

Operating Systems and Networks

Pierre Kleberger

pierre.kleberger@chalmers.se

Computer Science & Engineering

Adaption of slides by Andreas Larsson and Anders Gidenstam

With selected slides from:

- Kurose & Ross, "Computer Networking"

Roadmap

- Operating Systems
 - What is an Operating System
 - OS evolution
 - OS details
- Networking
 - The Internet
 - Network protocols
 - Security

What is an Operating System?

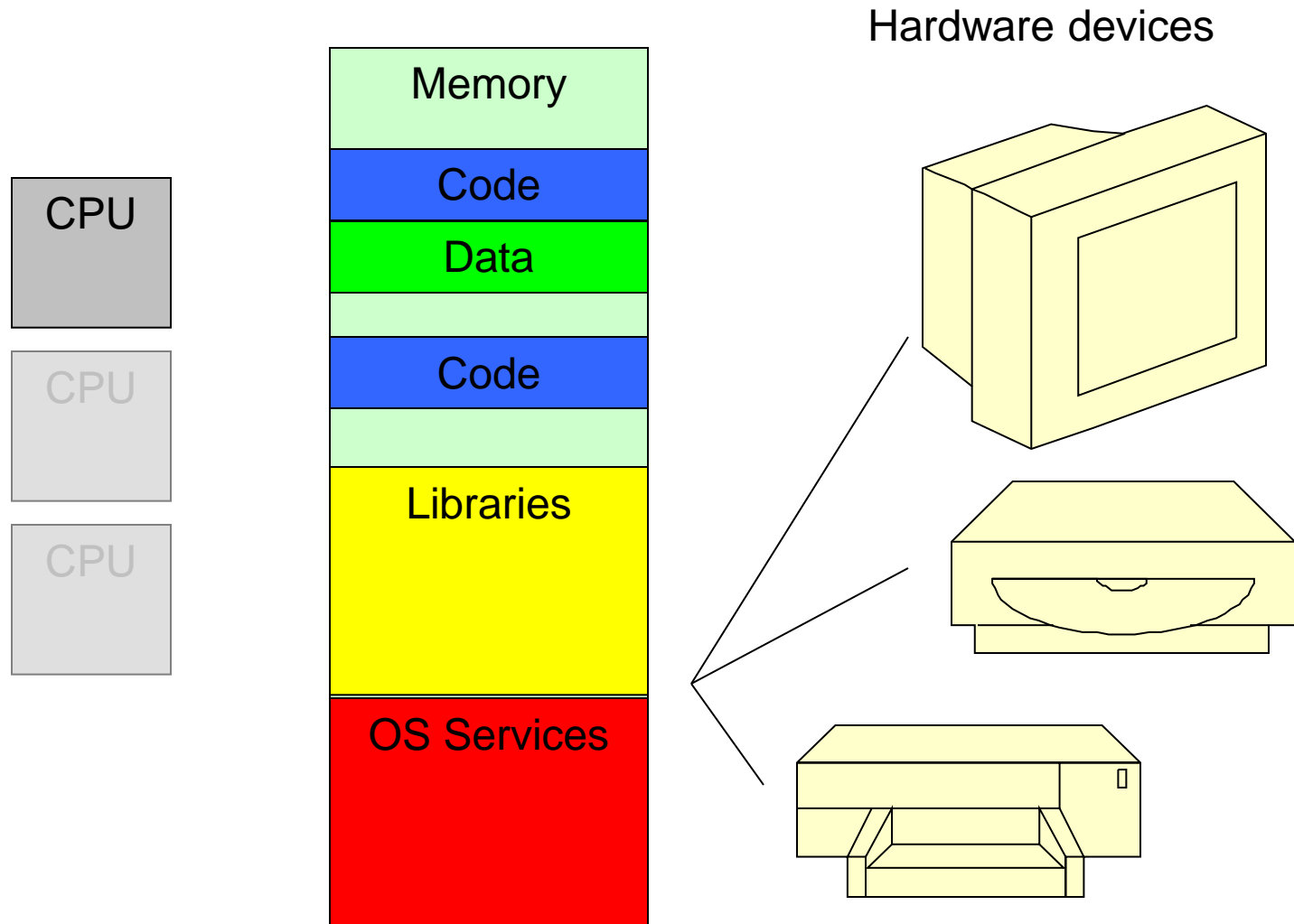
- Intermediary between the user and the hardware.
- Controls the execution of application programs
- Is an interface between applications and hardware
- Operating system goals:
 - Execute user programs
 - Facilitate problem solving for the users
 - Make the computer system convenient to use
 - Use the computer hardware in an efficient manner

The Computer: End-user's view

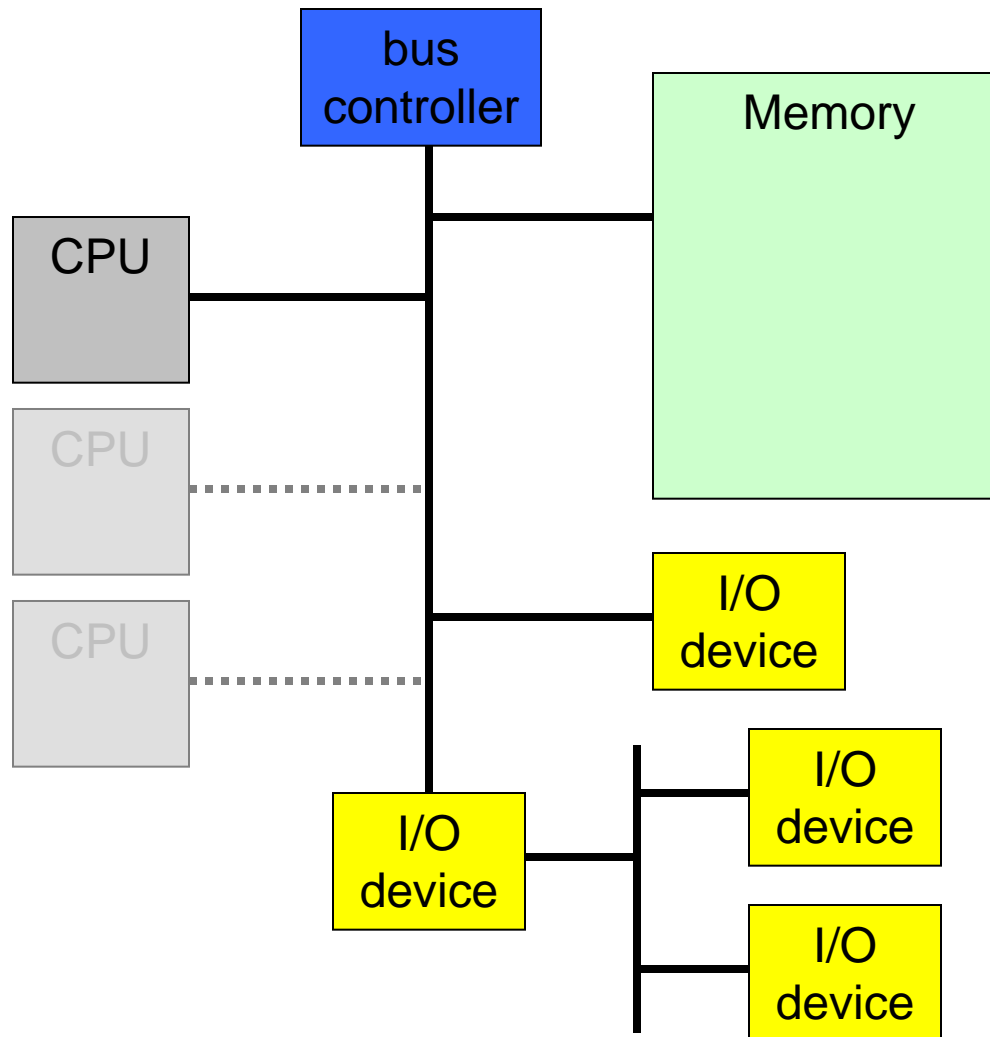


The Computer:

Application programmer's view

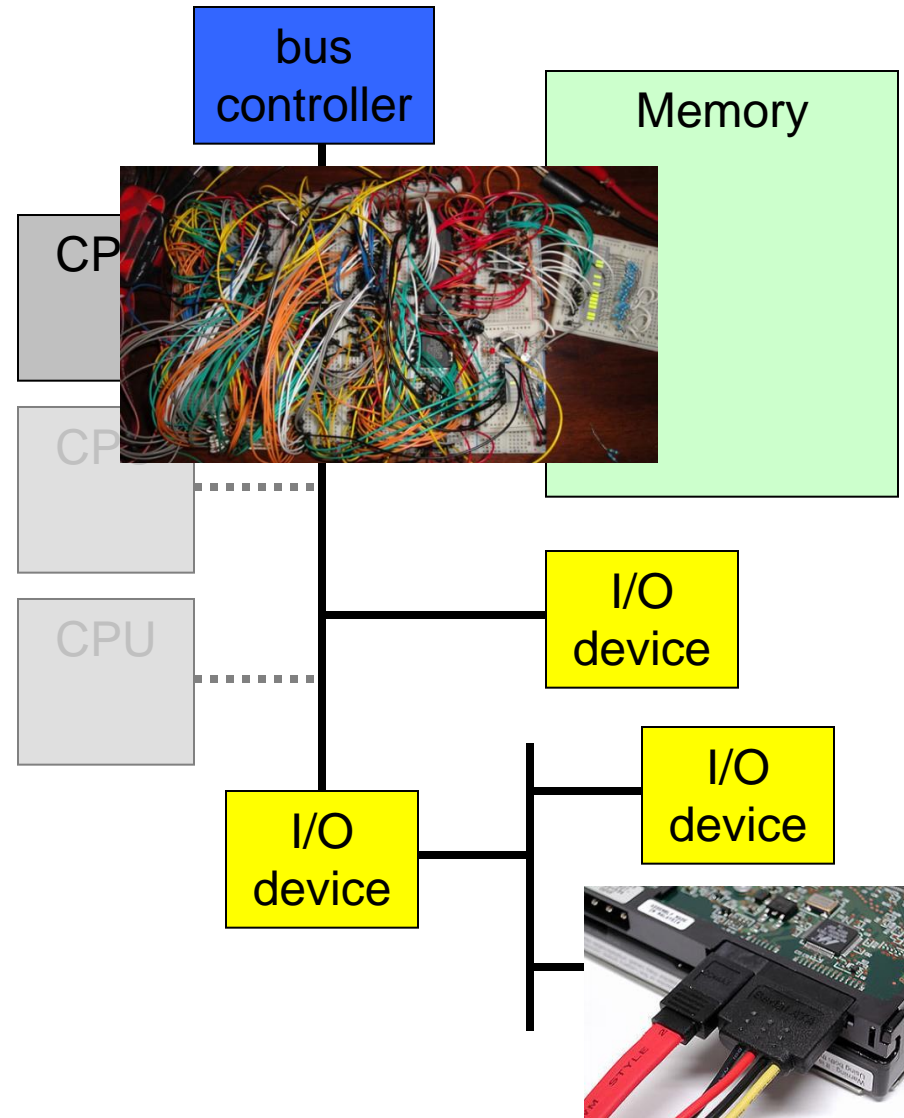


The Computer: OS programmer's view



Computer Hardware

- Processors
- Main Memory
 - Primary ("real") memory
 - Volatile
- I/O devices
 - secondary memory devices
 - communications devices
 - Screen, keyboard, network
- System bus
 - communication among processors, memory, and I/O modules

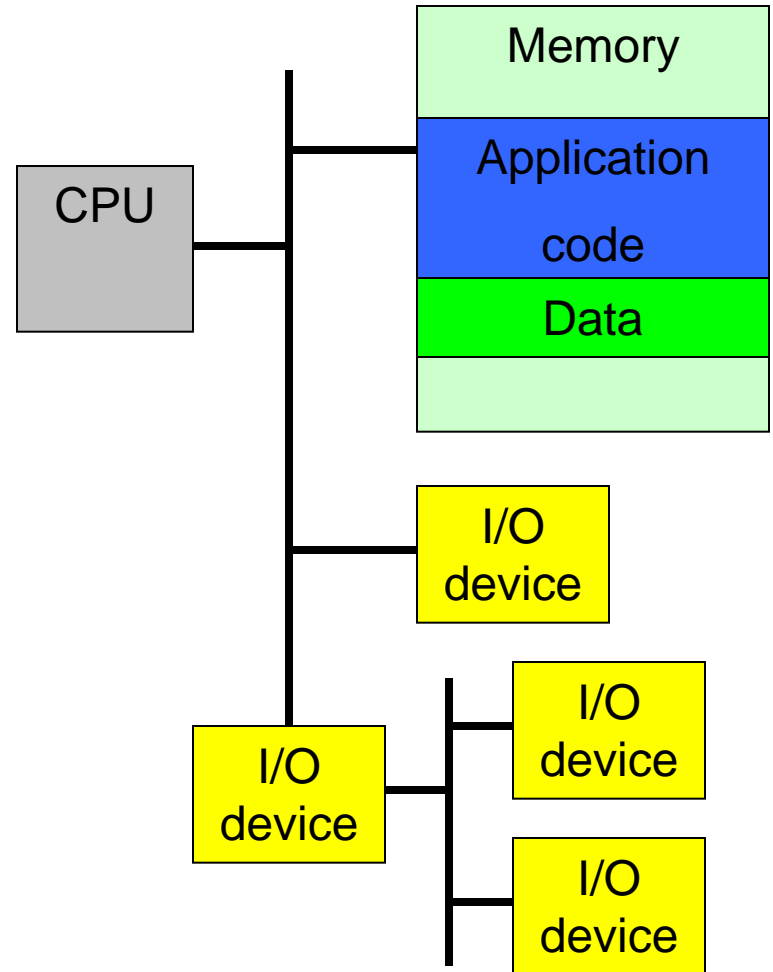


Introduction

- Operating Systems
 - What is an Operating System
 - OS evolution
 - OS details
- Networking
 - The Internet
 - Network protocols
 - Security

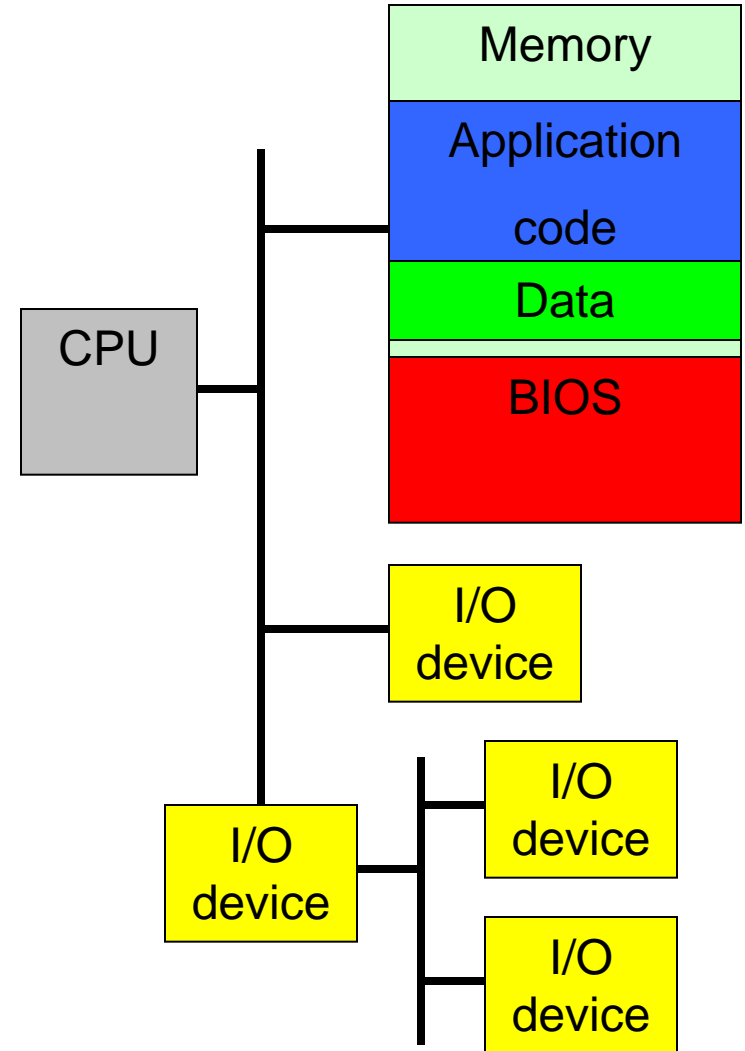
The evolution of operating systems

- The beginning
 - No OS
 - Every application had to do everything by itself
 - One program at a time
- Surely, still not so today?
 - microcontrollers



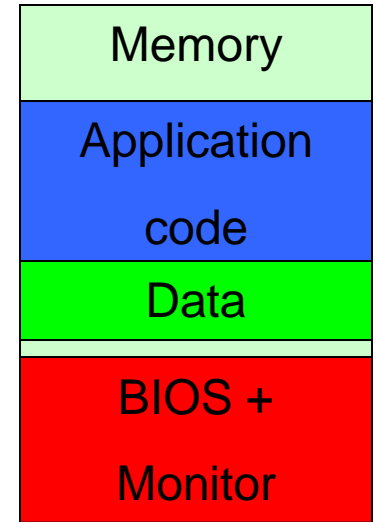
The evolution of operating systems

- BIOS
 - Basic Input Output System
 - In Read Only Memory (ROM)
 - Provides interface routines for accessing the hardware
- Still, only one program at a time



Batch processing

- In the 50s computers were expensive and rare, so efficient utilization was important
- Simple Batch Systems
 - Queue of jobs, run one at the time
 - Monitor
 - Software that controls the running programs
 - Batch jobs together
 - Program branches back to monitor when finished
 - Resident monitor is in main memory and available for execution



Uni- and Multiprogramming (1)

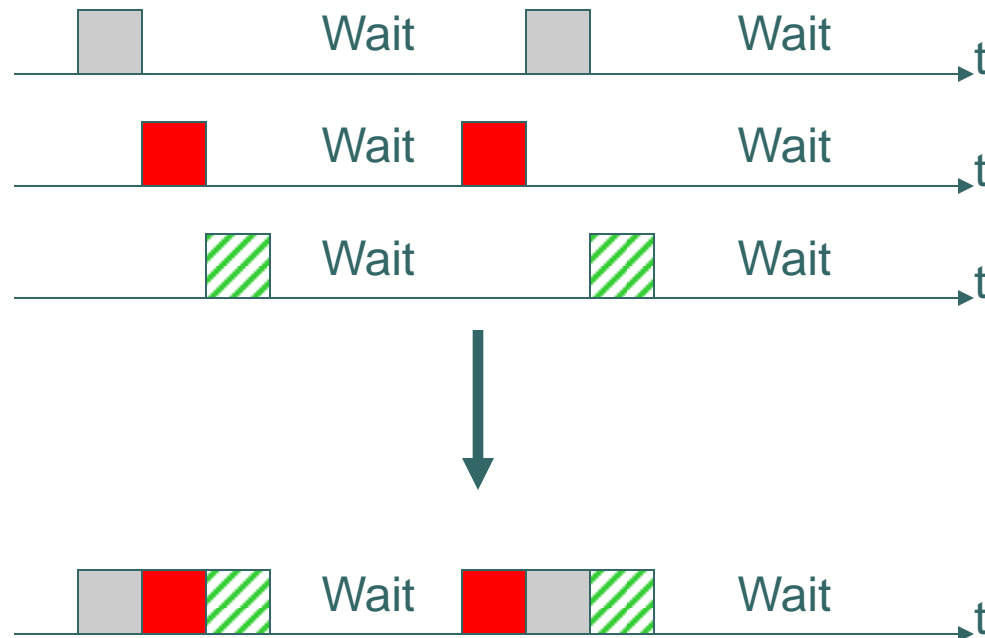
Uniprogramming

- One single program is running
- Processor must wait for I/O operations to complete before proceeding
- Leads to poor processor utilization



Uni- and Multiprogramming (2)

Multiprogramming

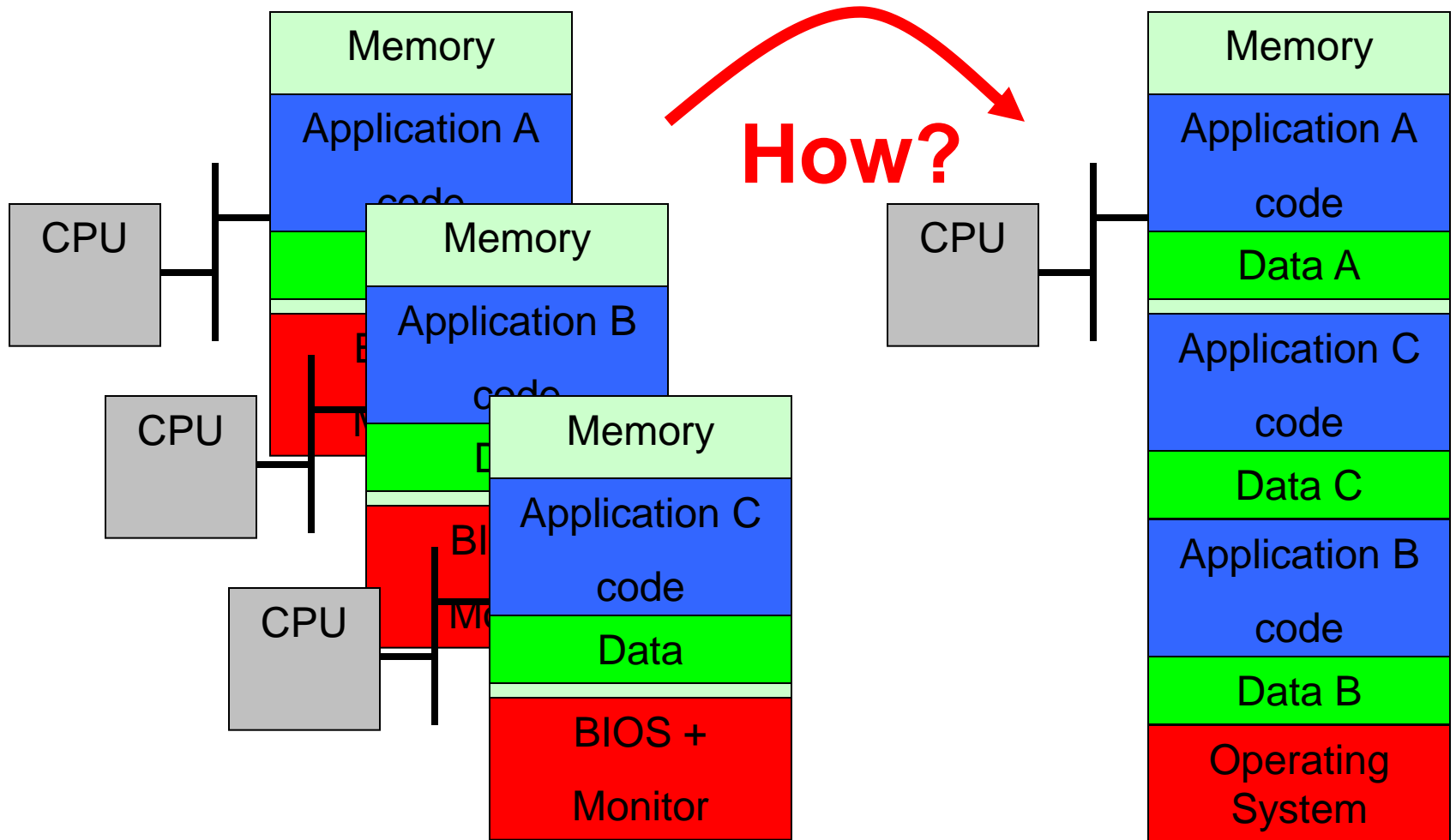


Uni- and Multiprogramming (3)

Multiprogramming

- Multiprogramming
 - Switch jobs at regular intervals
 - Benefits
 - Many applications running at the same time
 - Allows many simultaneous users
 - Interactive programs
 - “Real-time” interaction with user
 - Parallel/concurrent applications
 - Next step
 - Multiprocessor computers

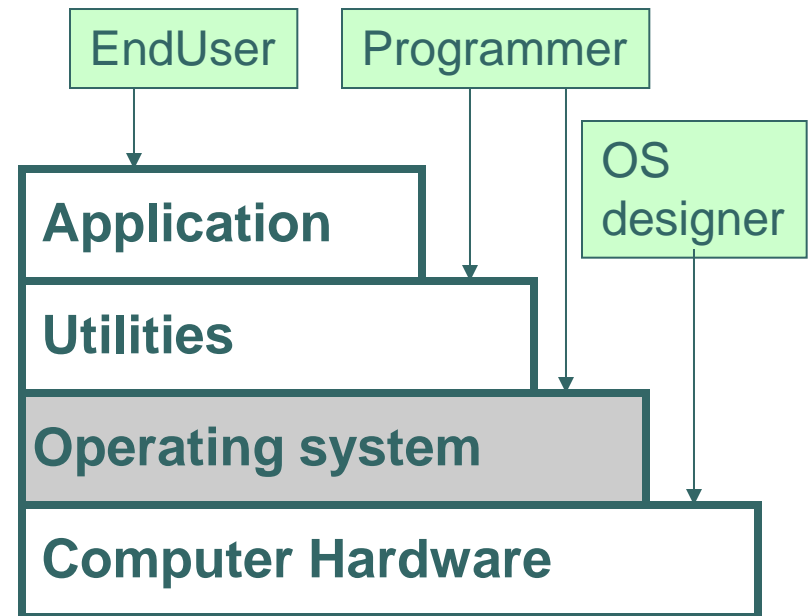
Multiprogramming – The Challenge



Operating System Architecture

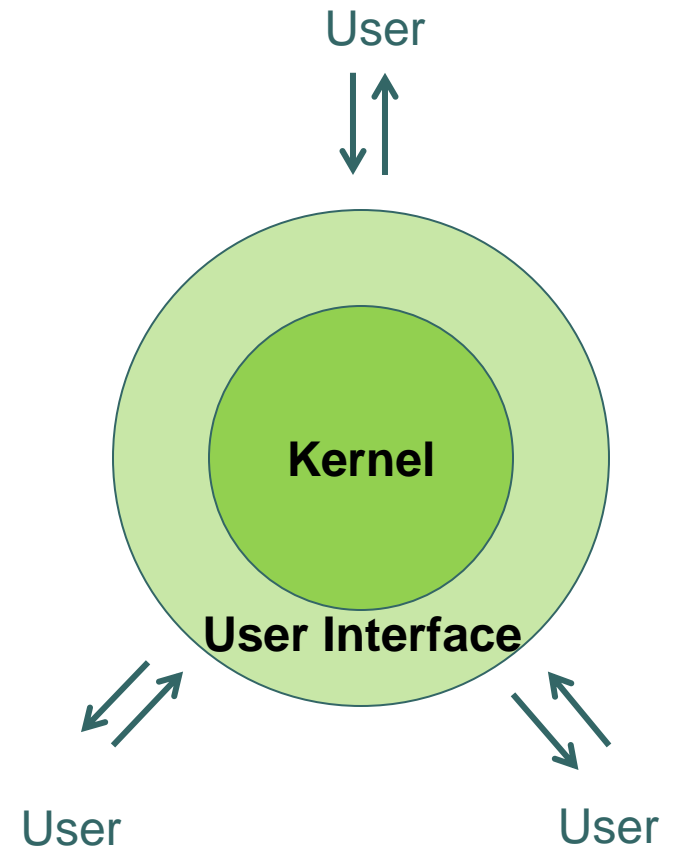
Software in the system

- Applications
- System software
 - Utilities
 - Compilers
 - Interpreters
 - Operating System
 - Shell
 - GUI
 - Command line
 - Kernel
 - The core of the OS



Services provided by the OS

- Program execution
- Access to I/O devices
- Controlled access to files
- Error detection
 - Hardware errors
 - Software errors
- Program development



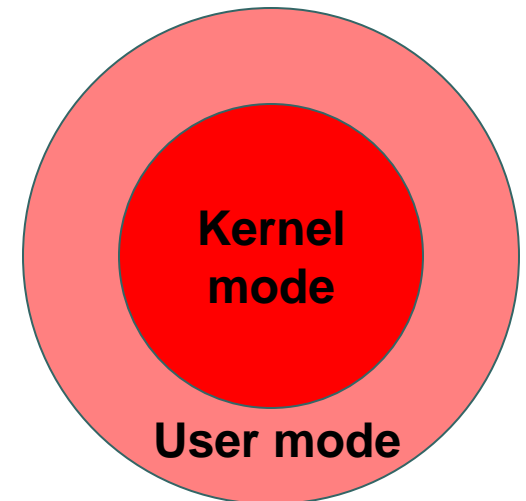
Kernel overview

- Portion of operating system that is in main memory
 - Contains most-frequently used functions
- Resource control
 - CPU Scheduling
 - Memory manager
 - File manager
 - Device drivers
- Bootstrap
 - Get the operating system running at system start



Kernel Security

- Privileged mode (Kernel mode)
 - Allowed to execute all CPU instructions
 - Access to all I/O devices
- Unprivileged mode (User mode)
 - Only a limited number of CPU instructions can be executed
 - e.g. access to memory and I/O devices are restricted



Roadmap

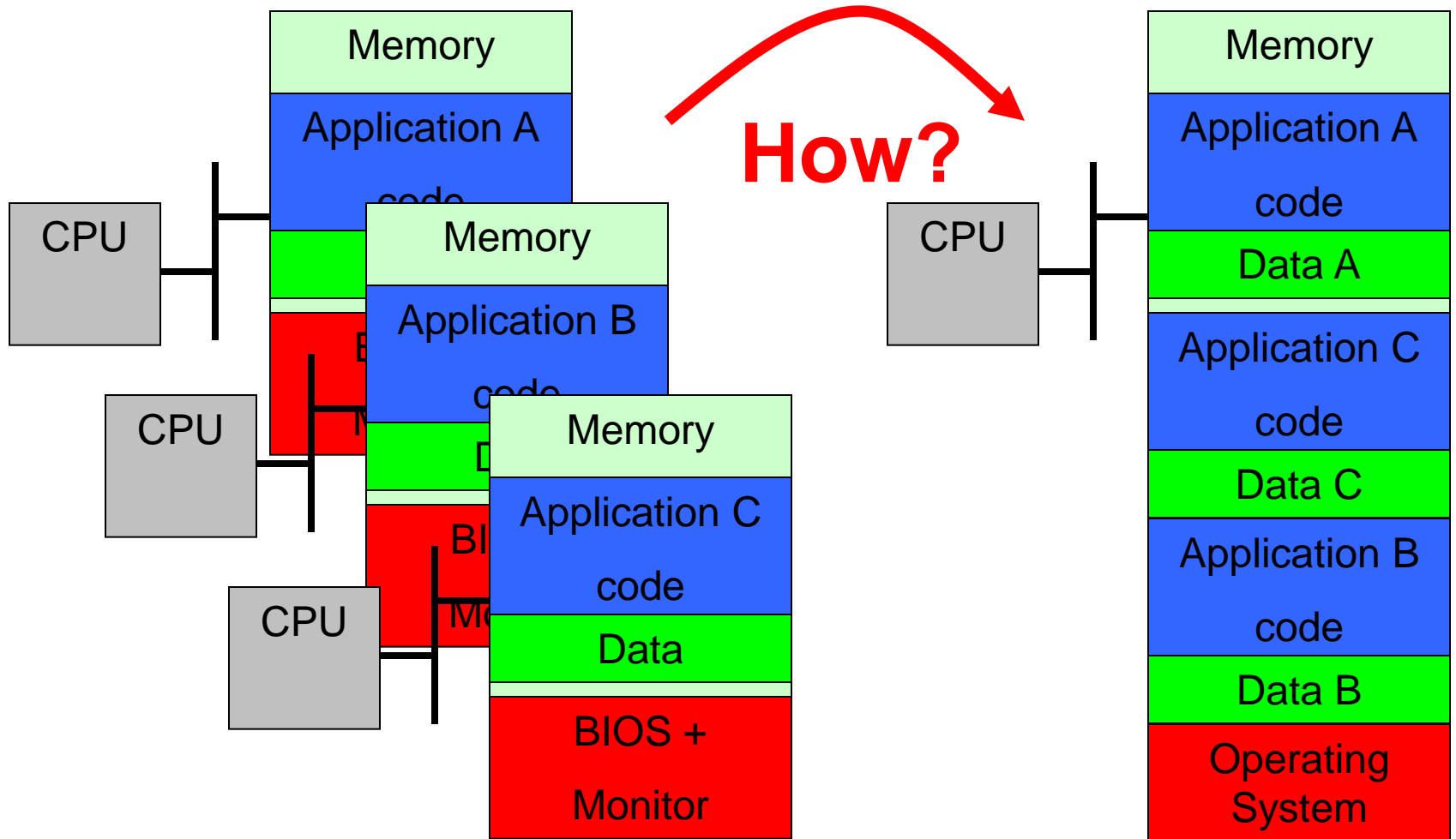
○ Operating Systems

- What is an Operating System
- OS evolution
- OS details
 - Context Switching
 - Virtual Memory
 - Resource Competition and Deadlock
 - File Systems
 - Interprocess Communication

○ Networking

Do you remember?

Multiprogramming – The Challenge



The answer:

Processes

P1: Word

P2: Chrome

P3: ...

- Process

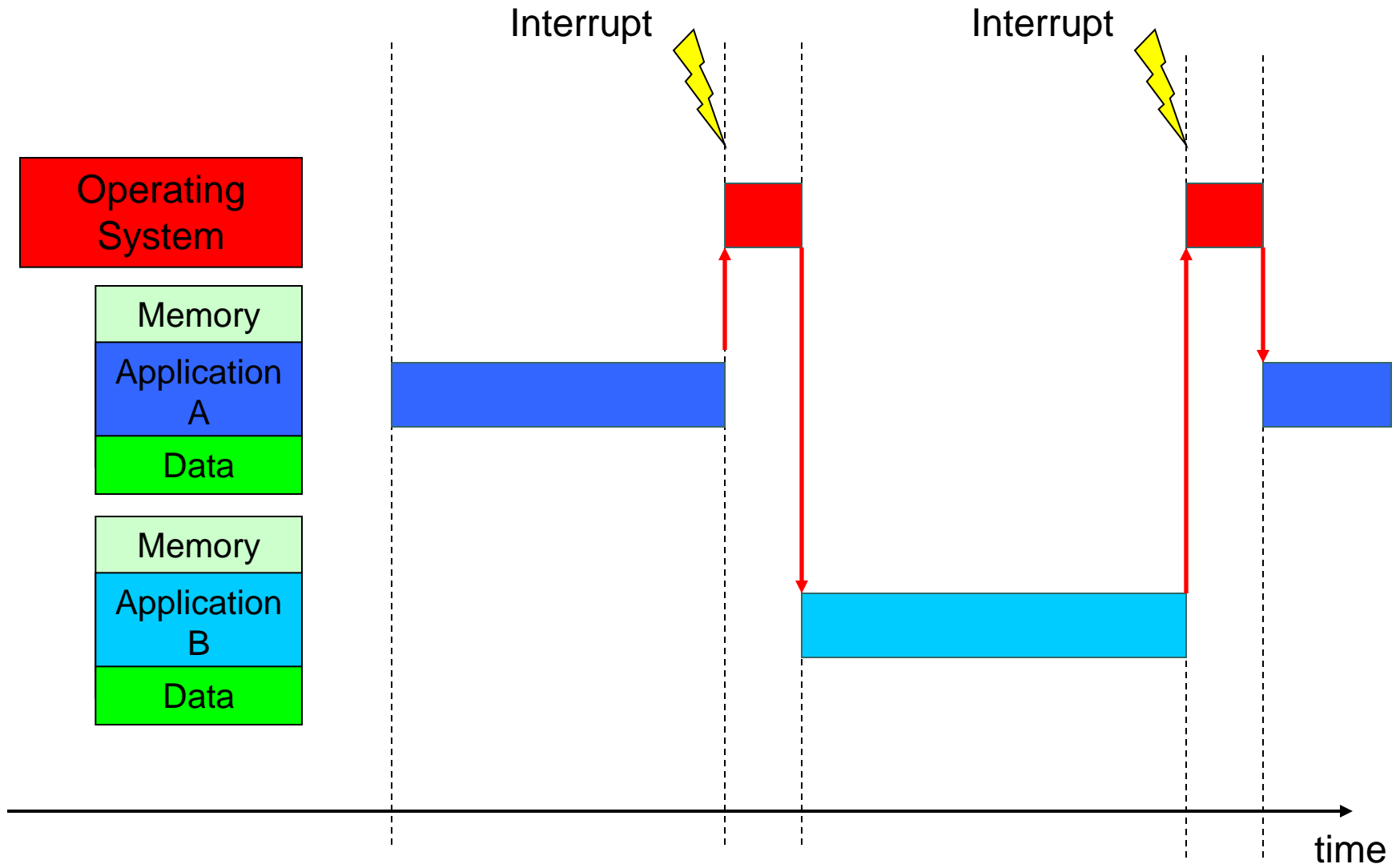
- A program in execution
- The OS presents a simpler "virtual" computer for exclusive use to the program/programmer

- A process includes:

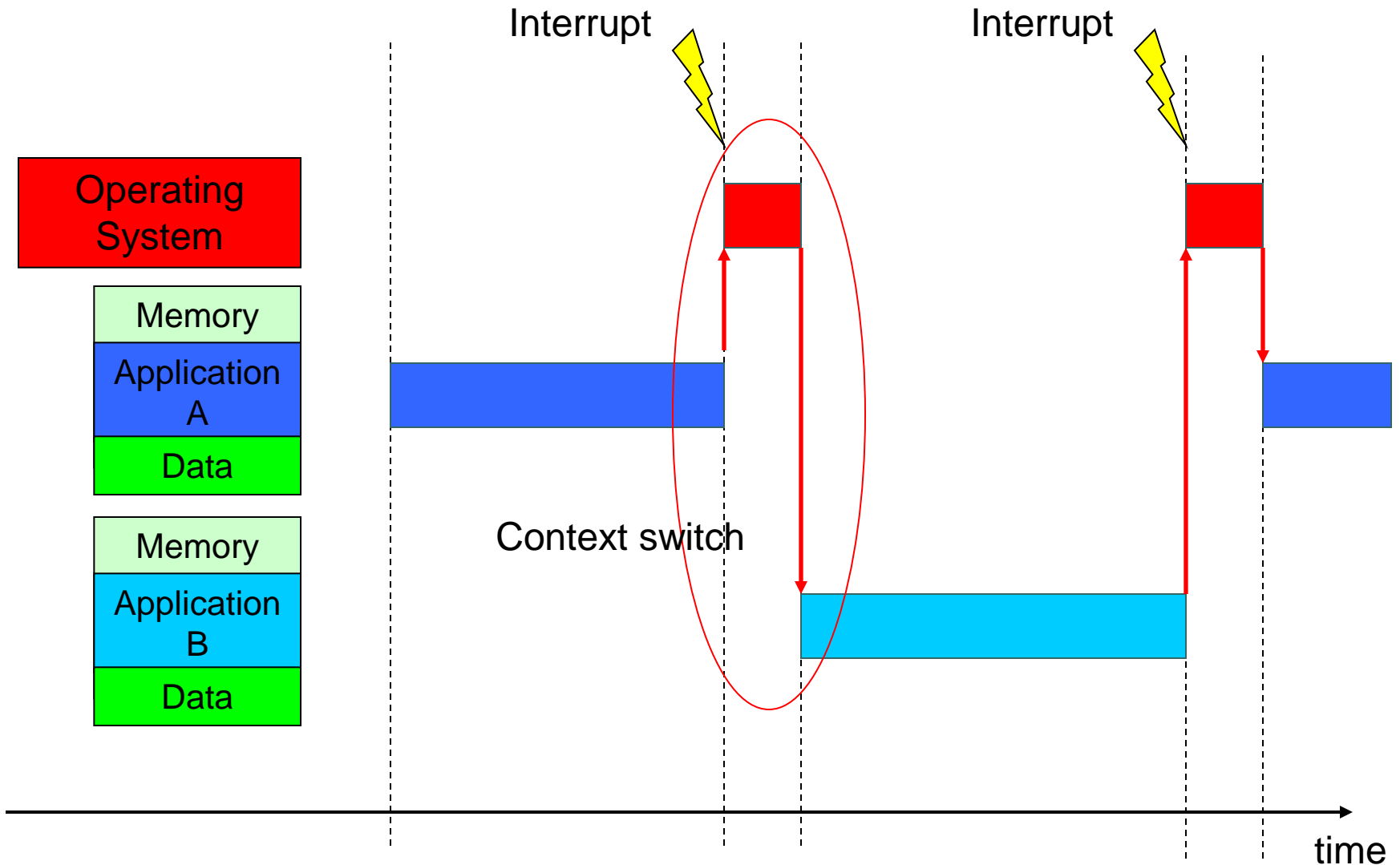
- Program code
- Program data and stack
 - The variables
- State
 - for context switches

} Snapshot of the state of the program in execution

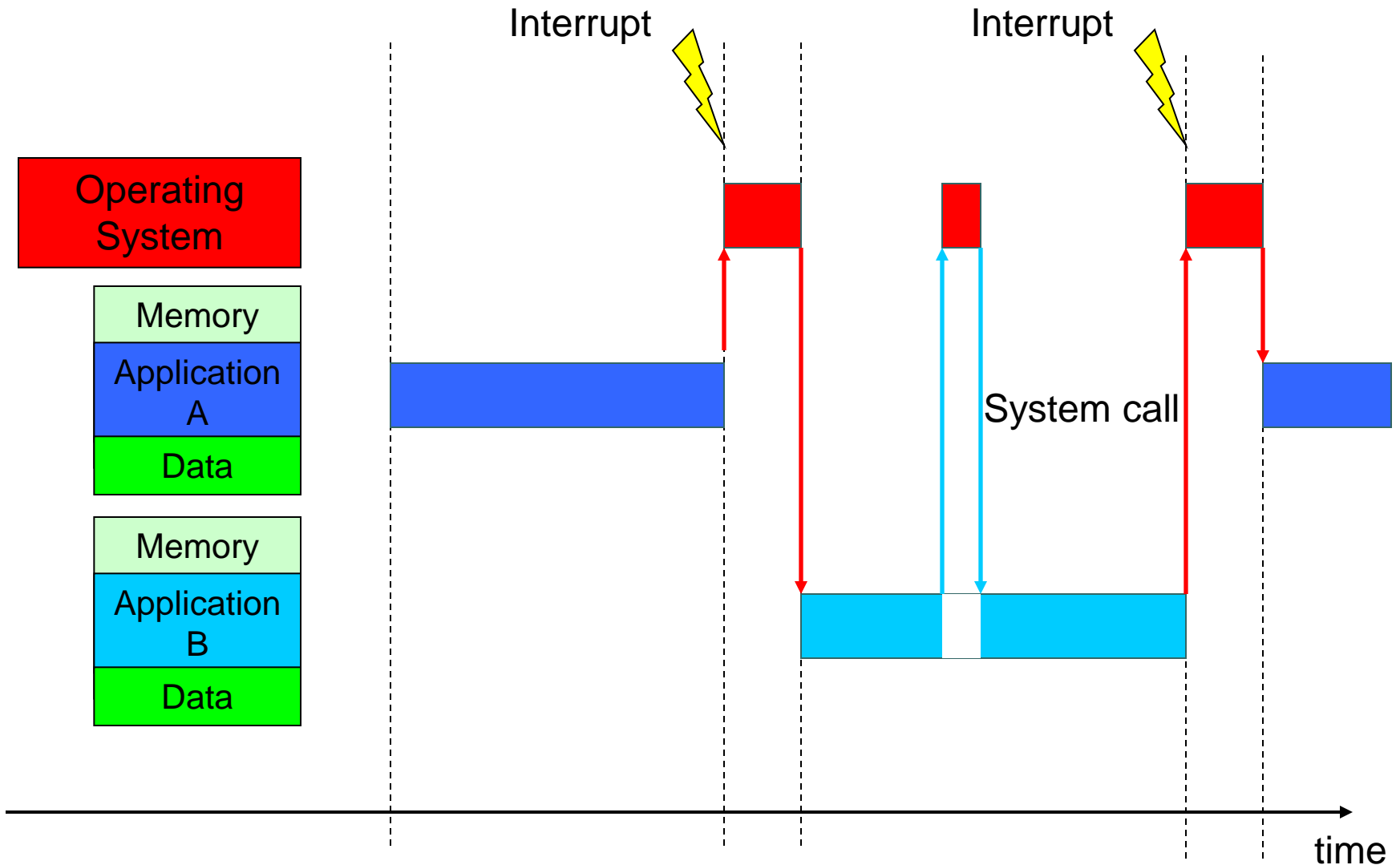
Sharing the processor



Sharing the processor



Sharing the processor



Context switch

Switching Process

- When switching to another process
 - Save state of old process
 - Load state of new process
- Reasons for switch
 - Interrupts
 - Blocking operations
 - I/O
 - Process synchronization
- Scheduling
 - Choosing the process to switch to

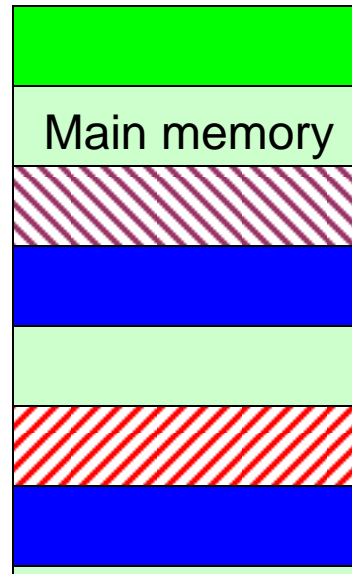
Virtual Memory

- The illusion of having
 - As much memory as are addressable
 - Probably more than the available physical memory.
 - 64-bit address => 16.8 million TB
 - All the memory by itself
- Allows the OS to move parts of processes on secondary memory
- Provides protection from other processes

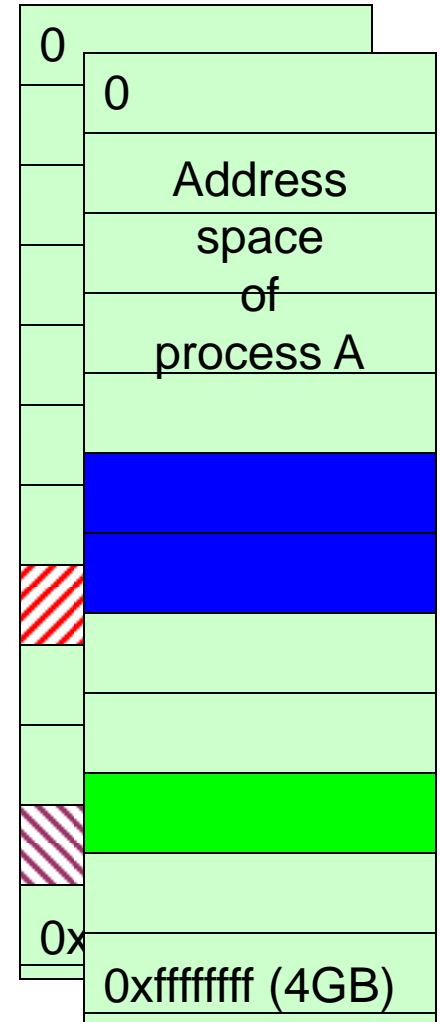
Paging

- Address space and main (physical) memory is divided into fixed-sized pages
- Each page may be located anywhere in main memory

Physical memory

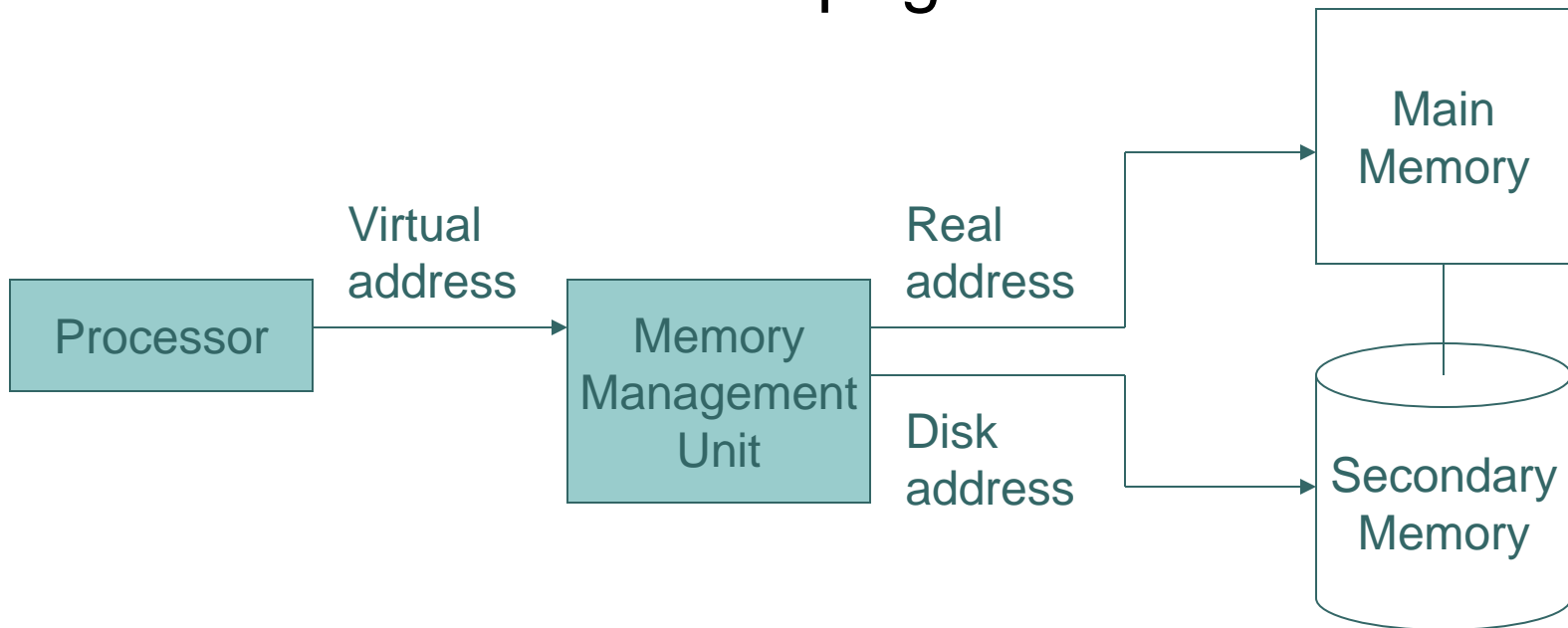


Virtual address spaces



Virtual memory addressing

- A virtual address is the combination of
 - A page number
 - An offset within the page



When physically memory filled

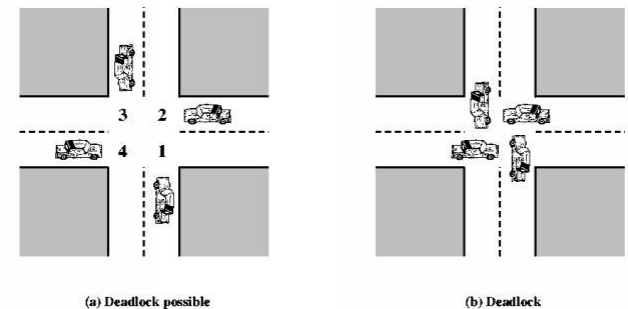
- Pages can be written to disk
 - They are "paged out"
 - Memory becomes available for other pages
 - Chosen pages should be seldomly used
 - If a process uses paged out memory
 - Needs to be read back into main memory
 - Probably ends up in another location

Competition for resources

- Processes use different resources
- Many resources cannot be shared at the same time between processes
 - Synchronization between processes are needed so it is used by one at a time
 - Locks are one tool to do this
- Problems can arise
 - Many processes want to use many resources at the same time
 - The order of acquiring them becomes important

Deadlock

- If a set of processes are all waiting on some other process in the set a deadlock has occurred.
- Example:
 - Two processes wants to transfer money between the same two accounts
 - Process A has got account 1
 - Process B has got account 2
 - Process A needs account 2 to do its transfer
 - Process B needs account 1 to do its transfer



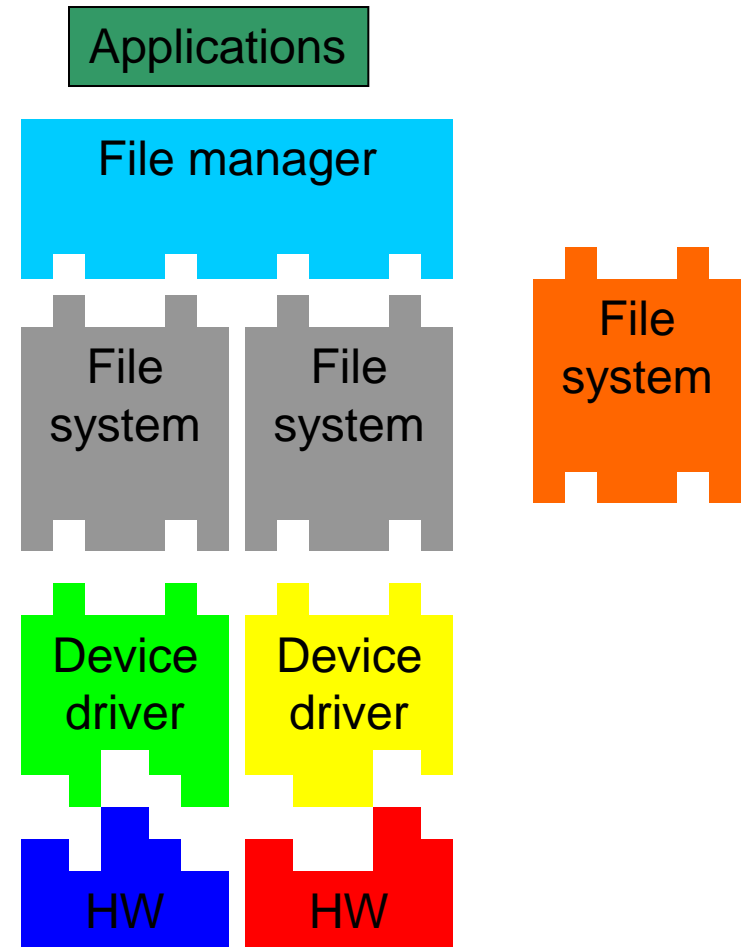
Deadlock:

Solution and Starvation

- Solve the deadlock: one process backs off
 - Release the resource
 - The other one can complete
 - Try again to acquire both resources
- Starvation can occur
 - If the same process backs off every time it might never finish
- Other solutions can avoid this

File systems

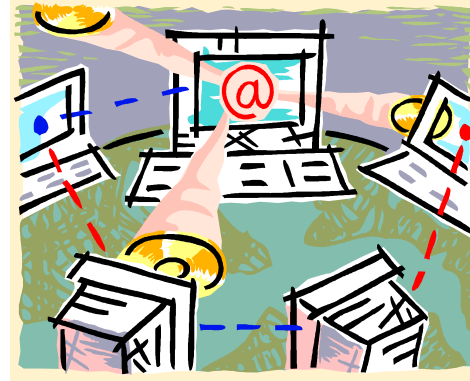
- An abstraction that provides
 - Long-term information storage in named files and directories
 - Allows hierarchical organization of data
 - Standard interface for applications
- The implementation is layered
 - Storage device
 - Device driver
 - File system implementation
 - OS file manager and application interface



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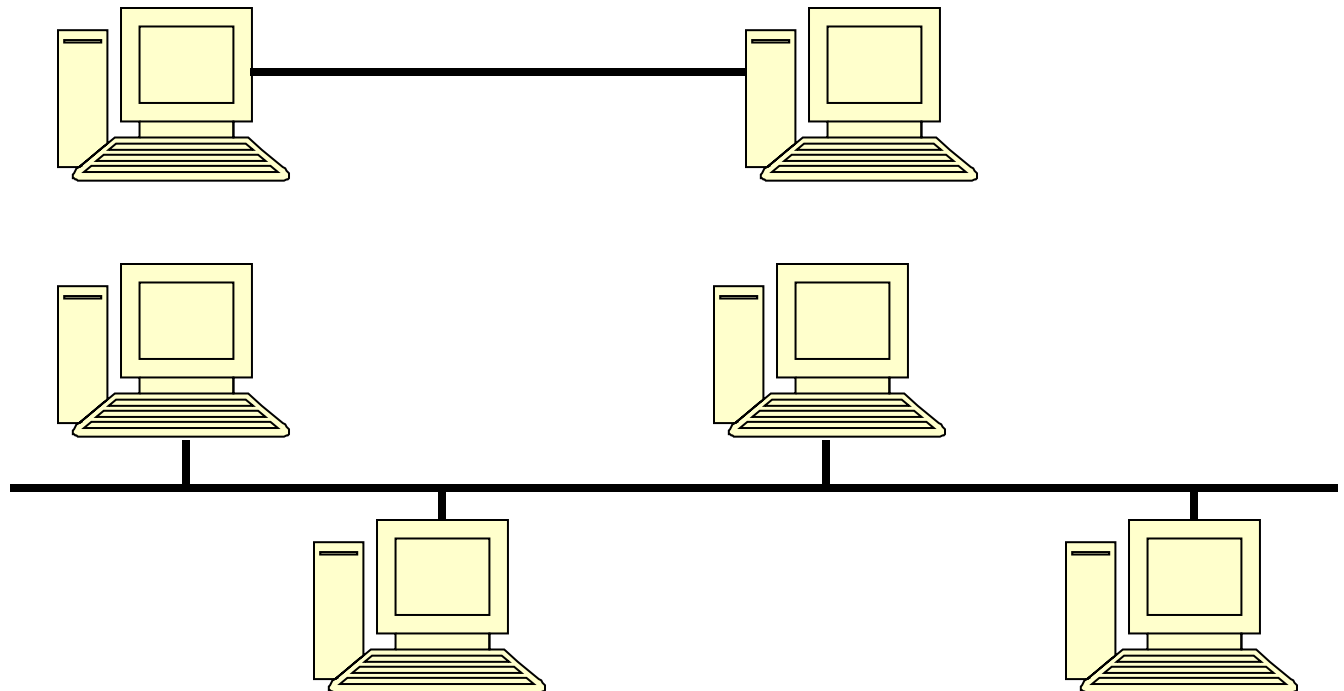
Networking



- Purpose
 - Allow applications (on different computers) to talk to each other

Networks - A bottom up view

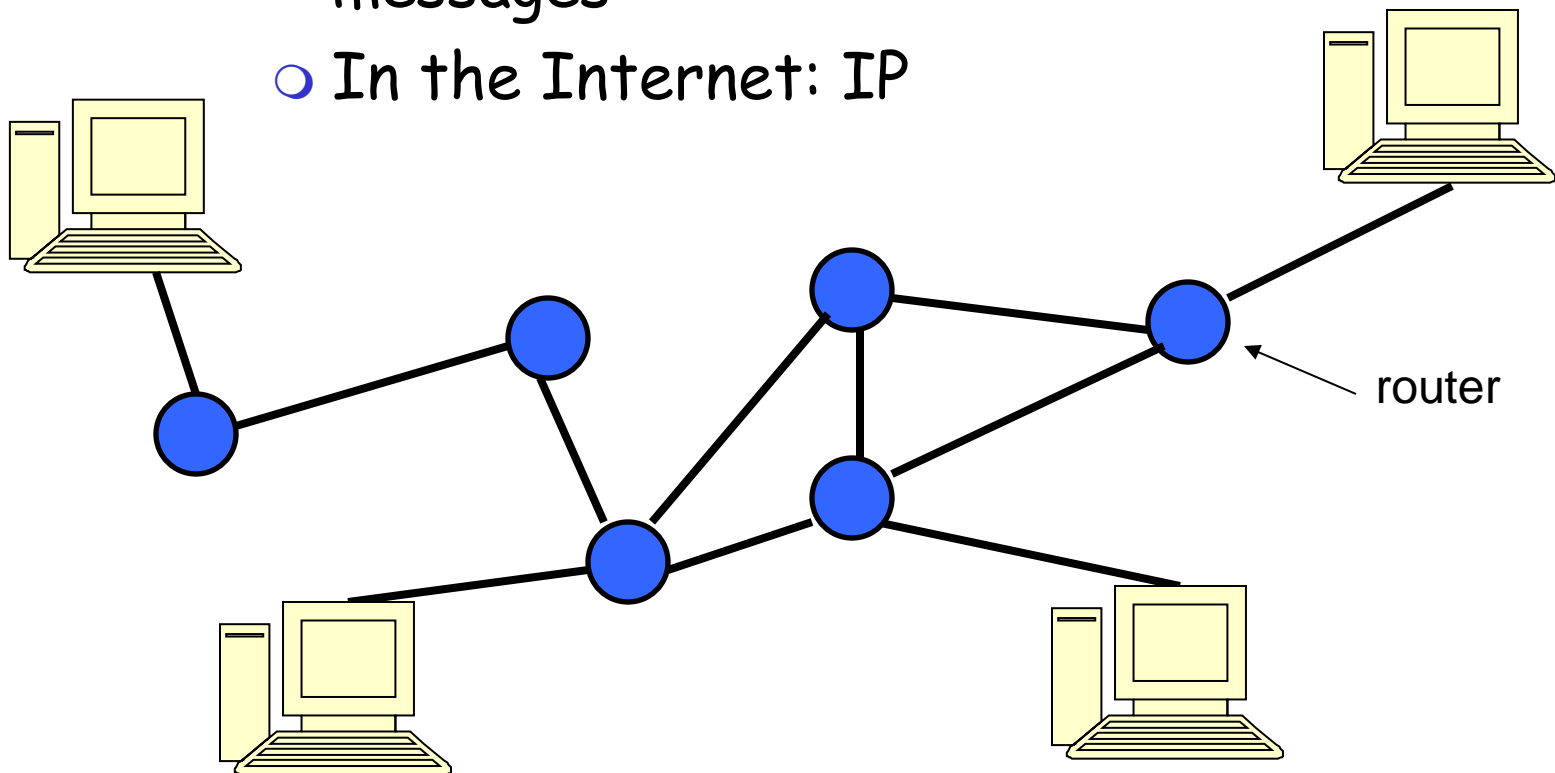
- ❑ The Link Layer - Computer to computer com.
 - Point-to-point
 - Shared medium (e.g. Ethernet)



Networks - A bottom up view

□ The Network Layer

- Host-to-host multihop routing of messages
- In the Internet: IP



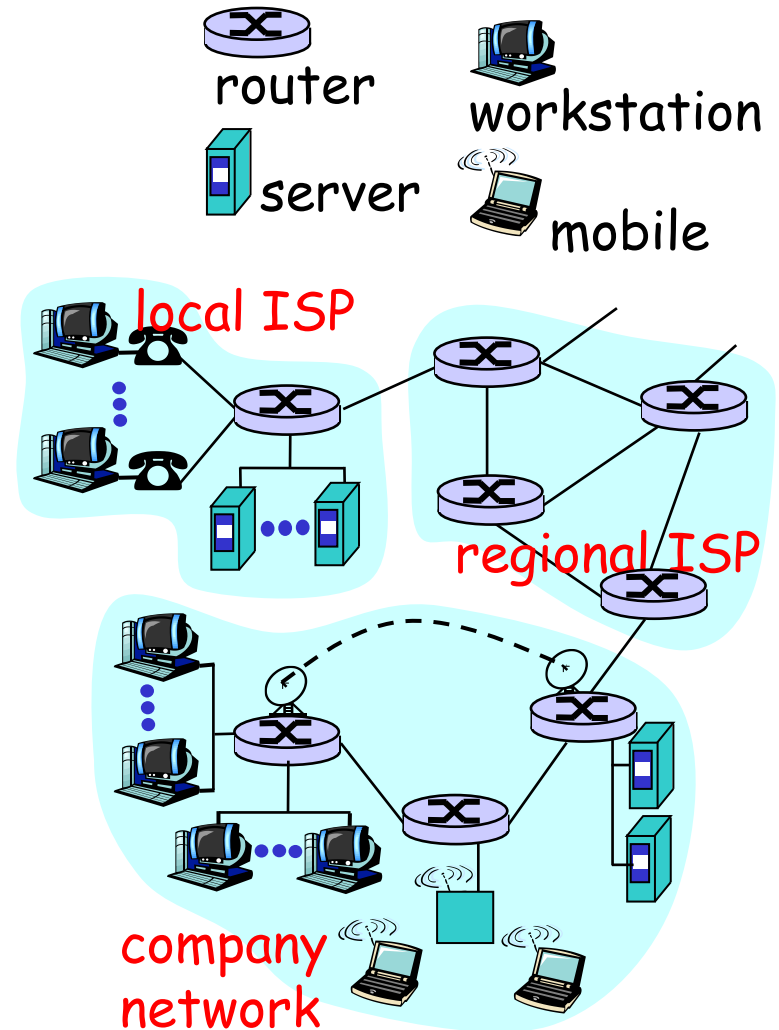
Networks - A bottom up view

□ The Transport Layer

- Application-to-application communication
- What kind of service?
 - Point-to-point or multicast?
 - Reliable (no messages are lost)?
 - Connection oriented?
- The Internet
 - TCP: connection oriented + reliable + point-to-point
 - UDP: connectionless + unreliable + point-to-point

What's the Internet: "nuts and bolts" view

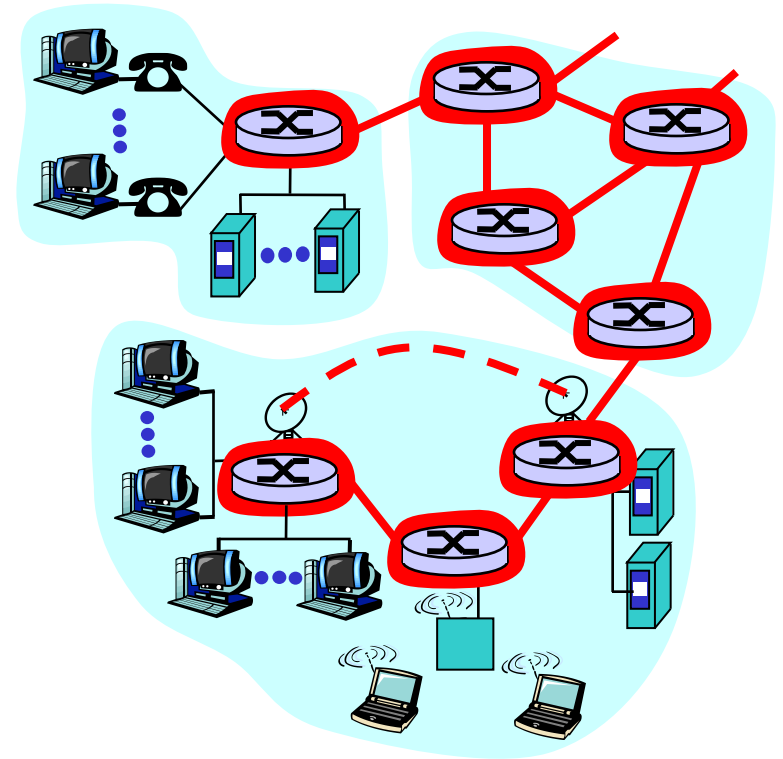
- ❑ millions of connected devices:
hosts, end-systems running
network apps
- ❑ *communication links*
- ❑ "network of networks"
 - connecting "devices" : *hubs, bridges, routers*
- ❑ *protocols*: control sending, receiving of msgs
- ❑ Internet standards
 - RFC: Request for comments
 - IETF: Internet Engineering Task Force



A closer look at network structure:

The network core:

- ❑ mesh of interconnected routers
- ❑ **fundamental question:** how is data transferred through net?
 - **circuit switching:** dedicated circuit per call: telephone net
 - **packet-switching:** data sent thru net in discrete "chunks"
 - hybrid form: **virtual circuits**



Network Core: Packet Switching

each end-end data stream
divided into *packets*

- ❑ user packets *share* network resources
- ❑ resources used *as needed*

store and forward:

- ❑ packets move one hop at a time
 - transmit over link
 - wait turn at next link

resource contention:

- ❑ aggregate resource demand can exceed amount available
- ❑ congestion: packets queue, wait for link use

Internet Addressing

- An Internet host is identified by
 - IP-address (IPv4)
 - A unique 32 bit id number
 - Hierarchical w.r.t. routing
 - Domain Name Service (DNS) name
 - Human readable name
 - Hierarchical w.r.t. country, organization etc.
 - Eg. zsh.chalmers.se
www.chalmers.se

IPv6

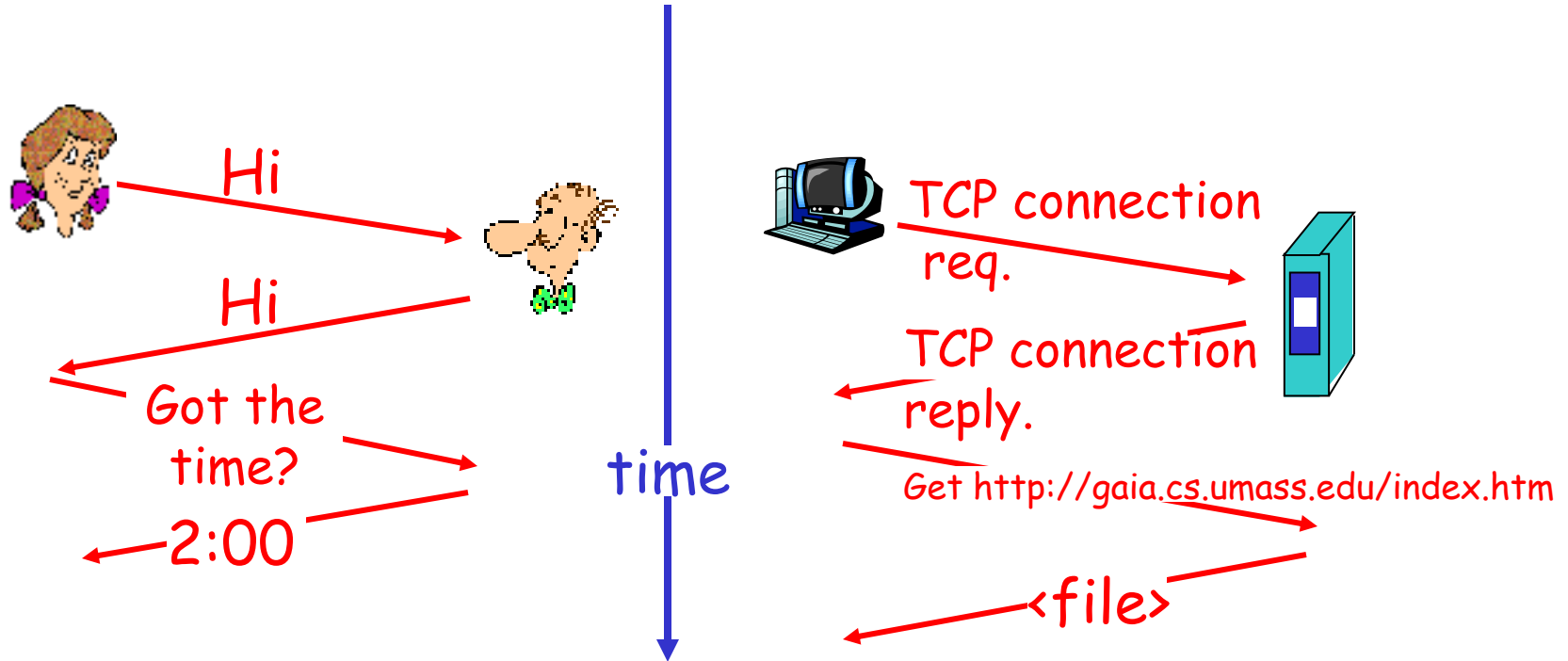
- “Normal” IP have 4 billion addresses
 - A lot of them are wasted
 - Chalmers alone have 65535 addresses
- We are running out of addresses
- Solution: IPv6
 - 128 bit addresses
 - 50 billion billion billion addresses / person
 - Allows solutions that is more hierarchical
 - We can waste address space freely

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What's a protocol?

a human protocol and a computer network protocol:



protocols define format, order of msgs sent and received among network entities and actions taken on msg transmission, receipt

Protocol "Layers"

Networks are complex!

- many "pieces":
 - hosts
 - routers
 - links of various media
 - applications
 - protocols
 - hardware, software

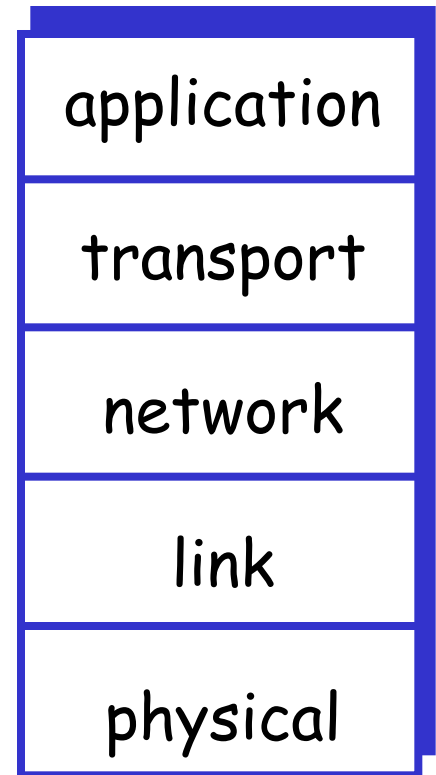
Question:

Is there any hope of *organizing*
structure of network?

Or at least our discussion of
networks

Internet protocol stack

- ❑ **application:** ftp, smtp, http, etc
- ❑ **transport:** tcp, udp, ...
- ❑ **network:** routing of datagrams from source to destination
 - ip, routing protocols
- ❑ **link:** data transfer between neighboring network elements
 - ppp, ethernet
- ❑ **physical:** bits "on the wire"



Why layering?

Dealing with complex systems:

- ❑ explicit structure allows identification, relationship of complex system's pieces
 - layered **reference model** for discussion
- ❑ **modularization eases maintenance**
 - change of implementation of layer's service transparent to rest of system
 - e.g., change in gate procedure doesn't affect rest of system

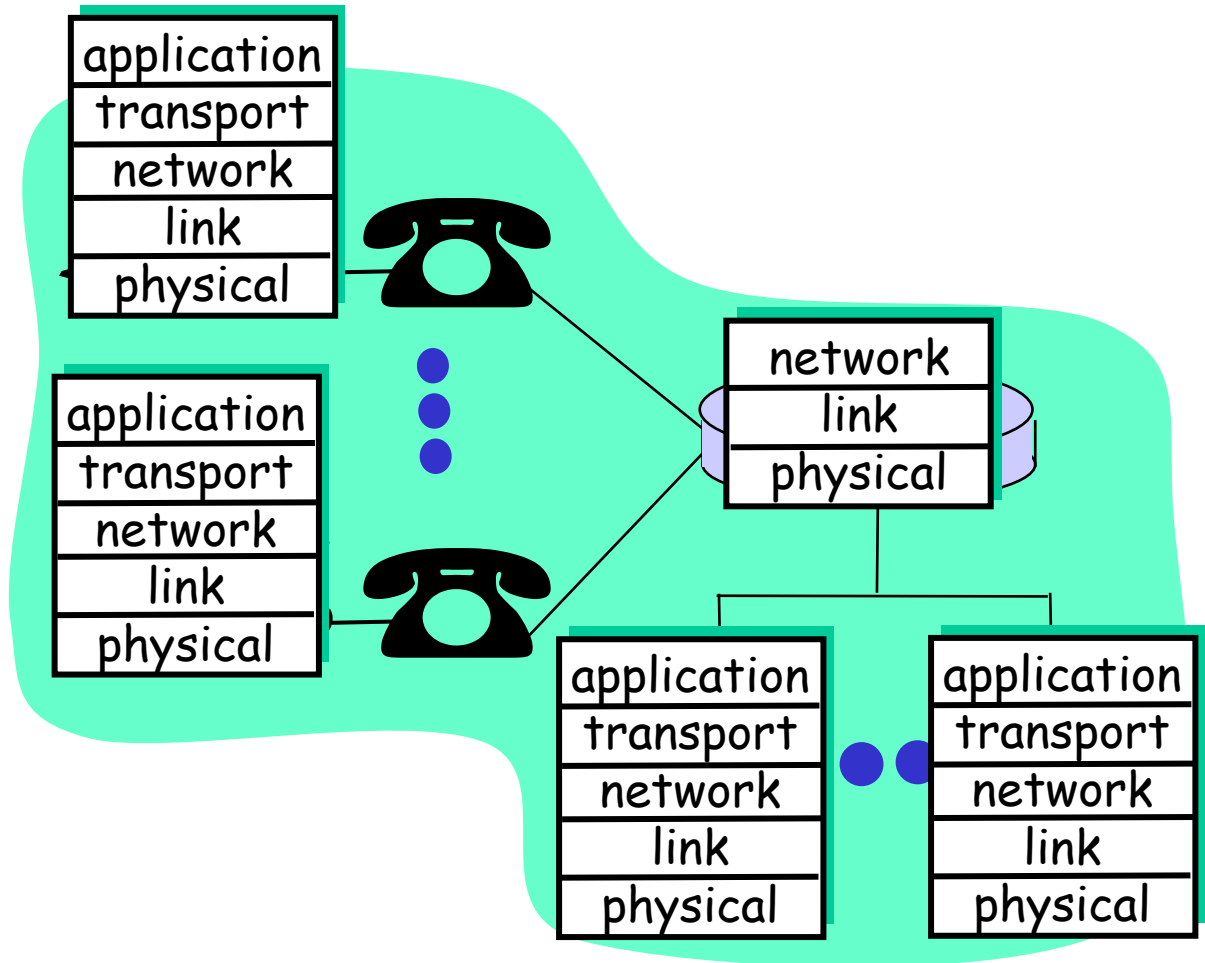
Terminology: Protocols, Interfaces

- ❑ Each **layer offers services** to the upper layers (shielding from the details how the services are implemented)
 - **service interface**: across layers in same host
- ❑ Layer n on a host carries a **conversation** with layer n on another host (data are not sent directly)
 - **host-to-host (aka peer-to-peer) interface**: defines messages exchanged with peer entity
- ❑ **Interfaces must be clean**
 - min info exchange
 - make it simple for protocol replacements
- ❑ **Network architecture** (set of layers, interfaces) vs **protocol stack** (protocol implementation)

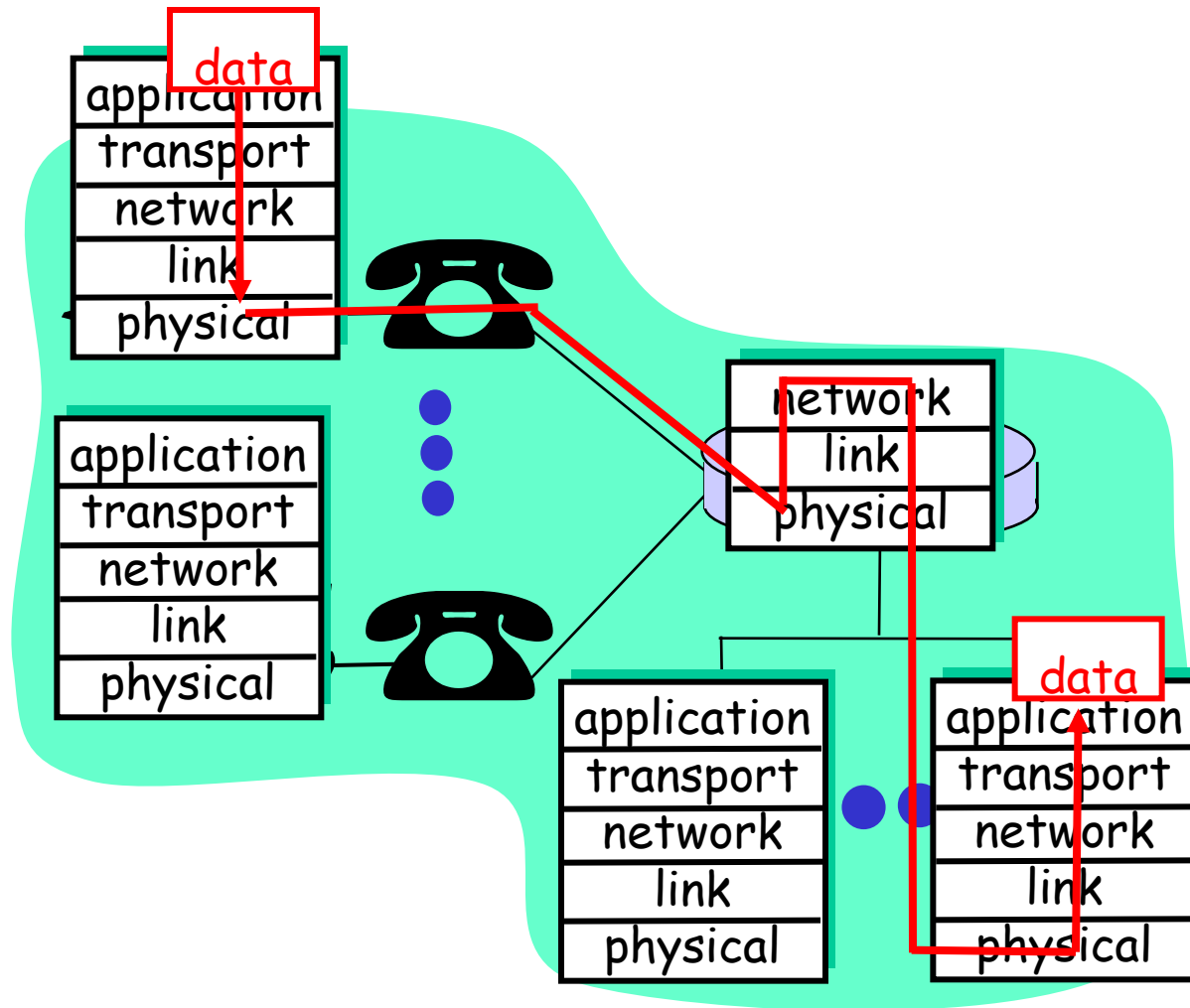
Layering: logical communication

Each layer:

- ❑ distributed
- ❑ "entities" implement layer functions at each node
- ❑ entities perform actions, exchange messages with peers

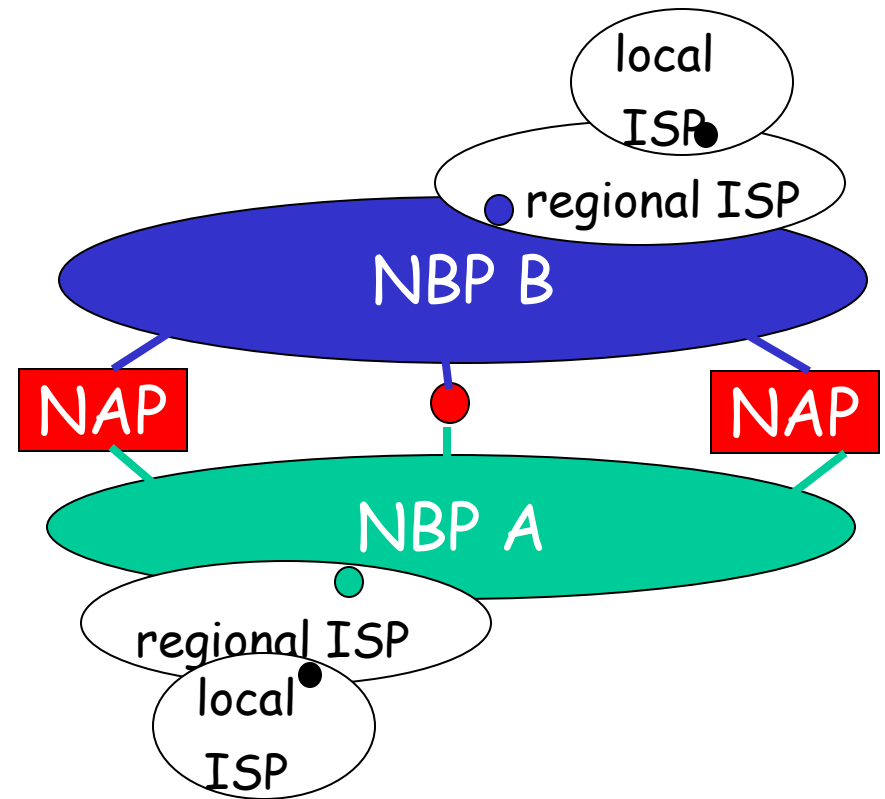


Layering: physical communication



Internet structure: network of networks

- ❑ roughly hierarchical
- ❑ national/international backbone providers (NBPs)
 - e.g. BBN/GTE, Sprint, AT&T, IBM, UUNet
 - interconnect (peer) with each other privately, or at public Network Access Point (NAPs: routers or (ATM) NWs of routers)
- ❑ regional ISPs
 - connect into NBPs
- ❑ local ISP, company
 - connect into regional ISPs



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What is network security?

Confidentiality: only sender, intended receiver should “understand” msg contents

- sender encrypts msg
- receiver decrypts msg

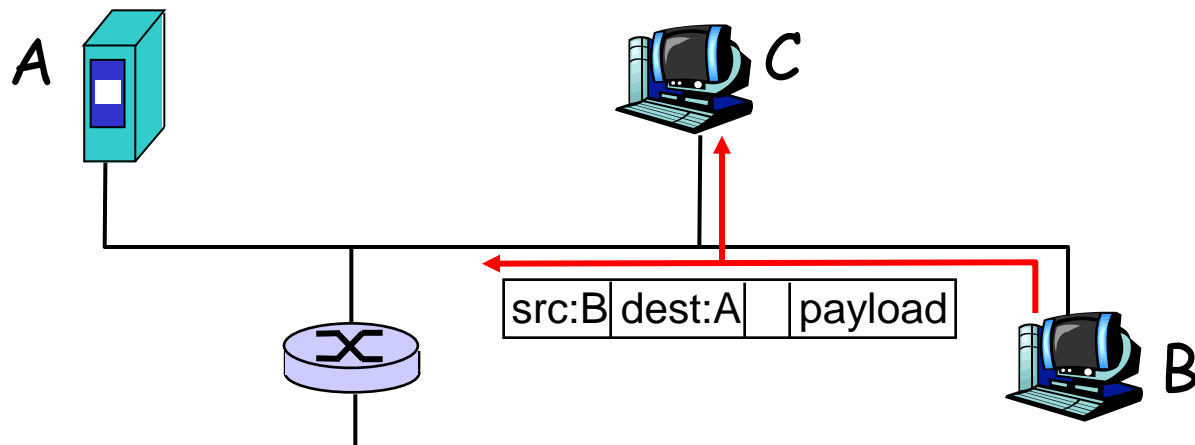
Integrity of Messages: sender, receiver want to ensure message not altered (in transit, or afterwards) without detection

Authentication: sender, receiver want to confirm identity of each other

Internet security threats

Packet sniffing:

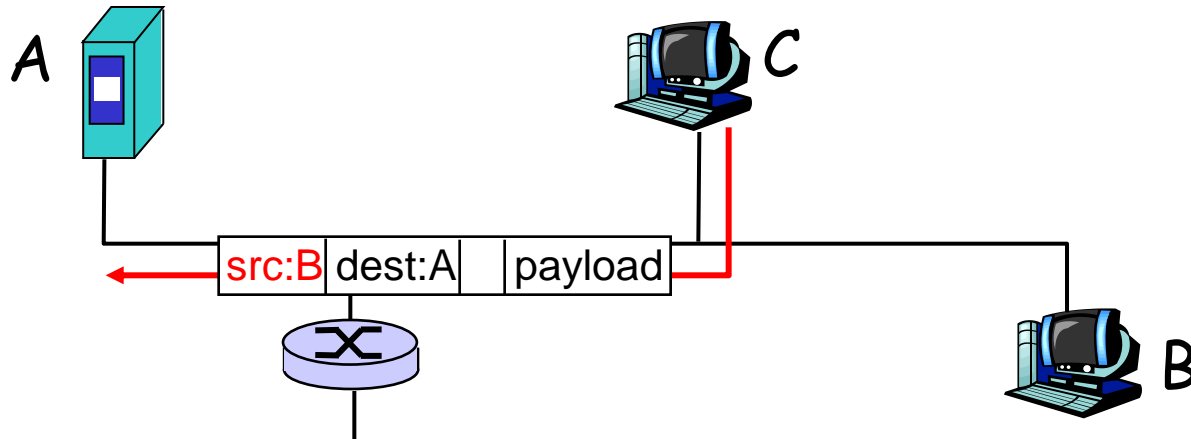
- broadcast media
- promiscuous NIC reads all packets passing by
- can **read all unencrypted data** (e.g. passwords)
- e.g.: C sniffs B's packets



Internet security threats

IP Spoofing:

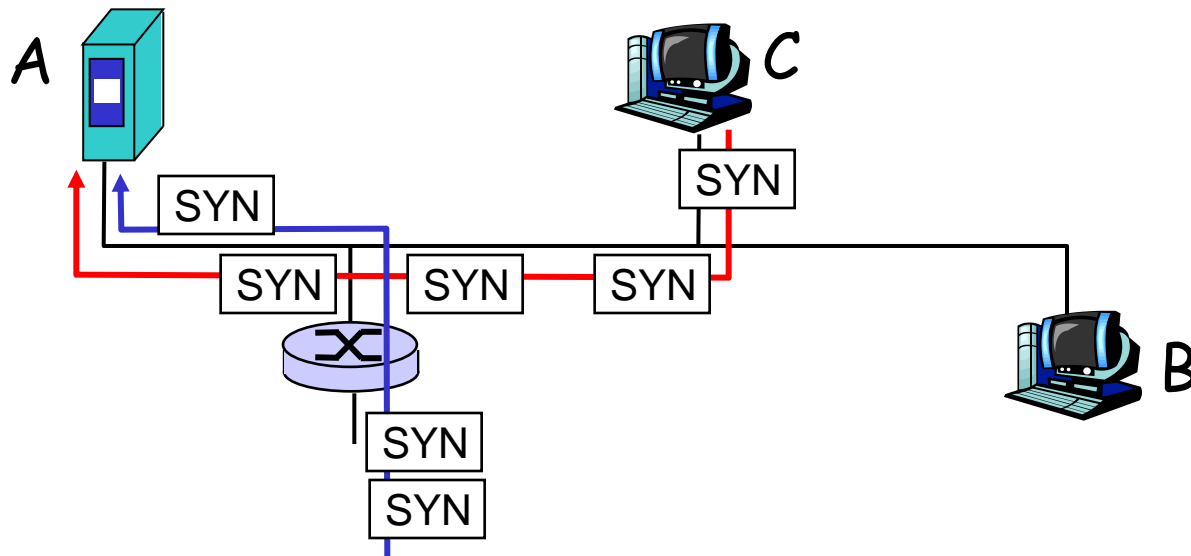
- can generate “raw” IP packets directly from application, putting **any value into IP source** address field
- receiver can't tell if source is spoofed
- e.g.: C pretends to be B



Internet security threats

Denial of service (DOS):

- flood of maliciously generated packets "swamp" receiver
- Distributed DOS (DDOS): multiple coordinated sources (or, rather, **spoofed packets**) swamp receiver
- e.g., C and remote host SYN-attack A



Encryption

- Symmetric
 - Encryption/Decryption with the same key
 - Key distribution a problem
- Public key encryption
 - Encryption with public key
 - Decryption only with private key
 - Key distribution still a problem

Viruses and Worms

- Virus
 - Malicious code fragment
 - Attaches itself to other programs on the host computer
- Worms
 - “Virus” that infect computers through the network.
 - The host has to run some vulnerable network service
- Trojan
 - Program with a hidden malicious agenda

Firewalls

firewall

isolates organization's internal net from larger Internet, allowing some packets to pass, blocking others.

Two firewall types:

- packet filter
- application gateways

To prevent denial of service attacks:

- SYN flooding: attacker establishes many bogus TCP connections. Attacked host alloc's TCP buffers for bogus connections, none left for "real" connections.

To prevent intruders from obtaining secret info.

- e.g., to monitor traffic going in/out from the network and discard sensitive information

Questions?

- Contact Information:

- Address:

- Pierre Kleberger
Computer Science & Engineering
Chalmers University of Technology

- Email:

- pierre.kleberger@chalmers.se

- Web:

- <http://www.chalmers.se/cse/EN/people/kleberger-pierre>