Operating Systems Security Some basics

Layers of a computer system

Applications

Services

Operating system (OS)

OS kernel

Hardware

 Where should the security of the system be placed?

 The security of a layer could normally be compromised by attacks from lower layers!

OS Protection Principles

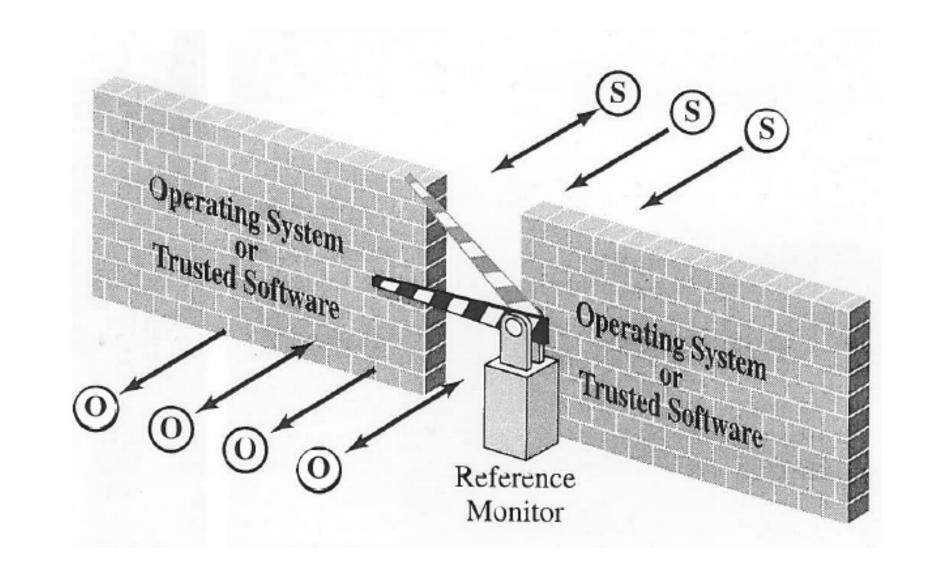
- The basis of OS protection is separation. The separation can be of four different kinds:
 - Physical: physical objects, such as CPU's, printers, etc.
 - Temporal: execution at different times
 - Logical: domains, each user gets the impression that she is "alone" in the system
 - Cryptographic: hiding data, so that other users cannot understand them
- "Computing is sharing and non-location security is separation"

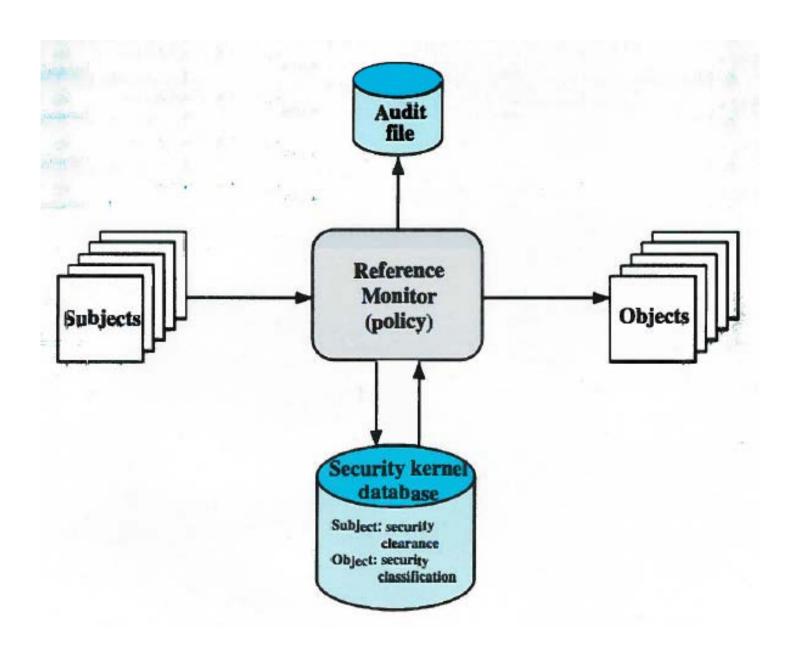
Protected objects

- In principle all objects in the OS need protection, but in particular those that are shareble, e.g.:
 - memory
 - I/O devices (disks, printers, tape drives, etc)
 - programs, procedures
 - data
 - hardware, such as
 - normal operating system mechanisms
 (e.g. file management logical, memory management physical)
 - bus control
 - interrupt control
 - status registers

Trusted operating system concepts

- There are a few basic concepts that are fundamental when dealing with trusted OS:
 - the kernel: is the part of the OS that performs the lowest-level functions
 - the security kernel: is responsible for enforcing the security mechanisms of the entire OS
 - the reference monitor (RM): is the part of the security kernel that controls access to objects
 - the trusted computing base (TCB): is everything in the trusted OS necessary to enforce the security policy





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Security policy and security model

- A **security policy** is a statement of the security we expect the system to enforce. The security can be expressed as a number of well-defined, consistent and implementable rules.
- A *security model* is a representation of the security policy for the OS.
- A formal security model is a mathematical description (formalisation) of the rules of the security policy. It could be used for formal proofs of security.

Development of a secure OS

- The development of secure OS can be made in six steps:
 - analyze of the system
 - choose/define a security policy
 - choose/create a security model (based on the policy)
 - choose implementation method
 - make a (conceptual) design
 - verify the correctness of the design
 - make an implementation
 - verify the implementation (?)
- There are feed-back loops between all of the above steps. Errors may occur in all above steps.

The two types of security costs

Make a trade-off between costs!

