Time triggered real time communication

Presentation overview

• Background

automotive electronics, an application area for time triggered communication.

• Time triggered protocols

TTPC, first commercial implementation. Originally from TU Vienna. Operational in civil aircrafts.

TTCAN, based on *Controller Area Network* (CAN) which is widely used in today's vehicular electronic systems.

FlexRay, based on BMW's "ByteFlight". Anticipated in next generation automotive electronic systems.

• Hybrid scheduling

combining static scheduling with fixed priority scheduling analysis.

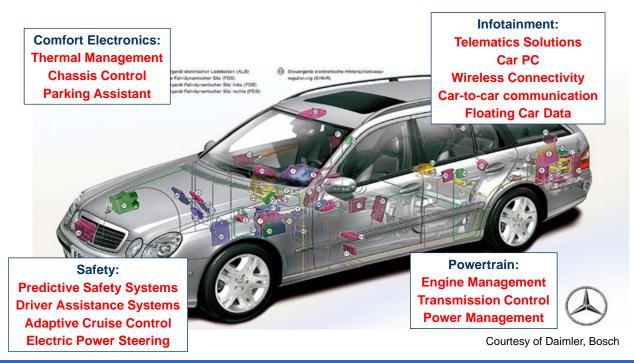
Time triggered real time communication

CHALMERS

Roger Johansson/2011

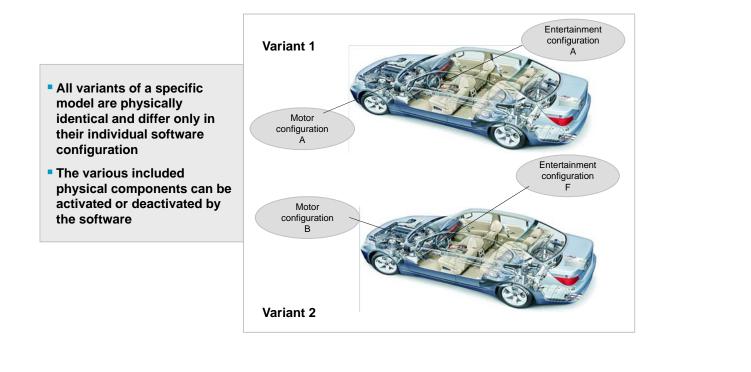
1

A premium passenger car is controlled and managed by 80+ Embedded Systems

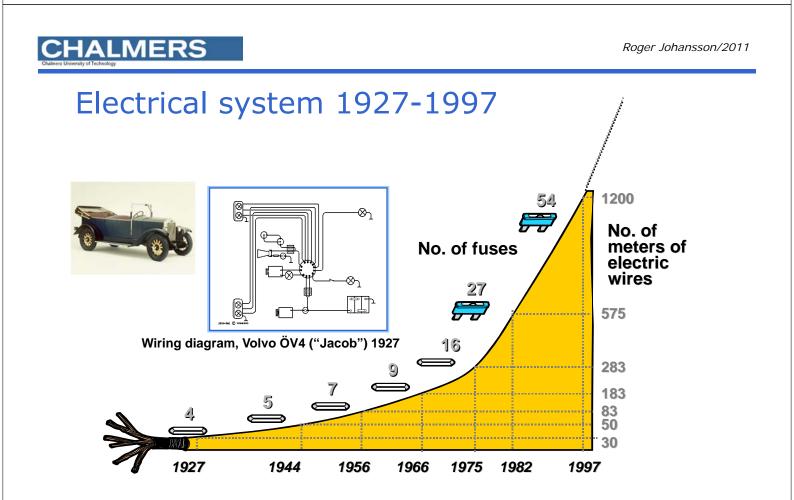


Time triggered real time communication

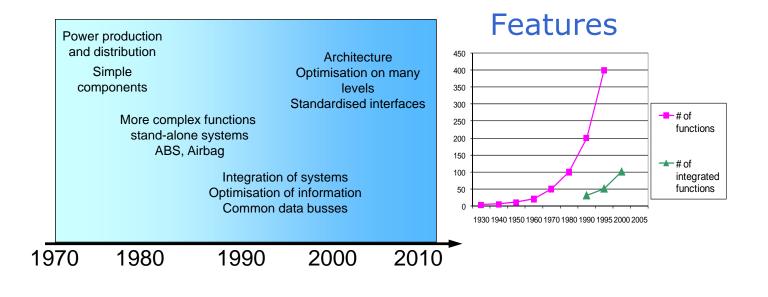
Virtual differentiation between variants



Time triggered real time communication



The evolution of the electrical system

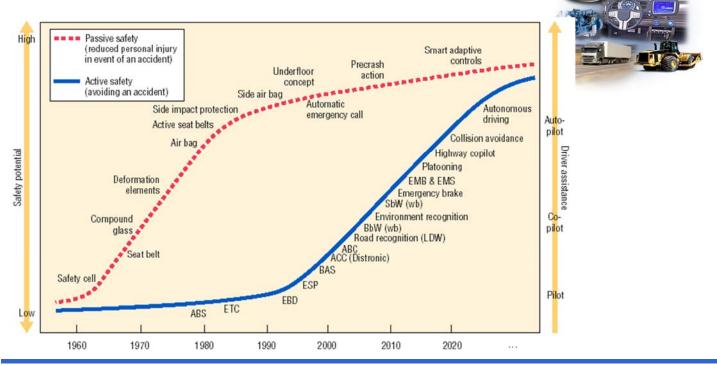


Time triggered real time communication

Roger Johansson/2011

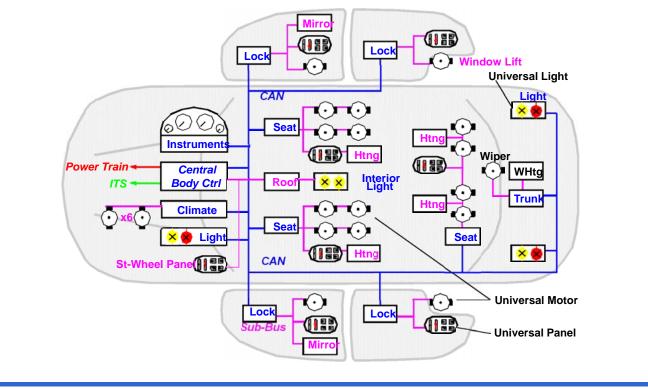
5

Automotive electronics roadmap



Time triggered real time communication

An electrical system...

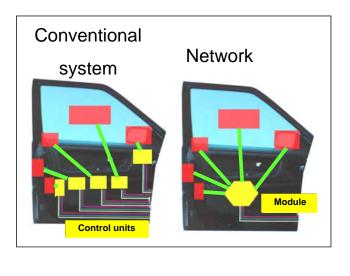


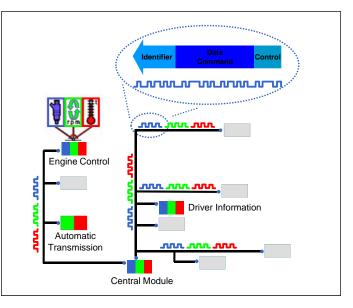
Time triggered real time communication

Roger Johansson/2011

7

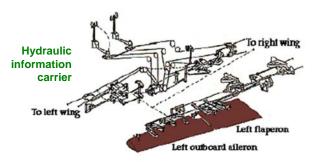
Multiplex Networks







By-wire control



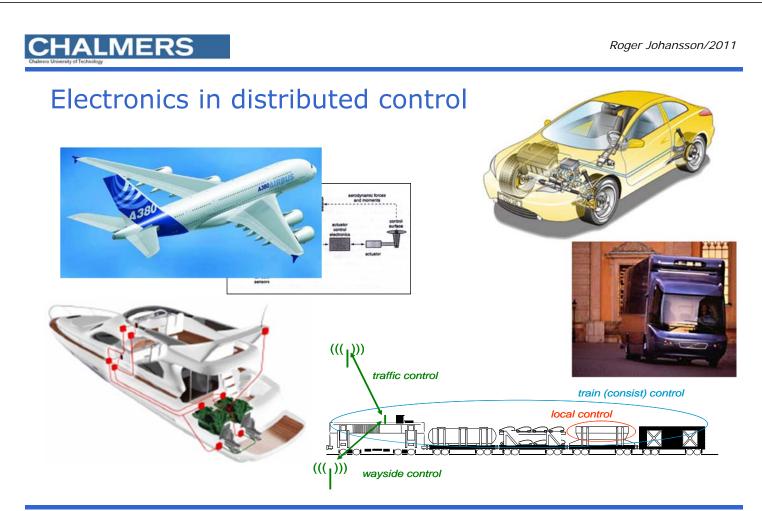
Electronic information carrier Fly-By-Wire Controls To right wing Electronics bay Left flaperon

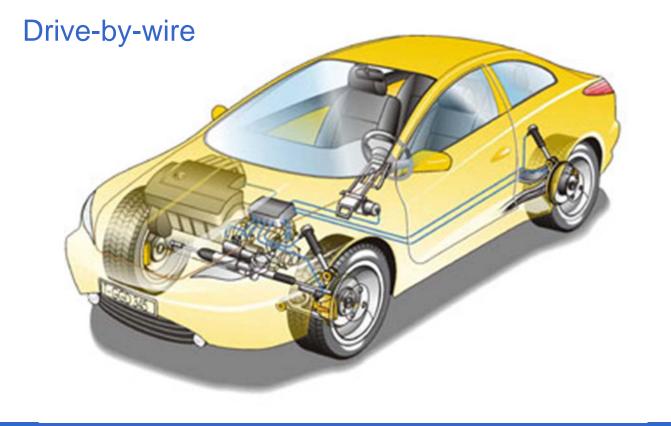
The F-8 Digital Fly-By-Wire (DFBW) fligh research project validated the principal concepts of all-electric flight control systems now used on nearly all modern high-performance aircraft and on military and civilian transports. The first flight of the 13-year project was on May 25, 1972.



Courtesy of Dryden Flight Research Center

Time triggered real time communication





Time triggered real time communication

Roger Johansson/2011

11

Control system implementation strategies

Local control

- Local information processing
- Independent control objects

Centralized global control

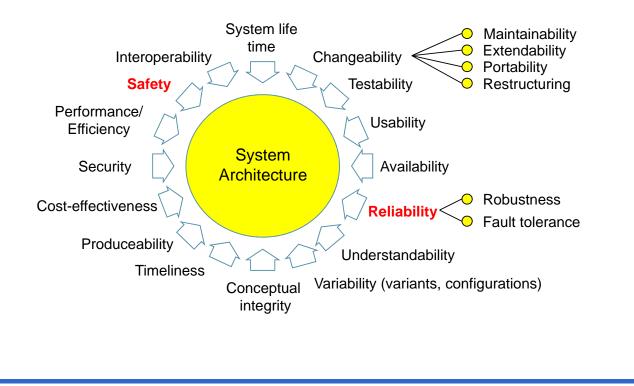
- Local and central information processing
- Interconnected control objects

Distributed global control

- Local and distributed information processing
- Interconnected control objects

CHALMERS

Non-functional requirements



Time triggered real time communication

Roger Johansson/2011

13

Tradeoffs from Safety/Reliability requirements

The extremes from reliability requirements leads to safety requirements.

Safety requirements implies redundancy, (Fail-Operational, Fail-Safe, etc).

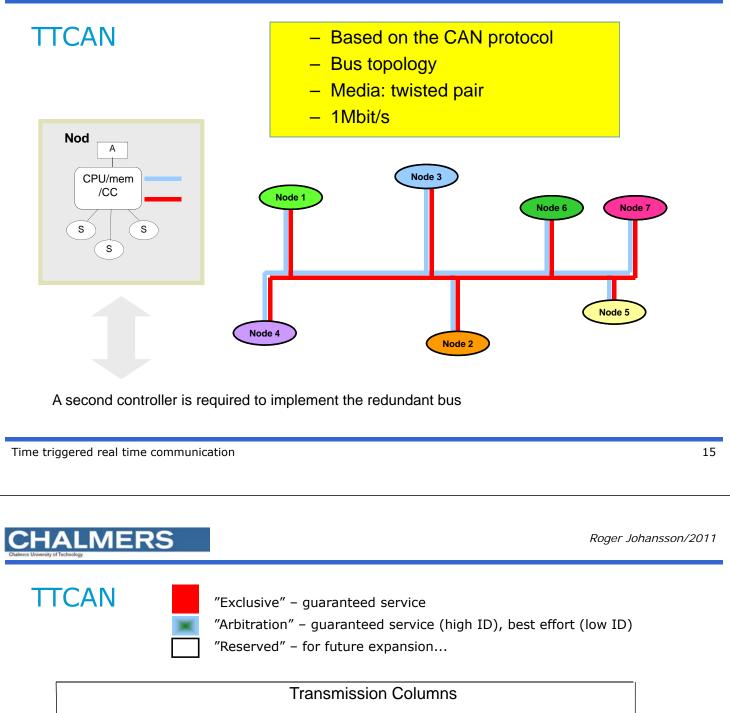
Safety requirements also demands predictability, we has to show, a priori, that the system will fulfill it's mission in every surrounding at every time.

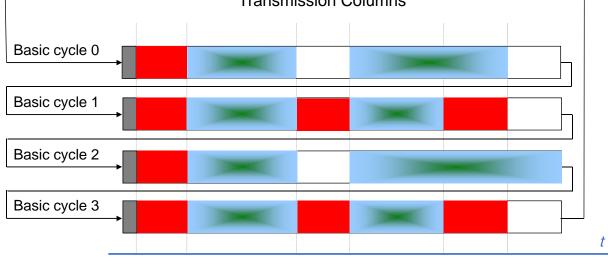
• In a distributed environment, only time triggered protocols and redundant buses can provide this safety. Contemporary TTP's are:

TTP/C, first commercial implementation. Originally from TU Vienna. Operational in civil aircrafts.

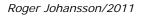
TTCAN, based on *Controller Area Network* (CAN) which is widely used in today's vehicular electronic systems.

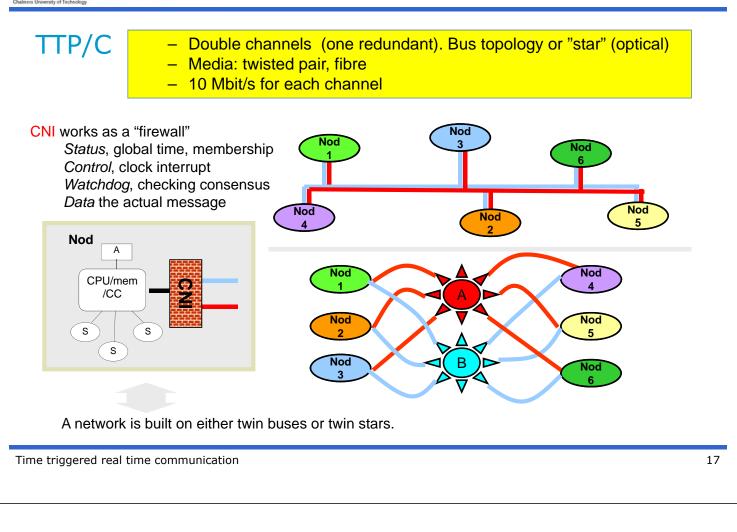
FlexRay, based on BMW's "ByteFlight". Operational in contemporary automotive electronic systems.





Time is global and measured in *network time units* (NTU's)



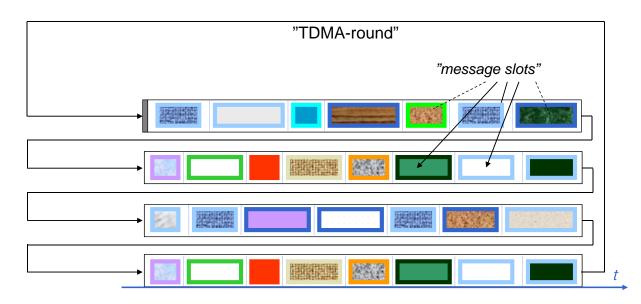




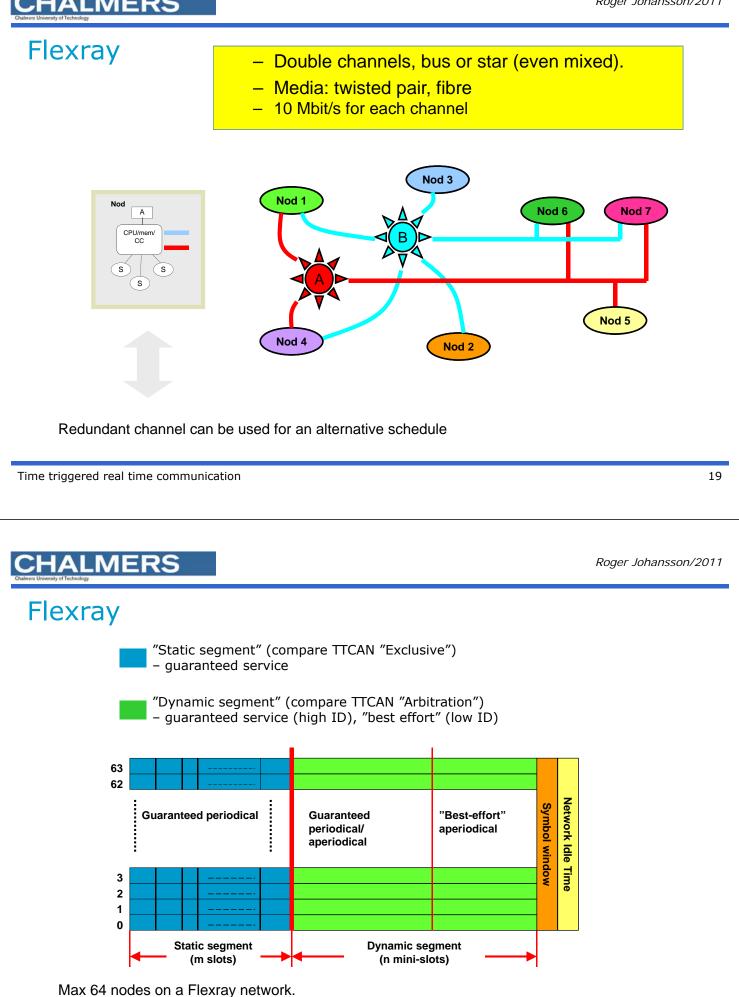
Roger Johansson/2011

TTP/C

All communication is statically scheduled Guaranteed service



Non periodical messages has to been fitted into static slots by the application



CHALMERS

Comparisons

All protocols targets real time applications.

TTCAN and Flexay combines time AND event triggered paradigms well.

All protocols are suitable for scheduling tools.

TTP/C has commercial production tools. Tools for TTCAN and Flexray are anticipated.

CAN, many years experiences, a lot of existing applications.

Implies migration of existing CAN applications into TTCAN.

TTP/C considered as complex.

Poor support for asynchronous events. High complexity, lacks second (or multiple) sources.

Flexray is the latest initiative.

Supported by most automotive suppliers.

Time triggered real time communication

CHALMERS

Combining time triggering with events: Example of Hybrid scheduling



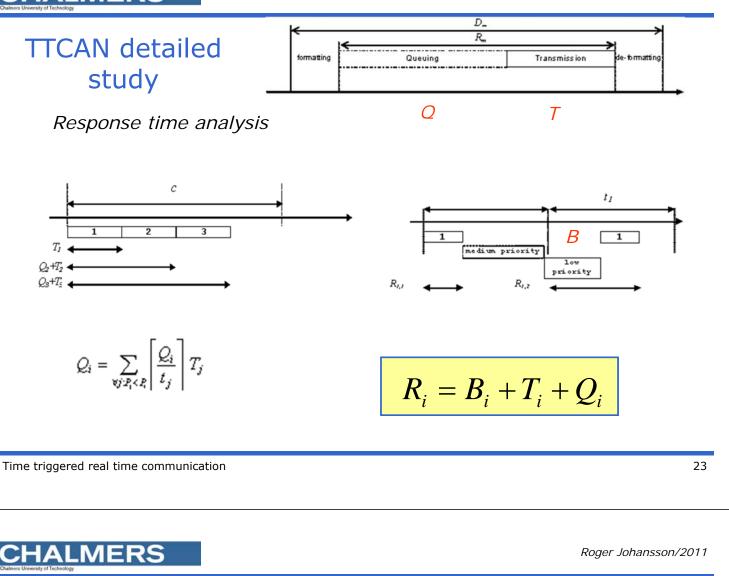


Roger Johansson/2011

Messages are sorted into three different categories:

- Hard real-time, for minimal jitter with guaranteed response time.
- Firm real-time, for guaranteed response time, but can tolerate jitter.
- Soft real-time, for "best effort" messages.

Time triggered real time communication



	Transr	nission C	olumns		
				time wi	ndows
Basic cycle 0			/		
Basic cycle 1				+	
Basic cycle 2					
Basic cycle 3					

After structuring:

 $M: \{M^h, M^f, M^s\}$, assume that at least M^h is defined. We now construct a matrix cycle. Due to protocol constraints, the schedule has to fulfil:

$$\mathbf{LCM}(M^h_p) = x \ 2^n$$

where:

- LCM is *least common multiple* period for the *M*^h message set;
- *x* is the preferred length of a basic cycle within LCM;
- *n* is the number of *basic cycles*.

Hardware constraints:

- Hwc1: $1 \le x \le 2^{y}$, has to be consistent with a hardware register, y bits
- Hwc2: $0 \le n \le k$, always a power of 2, constraint in hardware.
- Hwc3: # of triggers \leq *Tr*, columns in the matrix cycle. Limited by the number of available trigger registers.

Roger Johansson/2011

Multiple solutions satisfies the equation...

Choose a strategy:

Strategy 1:

Minimize number of *basic cycles*, requires a longer *basic cycle*, and more *triggers*.

Strategy 2:

Minimize length of *basic cycles*, increase probability of finding a feasible schedule for large message

Time triggered real time communication

Roger Johansson/2011

25

Persuing the strategies...

Construct a schedule for the following set:

 $M^{h} = (M1, M2, M3)$ with the following attributes (NTU): $M1_{p} = 1000, M1_{e} = 168$ $M2_{p} = 2000, M2_{e} = 184$ $M3_{p} = 3000, M3_{e} = 216$

It's obvious that:

LCM(M1, M2, M3) = 6000.

and:

 $6000 = x 2^n$

Strategy 1

MER

Minimizing number of basic cycles yields: $2^n = 1$, so n = 0 and x = 6000. Hwc1 and Hwc2 are fulfilled.

Total numbers of *triggers* for *N* messages in one *basic cycle* is:

$$\sum_{i=1}^{N} \frac{LCM(M)}{M^{i}}$$

in this case:

of triggers = $\frac{6000}{1000} + \frac{6000}{2000} + \frac{6000}{3000} = 11$

So, strategy 1, leads to a solution with:

- 1 basic cycle and 11 triggers.
- MAtrix cycle length is 6000 NTU.

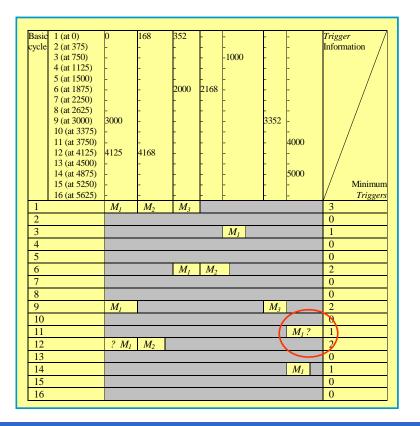
Bas	sic Cy	cle Trig	ggers								
0	168	352	1000	2000	2168	3000	3352	4000)4168	5000	
M_1	M_2	M_3	M_1	M_1	M_2	M_{I}	M_3	M_{I}	M_2	M_{I}	

Time triggered real time communication

Strategy 2

n = 0: $6000 = x 2^0$ $\Rightarrow x = 6000$ (same as strategy 1) n = 1: $6000 = x 2^1$ $\Rightarrow x = 3000$ n = 2: $6000 = x 2^2$ $\Rightarrow x = 1500$ n = 3: $6000 = x 2^3$ $\Rightarrow x = 750$ *n* = 4: $6000 = x 2^4$ $\Rightarrow x = 375$

n = 5: $6000 = x 2^5 \qquad \Rightarrow x = 187.5$



Roger Johansson/2011

Strategy 2

Avoid this conflict with the requirement that:

a basic cycle shall be at least as long as the shortest period in the message set.

Applying this restriction we get: n = 2, (x = 1500) which yields a feasible schedule:

Basic	1	0	168	352	-	-	-	1000	<i>Trigger</i> Information
cycle	2	-	-	-	2000	2168	-	-	Information
	3	3000	-	3352	-	-	4000	4168	Minimun
	4	-	-	-	5000	-	-	-	Trigger
1		M_1	M_2	M_3				M_1	4
2					M_1	M_2			2
3		M_1		M_3			M_1	M_2	4
4					M_1				1

Time triggered real time communication

Roger Johansson/2011

29

Verifying the events... (M)

	Grey slots are supposed to be allocated for <i>M</i> ^{<i>h</i>}										
Cycle	NTU-slots (Columns)										
1		q ₀									
2		q ₁ q ₂									
3.		q ₃		q ₄			q ₅				
2 ⁿ		q _{N-3}					q _{N-2}		q _{<i>N</i>-1}		

```
for each message m in M^{f}:

for message m = 1 up to last_m

for virtual message VM_{i} = 1 up to last_VM

if(Q_{m} + T_{m}) falls within (VM_{i,start}, VM_{i,completion})

Q_{m} = VM_{i,completion}

else

Q_{m} = \sum_{\forall j: P_{m} < P_{j}}^{1} \left[ \frac{Q_{m}}{t_{j}} \right] T_{j}

endif

end

end

end
```

CHALMERS

Conclusions

• Applicable real time communication protocols for future safety-critical applications has to provide strictly periodical (minimal jitter), periodical (jitter is negliable) and a-periodic communication to fully support control applications.

•Scheduling periodical and a-periodical events requires a new approach, *hybrid scheduling*.

• Hybrid scheduling is sparsely found in today's literature...



Thank you for your attention.