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# Course on Computer Communication and Networks

## Lecture 1 & part of lecture 2 Chapter 1: Introduction

EDA344/DIT 420, CTH/GU

Based on the book *Computer Networking: A Top Down Approach*, Jim Kurose, Keith Ross, Addison-Wesley.

# Roadmap





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- what's the Internet
- protocol layers
  - Communication through layers
- edge & core of any big network:
  - types of service, ways of information transfer, routing
- Internet layers & Logical vs physical communication
- Performance:
  - delays, loss
- Network/Internet structure complemented:
  - access net, physical media
  - backbones, NAPs, ISPs
- Security prelude

# the Internet: “nuts and bolts” view (1)





-  PC
-  server
-  wireless laptop
-  cellular handheld

- millions of connected computing devices: *hosts* = *end systems*
  - running *network apps*

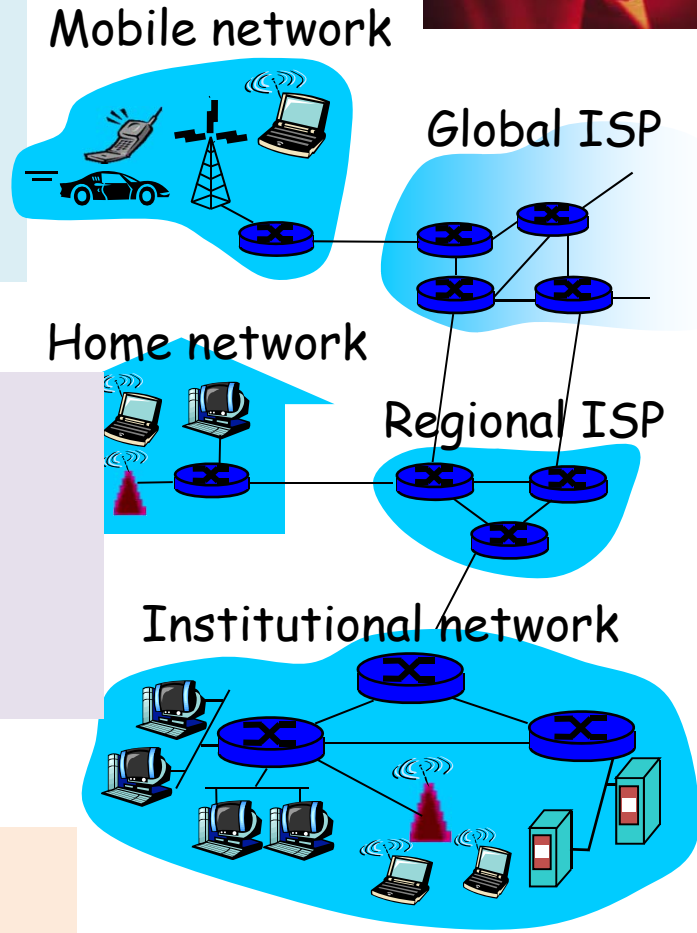
- *communication links*

- fiber, copper, radio, satellite
- transmission rate = *bandwidth*

-  access points
-  wired links

 router

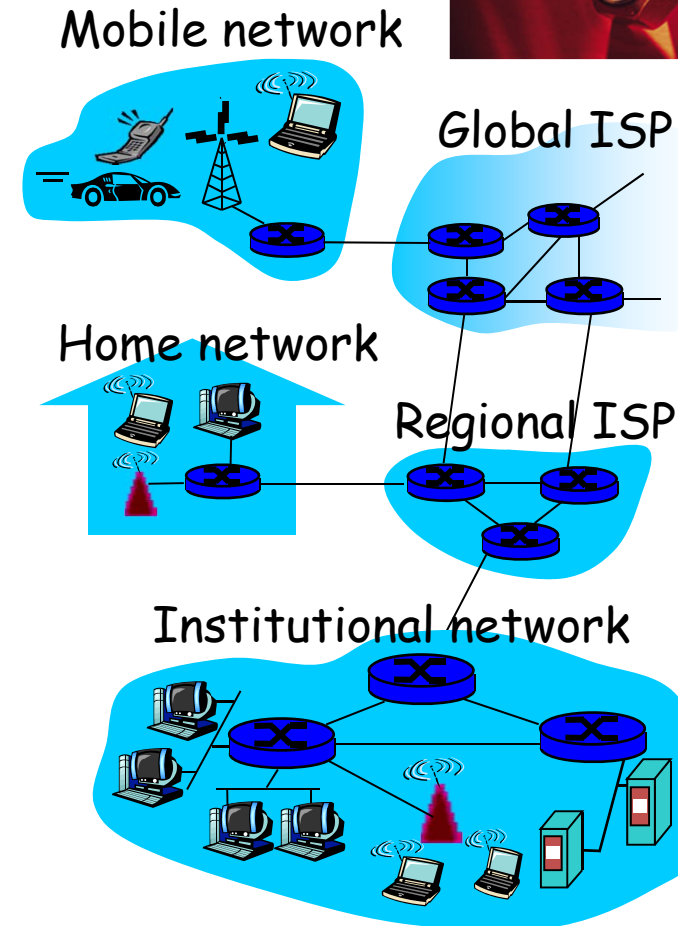
- *routers*: forward packets (chunks of data)



# the Internet: “nuts and bolts” view (2)

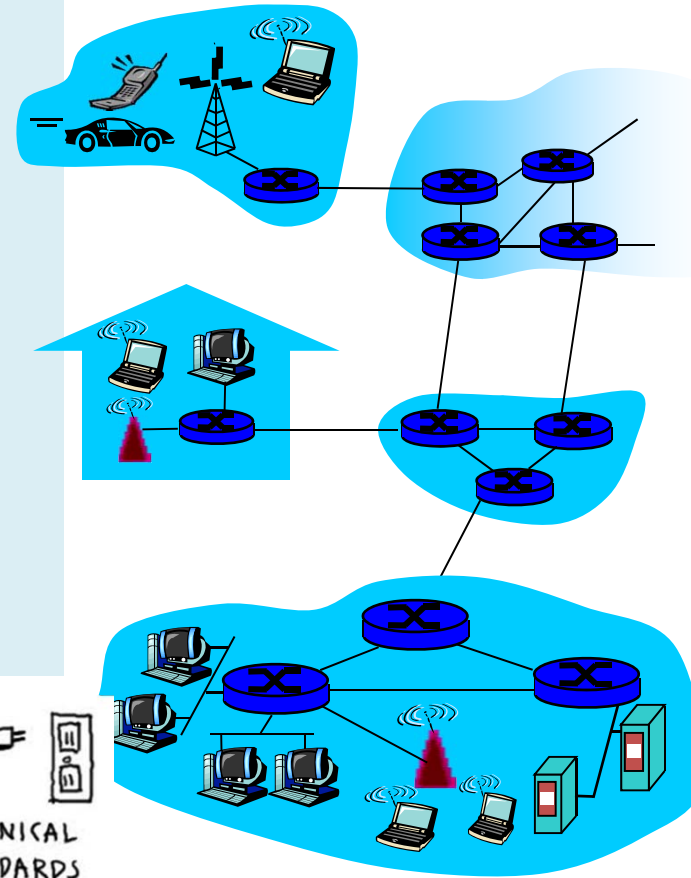


- *protocols* control sending, receiving of msgs
  - e.g., TCP, IP, HTTP, Skype, Ethernet
- *Internet: “network of networks”*
  - loosely hierarchical
  - public Internet versus private intranet



# the Internet: service view

- **communication *infrastructure*** enables distributed applications:
  - Web, VoIP, email, games, e-commerce, file sharing
- **communication *services* provided to apps:**
  - reliable data delivery from source to destination
  - “best effort” (unreliable) data delivery



## Internet standards

- RFC: Request for comments
- IETF: Internet Engineering Task Force



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## Networks are complex

- many “pieces”:
  - hosts
  - routers
  - links of various media
  - applications
  - hardware, software

## Question:

Is there any hope of *organizing* structure, study, development of networks?

# Layers of abstraction

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## Dealing with complex systems:

- explicit structure allows identification, relationship of complex system's pieces
  - layered **reference model** for discussion
- **modularization eases maintenance/es**
  - change of implementation of layer's service transparent to rest of system
  - e.g., change in gate procedure doesn't affect rest of system



# Terminology: Protocols, Interfaces

- Each **layer offers services** to the upper layers (**shielding** from the implementation details)
  - **service interface**: across layers in same host
- Layer  $n$  on a host carries a **conversation** with layer  $n$  on another host
  - **host-to-host interface**: defines messages exchanged with peer entity
- **Network architecture** (set of layers, interfaces) vs **protocol stack** (protocol implementation)

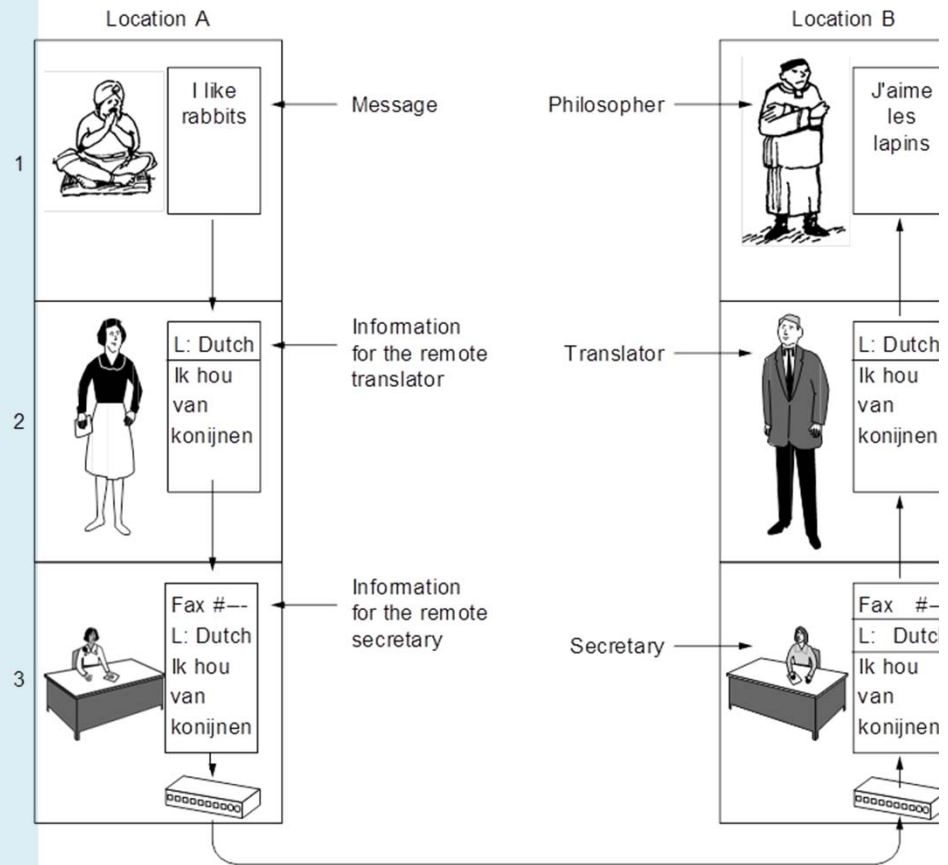
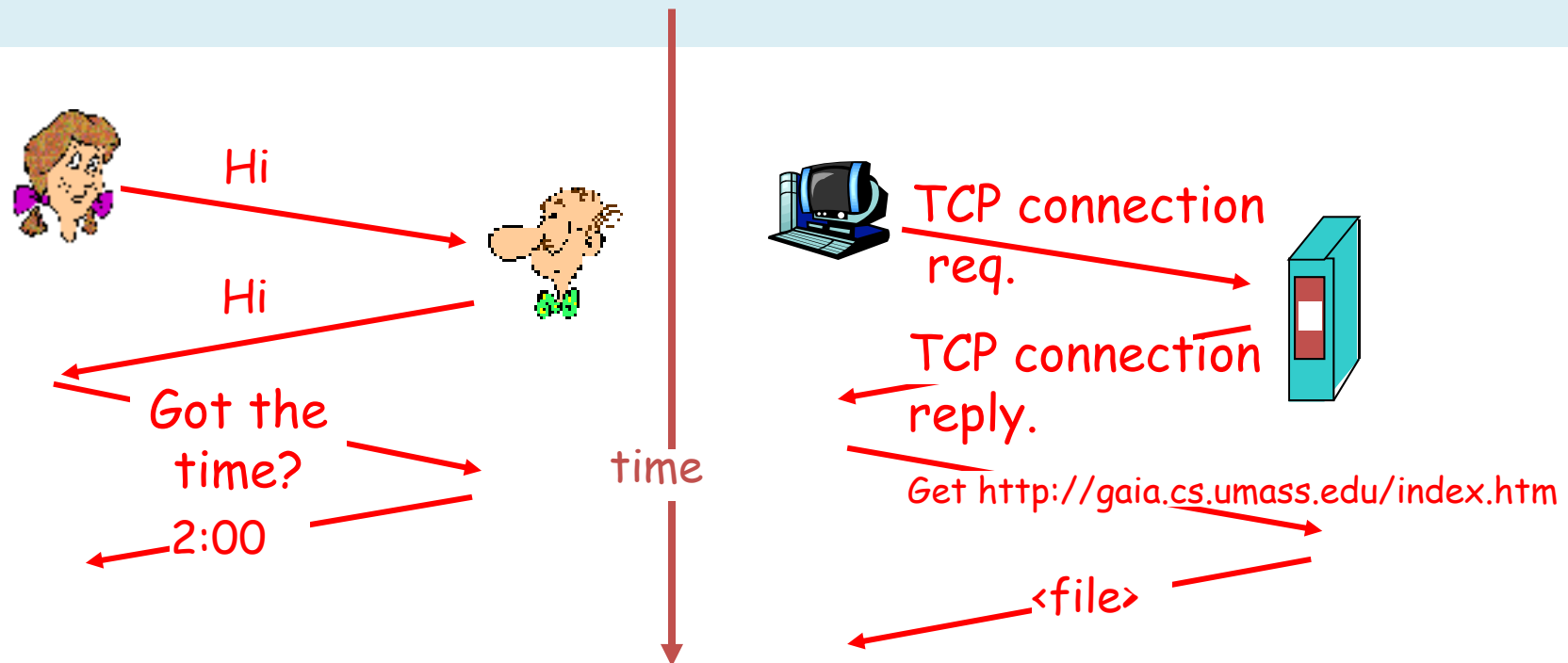


Fig. 1-10. The philosopher-translator-secretary architecture.

# What's a protocol?

a human protocol and a computer network protocol:



**host-to-host interface:** defines

- messages exchanged with peer entity: *format, order of msgs sent and received among network entities*
- *actions taken on msg transmission, receipt*

# Roadmap



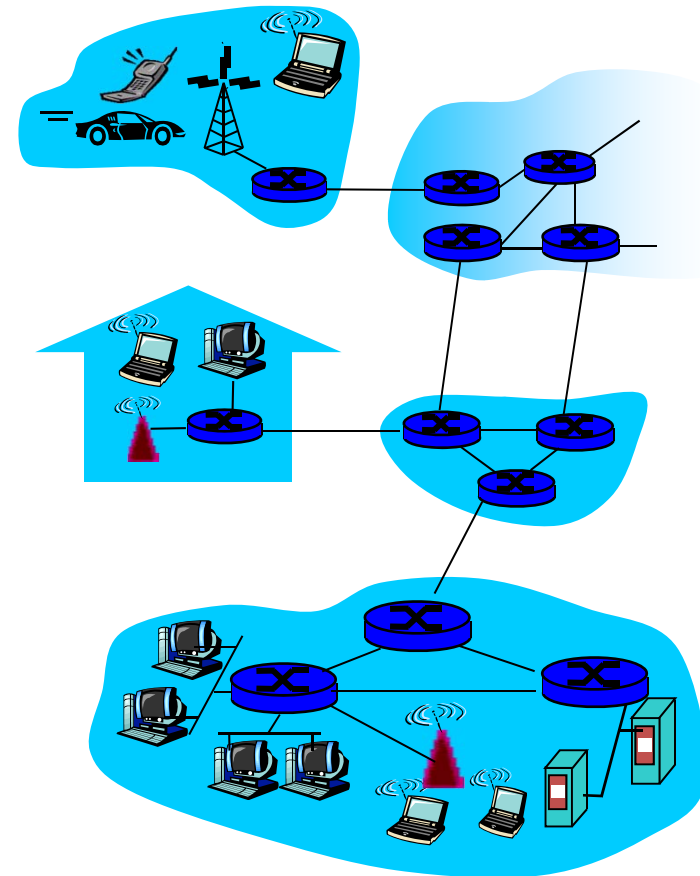
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# A closer look at (any big) network's structure:

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- **network edge:**  
applications and hosts
- **access networks,**  
physical media: wired,  
wireless  
communication links

- **network core:**
  - interconnected routers
  - network of networks



# The network edge:

## end systems (hosts):

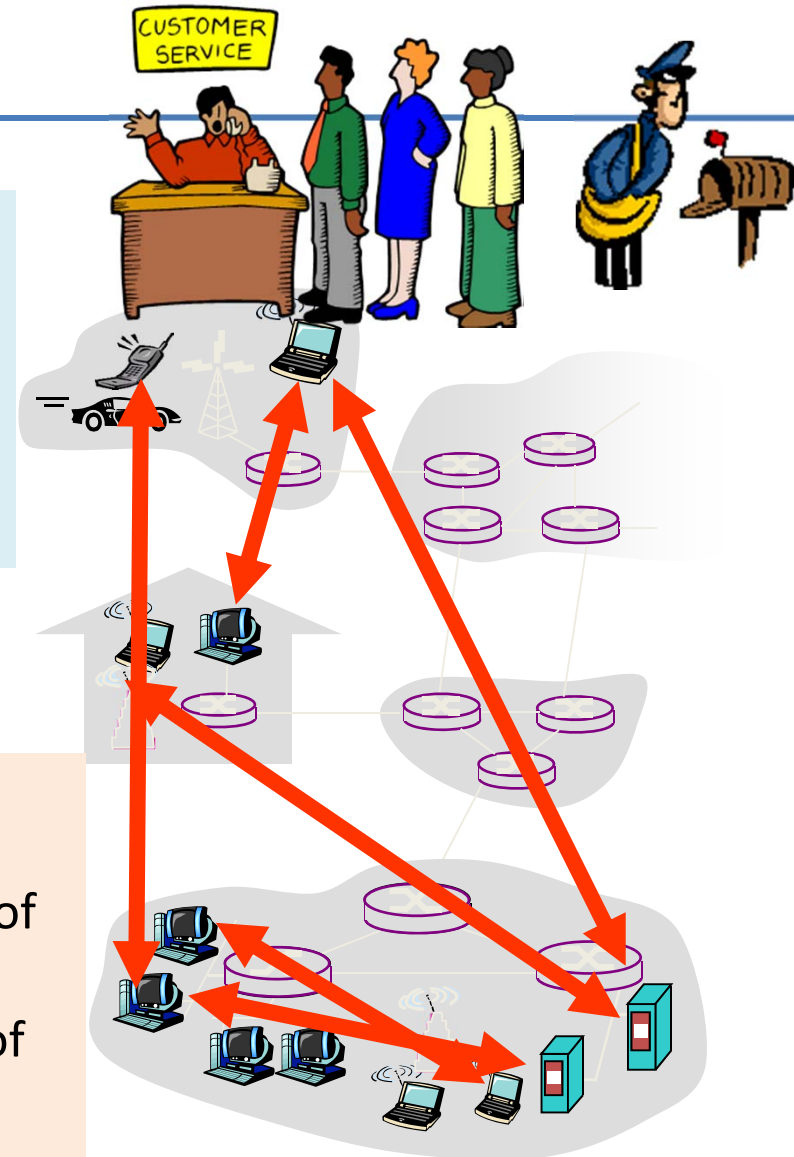
- ❑ run application programs e.g. in Internet Web, email, ...
- ❑ ... based on **network services** available at the edge

**Basic types of service** offered by the network to applications:

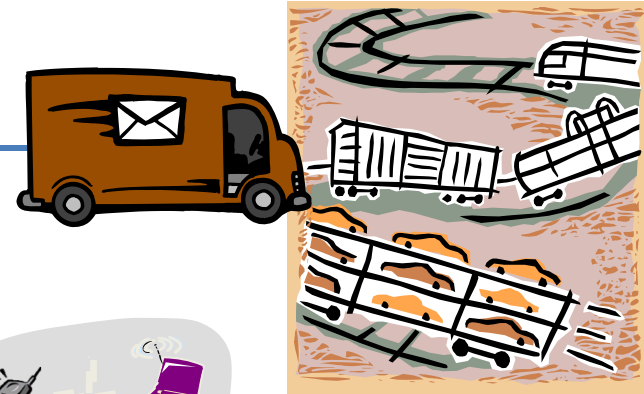
**connection-oriented**: reliable delivery of the data in the order they are sent

**connectionless**: “best effort” delivery of the data in arbitrary order

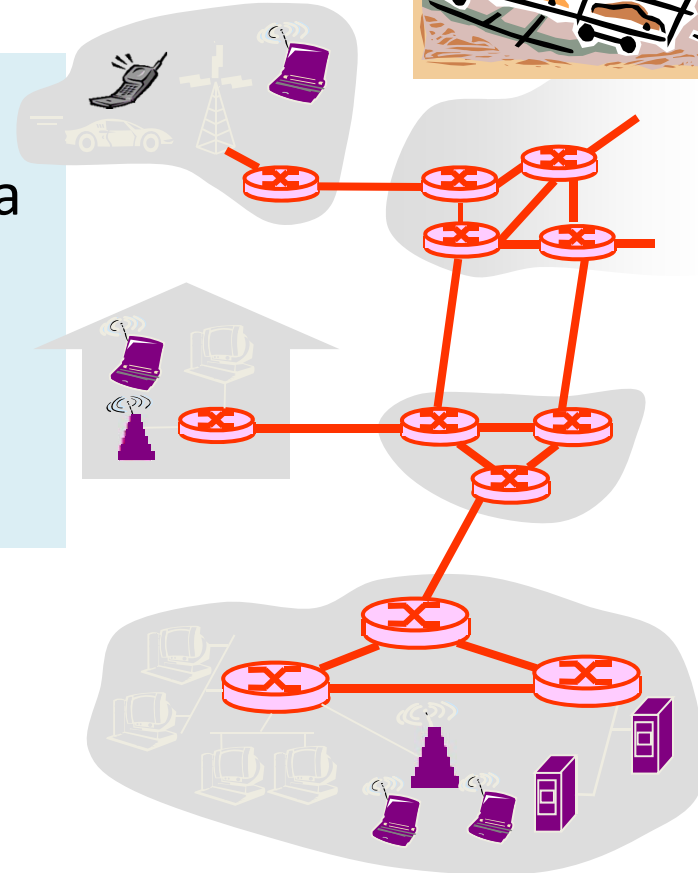
**Q: can we think of more types of service?**



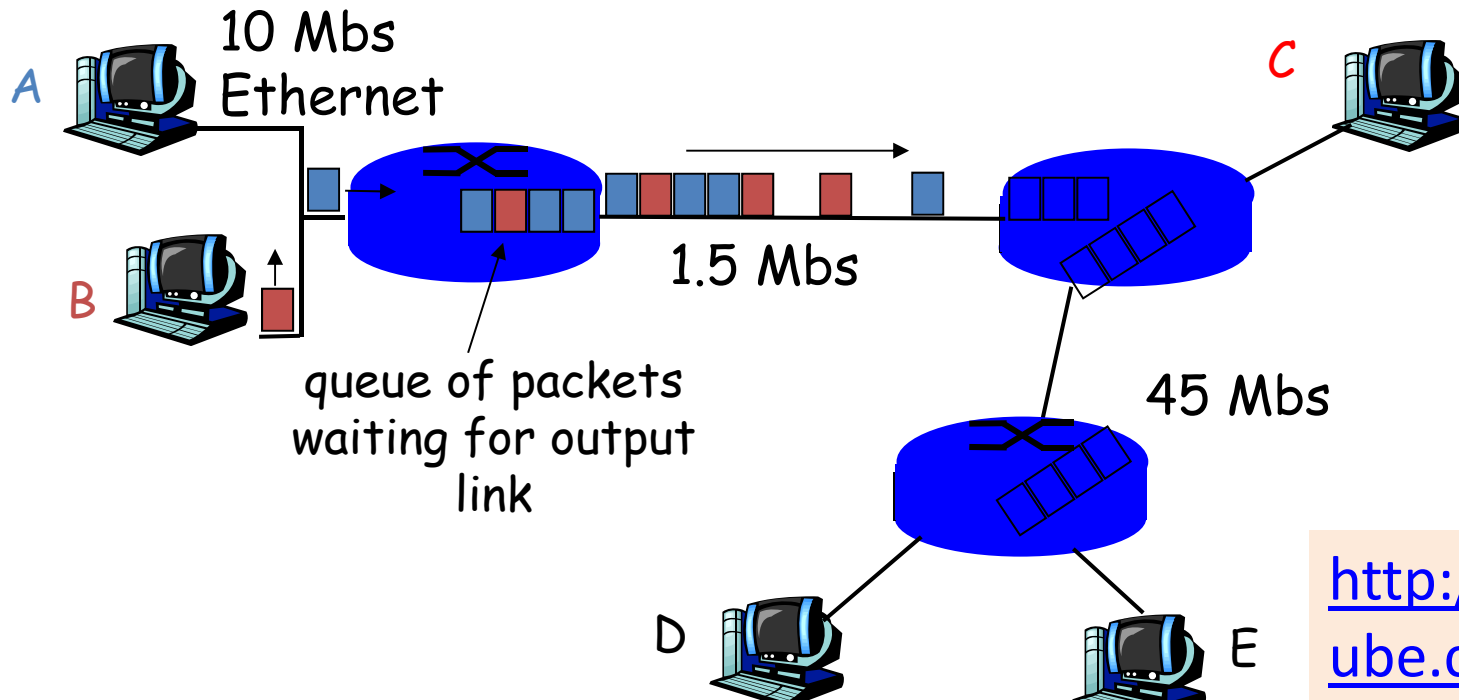
# The Network Core



- mesh of interconnected routers
- **fundamental question:** how is data transferred through net?
- **packet-switching:** data sent thru net in discrete “chunks”



# Network Core: Packet Switching



each end-end data stream divided into *packets*

- packets *share* network resources
- resources used *as needed*

**store and forward:**

- packets move one hop at a time
  - transmit over link; wait turn at next link

<http://www.youtube.com/watch?v=O7CuFIM4V54>

nice animation;  
note some of the  
terms in narration  
are not accurate  
(wrt protocol  
specifications)

# Packet-switched networks: routing

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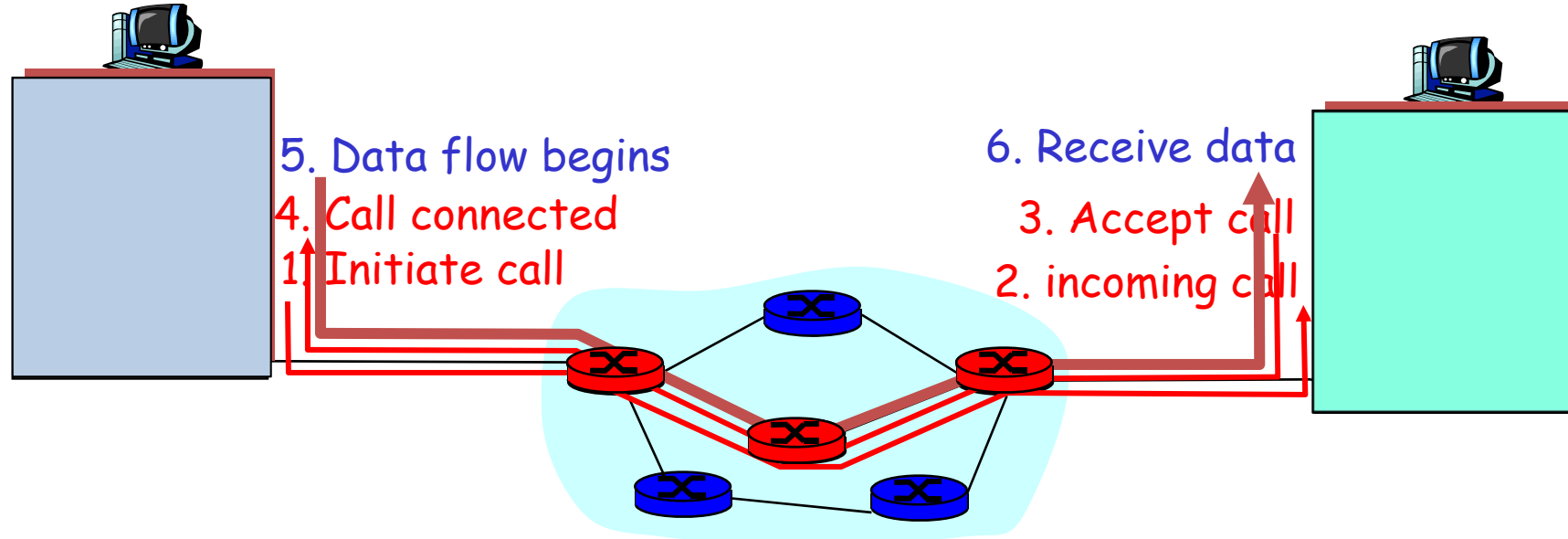
- What is routing's goal? find routes from source to destination
  - (Challenge 1) **path selection** algorithms
  - (Challenge2) Important **design issue/type of service offered:**
    - **datagram network:**
      - *destination address* determines next hop
      - routes may change during session
    - **virtual circuit network:**
      - fixed path determined at *call setup time*, remains fixed thru session
      - routers maintain per-call state



# Virtual circuits:

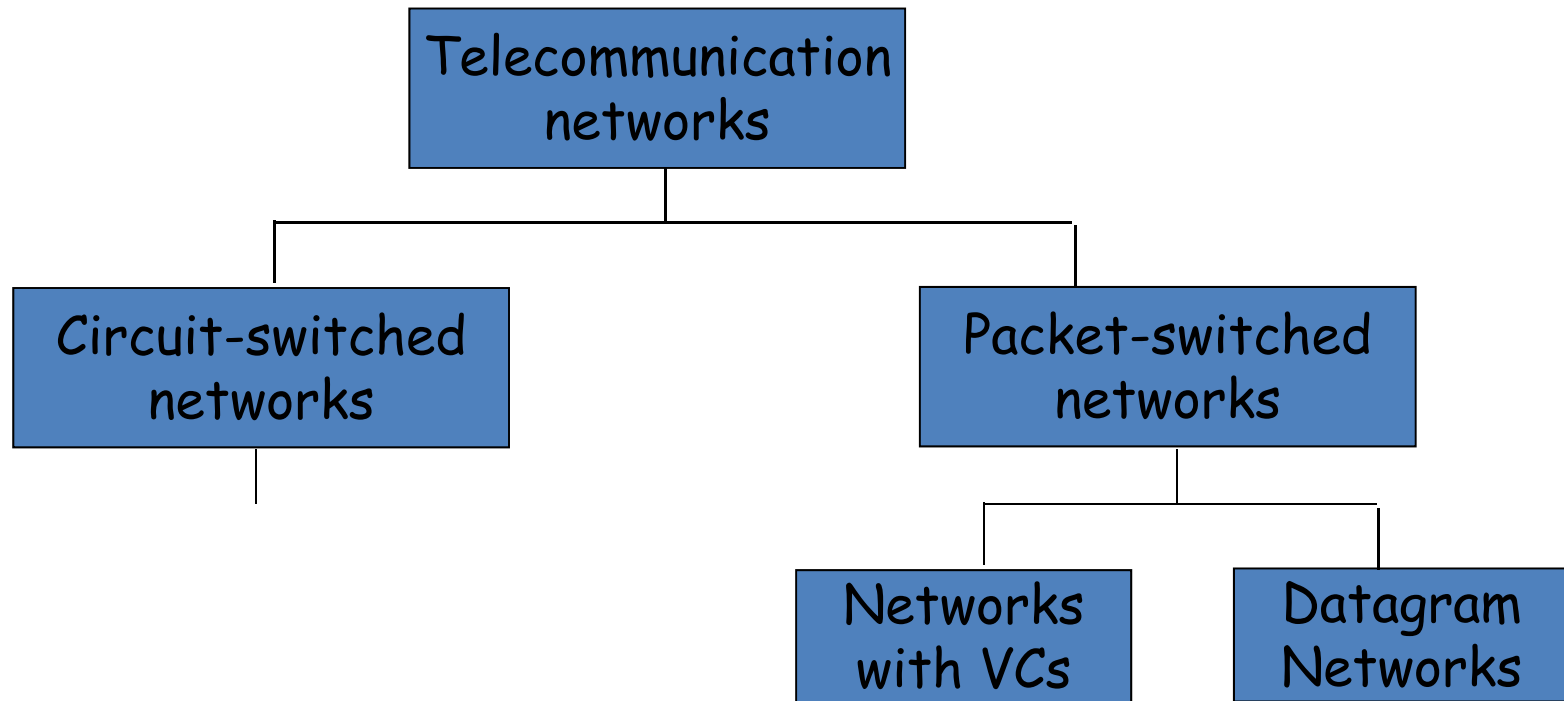
“source-to-dest path behaves almost like a physical circuit”

- call setup, teardown for each call *before* data can flow
  - signaling protocols to setup, maintain, teardown VC
- *every* router maintains “state” for *each* passing connection
- resources (bandwidth, buffers) may be *allocated* to VC



# Network Taxonomy

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- Datagram network (eg Internet) cannot be characterized either connection-oriented or connectionless.
  - Internet provides both connection-oriented (TCP) and connectionless services (UDP), at the network edge, to apps.

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# Layering – Some “history”:

## The OSI Reference Model

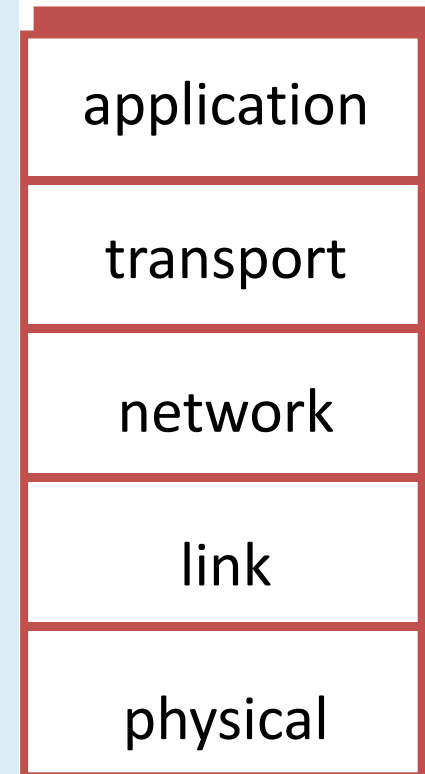
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- ISO (International Standards Organization) defined the OSI (Open Systems Interconnect) model to help vendors create interoperable network implementation
- Reduced the problem into smaller and more manageable problems: **7 layers**
  - a layer should be created where a different level of abstraction is needed; each layer should perform a well defined function)
  - The function of each layer should be chosen with an eye toward defining internationally **standardized** protocols
- “X dot” series (X.25, X.400, X.500) OSI model implementation (protocol stack)
- Did not really “fly”...

# Internet protocol stack

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- **application:** ftp, smtp, http, etc
- **transport:**
  - connection-oriented, reliable data delivery from source to destination (TCP);
  - connectionless, “best effort” (unreliable) data delivery (UDP)
- **network:** routing of datagrams from source to destination; best effort service
  - IP, routing protocols
- **link:** data transfer between neighboring network elements
  - Ethernet, WiFi, ...
- **physical:** bits “on the physical medium”



# Internet protocol stack

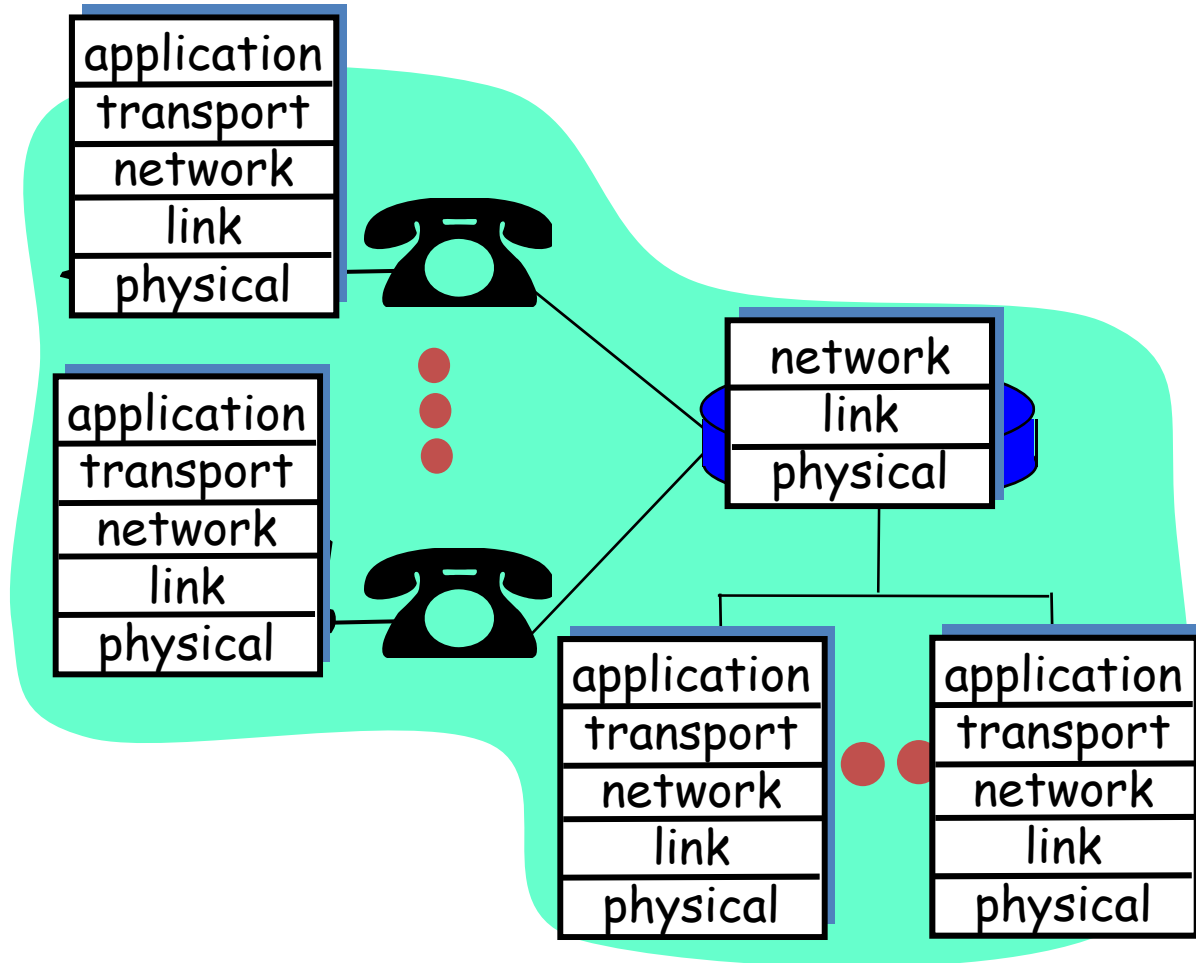
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- ❑ Architecture simple but not as thoroughly thought as OSI's
  - no clear distinction between interface-design and implementations;
  - hard to re-implement certain layers
  
- ❑ Successful protocol suite (**de-facto standard**)
  - was there when needed (OSI implementations were too complicated)
  - freely distributed with UNIX

# Layering: logical communication

Each layer:

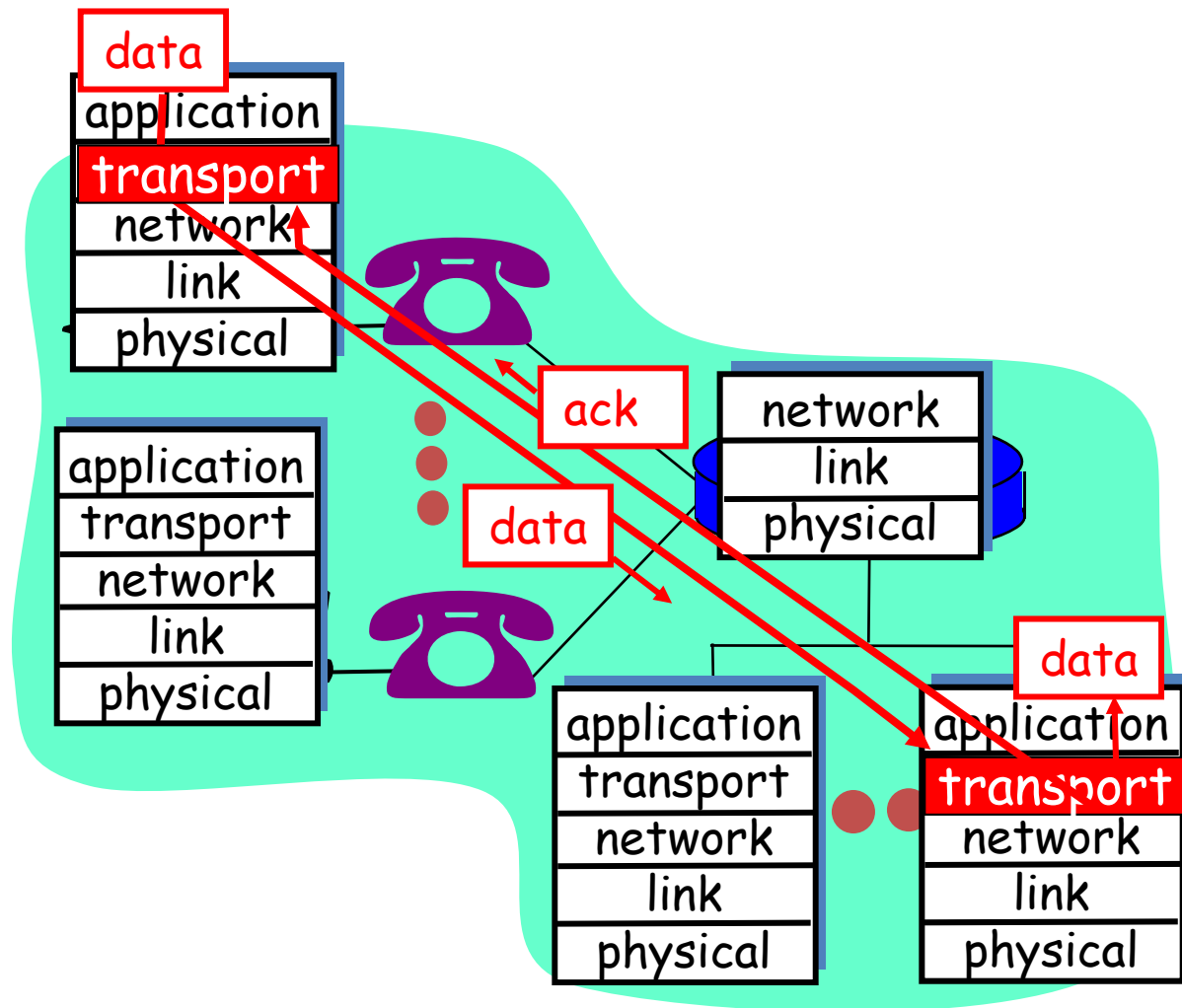
- distributed
- “entities” implement layer functions at each node
- entities perform actions, exchange messages with peers



# Layering: *logical* communication

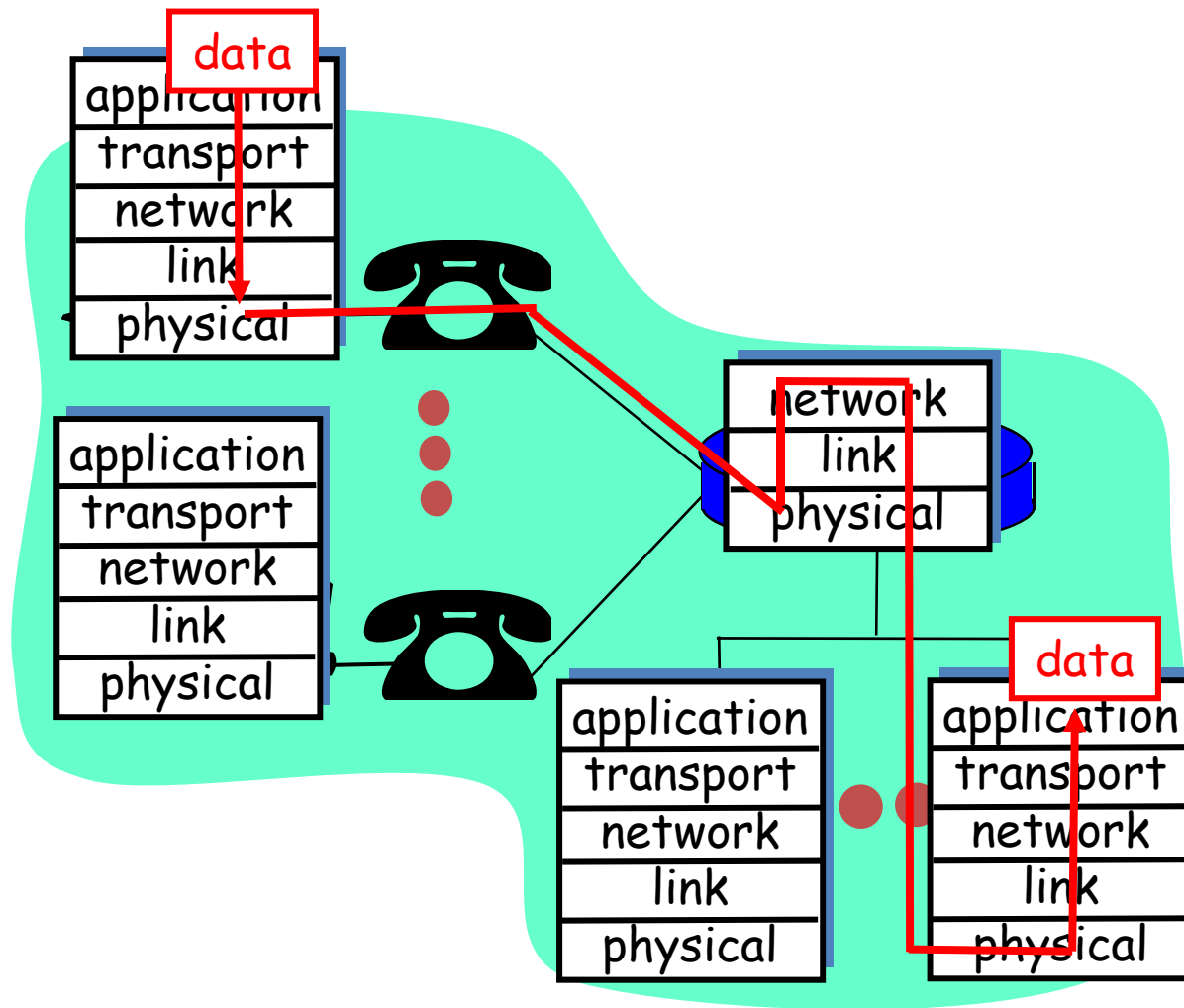
## E.g.: transport

- take data from app
- add addressing, form “datagram”
- send datagram to peer
- (possibly wait for peer to ack receipt)





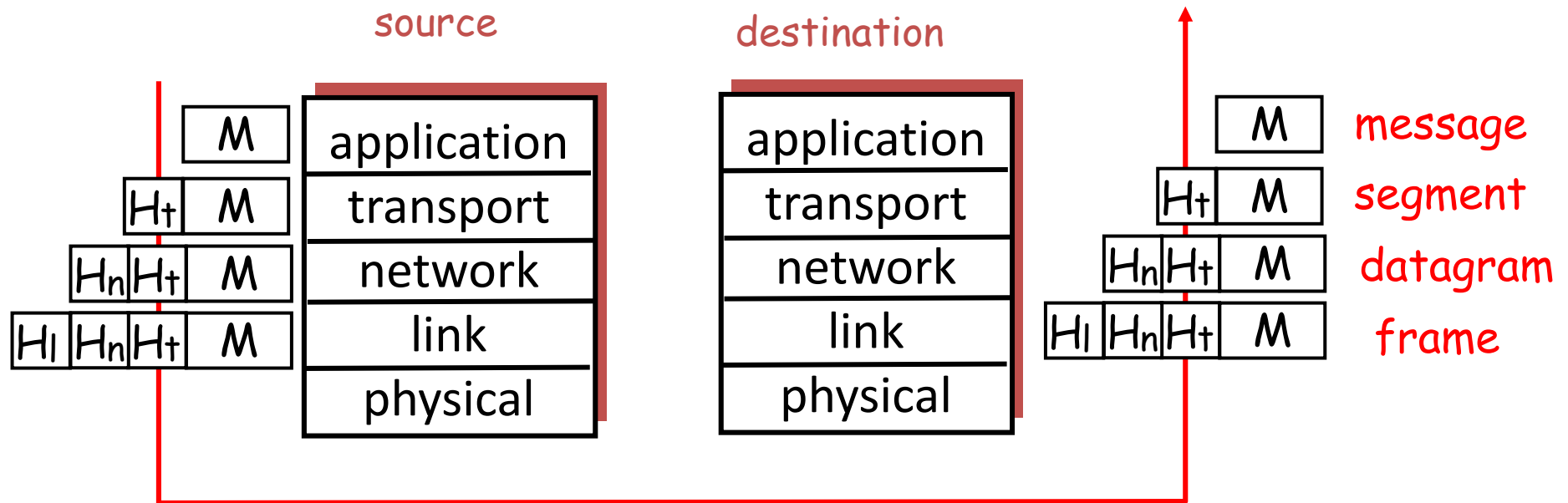
# Layering: physical communication



# Protocol layering and data

Each layer takes data from above

- adds header information to create new data unit
- passes new data unit to layer below



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# Delay in packet-switched networks

- **1. nodal processing:**

- check bit errors
- determine output link

- **2. queuing**

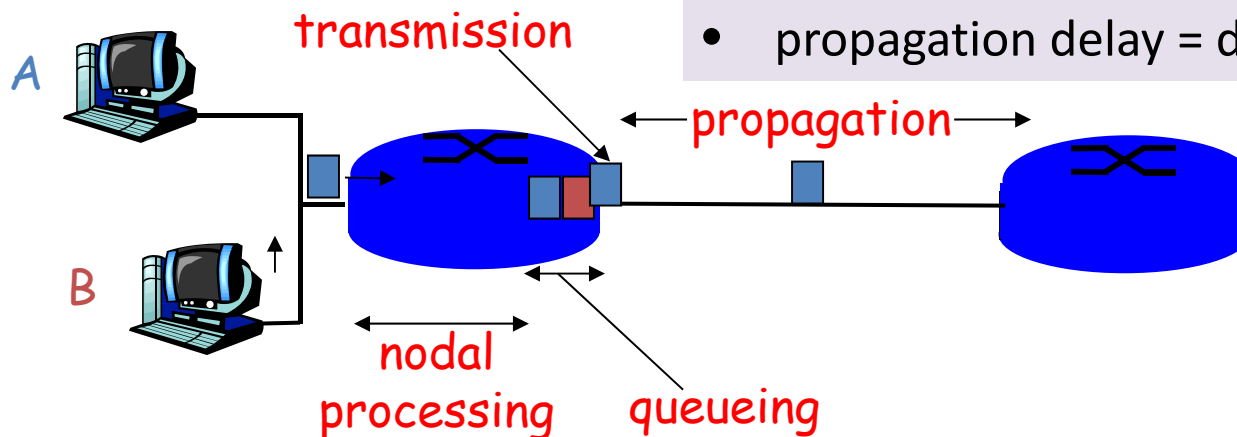
- time waiting at output link for transmission
- depends on congestion level of router

- **3. Transmission delay:**

- $R$  = link bandwidth (bps)
- $L$  = packet length (bits)
- time to send bits into link =  $L/R$

- **4. Propagation delay:**

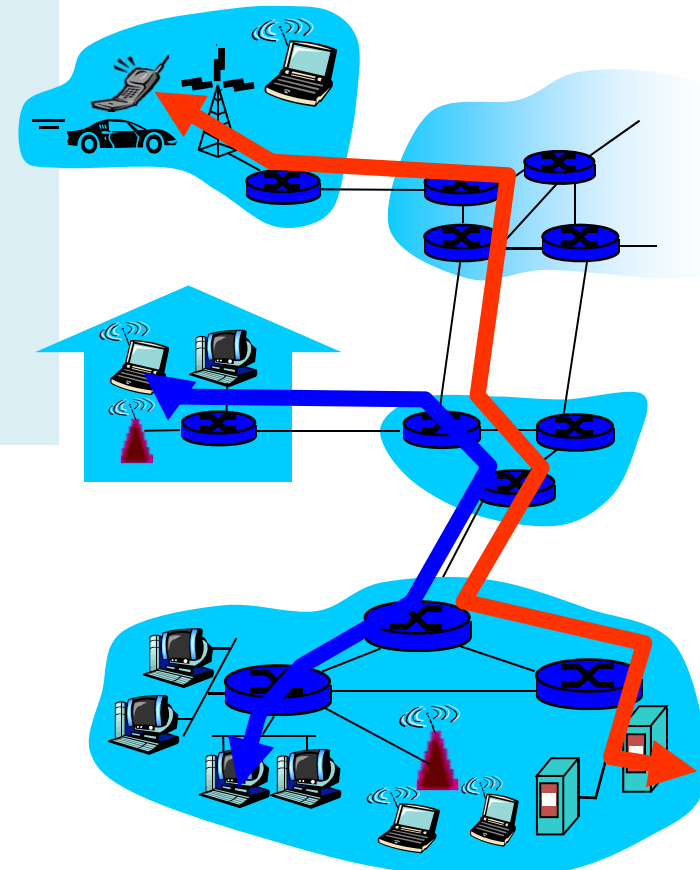
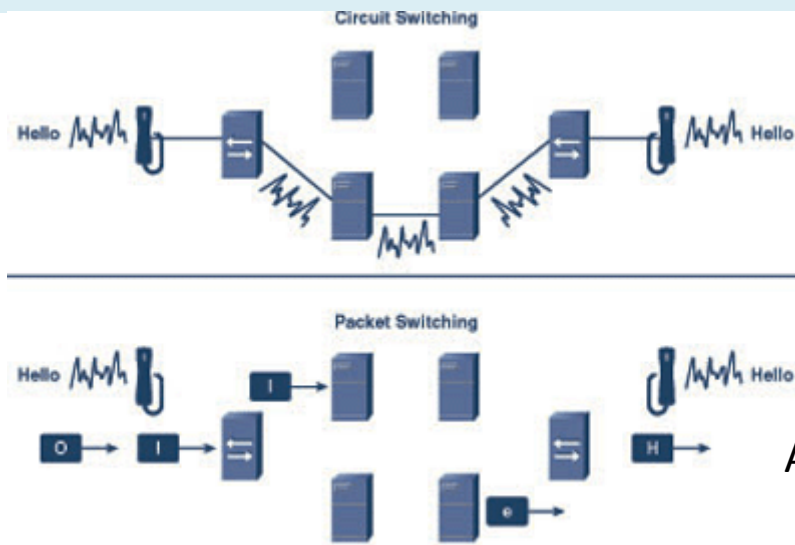
- $d$  = length of physical link
- $s$  = propagation speed in medium ( $\sim 2 \times 10^8$  m/sec)
- propagation delay =  $d/s$



# Network Core: Circuit Switching

End-end resources reserved/dedicated for “call” (analogue telephony)

- link bandwidth, switch capacity
- dedicated resources: no sharing
- circuit-like (guaranteed) performance
- call setup required



A nice video for Circuit vs packet switching  
<http://www.youtube.com/watch?v=Dq1zpiDN9k4&feature=related>

# Visualize delays: Circuit, message, packet switching

- store and forward behavior + other delays' visualization (fig. from "Computer Networks" by A. Tanenbaum,)

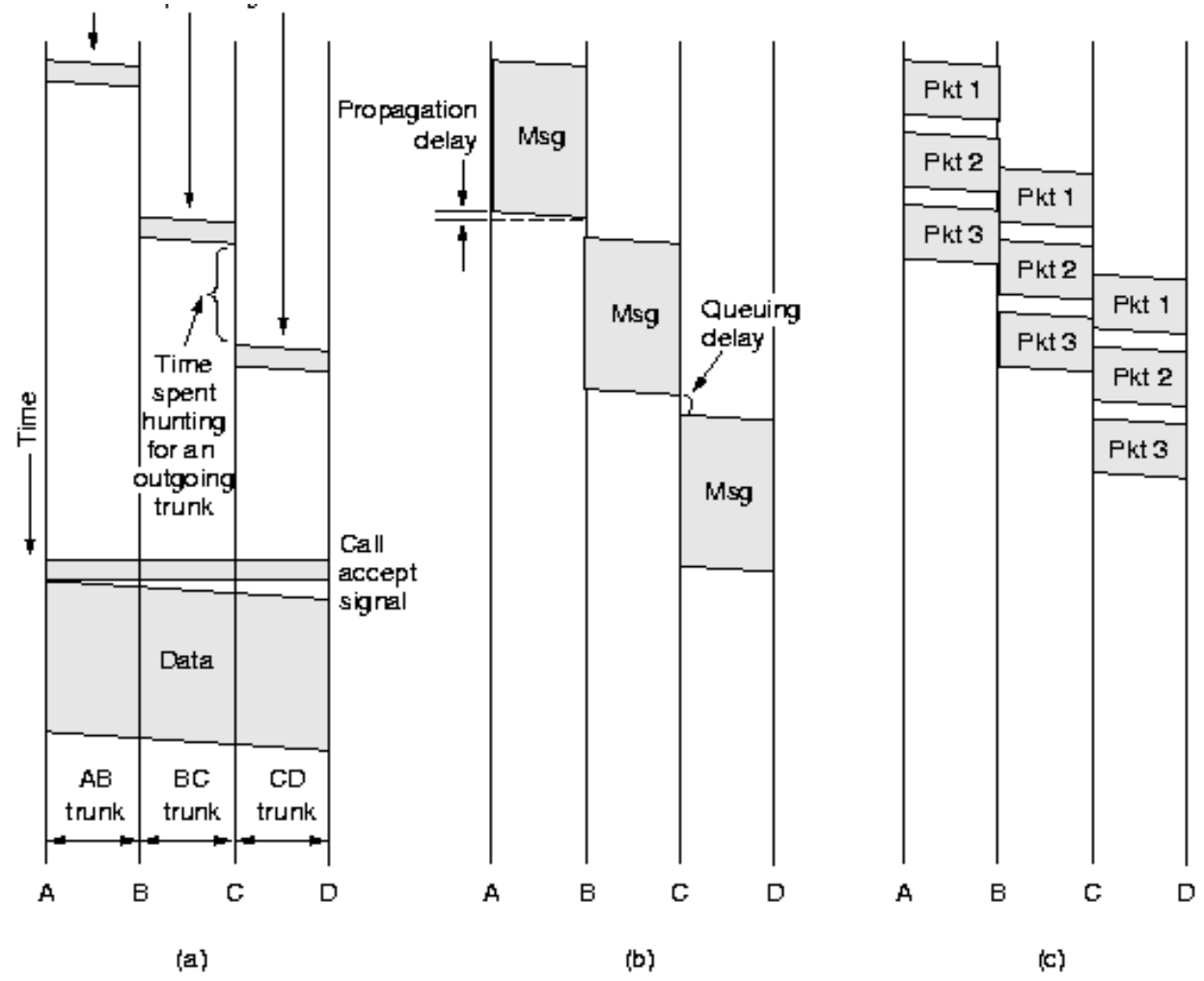
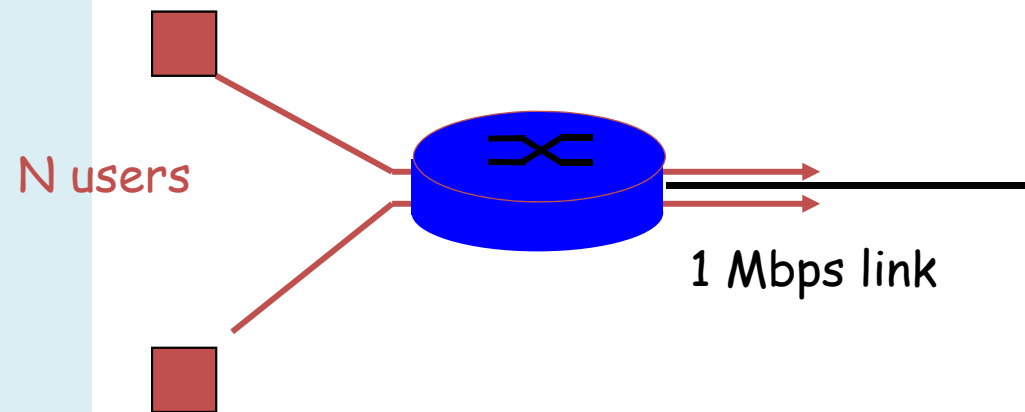


Fig. 2-35. Timing of events in (a) circuit switching, (b) message switching, (c) packet switching.

# Packet switching versus classical circuit switching

Packet switching allows more users to use the network!

- 1 Mbit link
- each user/connection:
  - 100Kbps when “active”
  - active 10% of time (bursty behaviour)
- circuit-switching how many users/connections?:
  - 10
- packet switching:
  - with 35 users, probability > 10 active less than 0.0004 ( $\Rightarrow$  almost all of the time, same the queuing behaviour is as in circuit switching)



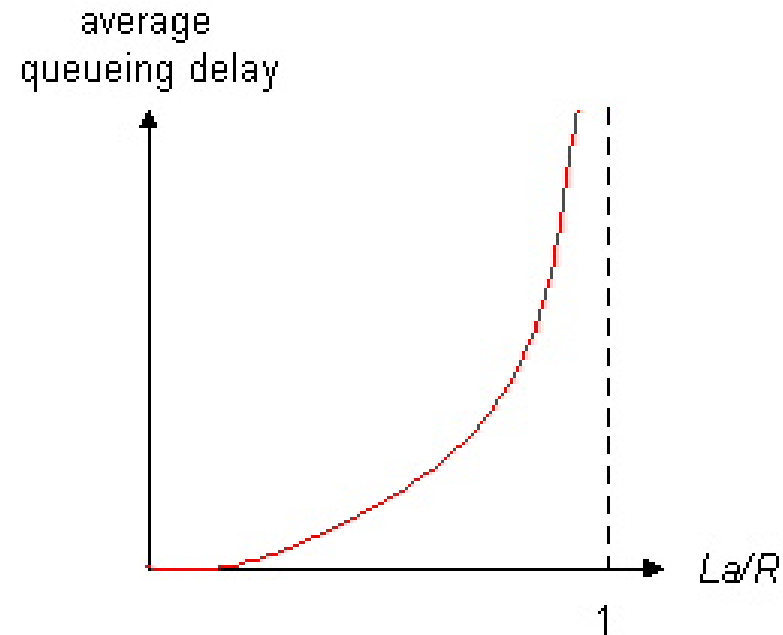
Hint: The probability of  $k$  out of  $n$  users active ( $p=0.1$  in our example)

$$f(k; n, p) = \Pr(X = k) = \binom{n}{k} p^k (1 - p)^{n-k}$$

# Queueing delay (revisited) ...

- $R$ =link bandwidth (bps)
- $L$ =packet length (bits)
- $a$ =average packet arrival rate

traffic intensity =  $La/R$



- $La/R \sim 0$ : average queueing delay small
- $La/R \rightarrow 1$ : delays become large
- $La/R > 1$ : more "work" arriving than can be serviced, average delay infinite! **Queues may grow unlimited**, packets can be **lost**



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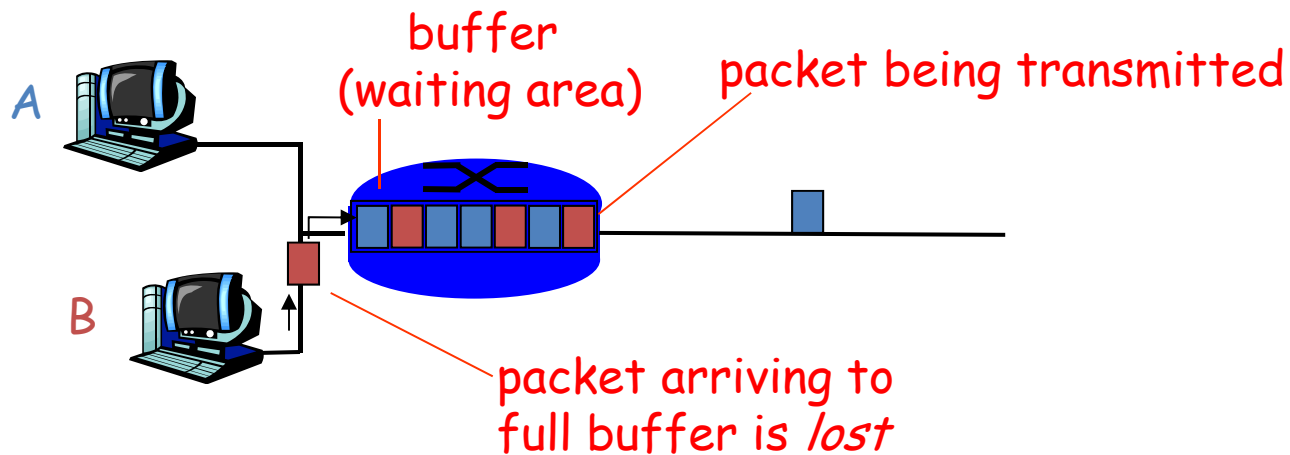
## Related with delays...

- packet loss
- throughput

# Packet loss

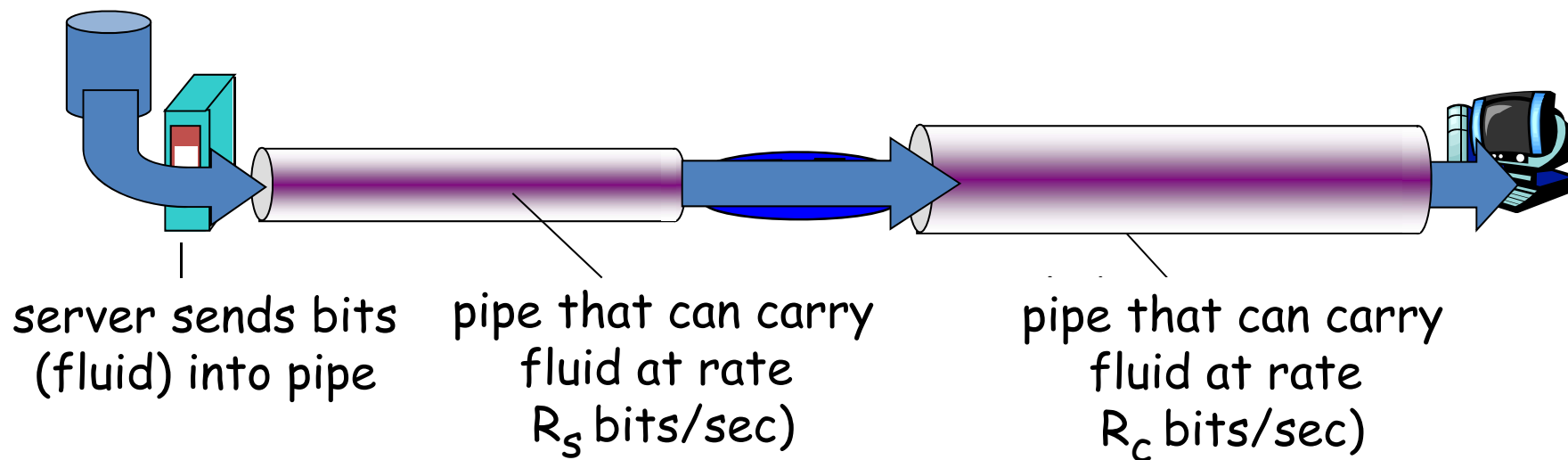
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- queue (aka buffer) preceding link has finite capacity
- packet arriving to full queue dropped (aka lost)
- lost packet may be retransmitted by previous node, by source end system, or not at all



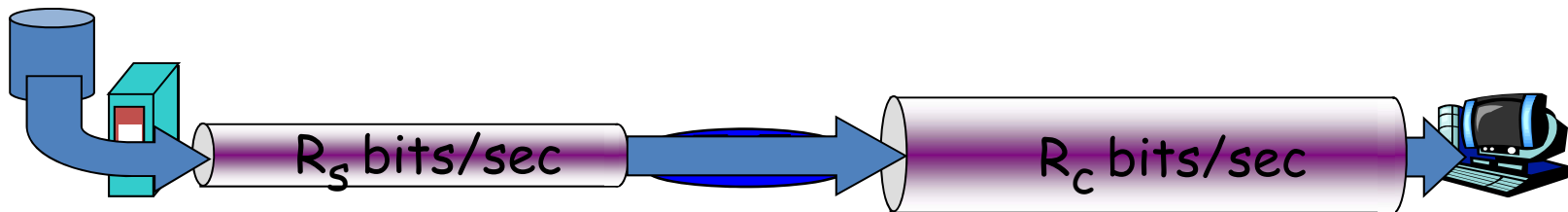
# Throughput

- *throughput*: rate (bits/time unit) at which bits transferred between sender/receiver
  - *instantaneous*: rate at given point in time
  - *average*: rate over longer period of time

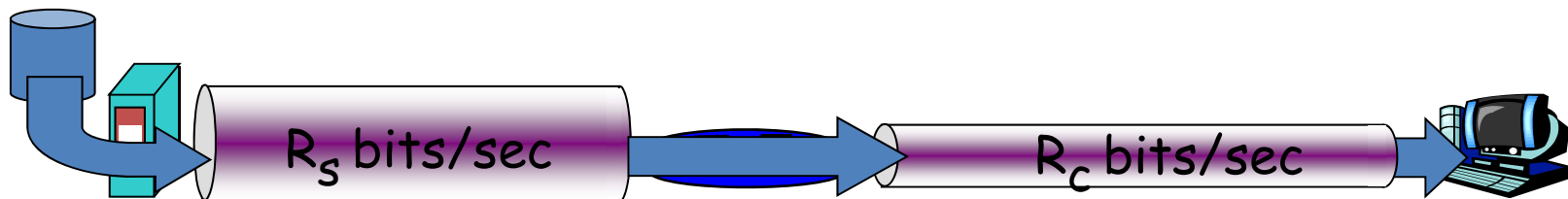


# Throughput (more)

- $R_s < R_c$  What is average end-end throughput?



- $R_s > R_c$  What is average end-end throughput?

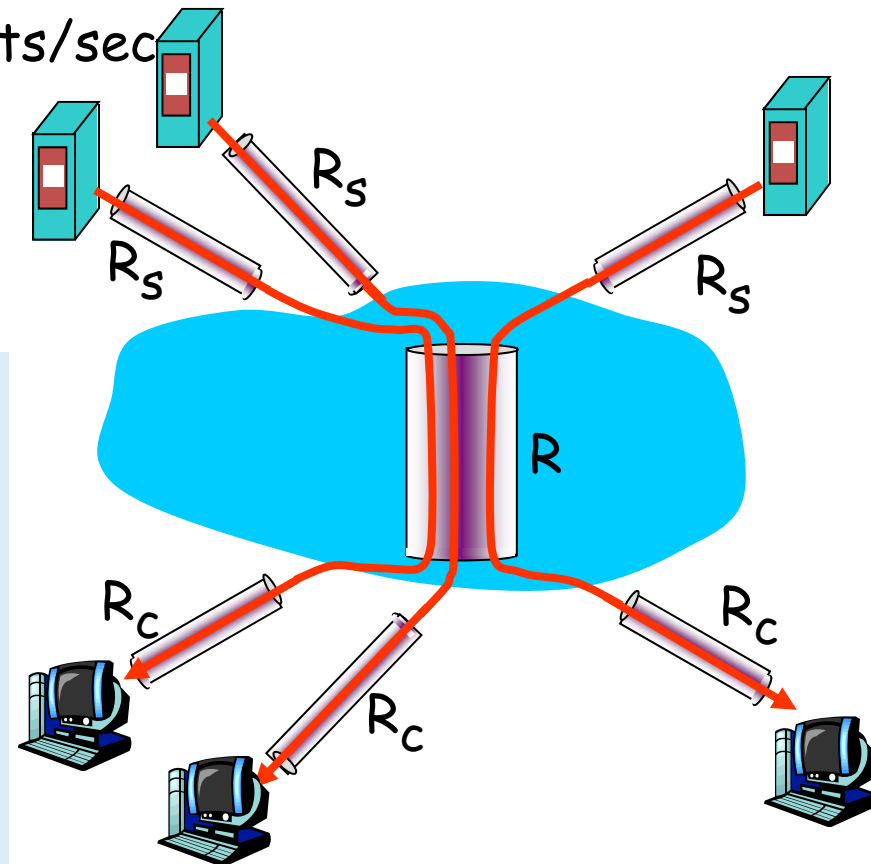


*bottleneck link*

link on end-end path that constrains end-end throughput

# Throughput: Internet scenario

10 connections (fairly) share  
backbone bottleneck link of  $R$  bits/sec

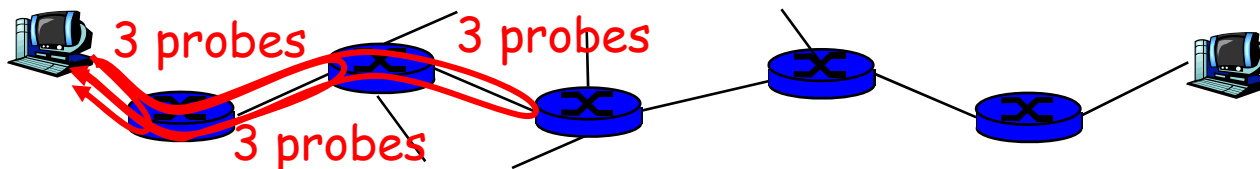


- per-connection end-end throughput:  $\min(R_c, R_s, R/10)$  (if fair)
- in practice:  $R_c$  or  $R_s$  is often bottleneck

## ... “Real” Internet delays and routes (1)...

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- What do “real” Internet delay & loss look like?
- Traceroute program: provides delay measurement from source to router along end-end Internet path towards destination. For all  $i$ :
  - sends three packets that will reach router  $i$  on path towards destination
  - router  $i$  will return packets to sender
  - sender times interval between transmission and reply.



## ...“Real” Internet delays and routes (2)...

**traceroute:** gaia.cs.umass.edu to www.eurecom.fr

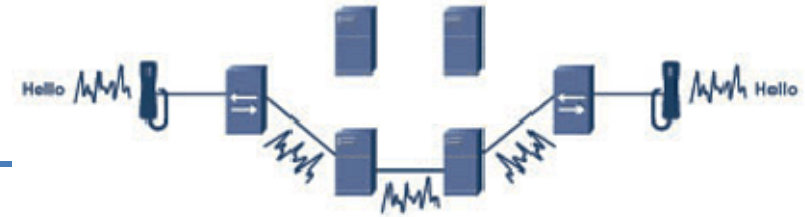
Three delay measurements from  
gaia.cs.umass.edu to cs-gw.cs.umass.edu

1	cs-gw (128.119.240.254)	1 ms	1 ms	2 ms
2	border1-rt-fa5-1-0.gw.umass.edu (128.119.3.145)	1 ms	1 ms	2 ms
3	cht-vbns.gw.umass.edu (128.119.3.130)	6 ms	5 ms	5 ms
4	jn1-at1-0-0-19.wor.vbns.net (204.147.132.129)	16 ms	11 ms	13 ms
5	jn1-so7-0-0-0.wae.vbns.net (204.147.136.136)	21 ms	18 ms	18 ms
6	abilene-vbns.abilene.ucaid.edu (198.32.11.9)	22 ms	18 ms	22 ms
7	nycm-wash.abilene.ucaid.edu (198.32.8.46)	22 ms	22 ms	22 ms
8	62.40.103.253 (62.40.103.253)	104 ms	109 ms	106 ms
9	de2-1.de1.de.geant.net (62.40.96.129)	109 ms	102 ms	104 ms
10	de.fr1.fr.geant.net (62.40.96.50)	113 ms	121 ms	114 ms
11	renater-gw.fr1.fr.geant.net (62.40.103.54)	112 ms	114 ms	112 ms
12	nio-n2.cssi.renater.fr (193.51.206.13)	111 ms	114 ms	116 ms
13	nice.cssi.renater.fr (195.220.98.102)	123 ms	125 ms	124 ms
14	r3t2-nice.cssi.renater.fr (195.220.98.110)	126 ms	126 ms	124 ms
15	eurecom-valbonne.r3t2.ft.net (193.48.50.54)	135 ms	128 ms	133 ms
16	194.214.211.25 (194.214.211.25)	126 ms	128 ms	126 ms
17	* * *			
18	* * *			
19	fantasia.eurecom.fr (193.55.113.142)	132 ms	128 ms	136 ms

trans-oceanic link

\* means no reponse (probe lost, router not replying)

# Packet switching properties

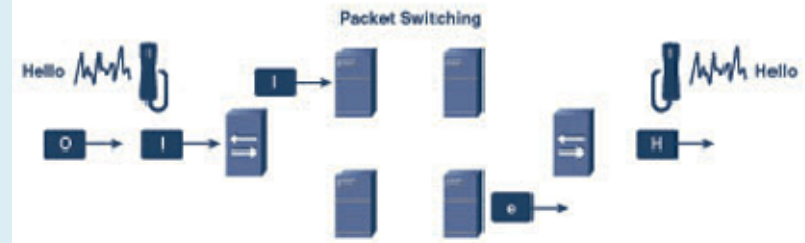


- **PS: Good:** Great for bursty data

- resource sharing
- no call setup

- **PS: Not so good? Excessive congestion:** packet delay and loss

- protocols needed for reliable data transfer, congestion control



- **Q: How to provide circuit-like behavior?**

- bandwidth guarantees are needed for some apps
- Some **routing policies** can help???

  - Cf virtual circuit



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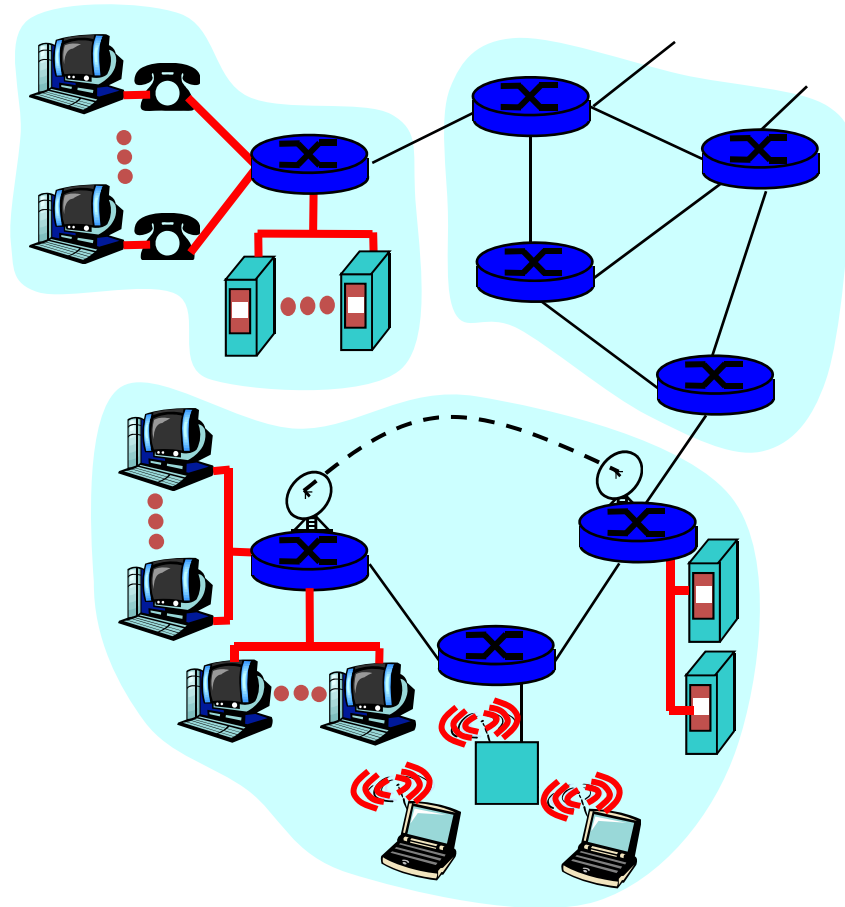
# Access networks and physical media

*Q: How to connect end systems to edge router?*

- residential access nets
- institutional access networks (school, company)
- mobile access networks

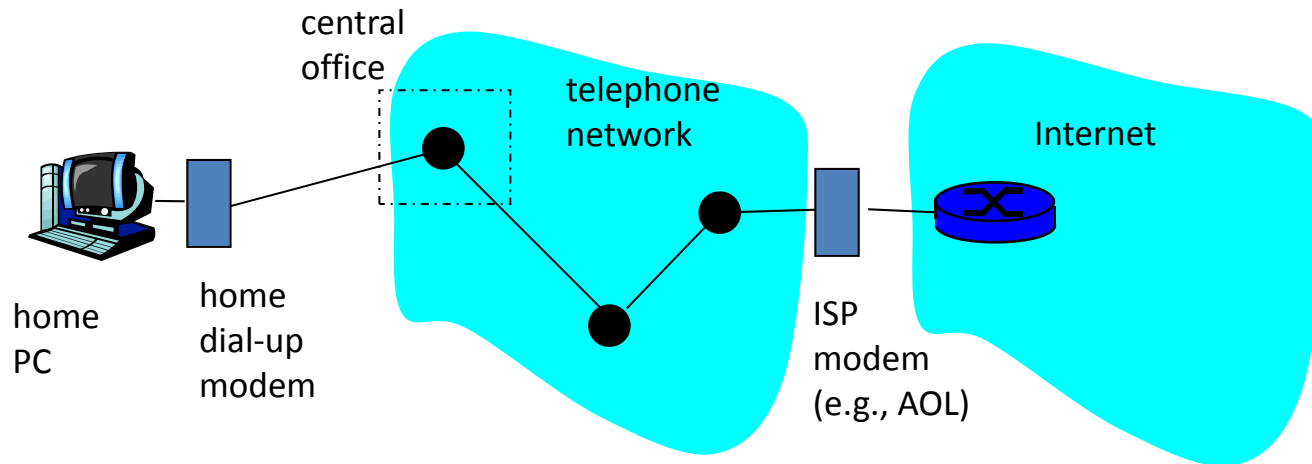
*Considerations:*

- bandwidth (bits per second) of access network?
- shared or dedicated?



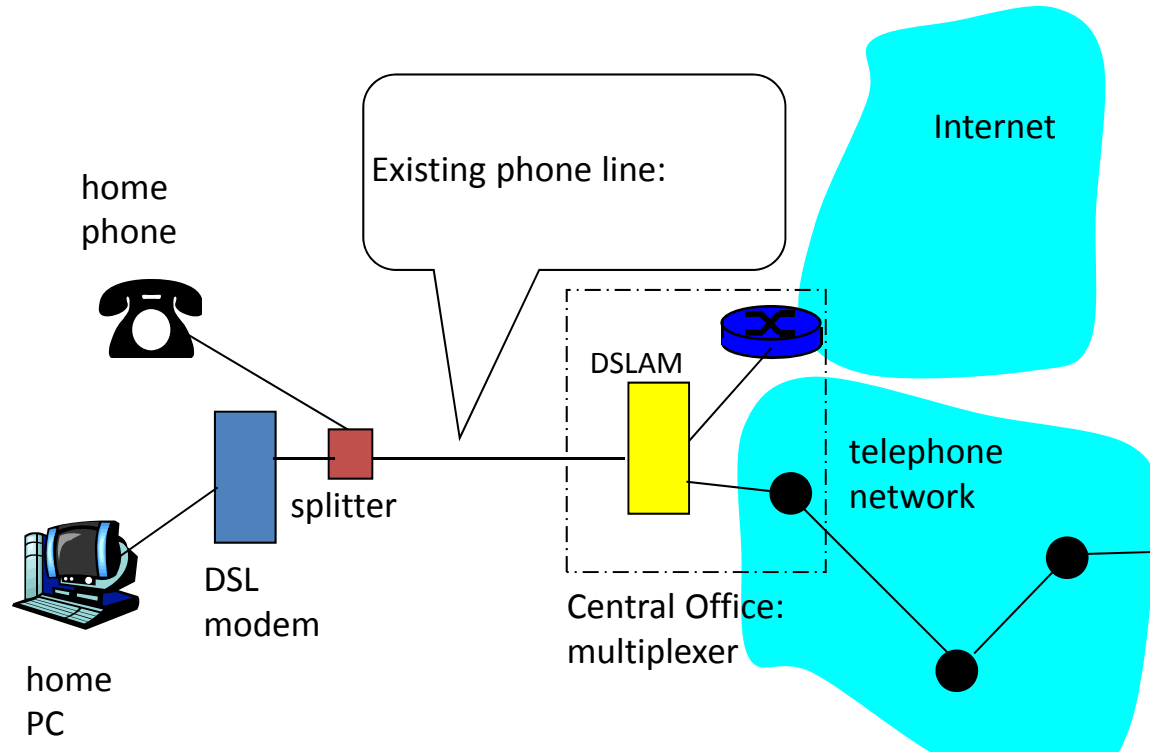
# Dial-up Modem

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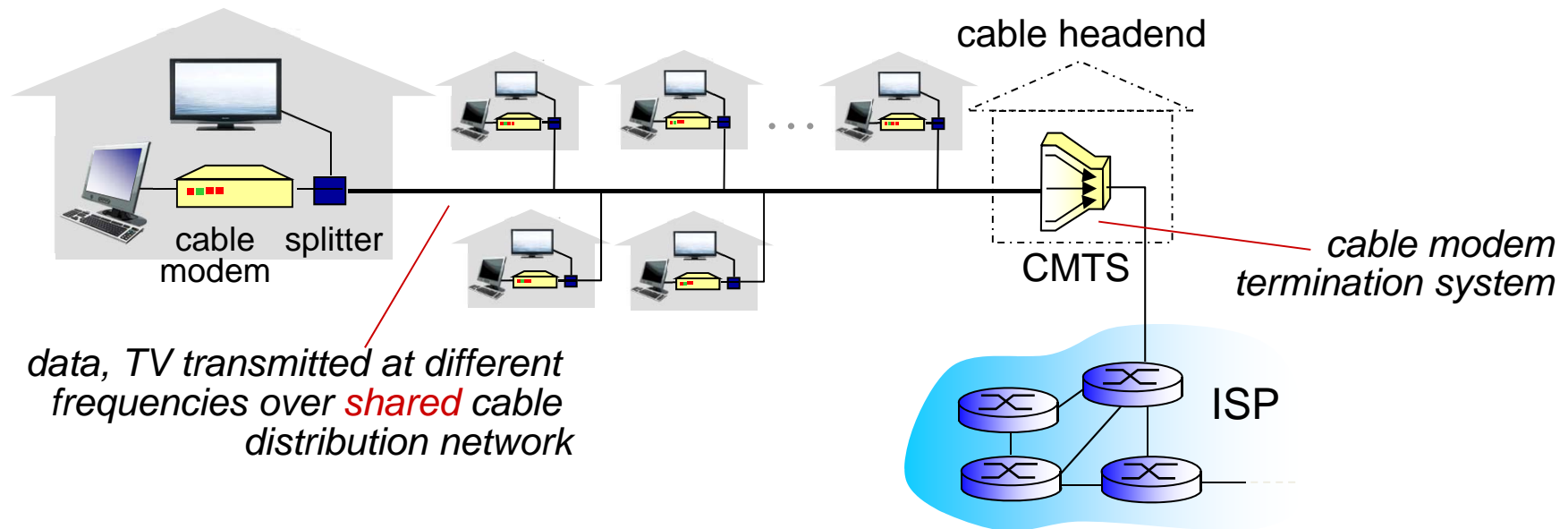
- ❖ Uses existing telephony infrastructure
  - ❖ Home is connected to **central office**
- ❖ up to 56Kbps direct access to router (often less)
- ❖ Can't surf and phone at same time: not **"always on"**

# Digital Subscriber Line (DSL)



- ❖ Also uses existing telephone infrastructure
  - ❖ dedicated physical line to telephone central office
- ❖ Asymmetric: commonly up to 2.5 Mbps upstream; up to 24 Mbps downstream

# Access net: cable network

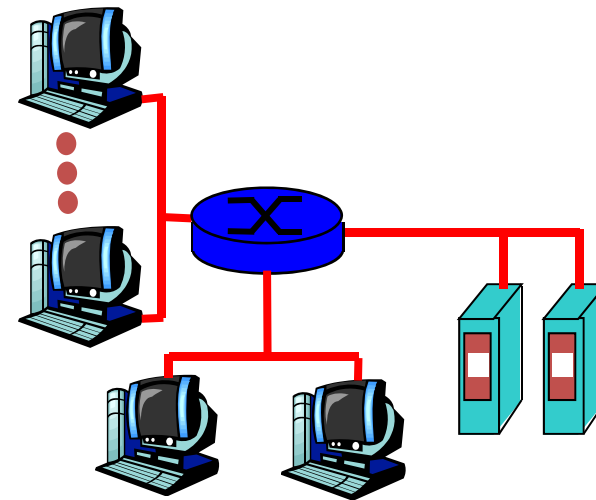


- ❖ **HFC: hybrid fiber coax**
  - asymmetric: up to 30Mbps downstream transmission rate, 2 Mbps upstream transmission rate
- ❖ **network** of cable, attaches homes to ISP router
  - homes *share access network* to cable headend
  - unlike DSL, which has dedicated access to central office
  - Related: **fiber to the home**

# Institutional access: local area networks

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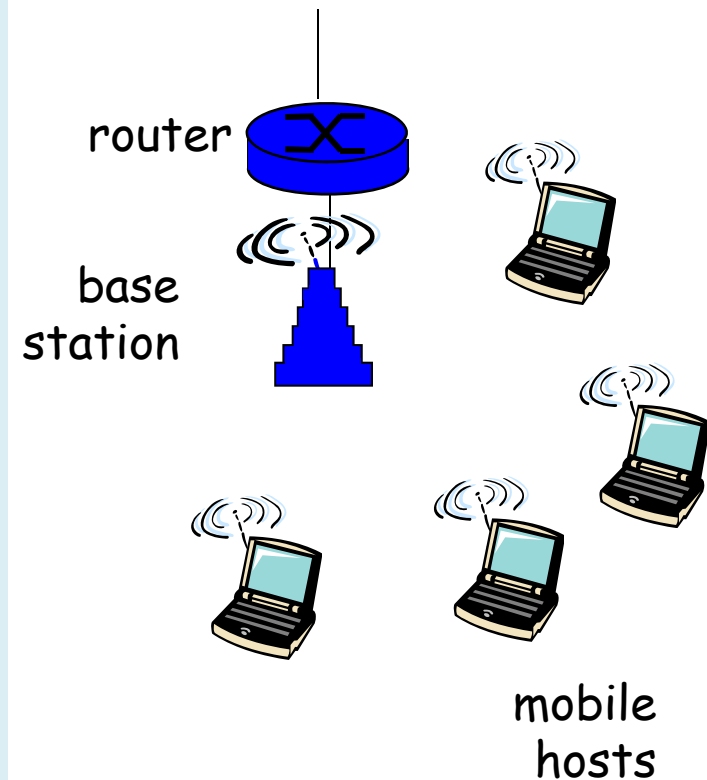
- **local area network (LAN)** connects end system to edge router
- **E.g. Ethernet:**
  - 10 Mbps, 100Mbps, Gigabit Ethernet
- **deployment:** institutions, home LANs



(a lot more on this later in the course)

# Wireless access networks

- shared *wireless* access network connects end system to router
  - via base station aka “access point”, or “ad hoc”
- **wireless LANs:**
  - 802.11b/g (WiFi): 11 or 54 Mbps
- **wider-area wireless access**
  - provided by telco operator
  - ~1Mbps over cellular system
  - WiMAX (10’s Mbps) over wide area

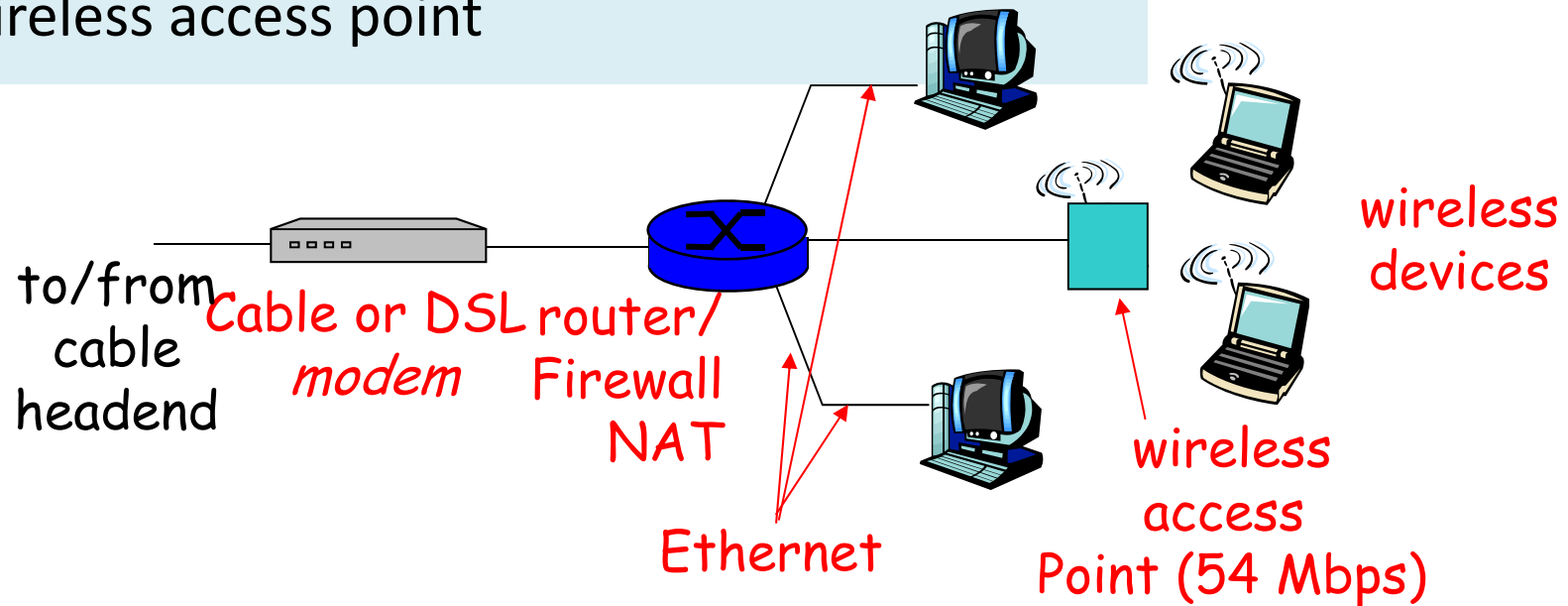


(a lot more on this later in the course)

# Home networks

## Typical home network components:

- DSL or cable modem or Fiber to the home
- router/firewall/NAT
- Ethernet
- wireless access point





# Physical Media

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- **physical link:** transmitted data bit propagates across link
  - **guided media:**
    - signals propagate in solid media: copper, fiber
  - **unguided media:**
    - signals propagate freely e.g., radio

# Physical media: wireless

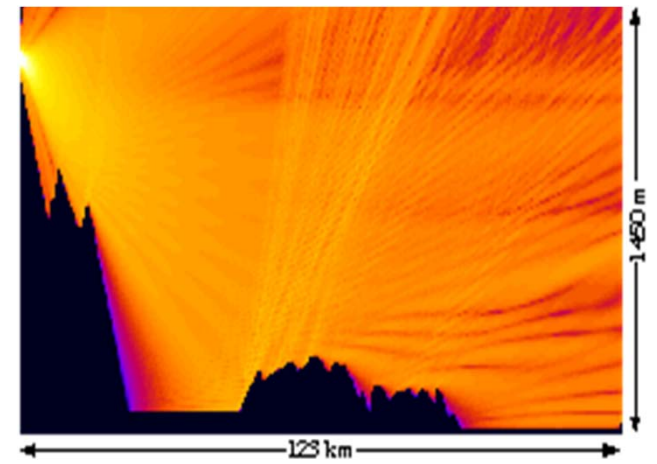
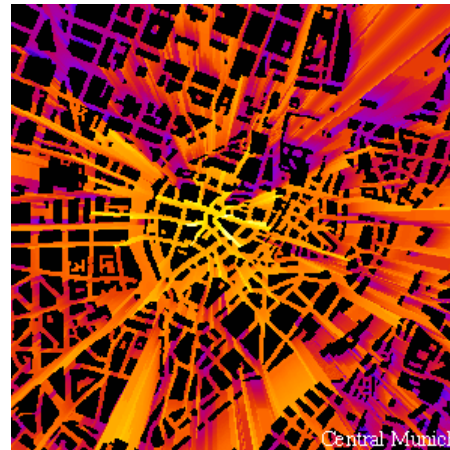
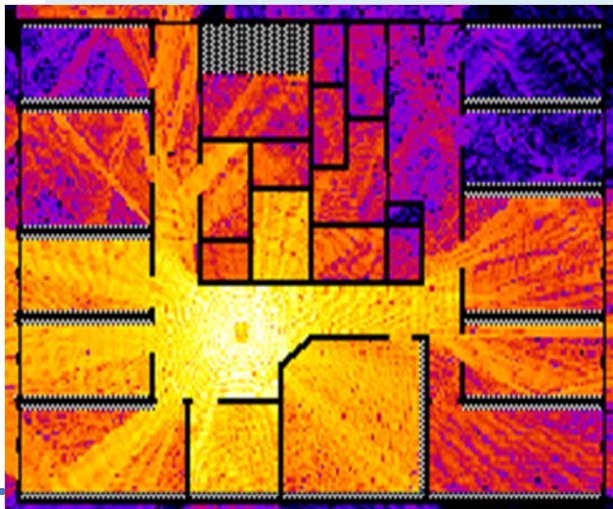
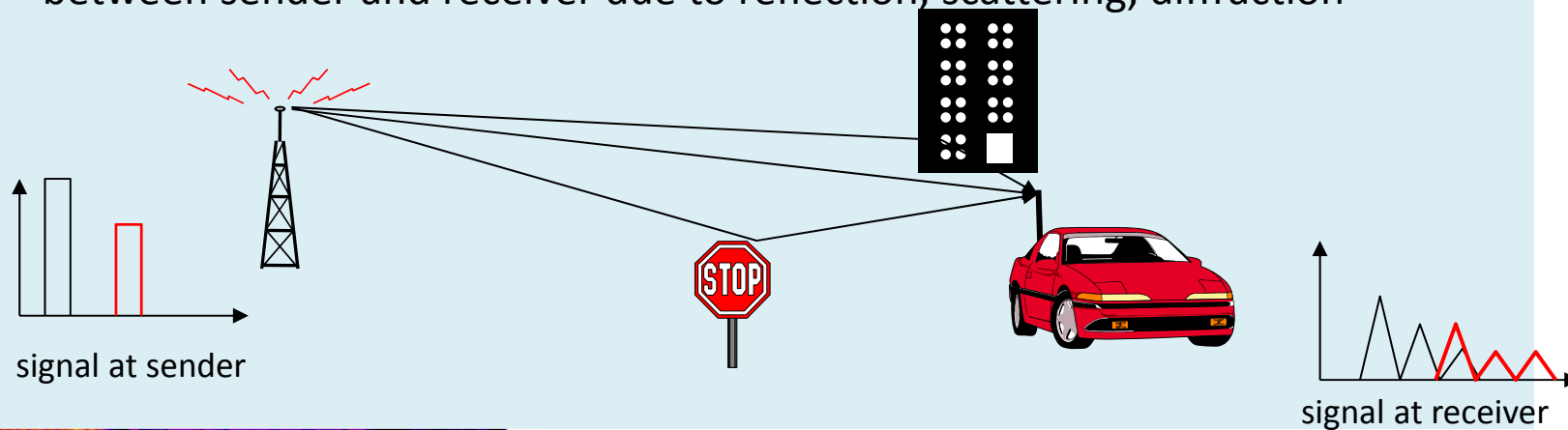
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- signal carried in electromagnetic spectrum
- **Omnidirectional**: signal spreads, can be received by many antennas
- **Directional**: antennas communicate with focused el-magnetic beams and must be aligned (requires higher frequency ranges)
- propagation environment effects:
  - reflection
  - obstruction by objects
  - interference

# Wireless: properties

## Attenuation, Multipath propagation

Signal can fade with distance, can get obstructed, can take many different paths between sender and receiver due to reflection, scattering, diffraction



# Physical Media: coax, fiber, twisted pair

## Coaxial cable:

- wire (signal carrier) within a wire (shield)
  - baseband: single channel on cable (common use in 10Mbps Ethernet)
  - broadband: multiple channels multiplexed on cable (commonly used for cable TV)



## Fiber optic cable:

- glass fiber carrying light pulses
- low attenuation
- high-speed operation:
  - 100Mbps Ethernet
  - high-speed point-to-point transmission (e.g., 5 Gps)
- low error rate



## Twisted Pair (TP)

- two insulated copper wires
  - Category 3: traditional phone wires, 10 Mbps Ethernet
  - Category 5 TP: more twists, higher insulation: high-speed Ethernet



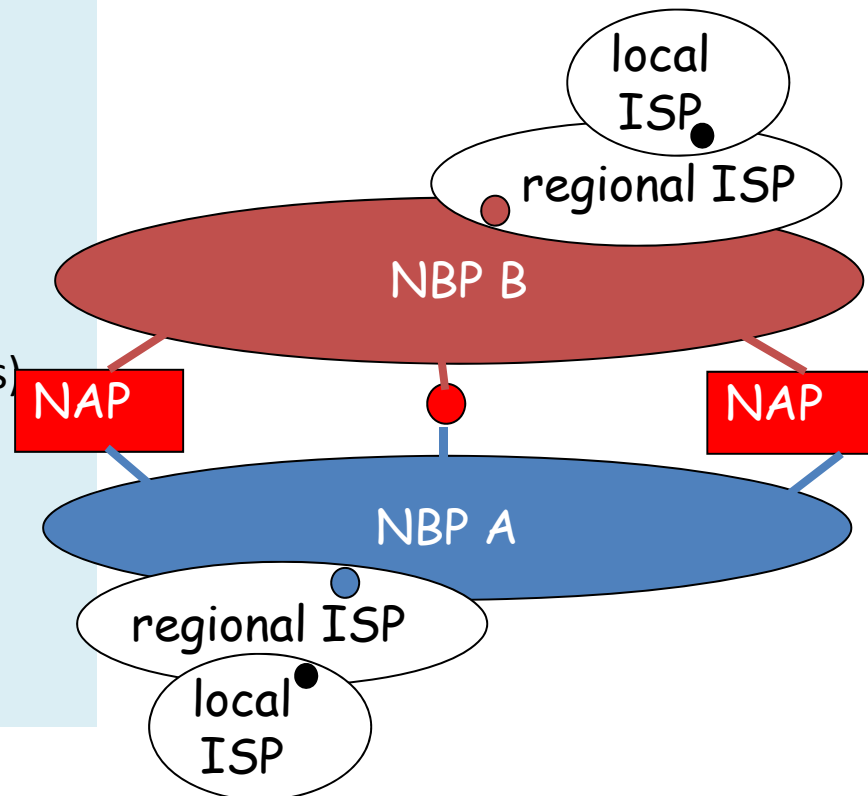
# Roadmap



- what's the Internet
- protocol layers
  - Communication through layers
- edge & core of any big network:
  - types of service, ways of information transfer, routing
- Internet layers & Logical vs physical communication
- Performance:
  - delays, loss
- Network/Internet structure complemented:
  - access net, physical media
  - backbones, NAPs, ISPs
- Security prelude

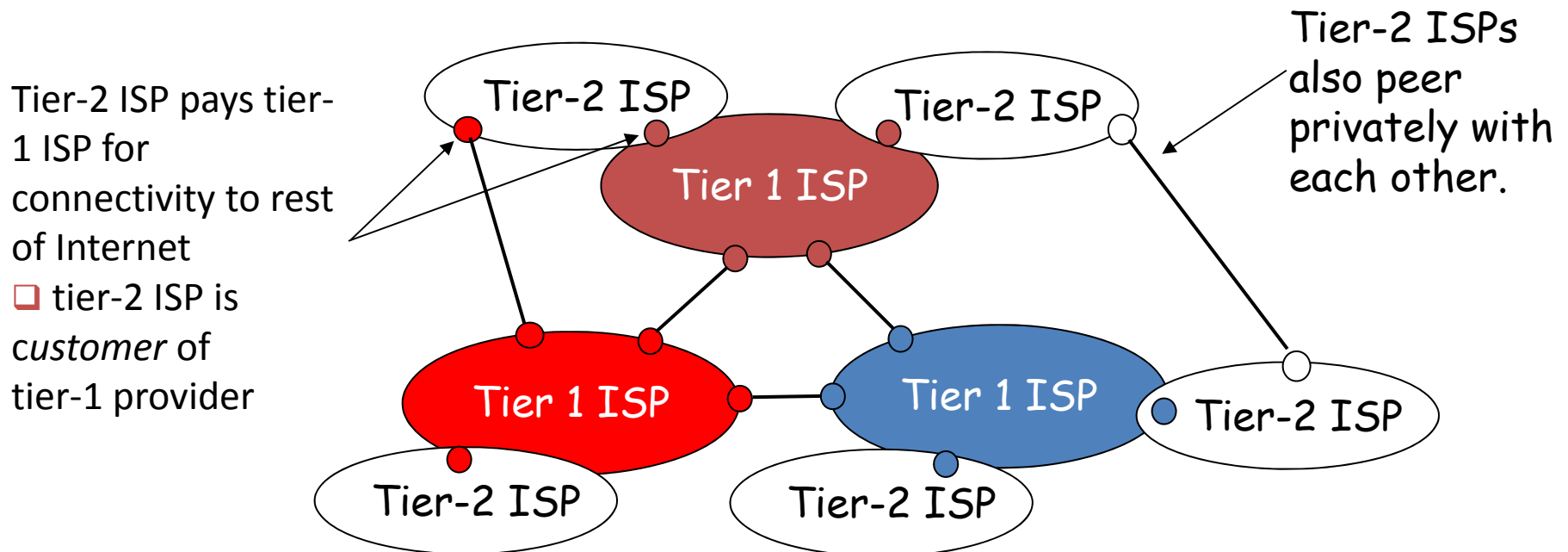
# Internet structure: network of networks

- roughly hierarchical
- national/international backbone providers (NBPs)- tier 1 providers
  - e.g. BBN/GTE, Sprint, AT&T, IBM, UUNet/Verizon, TeliaSonera
  - interconnect (peer) with each other privately, or at public Network Access Point (NAPs: routers or NWs of routers)
- regional ISPs, tier 2 providers
  - connect into NBPs; e.g. Tele2
- local ISP, company
  - connect into regional ISPs



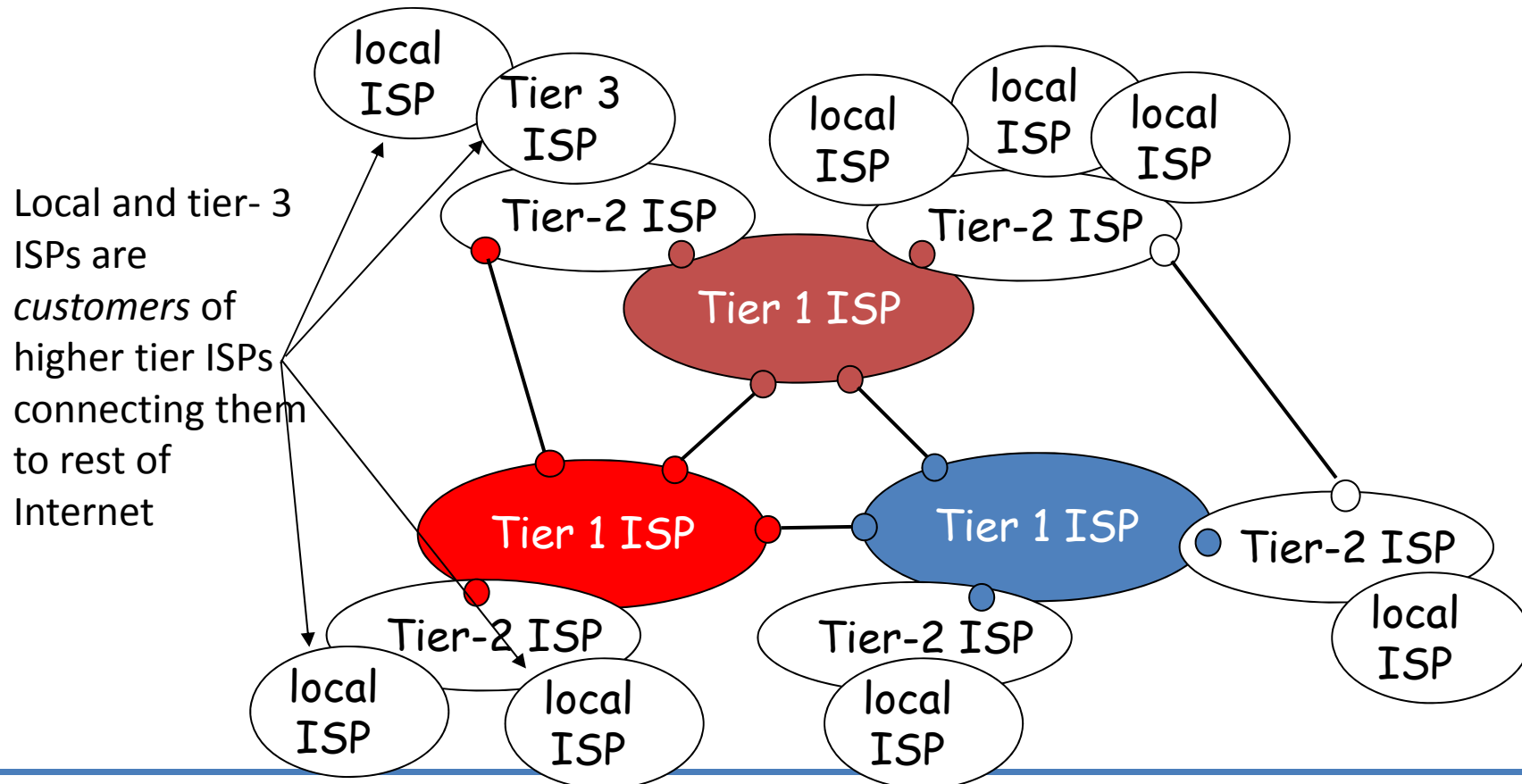
# Internet structure: network of networks

- “Tier-2” ISPs: smaller (often regional) ISPs
  - Connect to one or more tier-1 ISPs, possibly other tier-2 ISPs



# Internet structure: network of networks

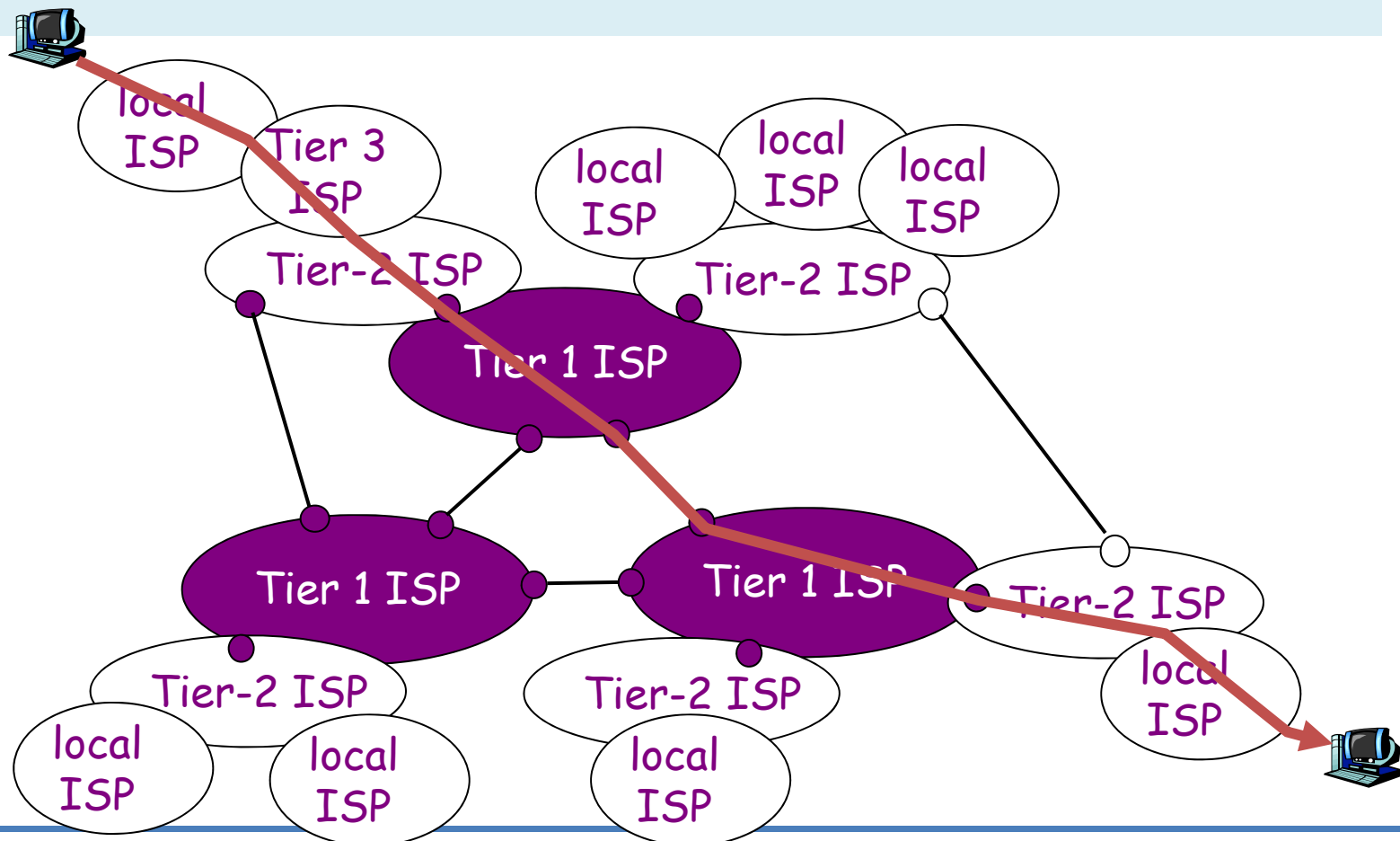
- “Tier-3” ISPs and local ISPs
  - last hop (“access”) network (closest to end systems)





# Internet structure: network of networks

- a packet passes through many networks



# Roadmap

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# Network Security

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- **The field of network security is about:**
  - how adversaries can attack computer networks
  - how we can defend networks against attacks
  - how to design architectures that are immune to attacks
- **Internet not originally designed with (much) security in mind**
  - *original vision*: “a group of mutually trusting users attached to a transparent network” 😊
  - Internet protocol designers playing “catch-up”
  - Security considerations in all layers!

# Bad guys can put malware into hosts via Internet

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- Malware can get in host from a **virus, worm, or trojan horse**.
- **Spyware malware** can record keystrokes, web sites visited, upload info to collection site.
- Infected host can be enrolled in a **botnet**, used for spam and DDoS attacks.
- Malware is often **self-replicating**: from an infected host, seeks entry into other hosts

# Example types of malware

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- **Trojan horse**

- Hidden part of some otherwise useful software
- Today often on a Web page (Active-X, plugin)

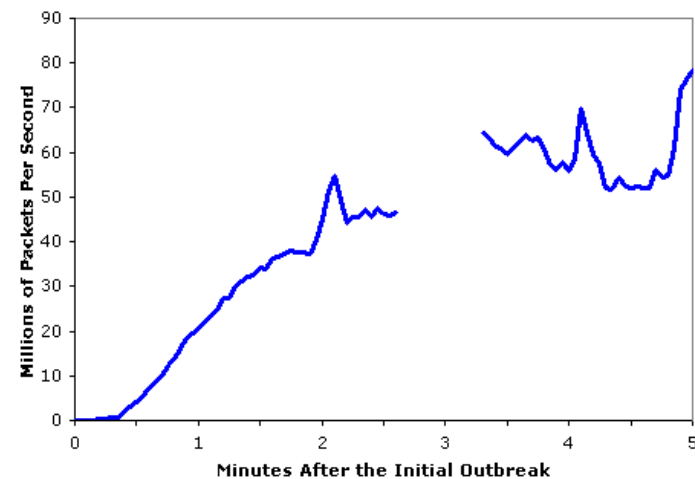
- **Virus**

- infection by receiving object (e.g., e-mail attachment), actively executing
- self-replicating: propagate itself to other hosts, users

- **Worm:**

- ❖ infection by passively receiving object that gets itself executed
- ❖ self-replicating: propagates to other hosts, users

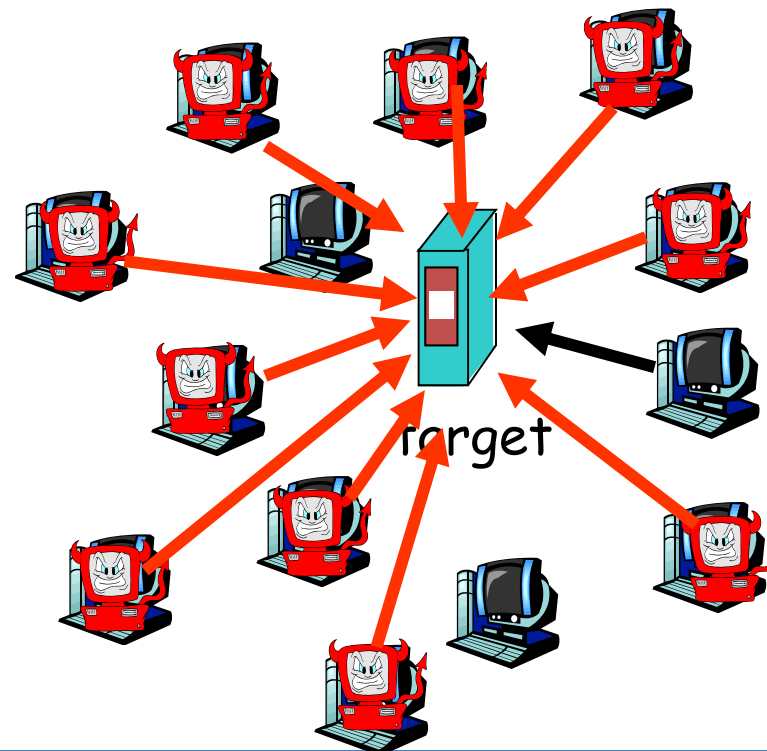
Sapphire Worm: aggregate scans/sec in first 5 minutes of outbreak (CAIDA, UWisc data)



# Bad guys can attack servers and network infrastructure

- Denial of service (DoS): attackers make resources (server, bandwidth) unavailable to legitimate traffic by overwhelming resource with bogus traffic

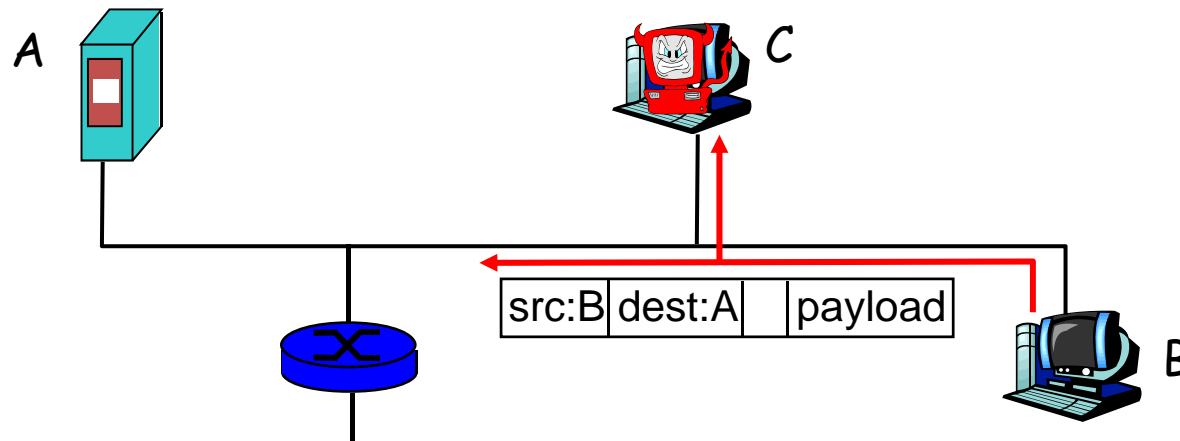
1. select target
2. break into hosts around the network (see botnet)
3. send packets toward target from compromised hosts



# The bad guys can sniff packets

## *Packet sniffing:*

- broadcast media (shared Ethernet, wireless)
- promiscuous network interface reads/records all packets (e.g., including passwords) passing by

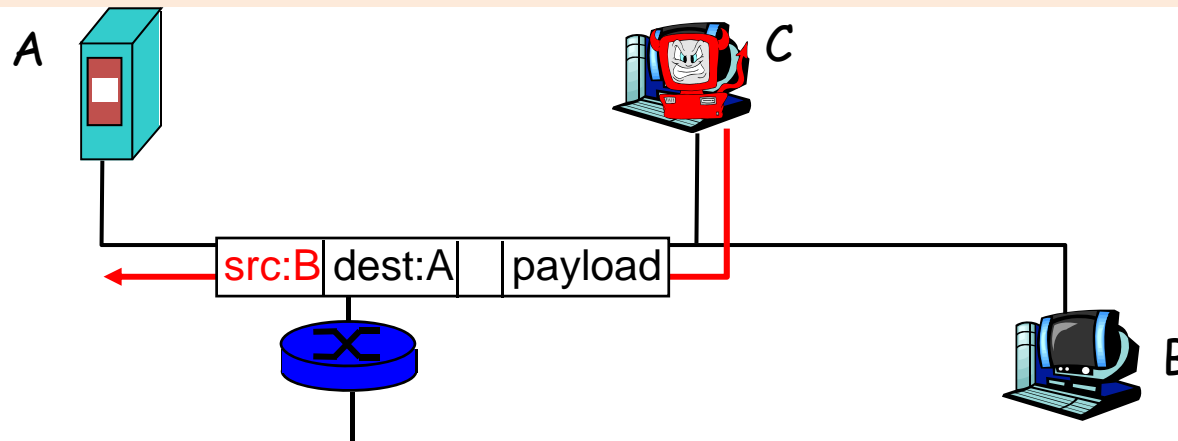


- ❖ Wireshark software used for end-of-chapter labs is a (free) packet-sniffer

# The bad guys can use false source addresses

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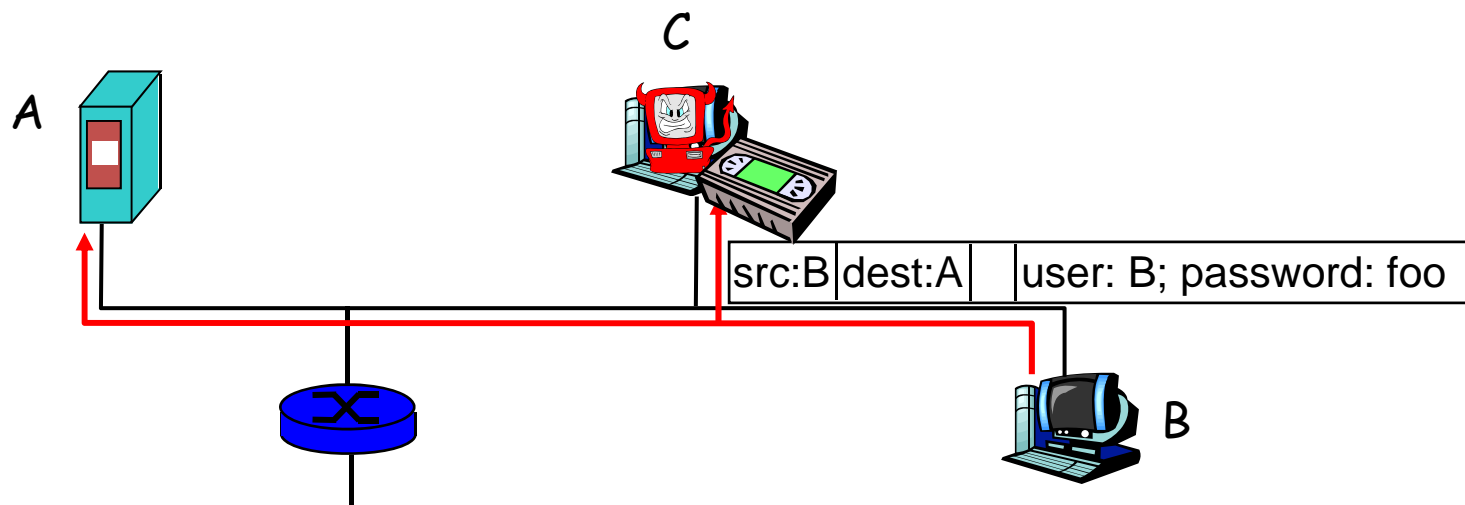
- *IP spoofing*: send packet with false source address





# The bad guys can record and playback

- *record-and-playback*: sniff sensitive info (e.g., password), and use later
  - password holder *is* that user from system point of view



# Roadmap

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# Chapter 1: Summary

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## Covered a “ton” of material!

- what’s the Internet
- what’s a protocol?
- protocol layers, service models
- network edge (types of service)
- network core (ways of transfer, routing)
  
- performance, delays, loss
- access net, physical media
- backbones, NAPs, ISPs
- Security concerns
  
- (history: read more corresponding section, interesting & fun 😊)

## You now hopefully have:

- context, overview, “feel” of networking
- A point of reference for context in the focused discussions to come

# Reading instructions

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## 1. Kurose Ross book

### Careful

4/e,5/e,6/e: 1.3, 1.4, 1.5

### Quick

4/e,5/e,6/e: the rest

Extra Reading (optional)

Computer and Network Organization: An Introduction,  
by Maarten van Steen and Henk Sips, Prentice Hall  
(very good introductory book for non-CSE students)

# Review questions

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Review questions from Kurose-Ross book, chapter 1 (for basic study)

- R11, R12, R13, R16, 17, R18, R19, R20, R21, R22, R23, R24, R25, R28.

Extra questions, for further study: delay analysis in packet switched networks:

<http://www.comm.utoronto.ca/~jorg/teaching/ece466/material/466-SimpleAnalysis.pdf>