Software Engineering using Formal Methods Java Modeling Language

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Road-map

first half of the course:

Modelling of distributed and concurrent systems

second half of course:

Deductive Verification of JAVA source code

- 1. *specifying* JAVA programs
- 2. proving JAVA programs correct

What kind of Specifications

system level specifications (requirements analysis, GUI, use cases) important, but not subject of this course

instead:

unit specification – contracts among implementers on various levels:

- application level application level
- application level library level
- library level library level

In the object-oriented setting:

units to be specified are interfaces, classes, and their methods

first focus on methods

methods specified by *potentially* referring to:

- result value,
- initial values of formal parameters,
- pre-state and post-stateaccessible part of pre/post-state

to stress the different roles – obligations – responsibilities in a specification:

widely used analogy of the specification as a contract

"Design by Contract" methodology

contract between *caller* and *callee* of method

callee guarantees certain outcome provided caller guarantees prerequisites

Running Example: ATM.java

```
public class ATM {
```

```
// fields:
private BankCard insertedCard = null;
private int wrongPINCounter = 0;
private boolean customerAuthenticated = false;
```

```
// methods:
public void insertCard (BankCard card) { ... }
public void enterPIN (int pin) { ... }
public int accountBalance () { ... }
public int withdraw (int amount) { ... }
public void ejectCard () { ... }
```

very informal Specification of 'enterPIN (int pin)':

Enter the PIN that belongs to the currently inserted bank card into the ATM. If a wrong PIN is entered three times in a row, the card is confiscated. After having entered the correct PIN, the customer is regarded as authenticated.

Getting More Precise: Specification as Contract

Contract states what is guaranteed under which conditions.

- precondition card is inserted, user not yet authenticated, pin is correct postcondition user is authenticated
- precondition card is inserted, user not yet authenticated, wrongPINCounter < 2 and pin is incorrect wrongPINCounter is increased by 1 user is not authenticated

precondition card is inserted, user not yet authenticated, wrongPINCounter >= 2 and pin is incorrect card is confiscated user is not authenticated

Meaning of Pre/Post-condition pairs

Definition

A **pre/post-condition** pair for a method m is **satisfied by the implementation** of m if:

When m is called in any state that satisfies the precondition then in any terminating state of m the postcondition is true.

- 1. No guarantees are given when the precondition is not satisfied.
- 2. Termination may or may not be guaranteed.
- **3.** Terminating state may be reached by normal or by abrupt termination (cf. exceptions).

non-termination and abrupt termination \Rightarrow next lecture

What kind of Specifications

Natural language specs are very important.

but this course's focus:

"formal" specifications:

Describing contracts of units in a mathematically precise language.

Motivation:

- higher degree of precision
- eventually: automation of program analysis of various kinds:
 - static checking
 - program verification

Java Modeling Language (JML)

JML is a specification language tailored to JAVA.

General JML Philosophy

Integrate

- specification
- implementation

in one single language.

 \Rightarrow JML is not external to JAVA

JML is JAVA + FO Logic + pre/post-conditions, invariants + more...

JML Annotations

JML extends JAVA by annotations.

JML annotations include:

- preconditions
- postconditions
- class invariants
- ✓ additional modifiers
- ✗ 'specification-only' fields
- ✗ 'specification-only' methods
- loop invariants

✓ ... × ...

: in this course, X: not in this course

JML/JAVA integration

JML annotations are attached to JAVA programs by writing them directly into the JAVA source code files

not to confuse JAVA compiler:

JML annotations live in in special comments, ignored by JAVA, recognized by JML tools.

from the file ATM. java

```
/*@ public normal_behavior
  @ requires !customerAuthenticated;
  @ requires pin == insertedCard.correctPIN;
  @ ensures customerAuthenticated;
  @*/
public void enterPIN (int pin) {
    if ( ...
}
```

Everything between /* and */ is invisible for JAVA.

```
/*@ public normal_behavior
  @ requires !customerAuthenticated;
  @ requires pin == insertedCard.correctPIN;
  @ ensures customerAuthenticated;
  @*/
public void enterPIN (int pin) {
    if ( ...
But:
```

A JAVA comment with '@' as its first character it is *not* a comment for JML tools.

JML annotations appear in JAVA comments starting with @.

How about "//" comments?

- /*@ public normal_behavior
 - @ requires !customerAuthenticated;
 - @ requires pin == insertedCard.correctPIN;

```
@ ensures customerAuthenticated; @*/
```

equivalent to:

```
//@ public normal_behavior
//@ requires !customerAuthenticated;
//@ requires pin == insertedCard.correctPIN;
//@ ensures customerAuthenticated;
```

The easiest way to comment out JML? I.e. comment out the comment:

```
/*u@ public normal_behavior ... @*/
```

 $//_{\sqcup}$ [©] public normal_behavior

```
//_0 requires !customerAuthenticated;
```

. . .

```
/*@ public normal_behavior
  @ requires !customerAuthenticated;
  @ requires pin == insertedCard.correctPIN;
  @ ensures customerAuthenticated;
  @*/
public void enterPIN (int pin) {
    if ( ...
```

What about the intermediate '@'s?

Within a JML annotation, a '@' is ignored:

- ▶ if it is the first (non-white) character in the line
- if it is the last character before '*/'.

 \Rightarrow The blue '@'s are not *required*, but it's a *convention* to use them.

```
/*@ public normal_behavior
  @ requires !customerAuthenticated;
  @ requires pin == insertedCard.correctPIN;
  @ ensures customerAuthenticated;
  @*/
public void enterPIN (int pin) {
    if ( ...
```

This is a **public** specification case:

- 1. it is accessible from all classes and interfaces
- 2. it can only mention public fields/methods of this class
- 2. Can be a problem. Solution later in the lecture.

```
/*@ public normal_behavior
  @ requires !customerAuthenticated;
  @ requires pin == insertedCard.correctPIN;
  @ ensures customerAuthenticated;
  @*/
public void enterPIN (int pin) {
    if ( ...
```

Each keyword ending with **behavior** opens a 'specification case'.

normal_behavior Specification Case

The method guarantees to *not* throw any exception (on the top level), *if the caller guarantees all preconditions of this specification case.*

```
/*@ public normal_behavior
@ requires !customerAuthenticated;
@ requires pin == insertedCard.correctPIN;
@ ensures customerAuthenticated;
@*/
public void enterPIN (int pin) {
    if ( ...
```

This specification case has two preconditions (marked by requires)

- 1. !customerAuthenticated
- 2. pin == insertedCard.correctPIN

here:

preconditions are boolean JAVA expressions

in general:

preconditions are boolean JML expressions (see below)

SEFM: Java Modeling Language

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```
/*@ public normal_behavior
    @ requires !customerAuthenticated;
    @ requires pin == insertedCard.correctPIN;
    @ ensures customerAuthenticated;
    @*/
```

specifies only the case where both preconditions are true in pre-state

the above is equivalent to:

```
/*@ public normal_behavior
  @ requires !customerAuthenticated;
  @ requires pin == insertedCard.correctPIN;
  @ ensures customerAuthenticated;
  @*/
public void enterPIN (int pin) {
    if ( ....
```

This specification case has one postcondition (marked by ensures)

customerAuthenticated

here:

postcondition is boolean JAVA expressions

in general:

postconditions are boolean JML expressions (see below)

different specification cases are connected by 'also'.

```
/*@ public normal_behavior
```

- @ requires !customerAuthenticated;
- @ requires pin == insertedCard.correctPIN;

```
@ ensures customerAuthenticated;
```

```
0
```

```
@ also
```

```
0
```

```
@ public normal_behavior
```

```
@ requires !customerAuthenticated;
```

@ requires pin != insertedCard.correctPIN;

```
@ requires wrongPINCounter < 2;</pre>
```

```
@ ensures wrongPINCounter == \old(wrongPINCounter) + 1;
@*/
```

```
public void enterPIN (int pin) {
```

```
if ( ...
```

```
/*@ <spec-case1> also
@
@ public normal_behavior
@ requires !customerAuthenticated;
@ requires pin != insertedCard.correctPIN;
@ requires wrongPINCounter < 2;
@ ensures wrongPINCounter == \old(wrongPINCounter) + 1;
@*/
public void enterPIN (int pin) { ...</pre>
```

for the first time, JML expression not a JAVA expression

\old(*E***)** means: *E* evaluated in the pre-state of enterPIN.

E can be any (arbitrarily complex) (JML) expression.

```
/*@ <spec-case1> also <spec-case2> also
  0
  @ public normal_behavior
  @ requires insertedCard != null;
  @ requires !customerAuthenticated;
  @ requires pin != insertedCard.correctPIN;
  @ requires wrongPINCounter >= 2;
  @ ensures insertedCard == null;
  @ ensures \old(insertedCard).invalid:
 @*/
```

public void enterPIN (int pin) { ...

two postconditions state that:

'Given the above preconditions, enterPIN guarantees:

insertedCard == null and \old(insertedCard).invalid'

Question:

could it be

```
@ ensures \old(insertedCard.invalid);
instead of
  @ ensures \old(insertedCard).invalid;
```

??

Specification Cases Complete?

```
consider spec-case-1:
```

- @ public normal_behavior
- @ requires !customerAuthenticated;
- @ requires pin == insertedCard.correctPIN;
- @ ensures customerAuthenticated;

what does spec-case-1 not tell about post-state?

recall: fields of class ATM:

insertedCard customerAuthenticated wrongPINCounter

what happens with insertCard and wrongPINCounter?

completing spec-case-1:

- @ public normal_behavior
- @ requires !customerAuthenticated;
- @ requires pin == insertedCard.correctPIN;
- @ ensures customerAuthenticated;
- @ ensures insertedCard == \old(insertedCard);
- @ ensures wrongPINCounter == \old(wrongPINCounter);

completing spec-case-2:

- @ public normal_behavior
- @ requires !customerAuthenticated;
- @ requires pin != insertedCard.correctPIN;
- @ requires wrongPINCounter < 2;</pre>
- @ ensures wrongPINCounter == \old(wrongPINCounter) + 1;
- @ ensures insertedCard == \old(insertedCard);
- **@ ensures** customerAuthenticated
- @ == \old(customerAuthenticated);

Completing Specification Cases

```
completing spec-case-3:
```

- @ public normal_behavior
- @ requires insertedCard != null;
- @ requires !customerAuthenticated;
- @ requires pin != insertedCard.correctPIN;
- @ requires wrongPINCounter >= 2;
- @ ensures insertedCard == null;
- @ ensures \old(insertedCard).invalid;
- **@ ensures** customerAuthenticated
- @ == \old(customerAuthenticated);
- @ ensures wrongPINCounter == \old(wrongPINCounter);

Assignable Clause

unsatisfactory to add

```
@ ensures loc == \old(loc);
```

for all locations loc which do not change

instead:

add assignable clause for all locations which may change

Q assignable $loc_1, \ldots, loc_n;$

Meaning: No location other than loc_1, \ldots, loc_n can be assigned to. Special cases:

No location may be changed:

```
@ assignable \nothing;
```

Unrestricted, method allowed to change anything:

```
@ assignable \everything;
```

completing spec-case-1:

- @ public normal_behavior
- @ requires !customerAuthenticated;
- @ requires pin == insertedCard.correctPIN;
- @ ensures customerAuthenticated;
- @ assignable customerAuthenticated;

completing spec-case-2:

- @ public normal_behavior
- @ requires !customerAuthenticated;
- @ requires pin != insertedCard.correctPIN;
- @ requires wrongPINCounter < 2;</pre>
- @ ensures wrongPINCounter == \old(wrongPINCounter) + 1;
- @ assignable wrongPINCounter;

Specification Cases with Assignable

```
completing spec-case-3:
```

- @ public normal_behavior
- @ requires insertedCard != null;
- @ requires !customerAuthenticated;
- @ requires pin != insertedCard.correctPIN;
- @ requires wrongPINCounter >= 2;
- @ ensures insertedCard == null;
- @ ensures \old(insertedCard).invalid;
- @ assignable insertedCard,
- @ insertedCard.invalid,

You can specify groups of locations as assignable, using '*'.

example:

```
@ assignable o.*, a[*];
```

makes all fields of object o and all locations of array a assignable.

JML extends the JAVA modifiers by additional modifiers.

The most important ones are:

- spec_public
- ► pure

Aim: admitting more class elements to be used in JML expressions.

JML Modifiers: spec_public

in enterPIN example, pre/post-conditions made heavy use of class fields

But: public specifications can only talk about public fields.

Not desired: make all fields public.

one solution:

- keep the fields private/protected
- make those needed for specification spec_public

different solution: use specification-only fields (not covered in this course)

JML Modifiers: pure

It can be handy to use method calls in JML annotations. Examples:

- o1.equals(o2)
- li.contains(elem)
- li1.max() < li2.min()</pre>

allowed if, and only if method is guaranteed to have no side effects. In JML, you can specify methods to be '**pure**':

public /*@ pure @*/ int max() { ...

'**pure**' puts obligation on implementer, not to cause side effects, but allows to use method in annotations

'pure' similar to 'assignable \nothing;', but global to method

JML Expressions \neq JAVA Expressions

boolean JML Expressions (to be completed)

- each side-effect free boolean JAVA expression is a boolean JML expression
- if a and b are boolean JML expressions, and x is a variable of type t, then the following are also boolean JML expressions:

How to express the following?

- an array arr only holds values ≤ 2
- the variable m holds the maximum entry of array arr
- all Account objects in the array accountProxies are stored at the index corresponding to their respective accountNumber field
- all created instances of class BankCard have different cardNumbers

to be answered in the next lecture

Literature for this Lecture

essential reading:

in KeY Book A. Roth and Peter H. Schmitt: Formal Specification. Chapter 5 only sections 5.1, 5.3, In: B. Beckert, R. Hähnle, and P. Schmitt, editors. Verification of Object-Oriented Software: The KeY Approach, vol 4334 of LNCS. Springer, 2006. (e-version via Chalmers Library)

further reading, all available at

http://www.eecs.ucf.edu/~leavens/JML/documentation.shtml:

JML Reference Manual Gary T. Leavens, Erik Poll, Curtis Clifton, Yoonsik Cheon, Clyde Ruby, David Cok, Peter Müller, and Joseph Kiniry.

JML Reference Manual

JML Tutorial Gary T. Leavens, Yoonsik Cheon.

Design by Contract with JML

JML Overview Gary T. Leavens, Albert L. Baker, and Clyde Ruby. JML: A Notation for Detailed Design