

Lava II

Mary Sheeran, Thomas Hallgren
Chalmers University of Technology

Generating VHDL

- In the simplest case
`writeVhdl "fullAdder" fullAdder`
- Assigning names to the inputs
`writeVhdlInput "fullAdder" fullAdder (var "carryIn", (var "a", var "b"))`
- Assigning names also to the outputs
`writeVhdlInputOutput "fullAdder" fullAdder
 (var "carryIn", (var "a", var "b")) (var "sum", var "carryOut")`
- Generic circuits are not supported, so you need to pick a size
`writeVhdlInputOutput "rippleCarryAdder" rippleCarryAdder
 (var "carryIn", (varList 8 "a", varList 8 "b"))
 (varList 8 "sum", var "carryOut")`

Generating VHDL (better)

Above method generates silly VHDL for combinational circuits

importing file VhdlNew11.hs (on Schedule page)

allows gen. of clocked or unclocked VHDL netlists

Append Clk or NoClk to end of previous function names

see Lava2.hs

```

library ieee;

use ieee.std_logic_1164.all;

entity
    rippleCarryAdder
is
port
    (

        carryIn : in std_logic
    ; a_0 : in std_logic
    ; a_1 : in std_logic
    ; a_2 : in std_logic
    ; a_3 : in std_logic
    ; b_0 : in std_logic
    ; b_1 : in std_logic
    ; b_2 : in std_logic
    ; b_3 : in std_logic

    ; sum_0 : out std_logic
    ; sum_1 : out std_logic
    ; sum_2 : out std_logic
    ; sum_3 : out std_logic
    ; carryOut : out std_logic
    );
end rippleCarryAdder;

```

```

architecture
    structural
of
    rippleCarryAdder
is
    signal w1 : std_logic;
    signal w2 : std_logic;
    signal w3 : std_logic;
    signal w4 : std_logic;
    signal w5 : std_logic;

    ...      . . .

    signal w29 : std_logic;
begin
    c_w2      : entity work.wire port map (carryIn, w2);
    c_w4      : entity work.wire port map (a_0, w4);
    c_w5      : entity work.wire port map (b_0, w5);
    c_w29     : entity work.andG port map (w25, w26, w29);
    c_w27     : entity work.xorG port map (w28, w29, w27);

    . . .
    c_sum_0   : entity work.wire port map (w1, sum_0);
    c_sum_1   : entity work.wire port map (w6, sum_1);
    c_sum_2   : entity work.wire port map (w13, sum_2);
    c_sum_3   : entity work.wire port map (w20, sum_3);
    c_carryOut : entity work.wire port map (w27, carryOut);
end structural;

```

Generic circuits again

Lava.Arithmetic.hs contains

```
binAdder :: ([Signal Bool], [Signal Bool]) -> [Signal Bool]
```

Generic circuits again

Lava.Arithmetic.hs contains

```
binAdder :: ([Signal Bool], [Signal Bool]) -> [Signal Bool]
```

```
> simulate binAdder ([low,high,low], [high,low,high])  
[high,high,high,low]
```

Generic circuits again

Lava.Arithmetic.hs contains

```
binAdder :: ([Signal Bool], [Signal Bool]) -> [Signal Bool]
```

Let's check if it is commutative!

First attempt

```
prop_AdderCommutative (as,bs) = ok
  where
    out1 = binAdder (as,bs)
    out2 = binAdder (bs,as)
    ok    = out1 <==> out2
```


First attempt

```
prop_AdderCommutative (as,bs) = ok
  where
    out1 = binAdder (as,bs)
    out2 = binAdder (bs,as)
    ok    = out1 <==> out2
```

smv prop_AdderCommutative **does not work!**

Need to fix size

```
prop_AdderCommutative_ForSize n =  
  forAll (list n) $ \ as ->  
    forAll (list n) $ \ bs ->  
      prop_AdderCommutative (as,bs)
```

Need to fix size

```
prop_AdderCommutative_ForSize n =  
  forAll (list n) $ \ as ->  
    forAll (list n) $ \ bs ->  
      prop_AdderCommutative (as,bs)
```

smv (prop_AdderCommutative_ForSize 16) **works!**

See Chapter 4 in the Lava tutorial.

Same effect but easier

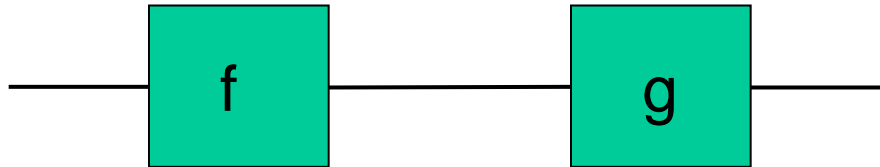
```
prop_AdderComm1 n  
  = prop_AdderCommutative (varList n "a", varList n "b")
```

```
fv_binAdd_Comm1 = smv (prop_AdderComm1 16)
```

works

Serial composition

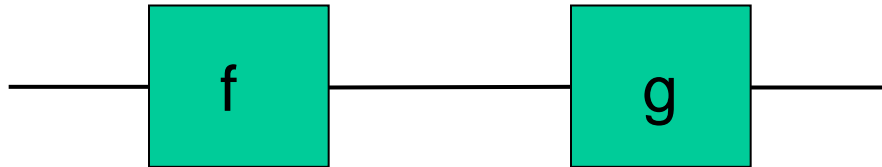
useful connection pattern



$f \rightarrow g$

Serial composition type

useful connection pattern

$$(->-) :: (a \rightarrow b) \rightarrow (b \rightarrow c) \rightarrow a \rightarrow c$$


Serial composition example

```
doubSum :: [Signal Int] -> Signal Int  
doubSum = map (*2) ->- sum
```

```
> simulate doubSum [1..8]  
72
```

Serial composition example

```
doubSum :: [Signal Int] -> Signal Int
doubSum = map (*2) ->- sum
```

could also have written

```
doubSum1 :: [Signal Int] -> Signal Int
doubSum1 as = sum (map double as)
  where
    double a = a * 2
```


Feedback and sequential circuits

First example

```
bad inp = out  
where  
  out = nand2(inp,out)
```

Feedback and sequential circuits

First example

```
bad inp = out  
where  
  out = nand2(inp,out)
```

> simulate bad low
high

> simulate bad high

*** Exception: combinational loop

Delay in VHDL

Signal assignments have no delay by default:

```
out <= a nand b;
```

Delay can be introduced explicitly:

```
out <= a nand b after 4ns;
```

Delay in Lava

The logical gates in the Lava library are "ideal" and have zero delay

Delay has to be modelled explicitly:

`delay init s`

delays the signal `s` by one time unit

The output during the first time unit is `init`

Delay in Lava

The Lava library does not care how long a time unit is.

It could be the gate delay, for analyzing the effect of delay in combinational circuits

But usually it is one clock cycle in a synchronously clocked sequential circuit.

Feedback and sequential circuits

Second example

nand2D = nand2 ->- delay low

good a = out

where

out = nand2D(a,out)

Feedback and sequential circuits

nand2D = nand2 ->- delay low

good a = out

where

out = nand2D(a,out)

*Main> simulate good high

*** Exception: evaluating a delay component



Need to use sequential simulation

Feedback and sequential circuits

nand2D = nand2 ->- delay low

good a = out

where

out = nand2D(a,out)

*Main> simulateSeq good [high,high,low,high]
[low,high,low,high]

Retiming

`nand2D = nand2 ->- delay low`

`delNand2 = delay (high,high) ->- nand2`

`sim0 = simulateSeq nand2D [(low,low),(high,low),(high,high),(low,low)]`

`sim1 = simulateSeq delNand2 [(low,low),(high,low),(high,high),(low,low)]`

`> sim0`

`[low,high,high,low]`

`> sim1`

`[low,high,high,low]`

Retiming

`nand2D = nand2 ->- delay low`

`delNand2 = delay (high,high) ->- nand2`

`sim0 = simulateSeq nand2`
`sim1 = simulateSeq delNand2`

Note that `delay` works on many types,
not just bits

`> sim0`

`[low,high,high,low]`

`> sim1`

`[low,high,high,low]`

FV

-- A general function for equivalence testing

propEQ circ1 circ2 inp = ok

where

out1 = circ1 inp

out2 = circ2 inp

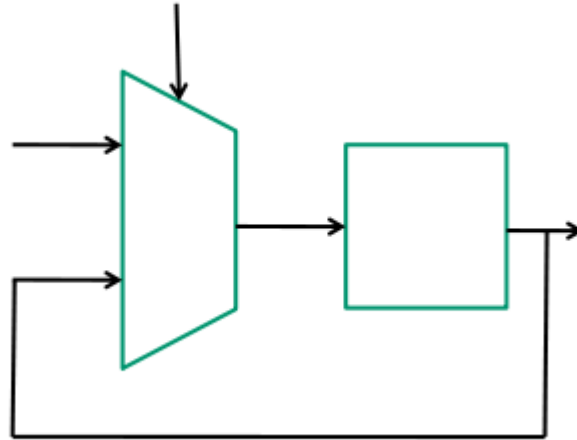
ok = out1 <==> out2

prop0 = propEQ nand2D delNand2

fv_prop0 = smv prop0

(on my laptop ca .1 sec, 60 BDD nodes allocated)

Register



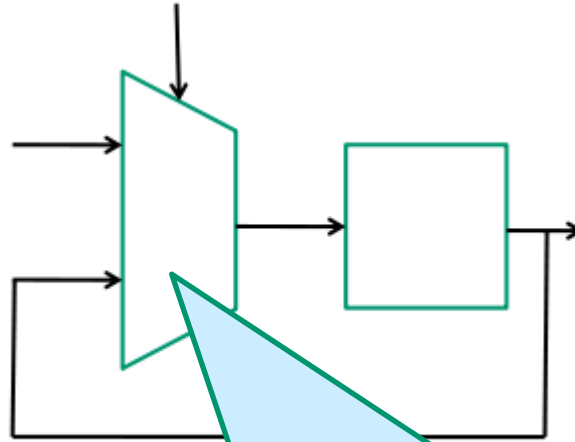
$\text{reg init } (w, \text{din}) = \text{dout}$

where

$\text{dout} = \text{delay init } m$

$m = \text{mux } (w, (\text{dout}, \text{din}))$

Register



multiplexer (also polymorphic)

```
reg init (w,din) =
```

```
  where
```

```
    dout = delay init m
```

```
    m     = mux (w,(dout,din))
```

```
  mux :: ... => (Signal Bool,(a,a)) -> a
```

using Haskell to generate inputs

```
-- infinite lists
```

```
lh :: [Bit]
```

```
lh = low : high : lh
```

```
ins :: Int -> [[Signal Int]]
```

```
ins n = map (replicate n) [1..]
```

```
regtst n = simulateSeq (reg (zeroList n)) (take 10 (zip lh (ins n)))
```

```
*Main> regtst 5
```

```
[[0,0,0,0,0],[0,0,0,0,0],[2,2,2,2,2],[2,2,2,2,2],[4,4,4,4,4],  
 [4,4,4,4,4],[6,6,6,6,6],[6,6,6,6,6],[8,8,8,8,8],[8,8,8,8,8]]
```

Questions?

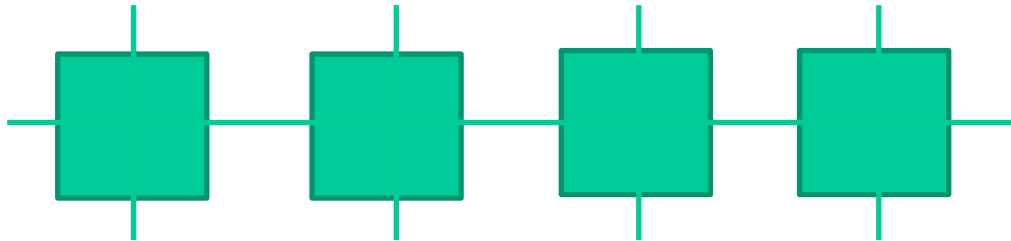
Connection patterns

Higher order functions that capture common ways
of plugging circuits together

Connection patterns

Higher order functions that capture common ways
of plugging circuits together

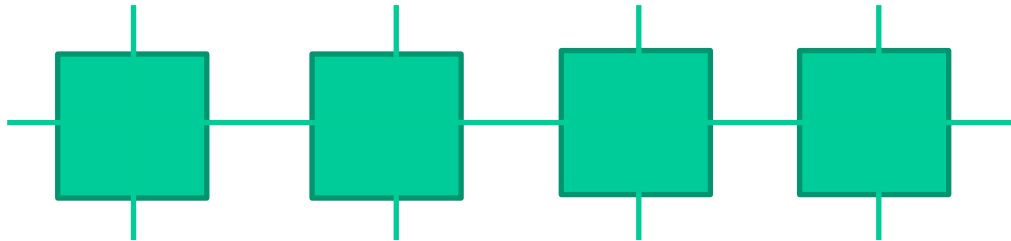
We saw `row`



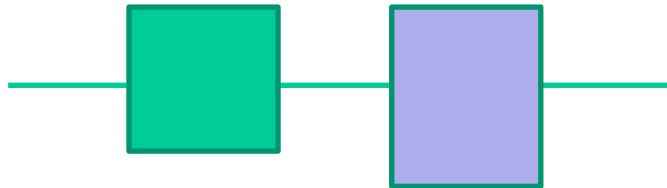
Connection patterns

Higher order functions that capture common ways of plugging circuits together

We saw row



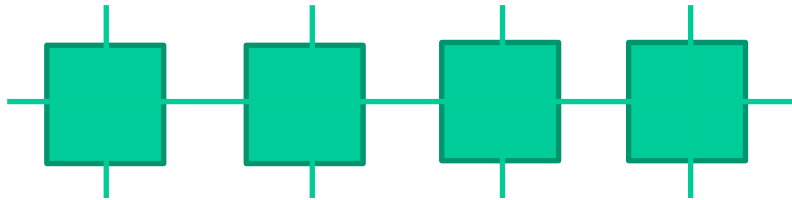
->-



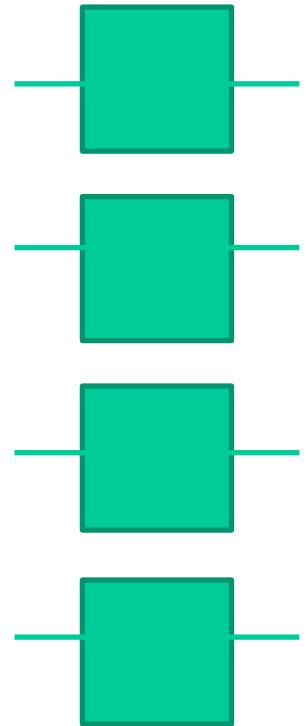
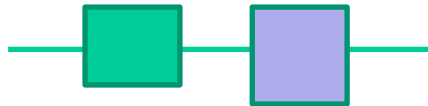
Connection patterns

Higher order functions that capture common ways of plugging circuits together

We saw row

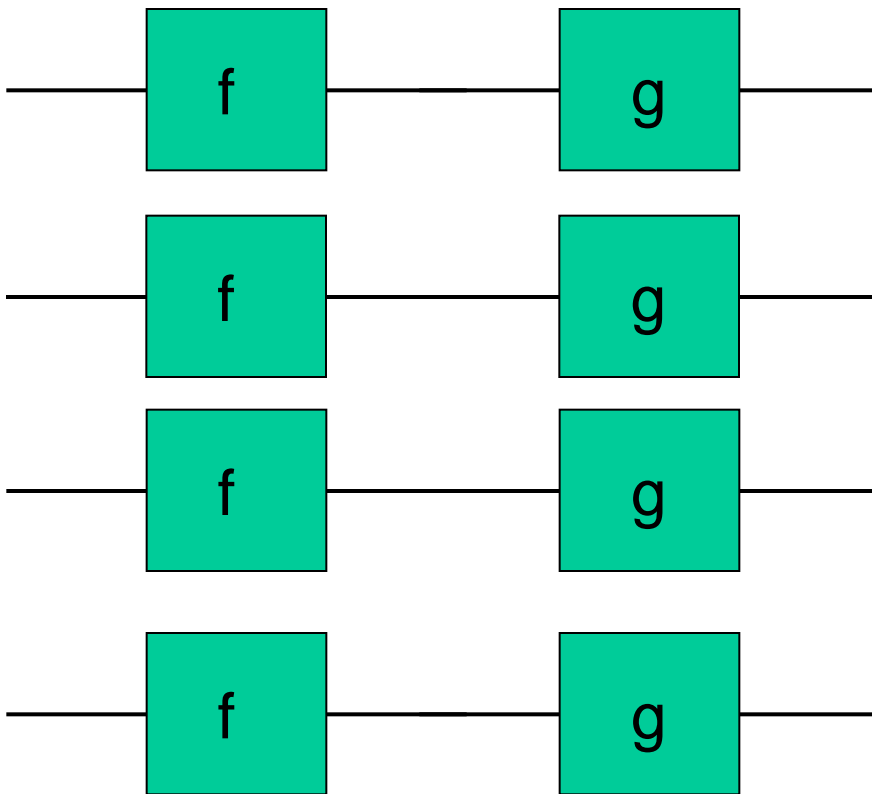


->-

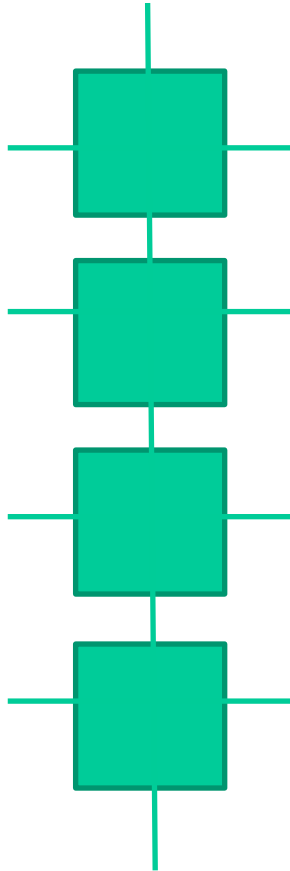


map

$\text{map } f \multimap \text{map } g = ??$



More connection patterns: column and grid



mirror circ $(a, b) = (c, d)$

where

$(d, c) = \text{circ } (b, a)$

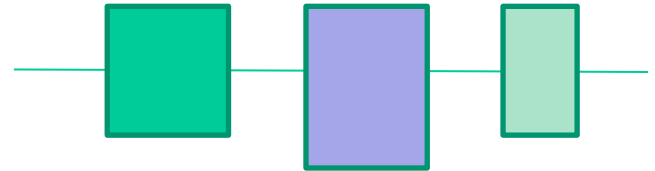
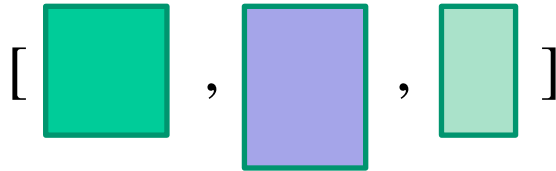
column circ = mirror (row (mirror circ))

grid circ = row (column circ)

(in Lava.Patterns)

could just define column recursively (exercise)

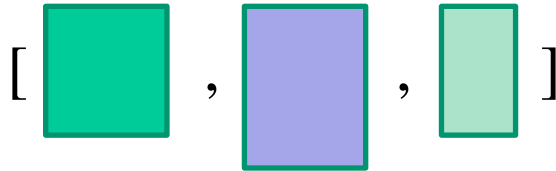
More connection patterns: compose



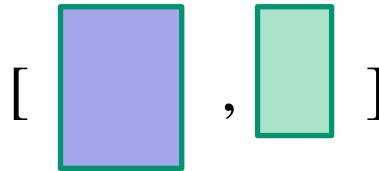
`compose :: [a -> a] -> a -> a`

(is in `Lava.Patterns`)

More connection patterns



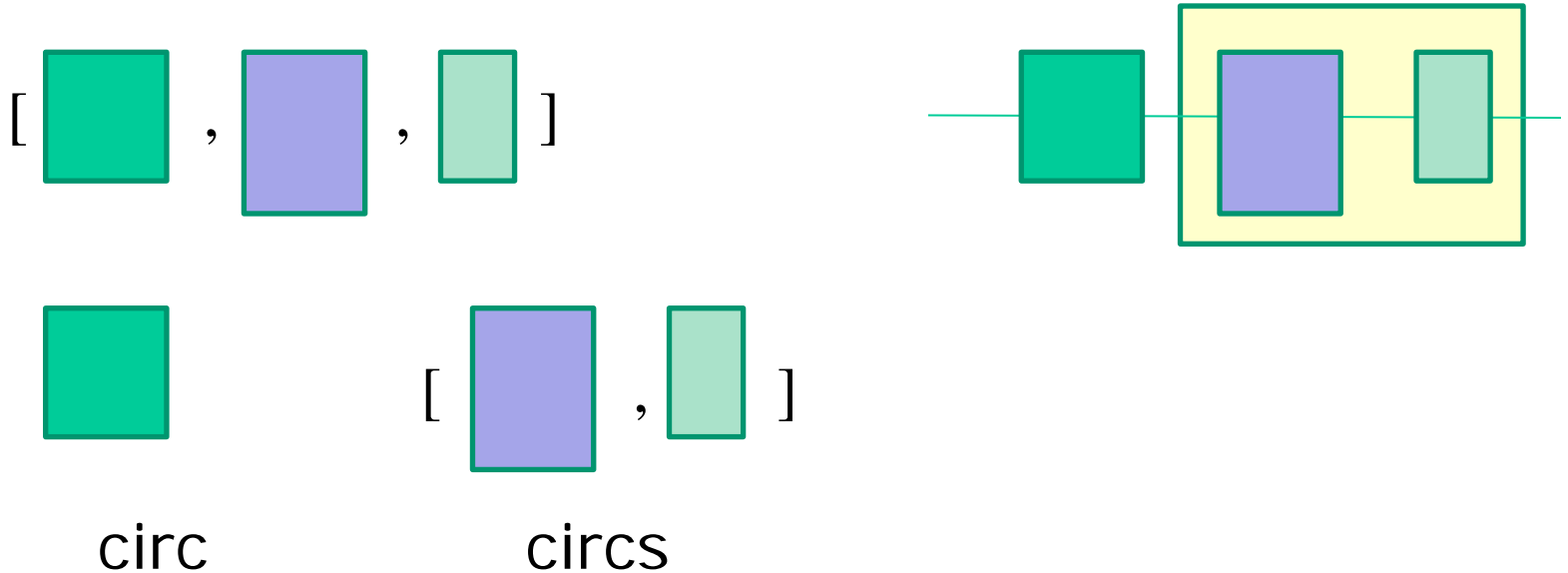
circ



circs

```
compose :: [a -> a] -> a -> a  
compose [] =  
compose (circ : circs) =
```

More connection patterns



```
compose :: [a -> a] -> a -> a
compose []           = id
compose (circ : circs) = circ ->- compose circs
```


compose n copies of function

```
composeN :: Int -> (a -> a) -> a -> a  
composeN n circ = compose (replicate n circ)
```

(in Lava.Patterns)

compose n copies of function

```
composeN :: Int -> (a -> a) -> a -> a  
composeN n circ = compose (replicate n circ)
```

```
doubN :: Int -> Signal Int -> Signal Int  
doubN n = composeN n (*2)
```

```
*Main> simulate (doubN 4) 1  
16
```

compose n copies of function

```
composeN :: Int -> (a -> a) -> a -> a  
composeN n circ = compose (replicate n circ)
```

Note that this is a **Haskell Int**
not a circuit level Int (Signal Int)

compose n copies of function

could also have defined this function recursively:

```
composeNR 0 circ = id  
composeNR n circ = circ -> composeNR (n-1) circ
```

compose n copies of function

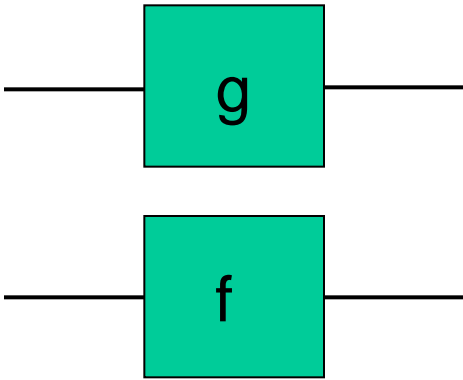
could also have defined this function recursively:

```
composeNR 0 circ = id  
composeNR n circ = circ -> composeNR (n-1) circ
```



This is a second very standard way to write recursive functions in Haskell

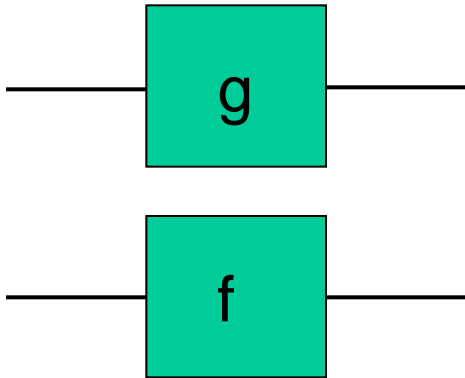
Par



$f \text{ -| - } g$

(in Lava.Patterns)

Par

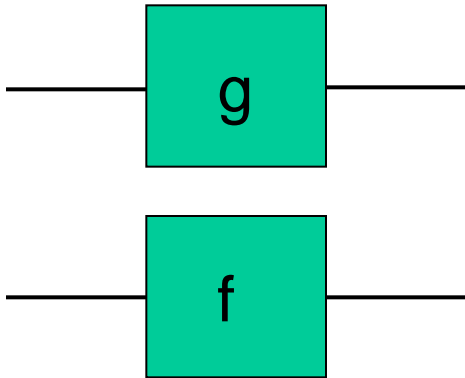


infixr 4 $-|-$

$\text{par circ1 circ2 } (a, b) = (\text{circ1 } a, \text{circ2 } b)$
 $\text{circ1 } -|- \text{ circ2} \quad \quad \quad = \text{par circ1 circ2}$

$f -|- g$

Par



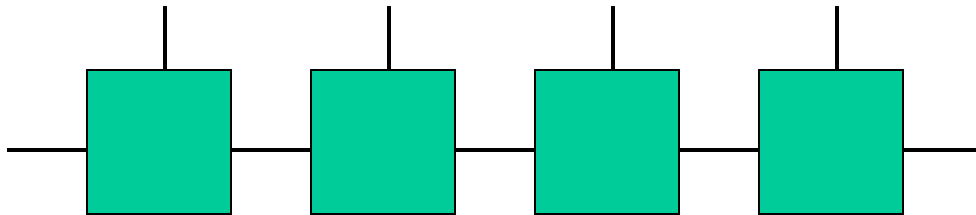
infixr 4 $-|-$

$\text{par circ1 circ2 } (a, b) = (\text{circ1 } a, \text{circ2 } b)$
 $\text{circ1 } -|- \text{ circ2} \quad \quad \quad = \text{par circ1 circ2}$

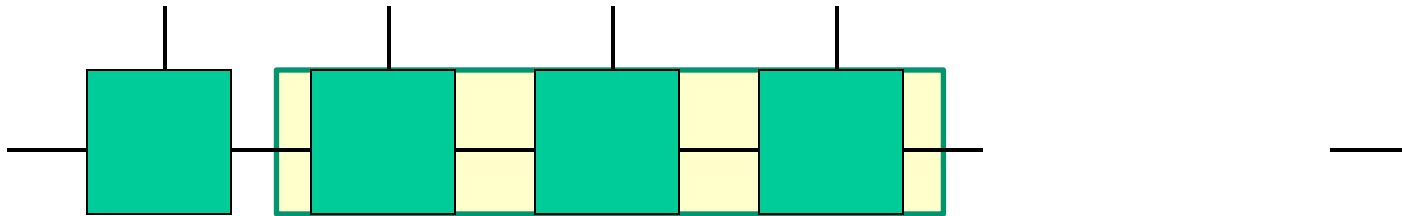
$f -|- g$

Q: What is the type of par ?

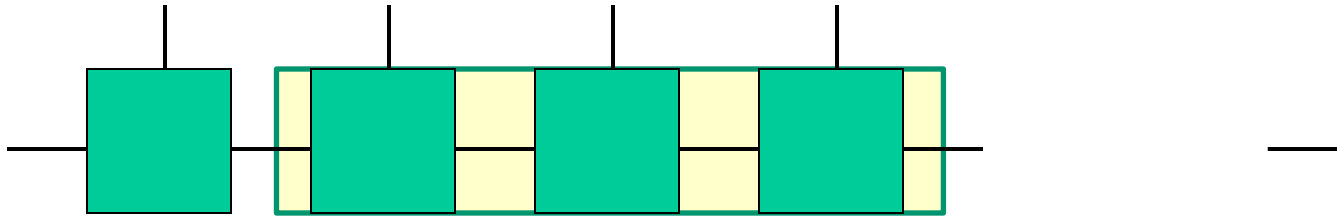
red f



red f

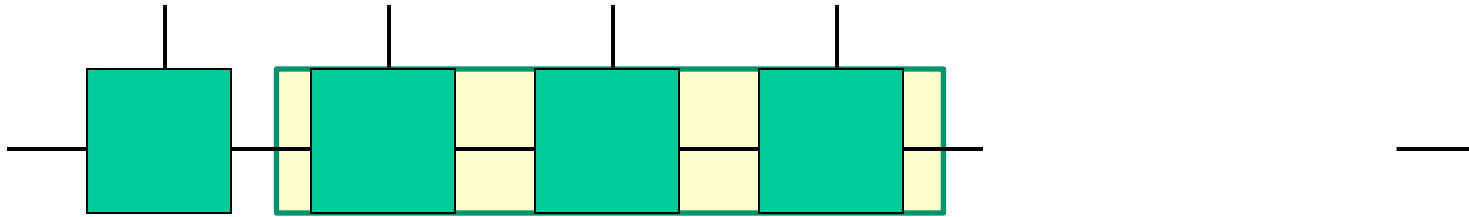


red f



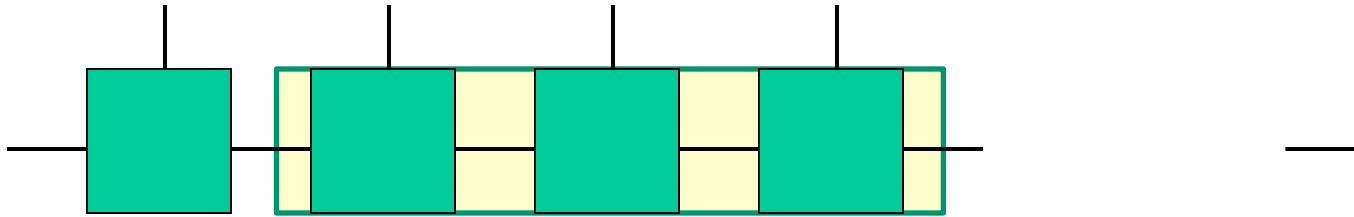
$\text{red} :: ((a,b) \rightarrow a) \rightarrow (a, [b]) \rightarrow a$

red f



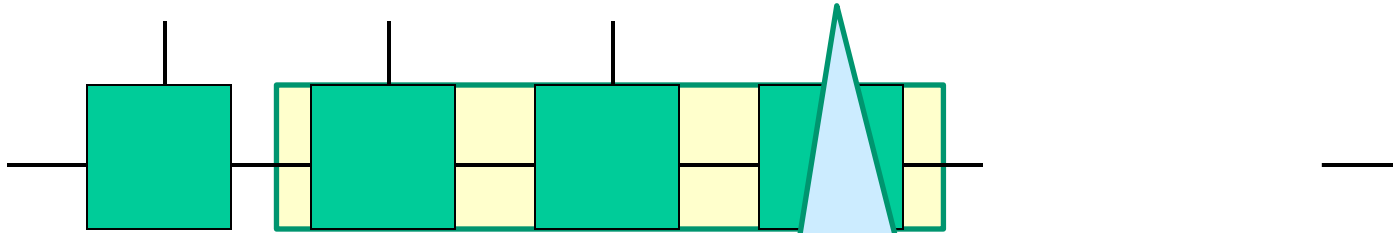
```
red :: ((a,b) -> a) -> (a, [b]) -> a
red f (a,[])      =
red f (a, (b:bs)) =
```

red f



```
red :: ((a,b) -> a) -> (a, [b]) -> a
red f (a,[])      = a
red f (a, (b:bs)) = red f (f(a,b), bs)
```

red f



```
red :: ((a,b) -> a) -> (a -> a)
red f (a,[])      = a
red f (a, (b:bs)) = red f (f a b, bs)
```

redtst = simulate

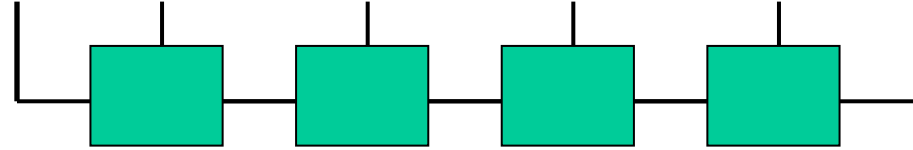
```
>redtst
18
```

This is very like `foldl` in Haskell
Here, I like to gather circuit inputs into one structure (pair). Makes it easier to see where the circuits are and to compose again with more connection patterns

lin (more like foldl1)

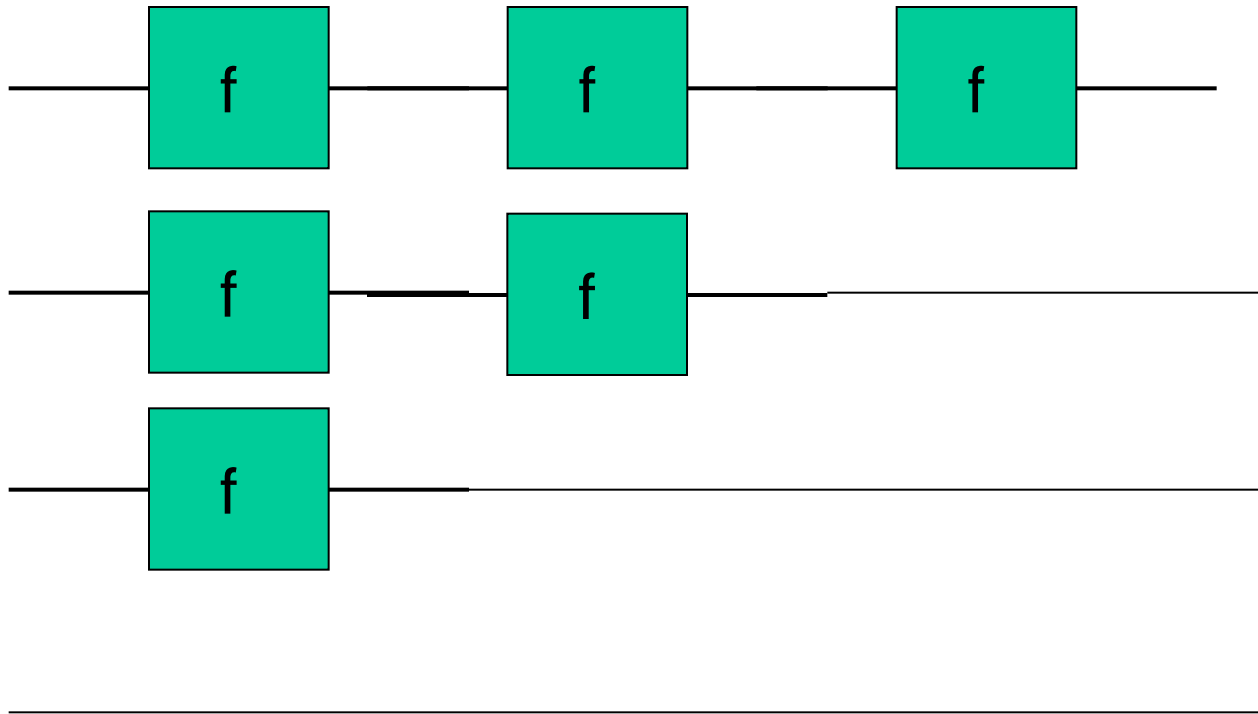
```
lin f (a:as) = red f (a,as)
lin _ []     = error "lin: empty list"
```

```
> simulate (lin plus) [1..5]
15
```



```
> simulate (lin plus) []
*** Exception: lin: empty list
```

tri f



Triangles (tri in Lava.Patterns)

```
tri circ [] =  
tri circ (inp:inps) =
```

Triangles (tri in Lava.Patterns)

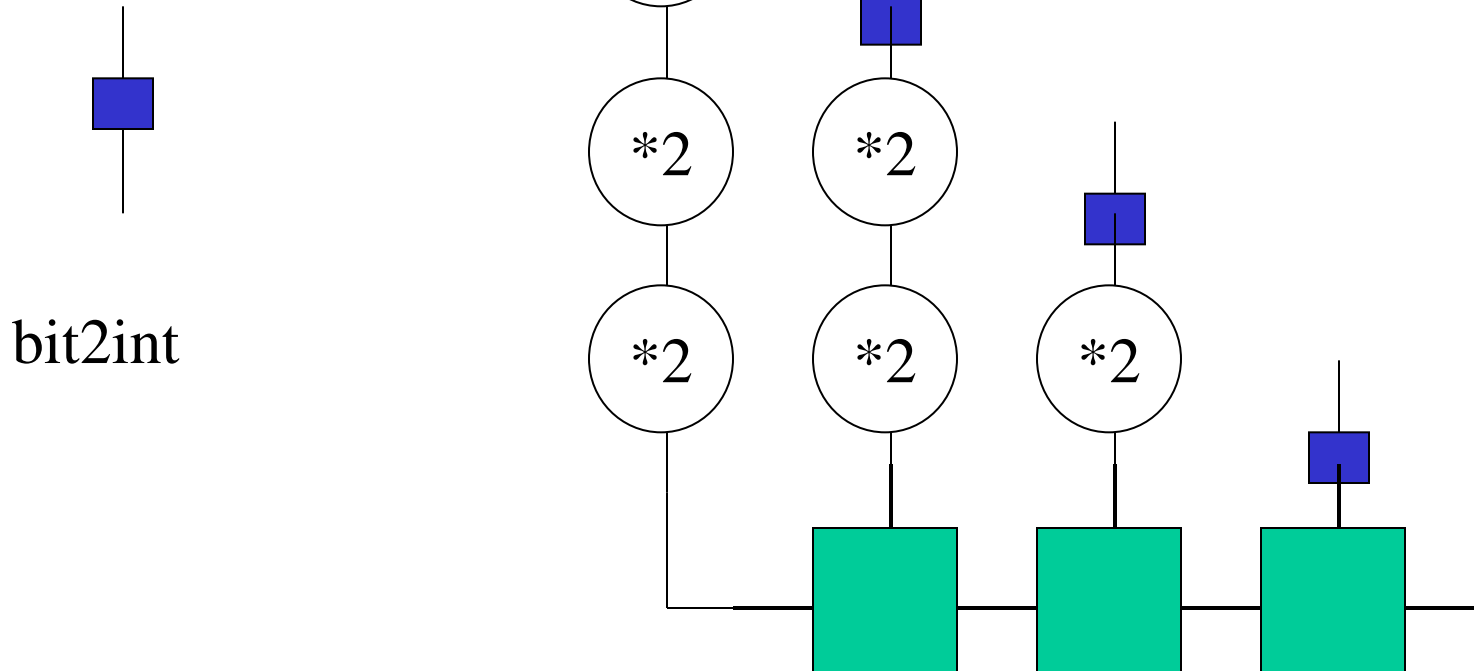
```
tri circ [] = []  
tri circ (inp:inps) = inp : (map circ ->- tri circ) inps
```

Triangles (tri in Lava.Patterns)

```
tri circ [] = []  
tri circ (inp:inps) = inp : (map circ ->- tri circ) inps
```

```
downtri f = reverse ->- tri f ->- reverse
```

Converting msb first binary to integer

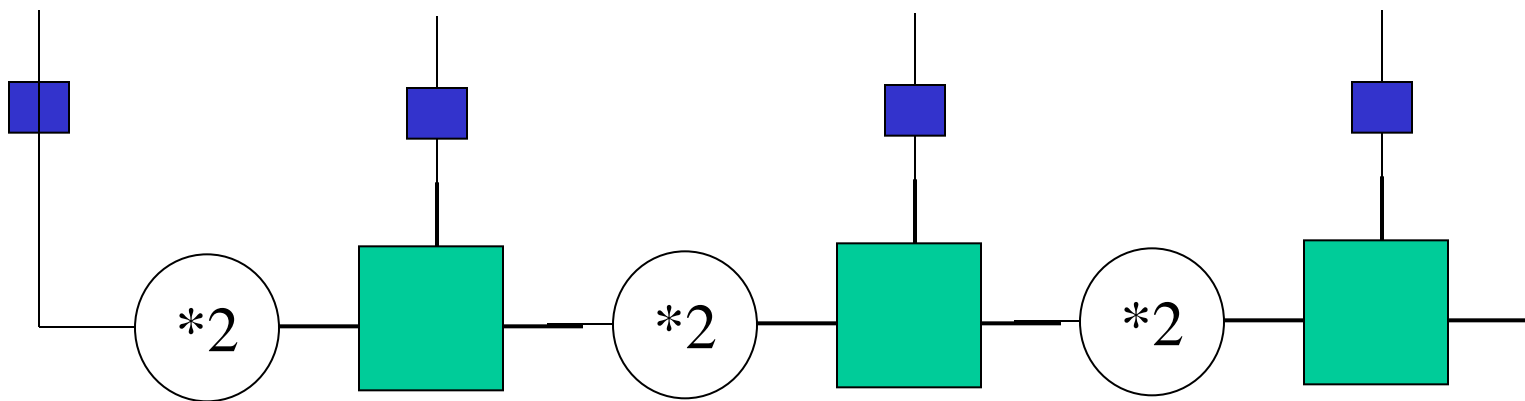


binary to integer

`msbbin2int = map bit2int ->- downtri (*2) ->- lin plus`

(`bin2int` is in Arithmetic module, lsb first)

Another way



binary to integer

```
msbbin2int' = map bit2int ->- lin cell  
  where  
    cell (a,b) = 2*a + b
```

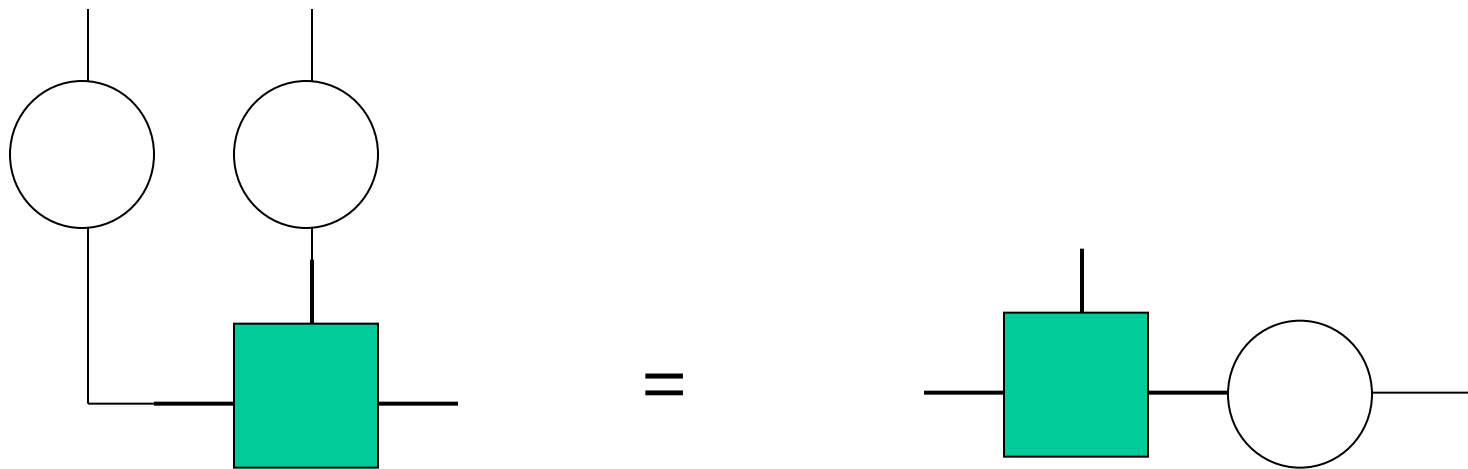
binary to integer

```
msbbin2int' = map bit2int ->- lin cell  
where  
  cell (a,b) = 2*a + b
```

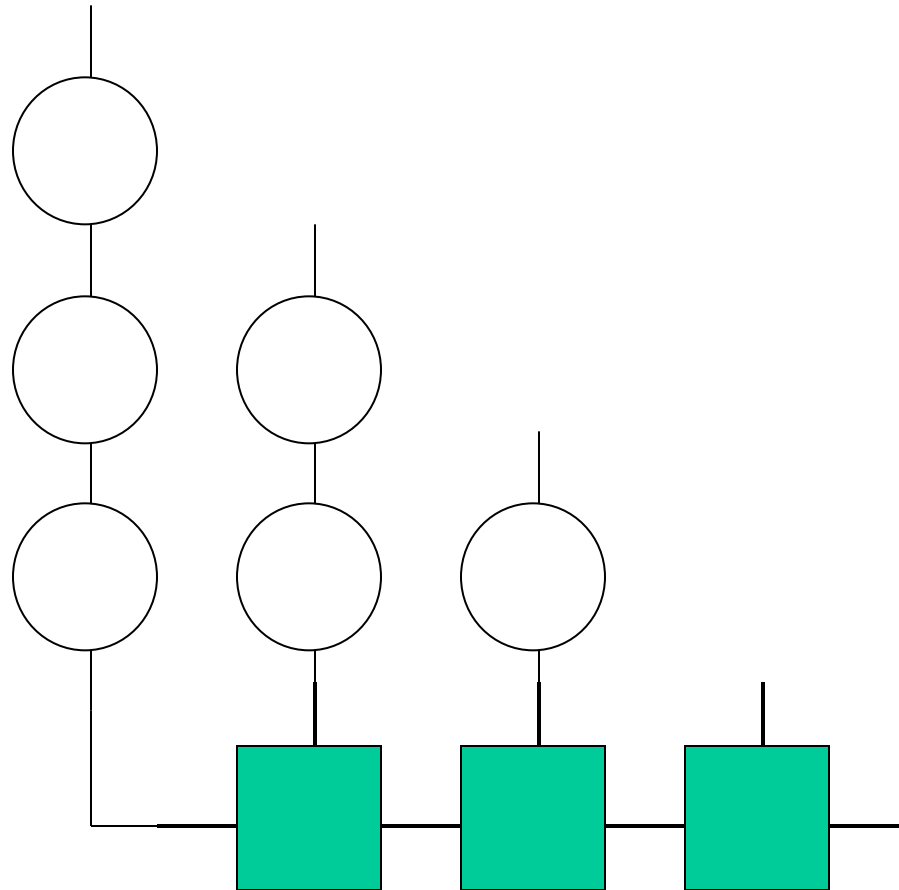
Note: we have no solver hooked up to
Lava that can do arithmetic ☹
So we can't formally verify equivalence
of the two different msb-bin to int
functions

On the bright side: a general rule!

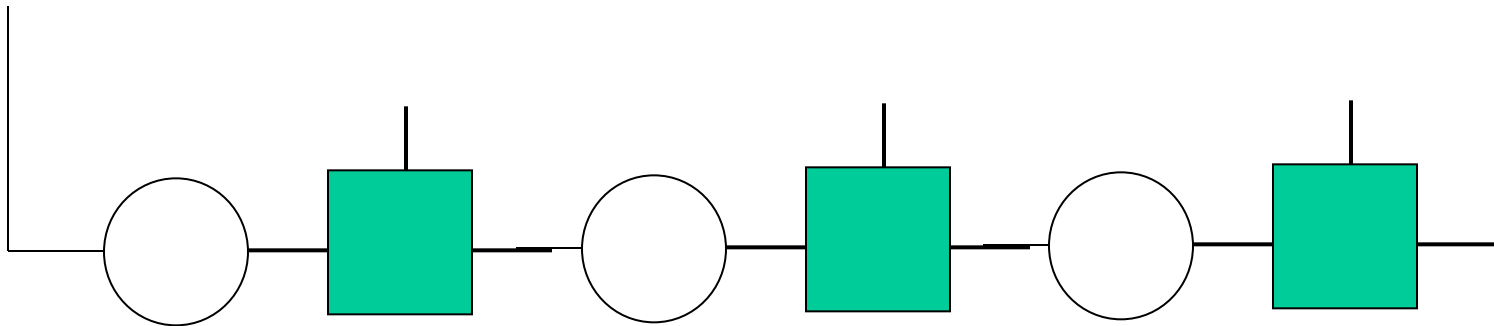
If



Then, this is the same as



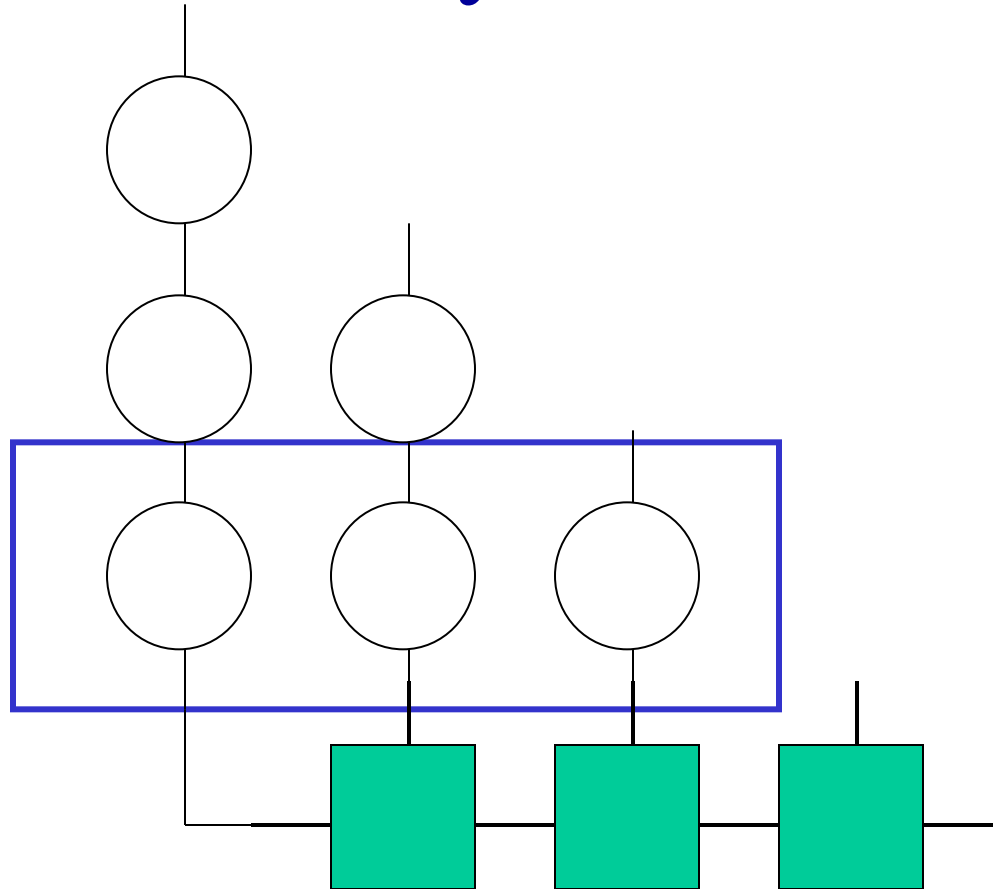
this, no matter what the
components are (and for any size)

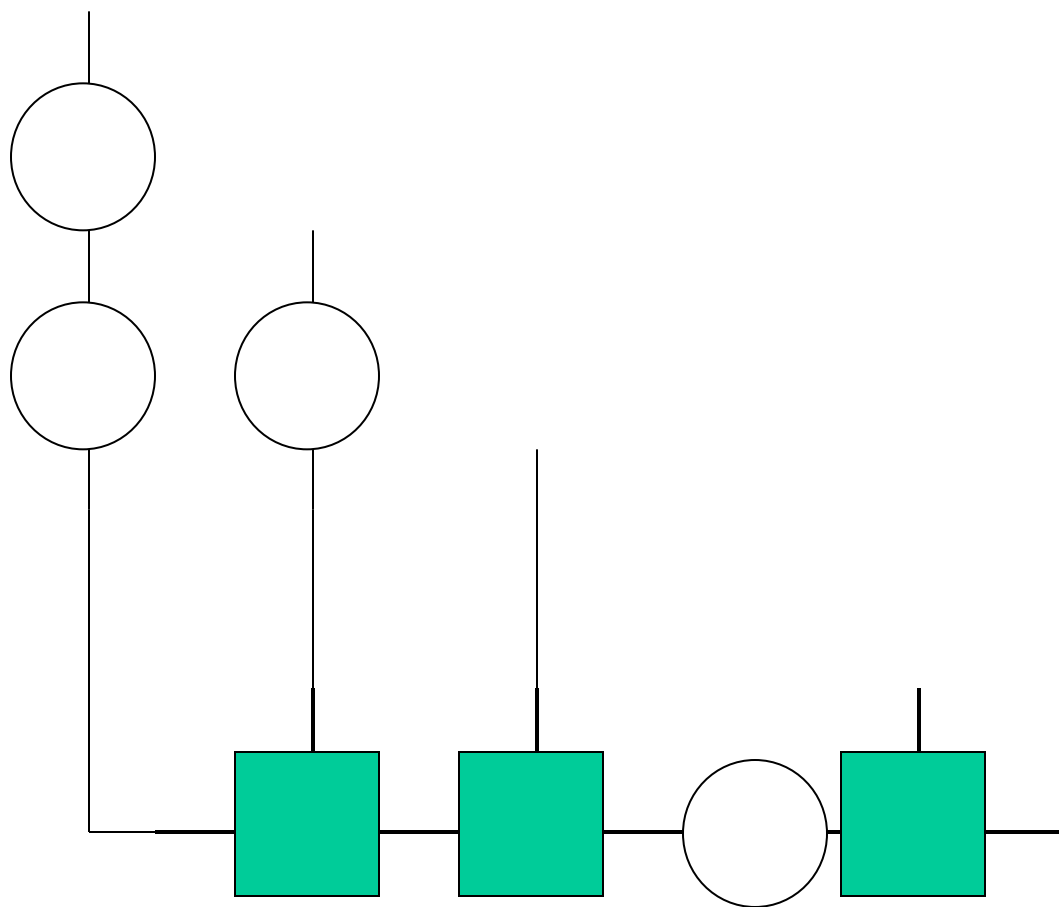


Why?

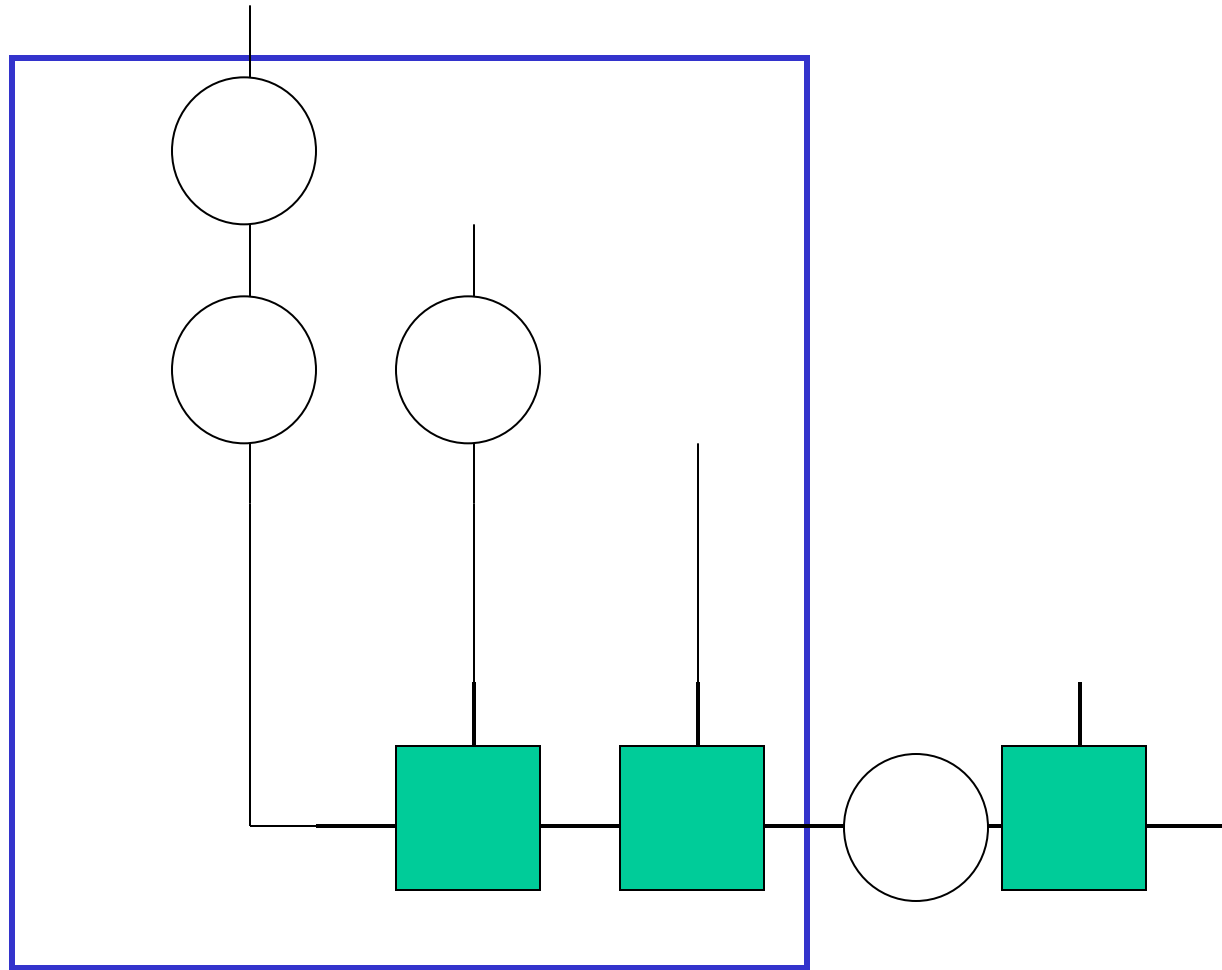
tri

map

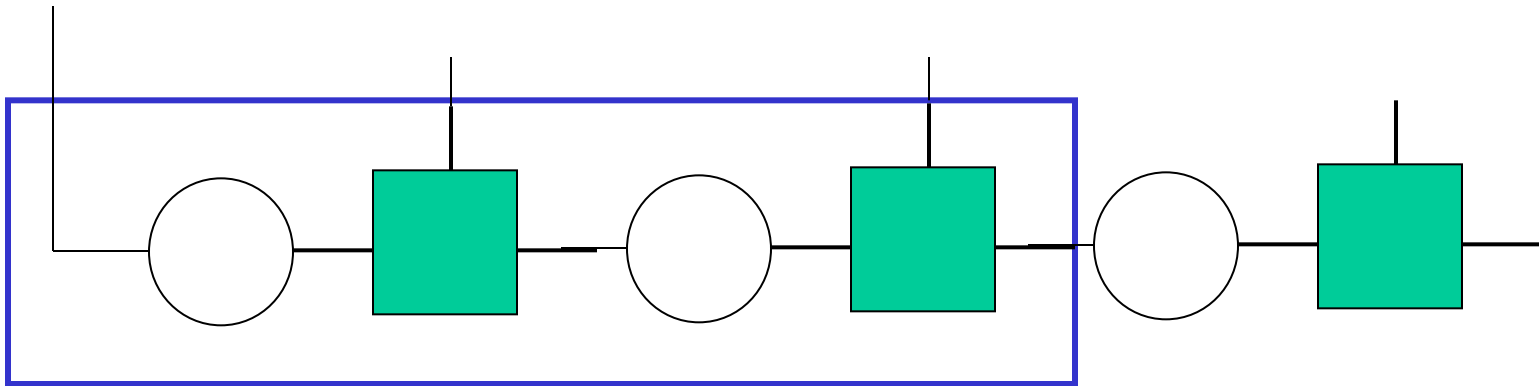




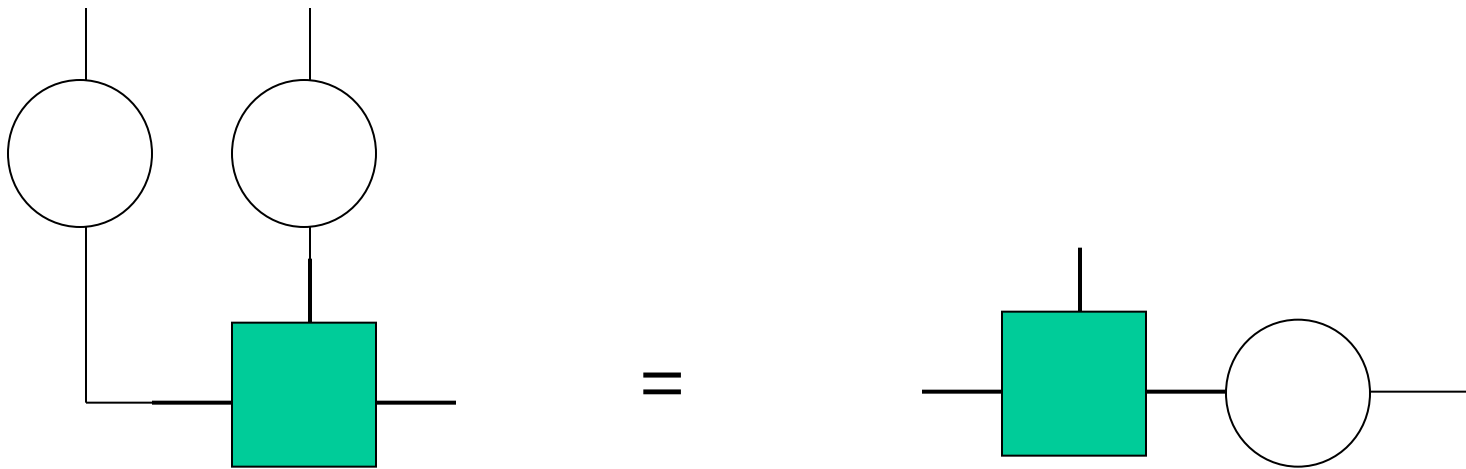
What's left?



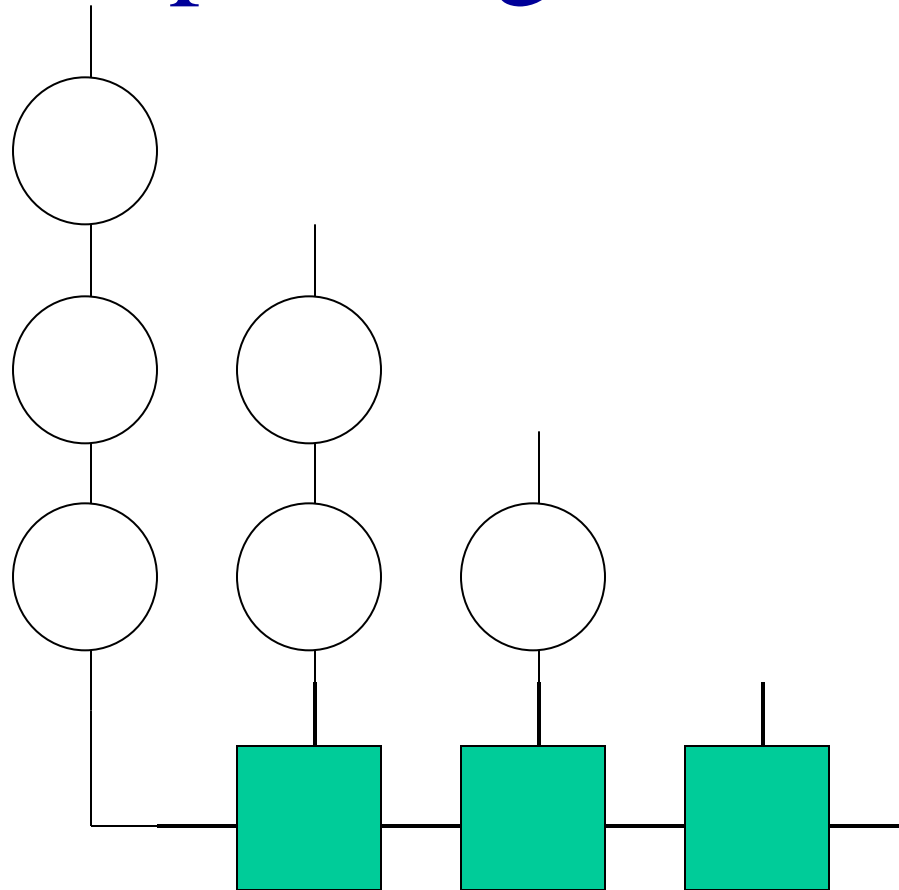
So, by induction



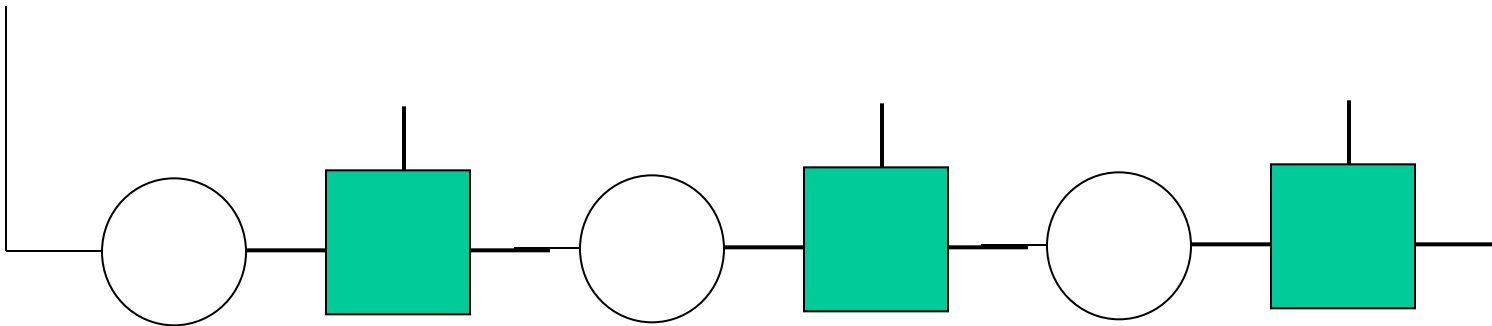
Does this look familiar?



Pipelining!



Equality holds for circle=delay



Checking equiv. of an instance

```
propEQ circ1 circ2 a = ok
  where
    out1 = circ1 a
    out2 = circ2 a
    ok = out1 <==> out2
```

```
propEQS circ1 circ2 n = propEQS circ1 circ2 (varList n "a")
```

```
withtri f = downtri (delay low) ->- lin f
```

```
piped f = lin cell
  where
    cell (a,b) = f (delay low a, b)
```

```
pipetst = smv (propEQS (withtri and2) (piped and2) 4)
```

Checking equiv. of an instance

```
propEQ circ1 circ2 a = ok
```

```
  where
```

```
    out1 = circ1 a
```

```
    out2 =
```

```
    ok =
```

```
propEQS
```

```
withtri f
```

```
pipet f = lin cell
```

```
  where
```

```
    cell (a,b) = f (delay low a, b)
```

```
pipetst = smv (propEQS (withtri and2) (pipet and2) 4)
```

On Mary's i5 laptop
Size 4 and 8 very quick
size 12 14 secs
size 16 didn't finish

Lava is good for stress-testing tools

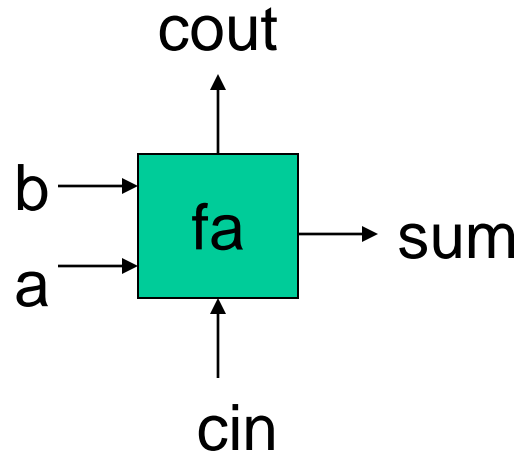
rList n "a")

Are connection patterns useful??

or do they just make programming harder?

Xilinx Lava provides part of the answer:

Full Adder in Xilinx Lava



$\text{fa}(\text{cin}, (a, b)) = (\text{sum}, \text{cout})$

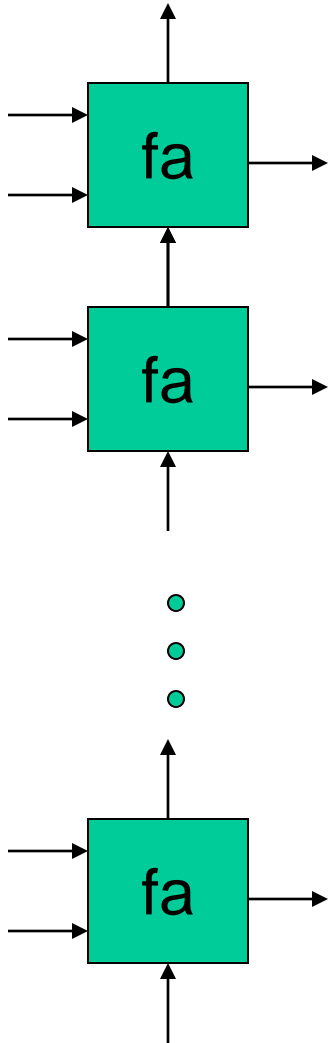
where

$\text{part_sum} = \text{xor}(a, b)$

$\text{sum} = \text{xorcy}(\text{part_sum}, \text{cin})$

$\text{cout} = \text{muxcy}(\text{part_sum}, (a, \text{cin}))$

Generic Adder



adder = **col** fa

Top Level

adder16Circuit

```
= do a <- inputVec "a" (bit_vector 15 downto 0)
      b <- inputVec "b" (bit_vector 15 downto 0)
      (s, carry) <- adder1 (a, b)
      sum <- outputVec "sum" (s++[carry])
              (bit_vector 16 downto 0)
```

? circuit2VHDL "add16" adder16Circuit

? circuit2EDIF "add16" adder16Circuit

? circuit2Verilog "add16" adder16Circuit

114 Lines of VHDL

```
library ieee ;
use ieee.std_logic_1164.all ;
entity add16 is
  port(a : in std_logic_vector (15 downto 0) ;
       b : in std_logic_vector (15 downto 0) ;
       c : out std_logic_vector (16 downto 0)
       ) ;
end entity add16 ;
```

```
library ieee, unisim ;
use ieee.std_logic_1164.all ;
use unisim.vcomponents.all ;
architecture lava of add16 is
  signal lava : std_logic_vector (0 to 80) ;
begin
  ...
  lut2_48 : lut2 generic map (init => "0110") port map (i0 => lava(5), i1 => lava(21), o => lava(48)) ;
  xorcy_49 : xorcy port map (li => lava(48), ci => lava(47), o => lava(49)) ;
  muxcy_50 : muxcy port map (di => lava(5), ci => lava(47), s => lava(48), o => lava(50)) ;
  lut2_51 : lut2 generic map (init => "0110") port map (i0 => lava(6), i1 => lava(22), o => lava(51)) ;
  xorcy_52 : xorcy port map (li => lava(51), ci => lava(50), o => lava(52)) ;
  muxcy_53 : muxcy port map (di => lava(6), ci => lava(50), s => lava(51), o => lava(53)) ;
  lut2_54 : lut2 generic map (init => "0110") port map (i0 => lava(7), i1 => lava(23), o => lava(54)) ;
  ...
end architecture lava;
```

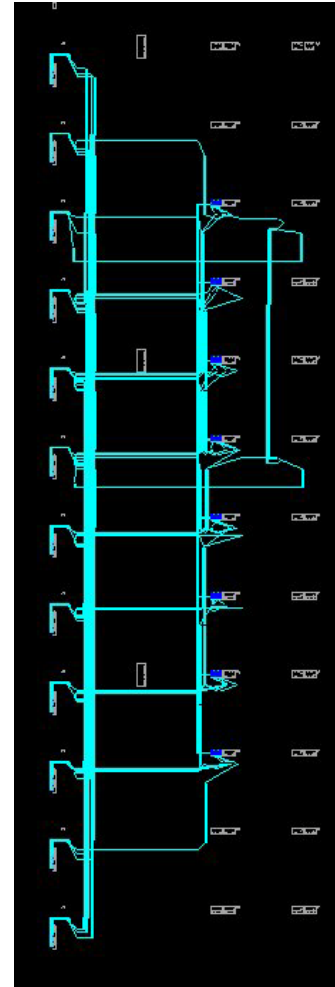
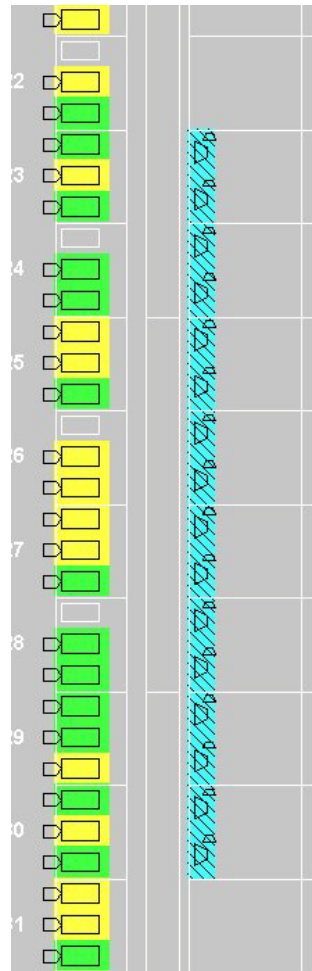
EDIF...

```
(edif add16
(edifVersion 2 0 0)
(edifLevel 0)
(keywordMap (keywordLevel 0))
(status
(written (timeStamp 2000 11 19 15 39 43)
(program "Lava" (Version "2000.14"))
(dataOrigin "Xilinx-Lava") (author "Xilinx Inc."))
)
)
...
(instance lut2_78
  (viewRef prim
    (cellRef lut2 (libraryRef lava_virtex_lib))
  )
  (property INIT (string "6"))
  (property RLOC (string "R-7C0.S1"))
)
...
(net lava_bit38
  (joined
    (portRef o (instanceRef muxcy_38))
    (portRef ci (instanceRef muxcy_41))
    (portRef ci (instanceRef xorcy_40))
  )
)
```

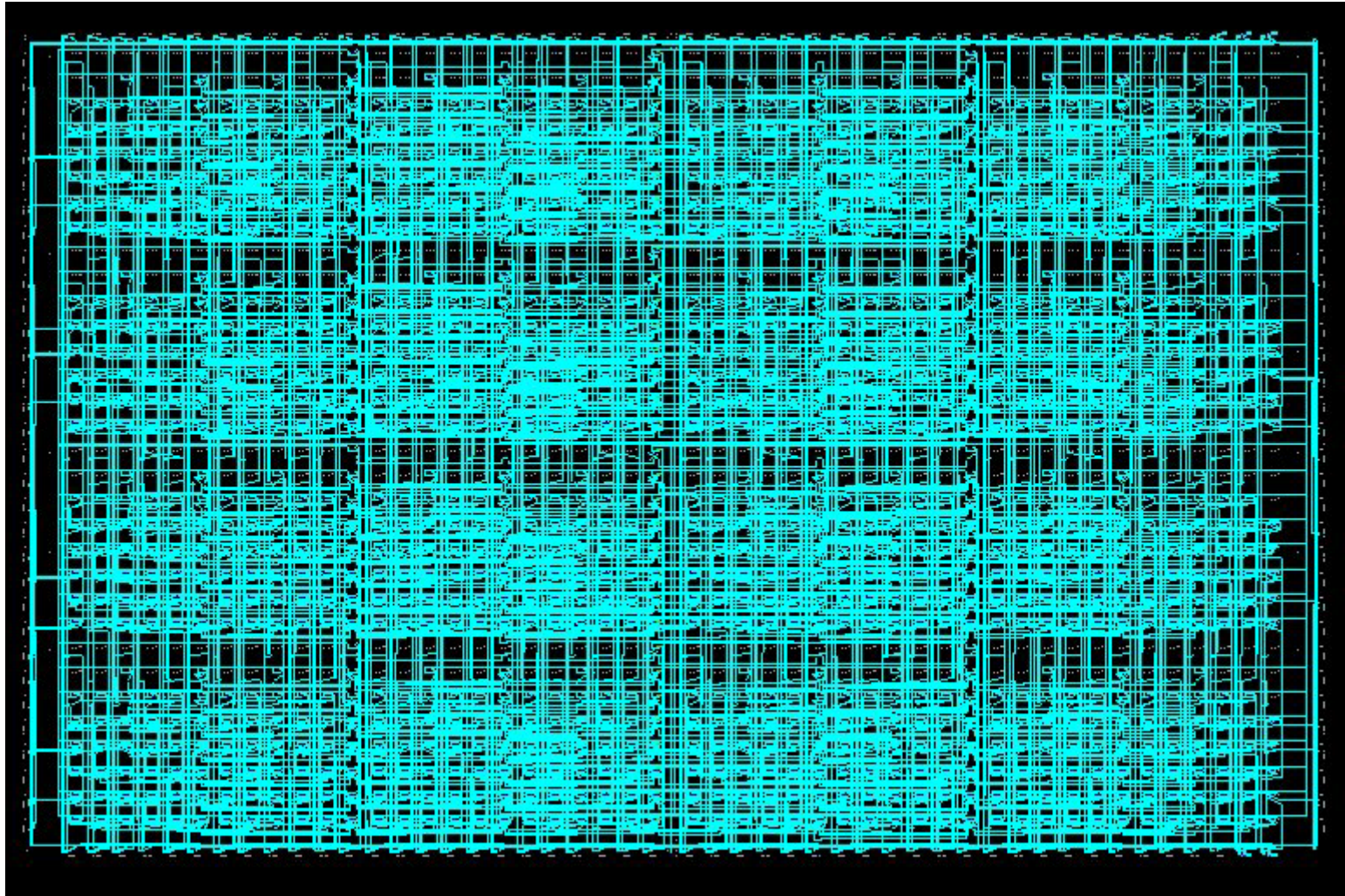
Xilinx FPGA Implementation

- 16-bit implementation on a XCV300 FPGA
- Vertical layout required to exploit fast carry chain
- No need to specify coordinates in HDL code

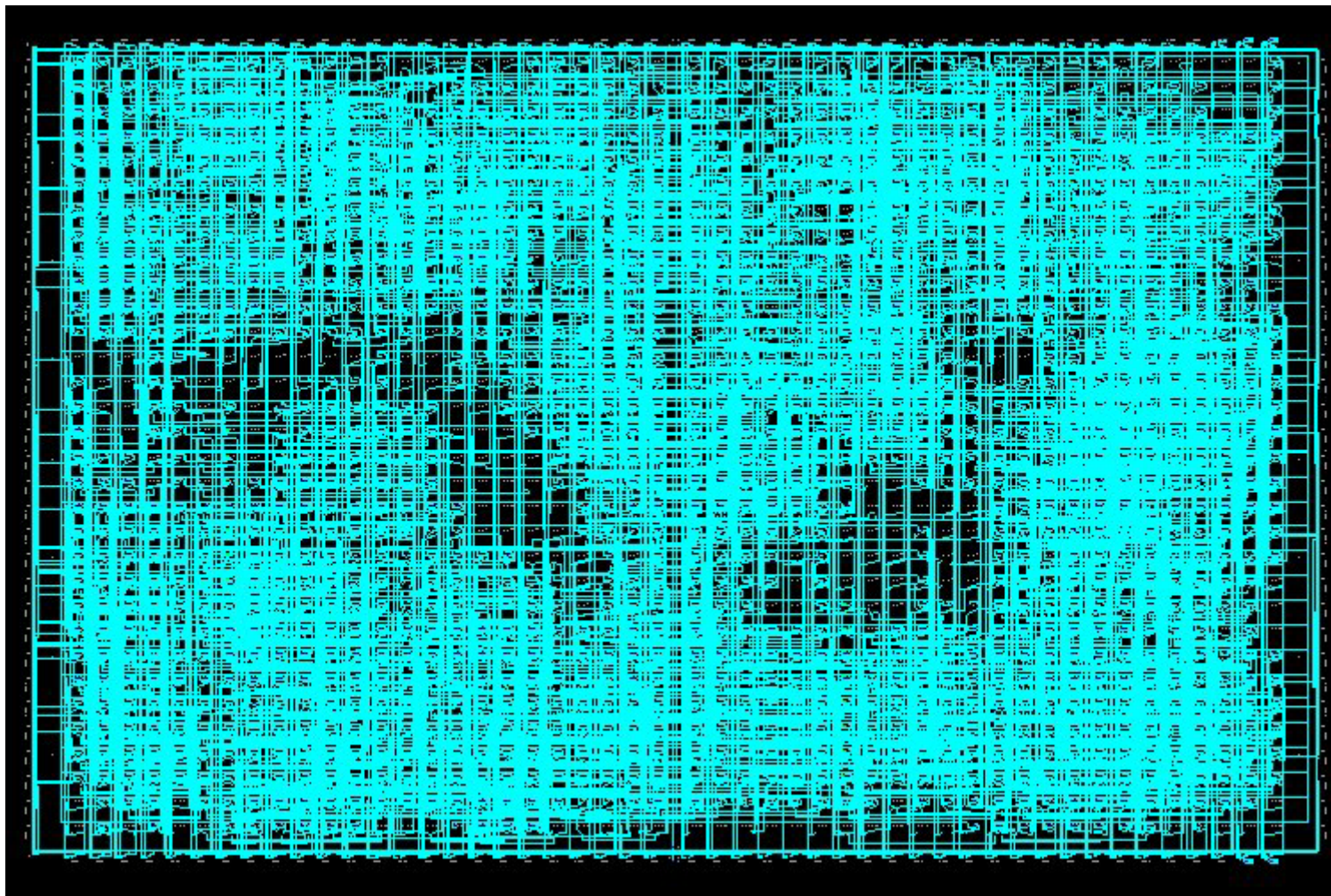
16-bit Adder Layout



Four adder trees



No Layout Information



Research questions: FPGA gen.

It is already quite hard to make use of fast carry chains on FPGAs (see papers on making fast adders, for instance)

How on earth can we make good use of DSPs, processors etc.??

One answer might be to use Functional Programming ☺
(This is what my latest grant proposal was about)

Another motivation for connection patterns and algebra

Work on Hawk for describing and reasoning about
processors showed really nice applications of transformations

See [John Matthew's slides](#)

Notes

Abstract circuits useful during design exploration (in simulation rather than formal verification)

Verify fixed size instances of generic circuits

Input to solvers (or VHDL netlist) generated by symbolic evaluation

Notes

No mechanical assistance in algebraic reasoning currently in Lava (but earlier research by Cachera using PVS was promising, 30 interactions down to 4 in mult. verif.)

Research question: what could we do with a first order logic theorem prover or SMT solver?

Exercise: Zero detection

- Define a generic circuit that
 - inputs a bit vector, and
 - outputs high if all bits are zero.

`zero_detect :: [Bit] -> Bit`

- Simple solution first
- Also think about circuit depth and delay