

# Software Engineering using Formal Methods

## Java Modeling Language

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# Road-map

first half of the course:

Modelling of distributed and concurrent systems

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first half of the course:

Modelling of distributed and concurrent systems

second half of course:

Deductive Verification of JAVA source code

1. *specifying* JAVA programs
2. proving JAVA programs correct

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*system level specifications*  
(requirements analysis, GUI, use cases)  
important, but  
*not subject of this course*

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*system level specifications*  
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important, but  
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instead:

*unit specification* – *contracts among implementers* on various levels:

- ▶ application level – application level
- ▶ application level – library level
- ▶ library level – library level

# Unit Specifications

**In the object-oriented setting:**

**units** to be specified are **interfaces**, **classes**, and their **methods**

first focus on methods

methods specified by *potentially* referring to:

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- ▶ pre-state and post-state

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methods specified by *potentially* referring to:

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- ▶ initial values of formal parameters,
- ▶ accessible part of pre/post-state

# Specifications as Contracts

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contract between *caller* and *callee* of method

*callee guarantees certain outcome* *provided* *caller guarantees prerequisites*

## Running Example: ATM.java

```
public class ATM {  
  
    // fields:  
    private BankCard insertedCard = null;  
    private int wrongPINCounter = 0;  
    private boolean customerAuthenticated = false;  
  
    // methods:  
    public void insertCard (BankCard card) { ... }  
    public void enterPIN (int pin) { ... }  
    public int accountBalance () { ... }  
    public int withdraw (int amount) { ... }  
    public void ejectCard () { ... }  
  
}
```

very informal Specification of 'enterPIN (**int** pin)':

*Enter the PIN that belongs to the currently inserted bank card into the ATM. If a wrong PIN is entered three times in a row, the card is confiscated. After having entered the correct PIN, the customer is regarded as authenticated.*

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user is not authenticated

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wrongPINCounter >= 2 and pin is incorrect

*postcondition*    card is confiscated  
user is not authenticated

# Meaning of Pre/Post-condition pairs

## Definition

A **pre/post-condition** pair for a method  $m$  is **satisfied by the implementation** of  $m$  if:

*When  $m$  is called in any state that satisfies the **precondition** then in any terminating state of  $m$  the **postcondition** is true.*

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non-termination and abrupt termination  $\Rightarrow$  next lecture

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## **“formal” specifications:**

Describing contracts of units in a mathematically precise language.

Motivation:

- ▶ higher degree of precision
- ▶ eventually: *automation* of program analysis of various kinds:
  - ▶ static checking
  - ▶ **program verification**

# Java Modeling Language (JML)

JML is a **specification language** tailored to **JAVA**.

## General JML Philosophy

Integrate

- ▶ specification
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in **one single language**.

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JAVA + **FO Logic** + **pre/post-conditions, invariants** + more. . .

# JML Annotations

JML **extends** JAVA by **annotations**.

## JML annotations include:

- ✓ preconditions
- ✓ postconditions
- ✓ class invariants
- ✓ additional modifiers
- ✗ 'specification-only' fields
- ✗ 'specification-only' methods
- ✓ loop invariants
- ✓ ...
- ✗ ...

✓: in this course, ✗: not in this course

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not to confuse JAVA compiler:

JML annotations live in in special comments,  
ignored by JAVA, recognized by JML tools.

# JML by Example

from the file ATM.java

```
⋮  
/*@ public normal_behavior  
    @ requires !customerAuthenticated;  
    @ requires pin == insertedCard.correctPIN;  
    @ ensures customerAuthenticated;  
    @*/  
public void enterPIN (int pin) {  
    if ( ...  
  
⋮
```

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```

Everything between `/*` and `*/` is invisible for JAVA.

# JML by Example

```
/*@ public normal_behavior
   @ requires !customerAuthenticated;
   @ requires pin == insertedCard.correctPIN;
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But:

A JAVA comment with '@' as its first character  
it is *not* a comment for JML tools.

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JML annotations appear in JAVA comments starting with @.

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How about “//” comments?

# JML by Example

```
/*@ public normal_behavior
  @ requires !customerAuthenticated;
  @ requires pin == insertedCard.correctPIN;
  @ ensures customerAuthenticated; @*/
```

*equivalent to:*

```
//@ public normal_behavior
//@ requires !customerAuthenticated;
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The easiest way to [comment out JML](#)? I.e. comment out the comment:

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/*@ public normal_behavior
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*equivalent to:*

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/*_@ public normal_behavior ... */
//_@ public normal_behavior
//_@ requires !customerAuthenticated;
...
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/*@ public normal_behavior
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What about the intermediate '@'s?



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What about the intermediate '@'s?

Within a JML annotation, a '@' is ignored:

- ▶ if it is the first (non-white) character in the line
- ▶ if it is the last character before '\*/'.

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- ▶ if it is the last character before '\*/'.

⇒ The blue '@'s are not *required*, but it's a *convention* to use them.

```
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   @ requires !customerAuthenticated;
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This is a **public** specification case:

1. it is accessible from all classes and interfaces
2. it can only mention public fields/methods of this class

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This is a **public** specification case:

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2. Can be a problem. Solution later in the lecture.

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Each keyword ending with **behavior** opens a 'specification case'.

## **normal\_behavior** Specification Case

The method guarantees to *not* throw any exception

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Each keyword ending with **behavior** opens a 'specification case'.

## **normal\_behavior** Specification Case

The method guarantees to *not* throw any exception (on the top level),

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Each keyword ending with **behavior** opens a 'specification case'.

## **normal\_behavior** Specification Case

The method guarantees to *not* throw any exception (on the top level),  
*if the caller guarantees all preconditions of this specification case.*

# JML by Example

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This specification case has two **preconditions** (marked by **requires**)

1. !customerAuthenticated
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here:

preconditions are *boolean JAVA expressions*

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here:

preconditions are *boolean JAVA expressions*

in general:

preconditions are *boolean JML expressions* (see below)

```
/*@ public normal_behavior
   @ requires !customerAuthenticated;
   @ requires pin == insertedCard.correctPIN;
   @ ensures customerAuthenticated;
   @*/
```

specifies only the case where **both** preconditions are true in pre-state  
the above is equivalent to:

```
/*@ public normal_behavior
   @ requires (      !customerAuthenticated
   @           && pin == insertedCard.correctPIN );
   @ ensures customerAuthenticated;
   @*/
```

# JML by Example

```
/*@ public normal_behavior
   @ requires !customerAuthenticated;
   @ requires pin == insertedCard.correctPIN;
   @ ensures customerAuthenticated;
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public void enterPIN (int pin) {
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This specification case has one **postcondition** (marked by **ensures**)

- ▶ `customerAuthenticated`

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here:

postcondition is *boolean JAVA expressions*

in general:

postconditions are *boolean JML expressions* (see below)

# JML by Example

different specification cases are connected by **'also'**.

```
/*@ public normal_behavior
   @ requires !customerAuthenticated;
   @ requires pin == insertedCard.correctPIN;
   @ ensures customerAuthenticated;
   @
   @ also
   @
   @ public normal_behavior
   @ requires !customerAuthenticated;
   @ requires pin != insertedCard.correctPIN;
   @ requires wrongPINCounter < 2;
   @ ensures wrongPINCounter == \old(wrongPINCounter) + 1;
   @*/

public void enterPIN (int pin) {
    if ( ...
```

# JML by Example

```
/*@ <spec-case1> also
   @
   @ public normal_behavior
   @ requires !customerAuthenticated;
   @ requires pin != insertedCard.correctPIN;
   @ requires wrongPINCounter < 2;
   @ ensures wrongPINCounter == \old(wrongPINCounter) + 1;
   @*/
public void enterPIN (int pin) { ...
```

for the first time, JML expression not a JAVA expression

**\old(*E*)** means: *E* evaluated in the pre-state of enterPIN.

*E* can be any (arbitrarily complex) (JML) expression.



# JML by Example

```
/*@ <spec-case1> also <spec-case2> also
   @
   @ public normal_behavior
   @ requires insertedCard != null;
   @ requires !customerAuthenticated;
   @ requires pin != insertedCard.correctPIN;
   @ requires wrongPINCounter >= 2;
   @ ensures insertedCard == null;
   @ ensures \old(insertedCard).invalid;
   @*/
public void enterPIN (int pin) { ...
```

two postconditions state that:

‘Given the above preconditions, enterPIN guarantees:

`insertedCard == null`    and    `\old(insertedCard).invalid`’

## Question:

could it be

```
@ ensures \old(insertedCard.invalid);
```

instead of

```
@ ensures \old(insertedCard).invalid;
```

??

# Specification Cases Complete?

consider spec-case-1:

```
@ public normal_behavior
@ requires !customerAuthenticated;
@ requires pin == insertedCard.correctPIN;
@ ensures customerAuthenticated;
```

what does spec-case-1 *not* tell about post-state?

# Specification Cases Complete?

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recall: fields of class ATM:

```
insertedCard
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recall: fields of class ATM:

```
insertedCard
customerAuthenticated
wrongPINCounter
```

what happens with insertCard and wrongPINCounter?

# Completing Specification Cases

completing spec-case-1:

```
@ public normal_behavior
@ requires !customerAuthenticated;
@ requires pin == insertedCard.correctPIN;
@ ensures customerAuthenticated;
@ ensures insertedCard == \old(insertedCard);
@ ensures wrongPINCounter == \old(wrongPINCounter);
```

# Completing Specification Cases

completing spec-case-2:

```
@ public normal_behavior
@ requires !customerAuthenticated;
@ requires pin != insertedCard.correctPIN;
@ requires wrongPINCounter < 2;
@ ensures wrongPINCounter == \old(wrongPINCounter) + 1;
@ ensures insertedCard == \old(insertedCard);
@ ensures customerAuthenticated
@      == \old(customerAuthenticated);
```

# Completing Specification Cases

completing spec-case-3:

```
@ public normal_behavior
@ requires insertedCard != null;
@ requires !customerAuthenticated;
@ requires pin != insertedCard.correctPIN;
@ requires wrongPINCounter >= 2;
@ ensures insertedCard == null;
@ ensures \old(insertedCard).invalid;
@ ensures customerAuthenticated
@      == \old(customerAuthenticated);
@ ensures wrongPINCounter == \old(wrongPINCounter);
```



# Assignable Clause

unsatisfactory to add

```
@ ensures loc == \old(loc);
```

for all locations *loc* which *do not* change

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instead:

add **assignable clause** for all locations which *may* change

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@ assignable loc1, ..., locn;
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@ assignable loc1, ..., locn;
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Meaning: No location other than *loc<sub>1</sub>, ..., loc<sub>n</sub>* can be assigned to.

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for all locations *loc* which *do not* change

instead:

add **assignable clause** for all locations which *may* change

```
@ assignable loc1, ..., locn;
```

Meaning: **No location other than  $loc_1, \dots, loc_n$  can be assigned to.**

Special cases:

**No** location may be changed:

```
@ assignable \nothing;
```

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Special cases:

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```

Unrestricted, method allowed to change anything:

```
@ assignable \everything;
```

# Specification Cases with Assignable

completing spec-case-1:

```
@ public normal_behavior
@ requires !customerAuthenticated;
@ requires pin == insertedCard.correctPIN;
@ ensures customerAuthenticated;
@ assignable customerAuthenticated;
```

# Specification Cases with Assignable

completing spec-case-2:

```
@ public normal_behavior
@ requires !customerAuthenticated;
@ requires pin != insertedCard.correctPIN;
@ requires wrongPINCounter < 2;
@ ensures wrongPINCounter == \old(wrongPINCounter) + 1;
@ assignable wrongPINCounter;
```

# Specification Cases with Assignable

completing spec-case-3:

```
@ public normal_behavior
@ requires insertedCard != null;
@ requires !customerAuthenticated;
@ requires pin != insertedCard.correctPIN;
@ requires wrongPINCounter >= 2;
@ ensures insertedCard == null;
@ ensures \old(insertedCard).invalid;
@ assignable insertedCard,
@           insertedCard.invalid,
```



# Assignable Groups

You can specify groups of locations as assignable, using '\*'.

example:

```
@ assignable o.*, a[*];
```

makes all fields of object o and all locations of array a assignable.

JML extends the JAVA modifiers by additional modifiers.

The most important ones are:

- ▶ `spec_public`
- ▶ `pure`

Aim: admitting more class elements to be used in JML expressions.

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different solution: use specification-only fields (not covered in this course)

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It can be handy to use method calls in JML annotations.

Examples:

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- ▶ `li.contains(elem)`
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**'pure'** similar to **'assignable \nothing;'**, but global to method

# JML Expressions $\neq$ JAVA Expressions

## **boolean** JML Expressions (to be completed)

- ▶ each **side-effect free** **boolean** JAVA expression is a **boolean** JML expression
- ▶ if *a* and *b* are **boolean** JML expressions, and *x* is a variable of type *t*, then the following are also **boolean** JML expressions:
  - ▶ *!a* (“not *a*”)
  - ▶ *a && b* (“*a* and *b*”)
  - ▶ *a || b* (“*a* or *b*”)

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  - ▶ a && b ("a and b")
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  - ▶ a ==> b ("a implies b")
  - ▶ a <==> b ("a is equivalent to b")
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to be answered in the next lecture



# Literature for this Lecture

*essential reading:*

**in KeY Book** A. Roth and Peter H. Schmitt: Formal Specification.  
Chapter 5 **only sections 5.1, 5.3**, In: B. Beckert, R. Hähnle, and  
P. Schmitt, editors. *Verification of Object-Oriented Software: The  
KeY Approach*, vol 4334 of *LNCS*. Springer, 2006.  
(e-version via Chalmers Library)

*further reading*, all available at

<http://www.eecs.ucf.edu/~leavens/JML/documentation.shtml>:

**JML Reference Manual** Gary T. Leavens, Erik Poll, Curtis Clifton,  
Yoonsik Cheon, Clyde Ruby, David Cok, Peter Müller, and  
Joseph Kiniry.

*JML Reference Manual*

**JML Tutorial** Gary T. Leavens, Yoonsik Cheon.  
*Design by Contract with JML*

**JML Overview** Gary T. Leavens, Albert L. Baker, and Clyde Ruby.  
*JML: A Notation for Detailed Design*