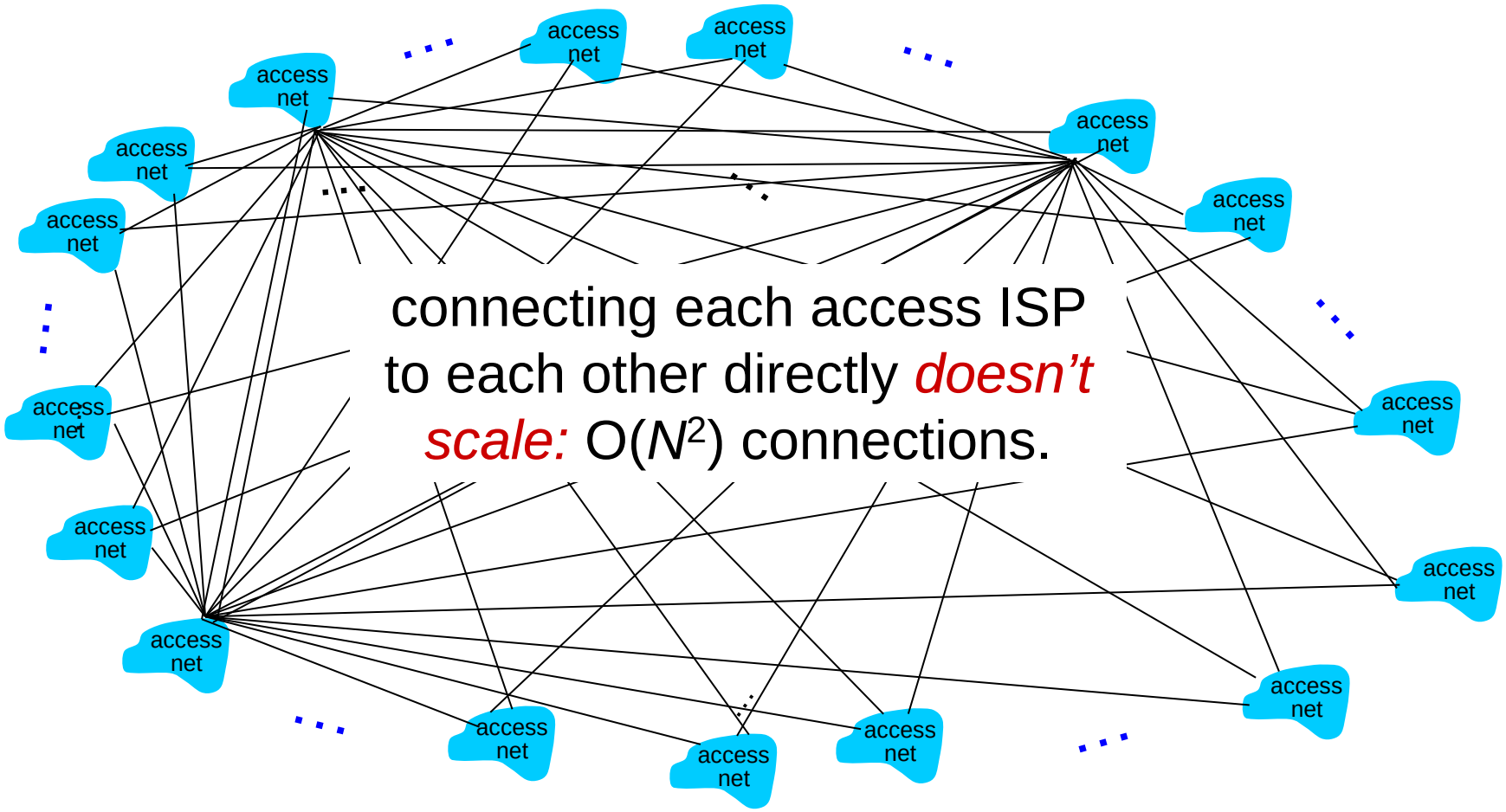
Internet structure: network of networks

Question: given *millions* of access ISPs, how to connect them together?



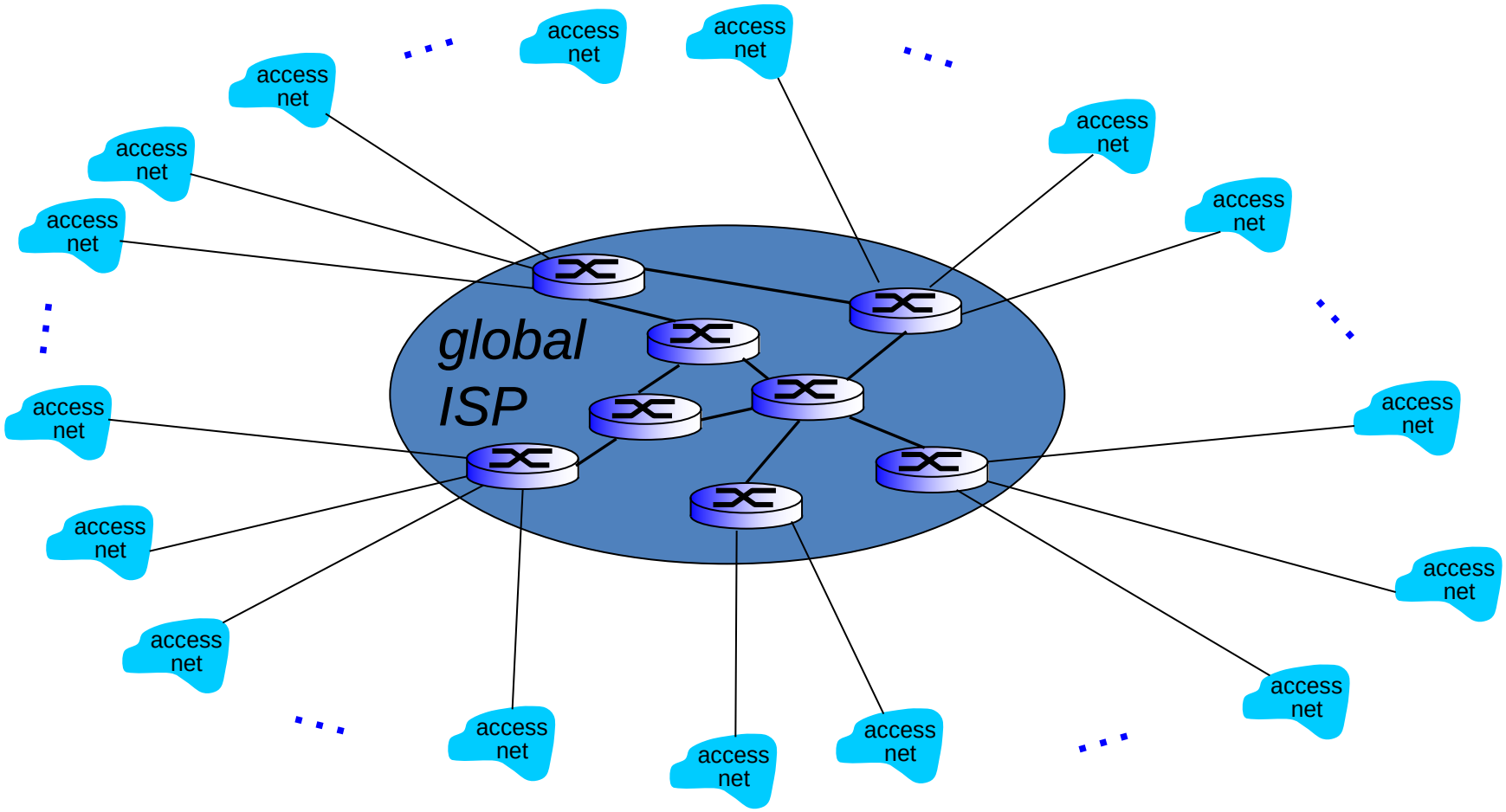
Internet structure: network of networks

Option: connect each access ISP to every other access ISP?



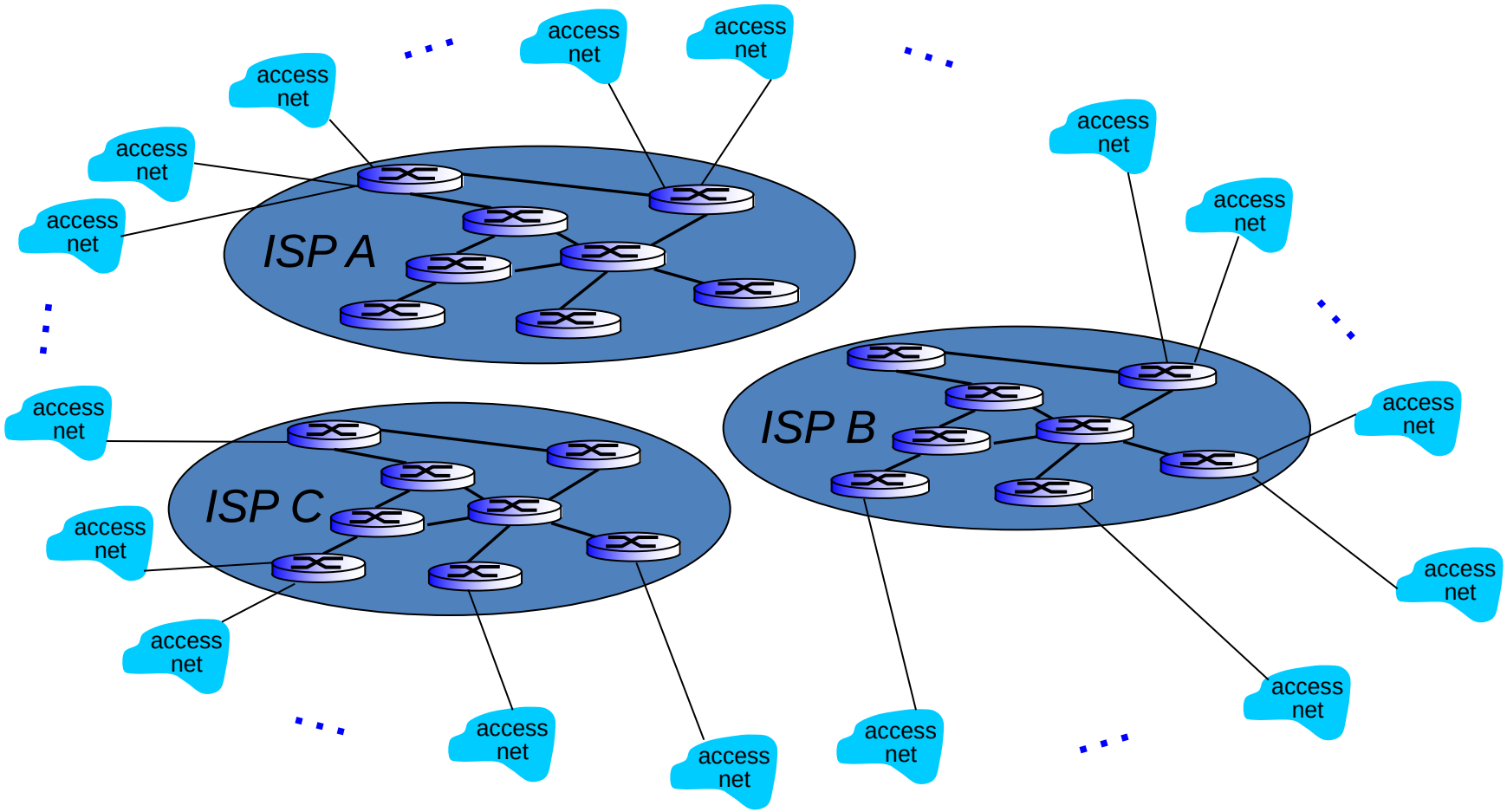
Internet structure: network of networks

Option: connect each access ISP to a global transit (imaginary) ISP?
Customer and provider ISPs have economic agreement.



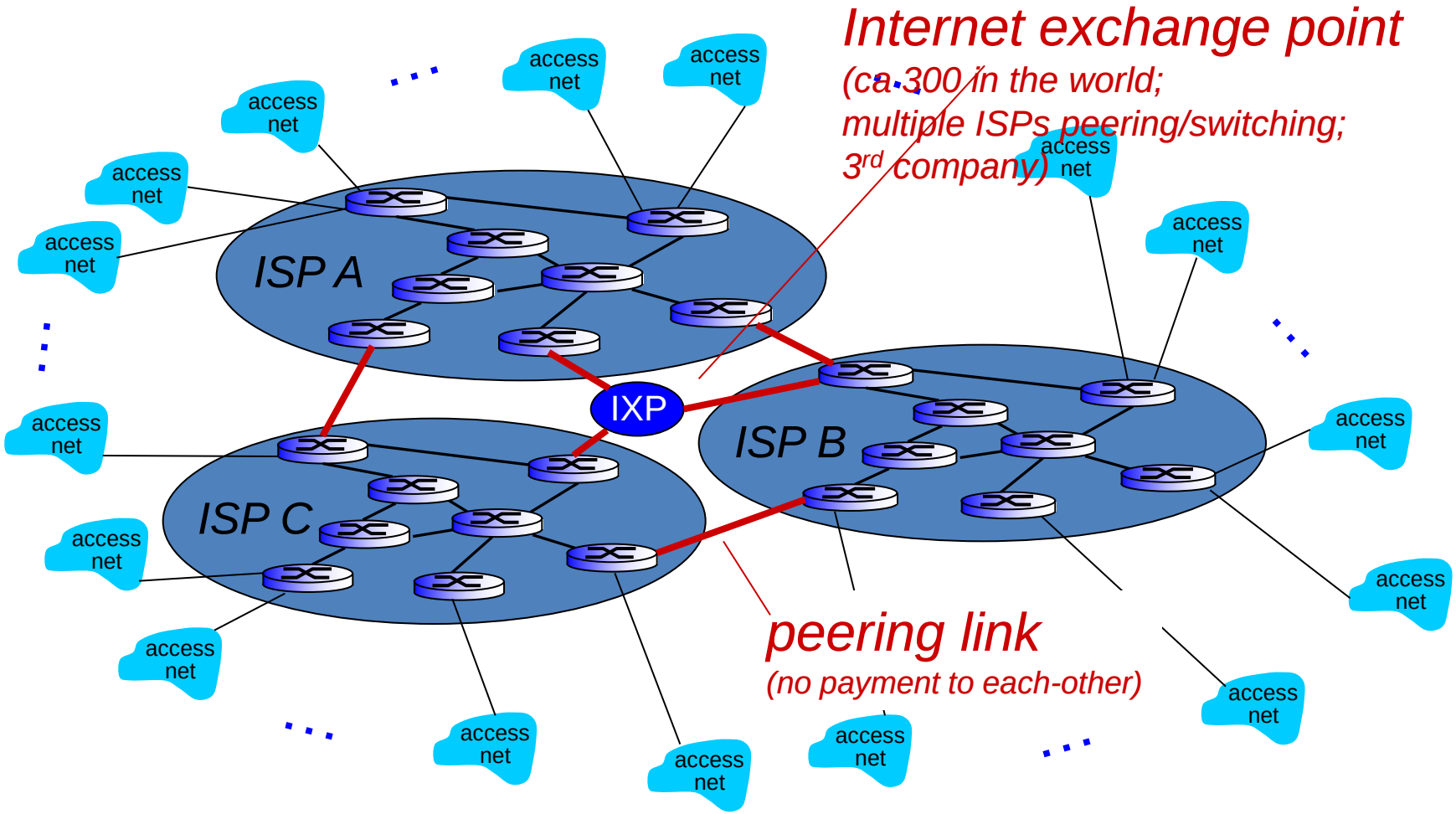
Internet structure: network of networks

But if one global ISP is viable business, there will be competitors ….



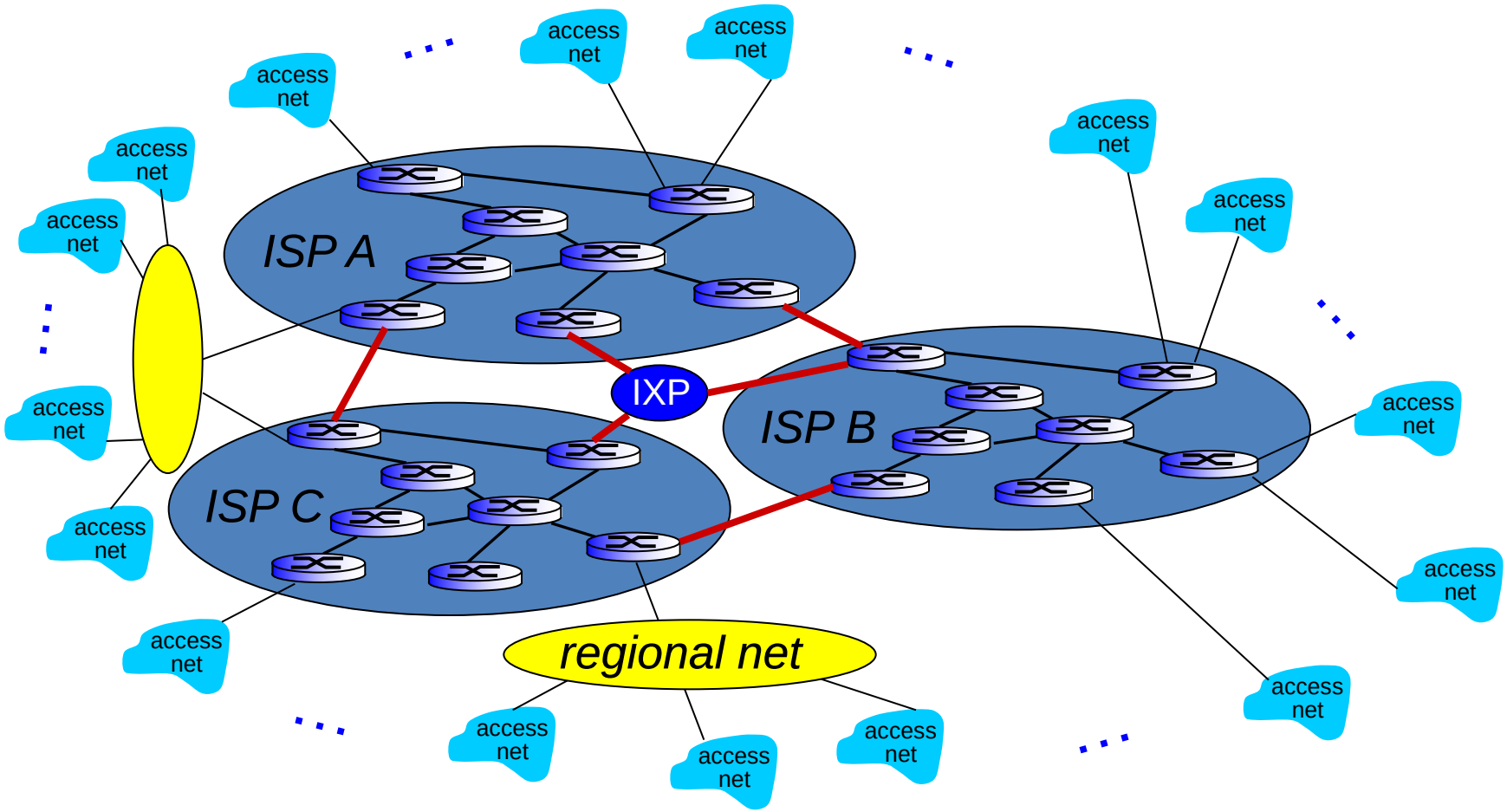
Internet structure: network of networks

But if one global ISP is viable business, there will be competitors …. which must be interconnected



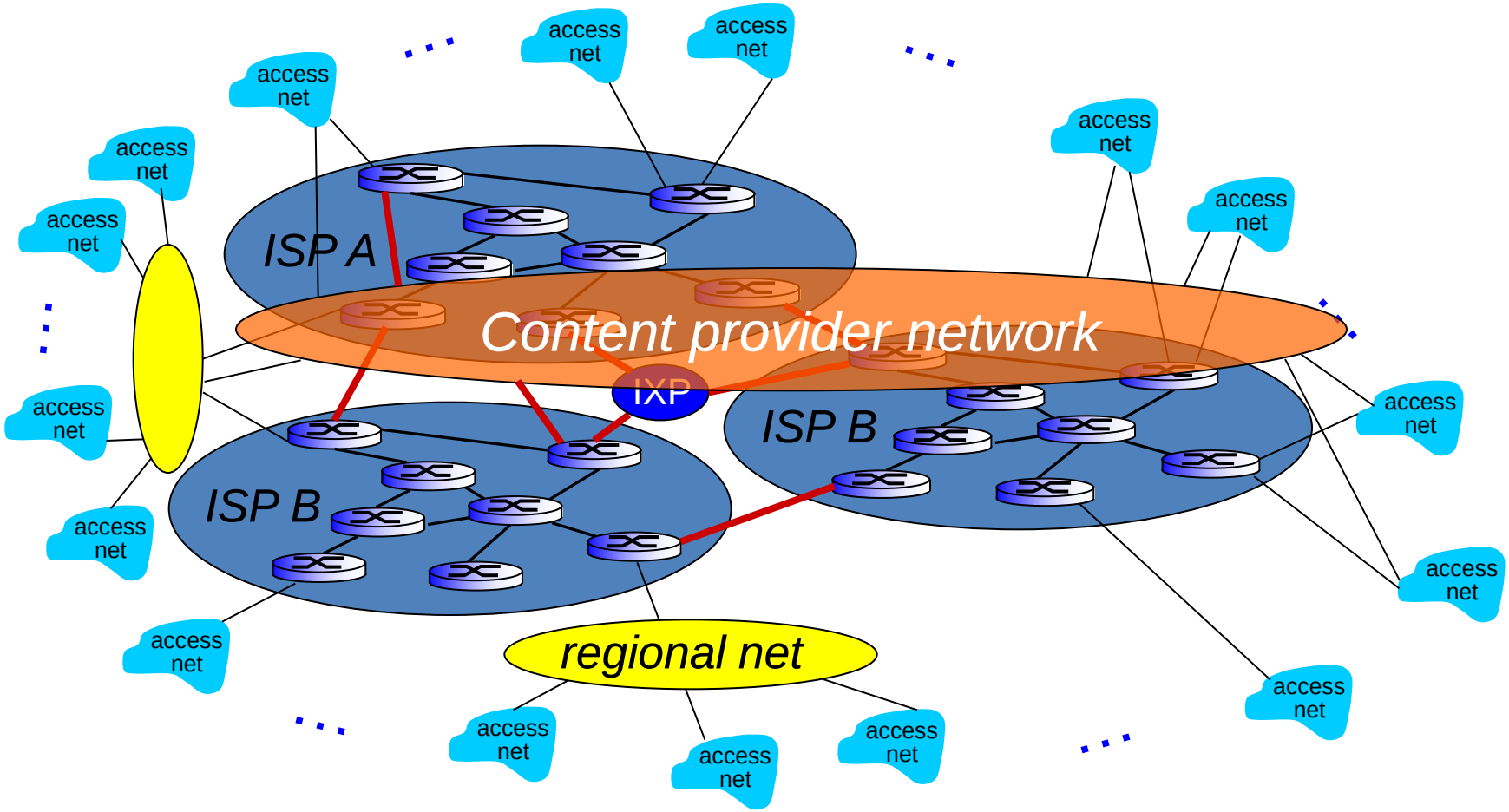
Internet structure: network of networks

… and regional networks may arise to connect access nets to ISPS

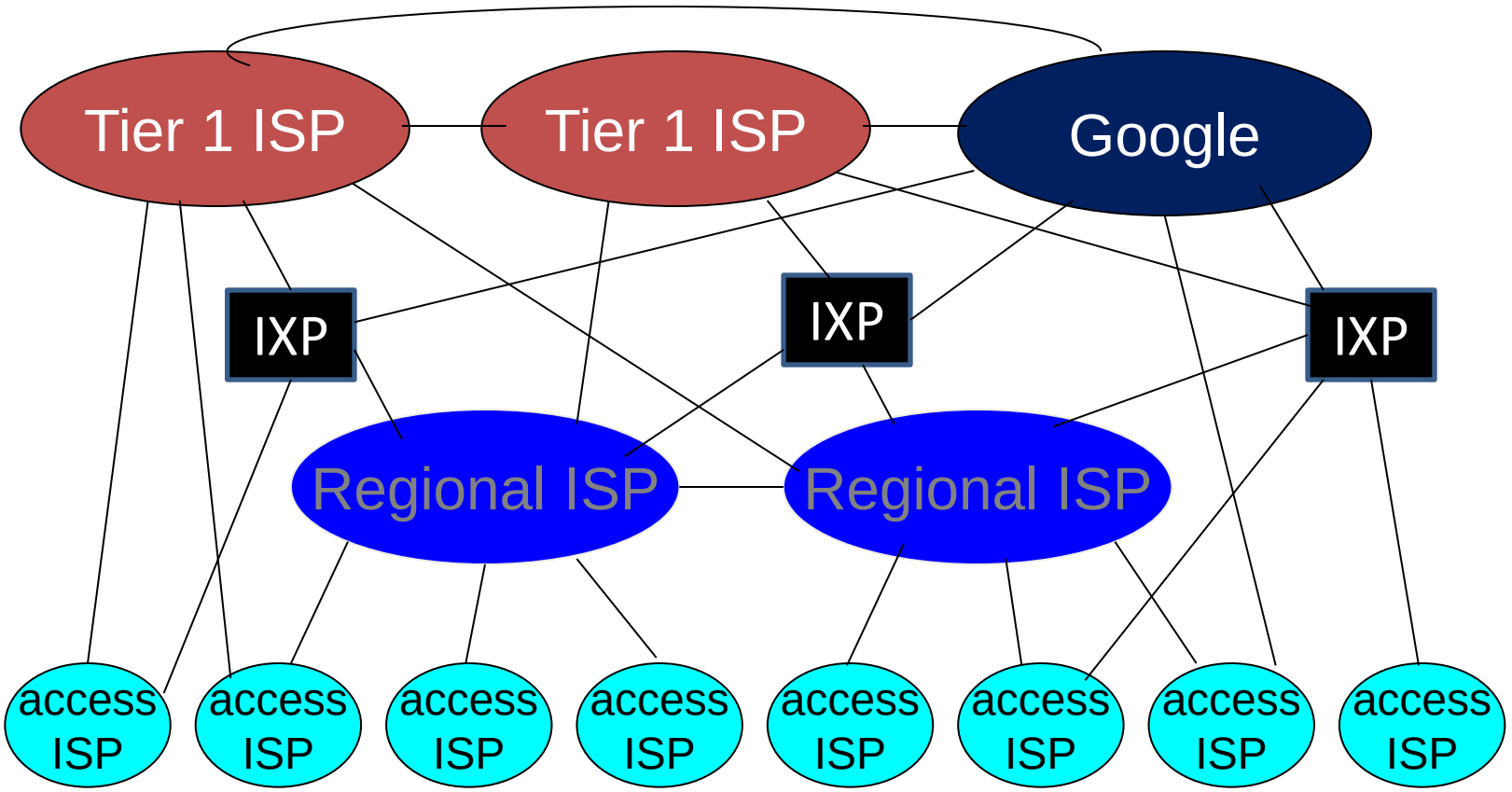


Internet structure: network of networks

… and content provider networks (e.g., Google, Microsoft, Akamai) may run their own network, to bring services, content close to end users



Internet structure: network of networks



- at center: small # of well-connected large networks
  - "tier-1" commercial ISPs (e.g., Level 3, Sprint, AT&T, NTT), national & international coverage
  - A new form of content provider network (e.g, Google): private network that connects it data centers to Internet, often bypassing tier-1, regional ISPs

End-of-recap….

Thank you

Recall, important for the exam:

When/where: wednesday March 18, 14.00-18.00, M

You may have with you:

- English-X dictionary
- no calculators, PDAs, etc (if/where numbers matter, do rounding)

To think during last, summary-study

Overview; critical eye; explain, ask yourselves: why is this so? / How does it work?

Good luck with all your efforts!!!

"If you hear a voice within you say 'you cannot paint,' then by all means paint, and that voice will be silenced." – Vincent Van Gogh