

Monads reference card

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	monad component	DSL application	
Syntax	$m :: * \rightarrow *$ $return :: a \rightarrow m a$ $(\gg=) :: m [a] \rightarrow ([a] \rightarrow m b) \rightarrow m b$	Expressions parameterized on return type constant expression bind an a returned by the lhs into the rhs	
Laws	name	law	a DSL aspect
“do”	$\text{do } x \leftarrow \alpha$ $y \leftarrow \beta$ γ	$return a \gg= (\lambda x. m x) \equiv m a$ $m \gg= (\lambda x. return x) \equiv m$ $(m \gg= f) \gg= g \equiv m \gg= (\lambda x. f x \gg= g)$	inlining/factorizing a constant removal/introduction of useless return extension/shrinking of scope
Comprehensions	$[\quad \gamma$ $ \quad x \leftarrow \alpha$ $, \quad y \leftarrow \beta$ $] \quad \text{return } \gamma$	$\alpha \gg= \lambda x.$ $\beta \gg= \lambda y.$	
	<ul style="list-style-type: none"> • parentheses are not needed • x may appear in γ 		
	<ul style="list-style-type: none"> • $\gg=$ can be used to “flatten” levels of the monad. • $join :: m(m a) \rightarrow m a$ • $join xs = xs \gg= id$ 		