

A (really) simple introduction to buffer overflows

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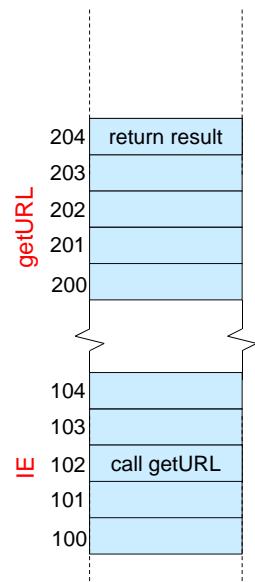


Exploits

- program has a security hole
- exploit = input that abuses the vulnerability
- In this module we will discuss an example:
the Buffer overflow

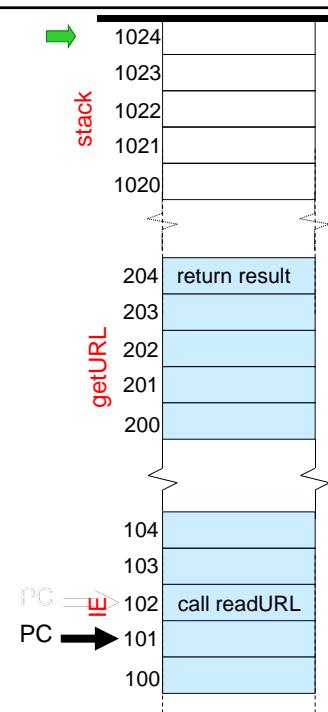
software

- sequence of instructions in memory
- logically divided in functions that call each other
 - function 'IE' calls function 'getURL' to read the corresponding page
- in CPU, the program counter contains the address in memory of the next instruction to execute
 - normally this is the next address (instruction 100 is followed by instruction 101, etc)
 - not so with function call



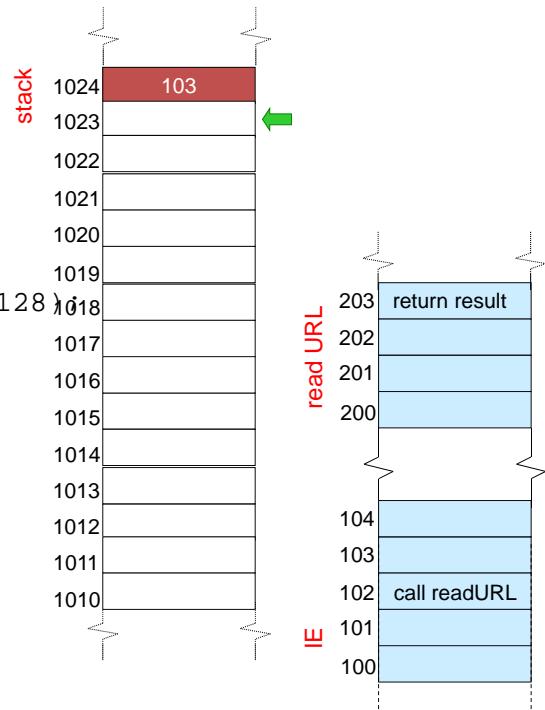
software

- so how does our CPU know where to return?
 - it keeps administration on a 'stack'



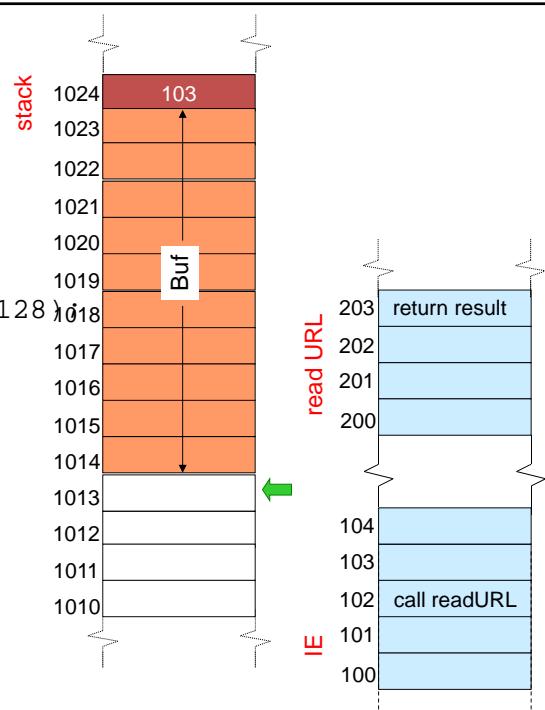
real functions have variables

```
getURL ()
{
    char Buf[10];
    read(keyboard,Buf,128)
    get_webpage (Buf);
}
IE ()
{
    getURL ();
}
```



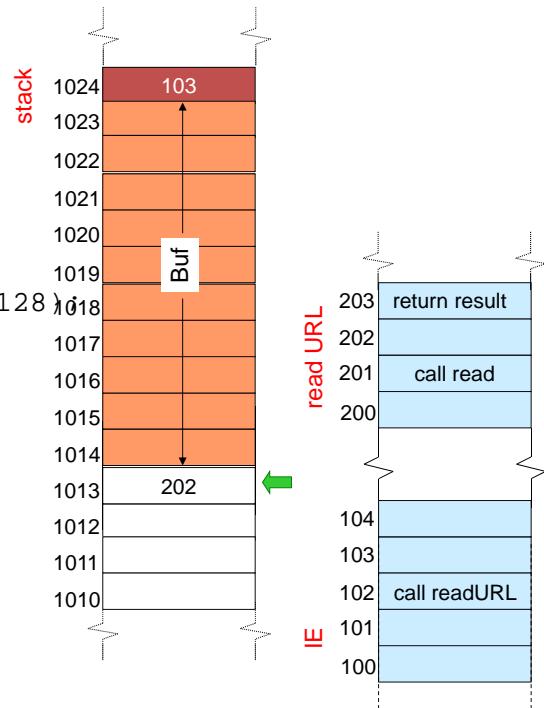
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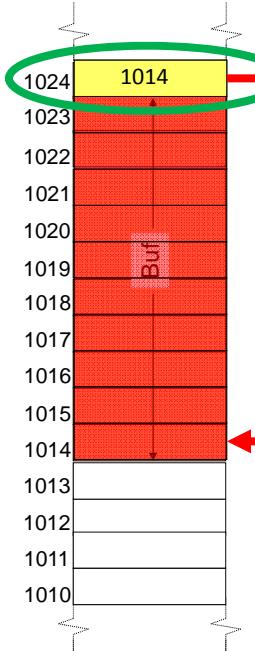


what is next?

- we have learned a lot
- but where are the vulnerabilities?
- and how do we exploit them?

Exploit

```
getURL ( )
{
    char Buf[10];
    → read(keyboard,Buf,128);
    get_webpage (Buf);
}
IE ( )
{
    getURL ( );
}
```



That is it, really

- all we need to do is stick our program in the buffer