



Real-Time Systems

Exercise #2

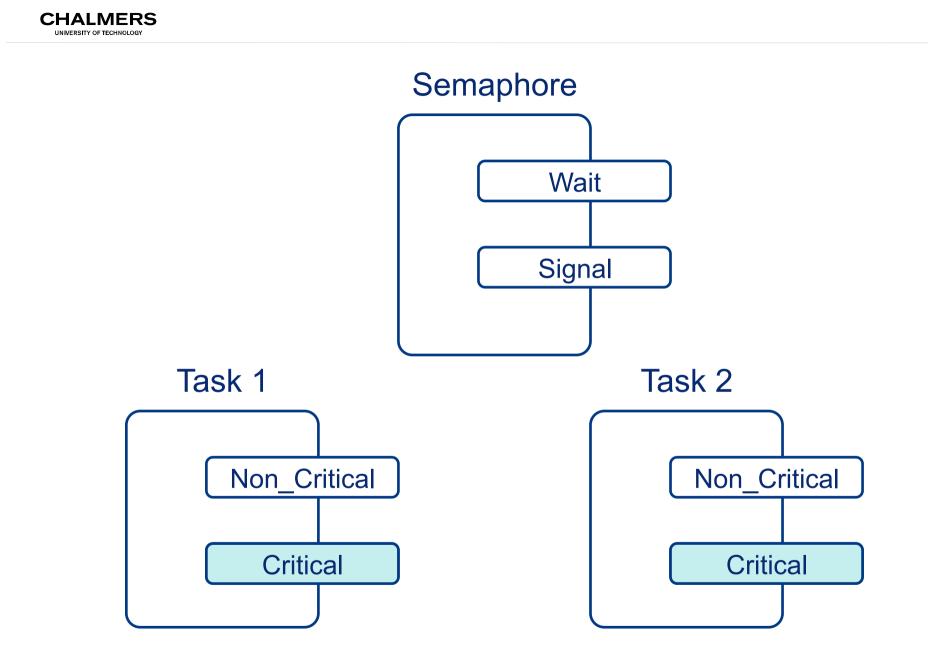
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Today: Revise key concepts

- SYNC
- Callbacks
- Queues
- Semaphores

Use a semaphore to synchronize two concurrent tasks in TinyTimber.



Call-back functionality in TinyTimber:

• TinyTimber guarantees that an object is handled like an exclusive resource during the execution of a method that belongs to the object <u>if</u> that method is called using SYNC().

If multiple concurrent calls, using SYNC(), are made to methods belonging to the same object, only one call will be granted access to the object. The other calls will be blocked (put in a waiting queue.)

When the object is available again, one of the blocked calls will be unblocked and the corresponding method is executed by means of a basic call-back functionality in TinyTimber.

Call-back functionality in TinyTimber (cont'd):

- Although this basic call-back functionality is sufficient in many cases, TinyTimber lacks one powerful property that protected objects, monitors and semaphores have:
 - The basic call-back functionality in TinyTimber cannot account for conditions relating to the <u>contents of an object</u>.

Note: this prevents us from implementing blocking versions of the Get/Put methods in the circular buffer example in an earlier lecture.

• Thus, in order to use advanced resource management with TinyTimber we must provide a <u>call-back functionality add-on</u>.

Call-back functionality add-on:

- A task requests access to a certain resource (object) with a call to a method belonging to that resource (object).
- If access is not granted (because a condition regarding the object state prevents this) the method call will be blocked.
- If the calling task used ASYNC() to request the resource the task itself is not blocked but continues executing code.
- Implementing call-back functionality means that a calling task supplies ASYNC() with information about a method to wake up (call back) when the resource becomes available.
- Since multiple tasks may want to request the resource, the provided call-back information must be stored in a queue.

Method parameter and return-value convention:

- TinyTimber uses a uniform approach to method definitions: all methods must have two parameters of specific types
 - The first parameter <u>must</u> be a pointer to an object of the class to which the method belongs. This pointer (often named 'self') allows the methods to access the state variables of the object.
 - The second parameter <u>must</u> be of type 'int' and can be used as an input parameter to the method (but can also be ignored).
- For this reason calls to method operations in the kernel (TINYTIMBER(), ASYNC(), SYNC(), AFTER(), ...) <u>must</u> include these parameters in addition to a method reference.
- The return value of a method must be of type 'int', unless no value is returned (in which case type 'void' is used).

Method parameter and return-value work-around:

- If an input parameter of type 'xxx' (different than 'int') is needed for the method, type casting the argument to type 'int' must be performed at call time; then the parameter is type-cast back to type 'xxx' within the method itself.
- If multiple input parameters are needed, they should be stored in a struct, and a pointer to the struct should be passed as the argument at call time (with appropriate type casting).
- This work-around is also applicable to return values.

A semaphore s is an integer variable with value domain ≥ 0 Atomic operations on semaphores:

<pre>Init(s,n):</pre>	assign s an initial value n
Wait(s):	<pre>if s > 0 then s := s - 1; else "block calling task";</pre>
Signal(s):	<pre>if "any task that has called Wait(s) is blocked" then "allow one such task to execute"; else s := s + 1;</pre>

Problem: Implement a class Semaphore in C using the TinyTimber kernel, to synchronize two concurrent tasks.

- The object should receive an initial value when it is created.
- The class should have two methods, Wait and Signal, that work in accordance with the definition of semaphores.
- The methods should have support for call-back functionality

Solution overview:

- 1. Define a **data type for call-back information**, that can also be stored as an element in a queue
- 2. Implement functions for manipulating **a queue** containing elements of the call-back information data type
- 3. Define a **class Semaphore with Wait and Signal** methods, as well as an initialization macro
- 4. Implement the Semaphore (Wait and Signal methods)
- 5. Create application code that **uses** the semaphore

// Define a data type for call-back information, that can also
// be used as an element in a queue

```
struct call_block;
typedef struct call_block *Caller;
```

```
typedef struct call_block {
   Caller next; // for use in linked lists
   Object *obj;
   Method meth;
} CallBlock;
```

#define initCallBlock() { 0, 0, 0 }

// Implement functions for manipulating a queue containing
// elements of the call-back information data type

```
void c enqueue(Caller c, Caller *queue) {
    Caller prev = NULL, q = *queue;
    while (q) { // find last element in queue
        prev = q;
        q = q - next;
    if (prev == NULL)
        *queue = c; // empty queue: put `c' first
    else
        prev->next = c; // non-empty queue: put `c' last
    c->next = NULL;
}
Caller c dequeue (Caller *queue) {
    Caller c = *queue;
    if (C)
        *queue = c->next; // remove first element in queue
    return c;
```

// Define a class Semaphore with Wait and Signal methods, // as well as an initialization macro

typedef struct {
 Object super;
 int value;
 Caller queue;

} Semaphore;

// Note that TinyTimber methods only accept type 'int' for the second // parameter. This means that, if we want to send a parameter of another // scalar type (i.e. a pointer), we will have to trick the system by // "type casting" to 'int' before a call, and then back to the original // type within the method.

void Wait(Semaphore*, int); void Signal(Semaphore*, int);

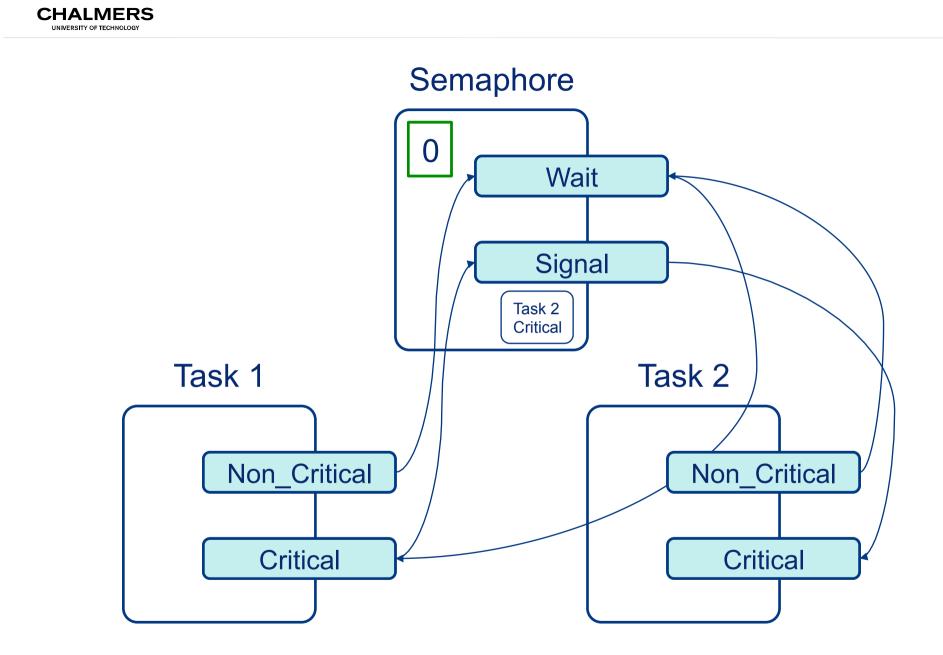
#define initSemaphore(n) { initObject(), n, 0 }

// Implement the methods Wait and Signal

```
void Wait(Semaphore *self, int c) {
   Caller wakeup = (Caller) c; // type-cast back from `int'
   if (self->value > 0) {
      self->value--;
      ASYNC(wakeup->obj, wakeup->meth, 0);
   }
   else
      c_enqueue(wakeup, &self->queue);
}
void Signal(Semaphore *self, int unused) {
   if (self->queue) {
      Caller wakeup = c_dequeue(&self->queue);
      ASYNC(wakeup->obj, wakeup->meth, 0);
   }
   else
      self->value++;
}
```

```
// Define two identical tasks using the same semaphore
Semaphore Sem = initSemaphore(1); // binary semaphore
typedef struct {
    Object super;
    CallBlock cb; // where call-back information is stored
} Task;
Task task1 = { initObject(), initCallBlock() };
Task task2 = { initObject(), initCallBlock() };
```

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}

Example: semaphores in C

```
void kickoff(Task *self, int unused) { // TinyTimber's first scheduled
ASYNC(&task1, Non_Critical, 0); // event
ASYNC(&task2, Non_Critical, 0); // spawn two identical tasks
}
int main() { // we enter here after system startup,
```