# CRDTS Data Types for EC Systems

## Distributed Systems Parallel Systems

# Problem?

### Eventual Consistency

Eventual consistency is a consistency model used in distributed computing that informally guarantees that, if no new updates are made to a given data item, eventually all accesses to that item will return the last updated value.

--Wikipedia



# Distributed

### Distributed System

A distributed system is one in which the failure of a computer you didn't even know existed can render your own computer unusable

-Leslie Lamport



## Scale Up

\$\$\$Big Iron (still fails)



### Scale Out

Commodity Servers
CDNs, App servers
Expertise



# 

# Tolerance

## 

# Latency

### Low Latency

Amazon found every 100ms of latency cost them 1% in sales.



## Low Latency

Google found an extra 0.5 seconds in search page generation time dropped traffic by 20%.



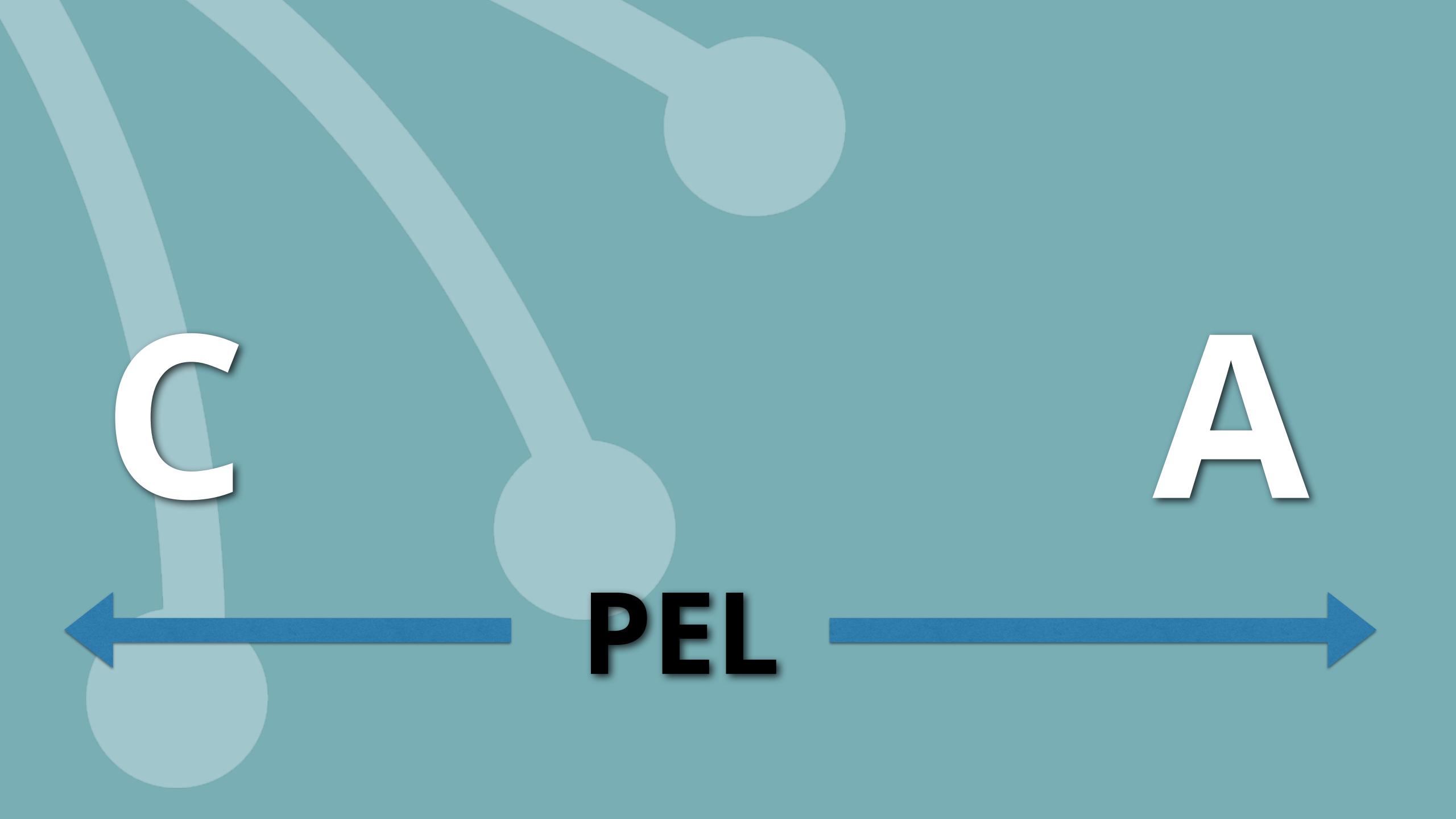
# Trade Off

# 

## http://aphyr.com/posts/288-the-network-is-reliable

# 





### Causal RYOW Session Monotonic Read



### Pick Your Own

Replicated Data Consistency Explained Through Baseball

http://research.microsoft.com/apps/pubs/default.aspx?id=157411 (Doug Terry)



# Who Pays?

# Developers But how?

### Google F1

"We have a lot of experience with eventual consistency systems at Google."

"We find developers spend a significant fraction of their time building extremely complex and error-prone mechanisms to cope with eventual consistency"



#### Google F1

"Designing applications to cope with concurrency anomalies in their data is very error-prone, time-consuming, and ultimately not worth the performance gains."



## Riak Overview

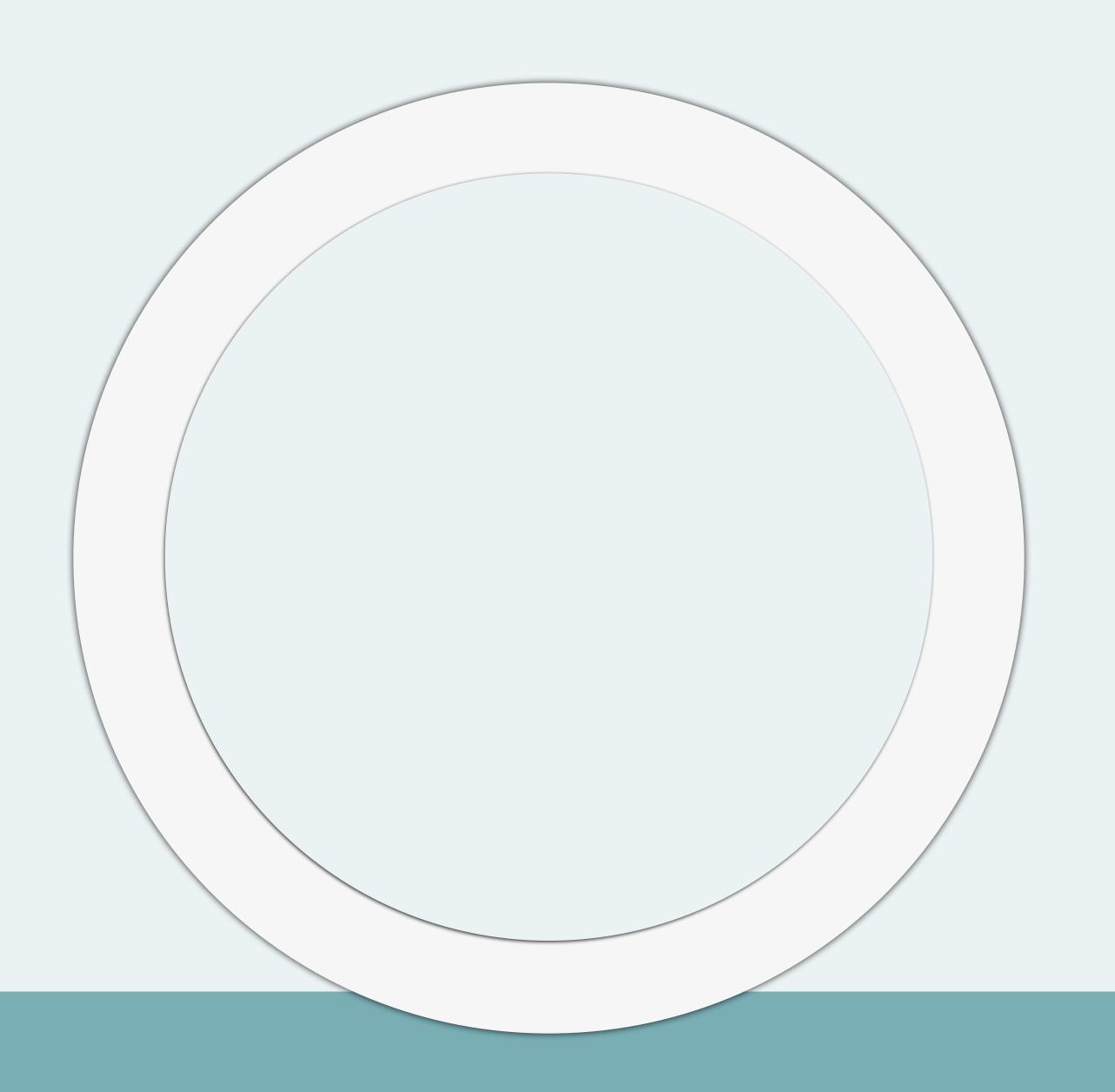
## Riak Overview

Erlang implementation of Dynamo

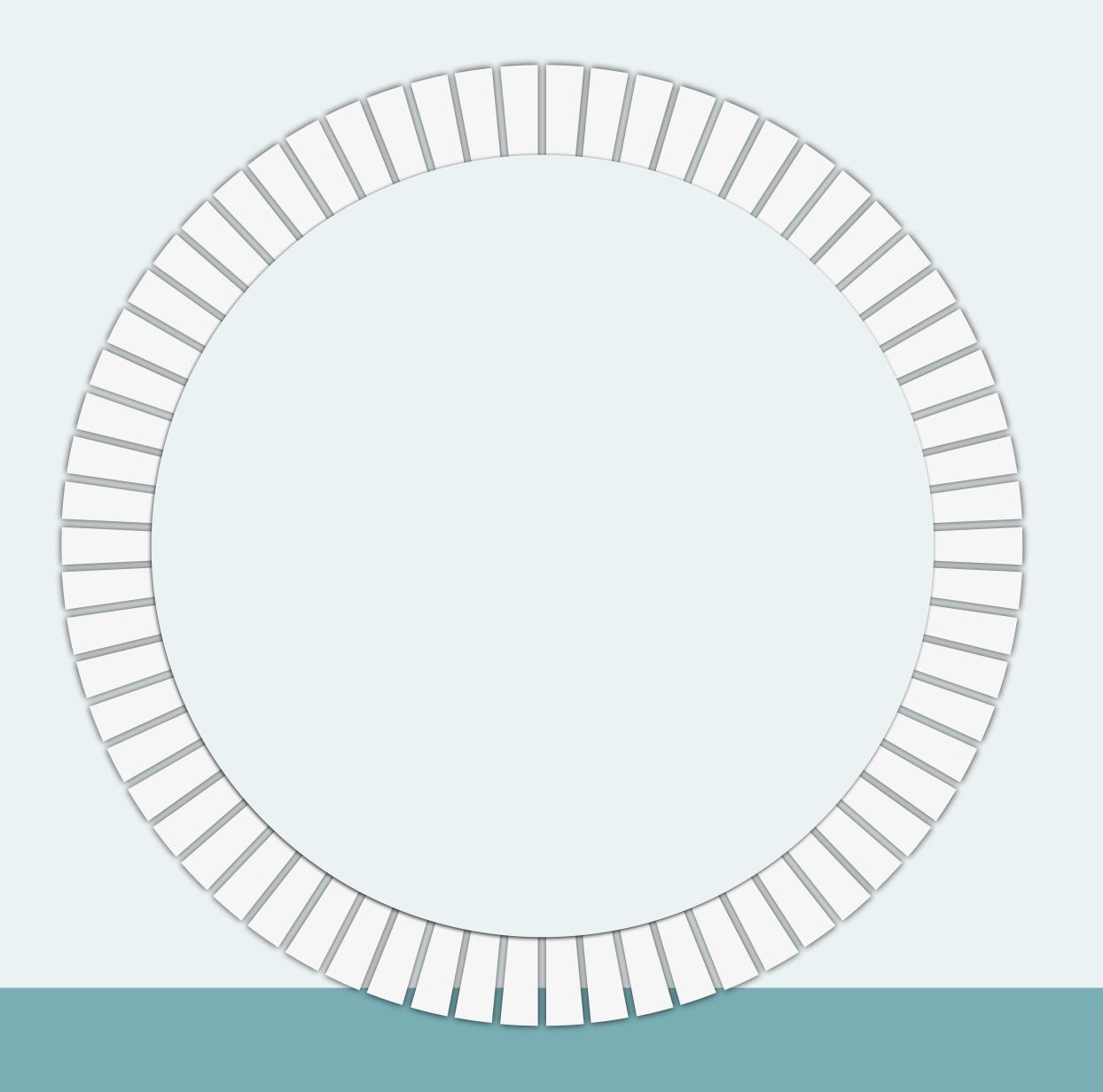
## {"key": "value"}



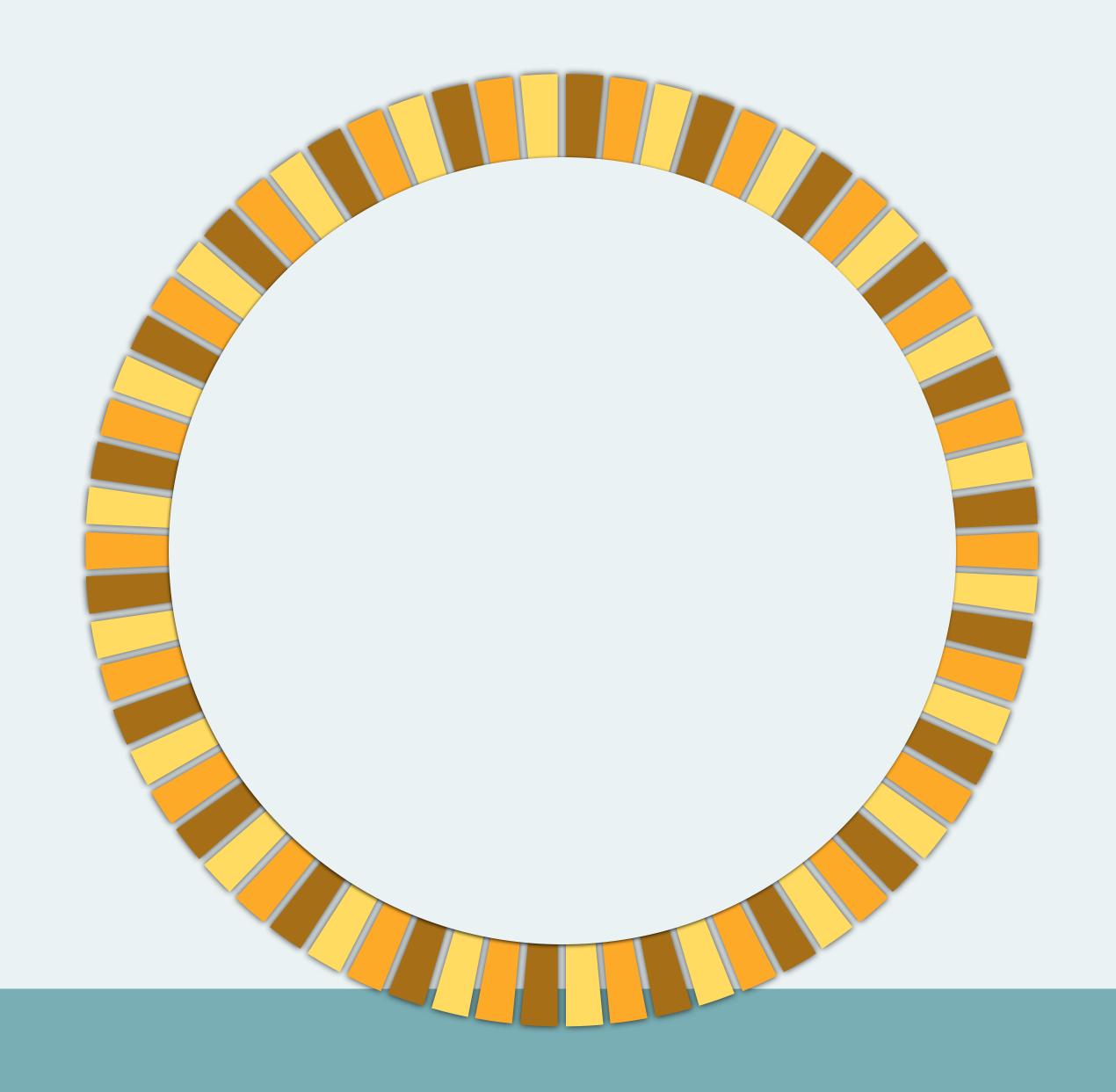
# RIAK OVERVIEW Consistent Hashing













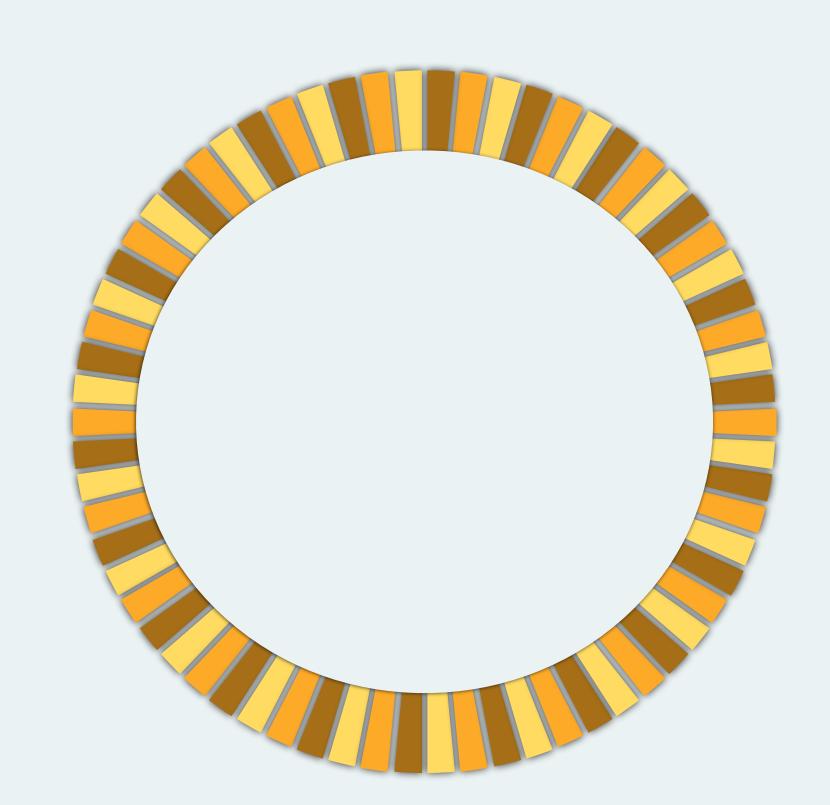
# Riak Overview Dynamic Membership

# RIAK OVERVIEW Replication factor

Replica

Replica

Replica



### Availability

Any non-failing node can respond to any request

--Gilbert & Lynch



# Riak Overview Two Writes: {Writer, Value, Time}

[{a, v1, a1}] [{b, v2, b1}] [{a, v1, a1}]



# Riak Overview Last Writer Wins Allow Mult

## Riak Overview

Last Writer Wins

[{b, v1, t2}]

[{b, v1, t2}]

[{b, v1, t2}]

## http://aphyr.com/posts/ 299-the-trouble-withtimestamps

### Riak Overview Allow Mult

[{a, v1, a1}, {b, v2, b1}]

[{a, v1, a1}, {b, v2, b1}]

[{a, v1, a1}, {b, v2, b1}]

# User specified IVIERSE

# Semantic Resolution

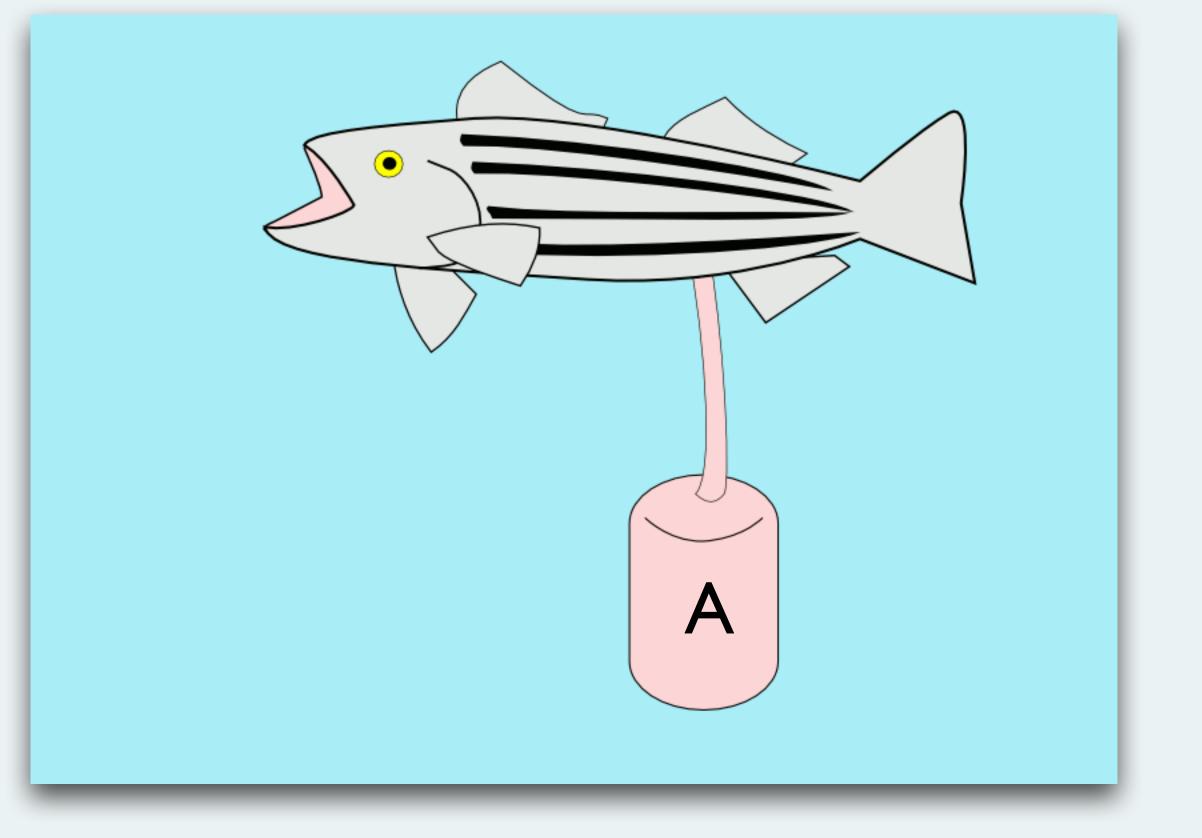


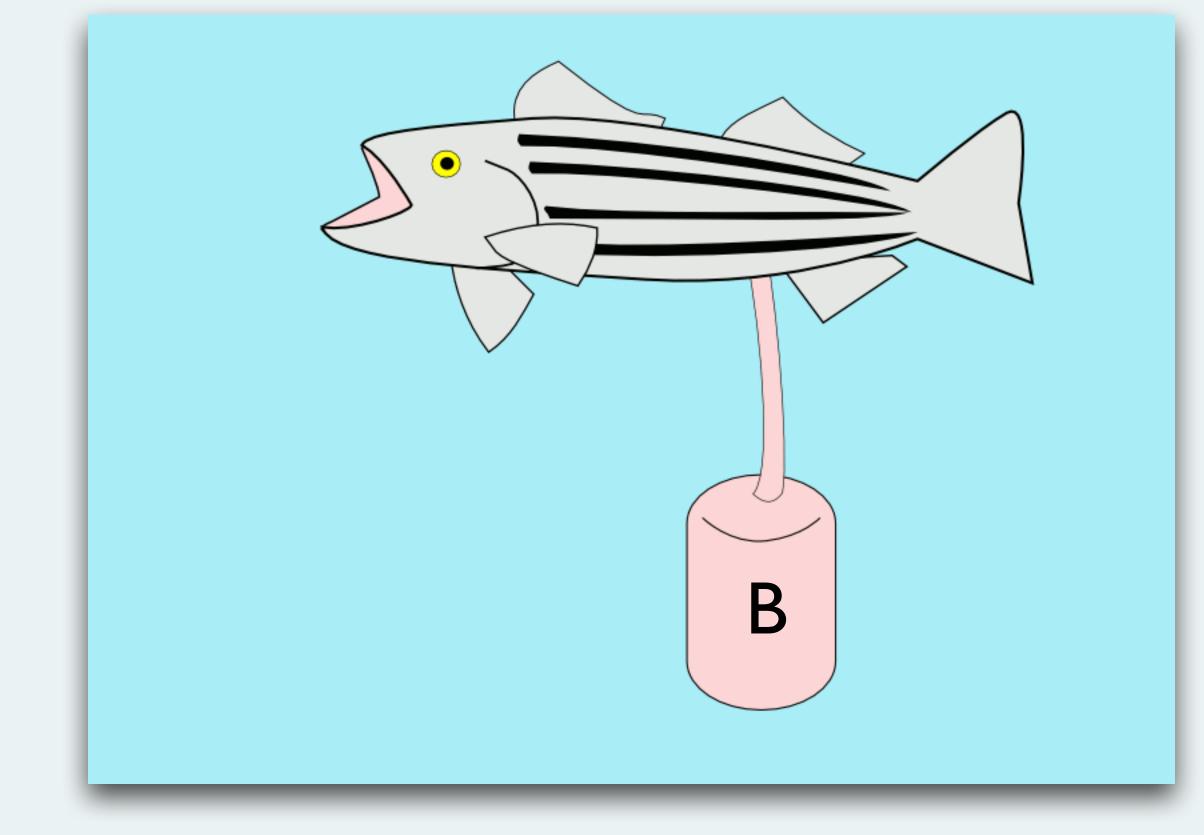
```
if (result.hasConflicts()) {
    // TODO: What should we do???
```





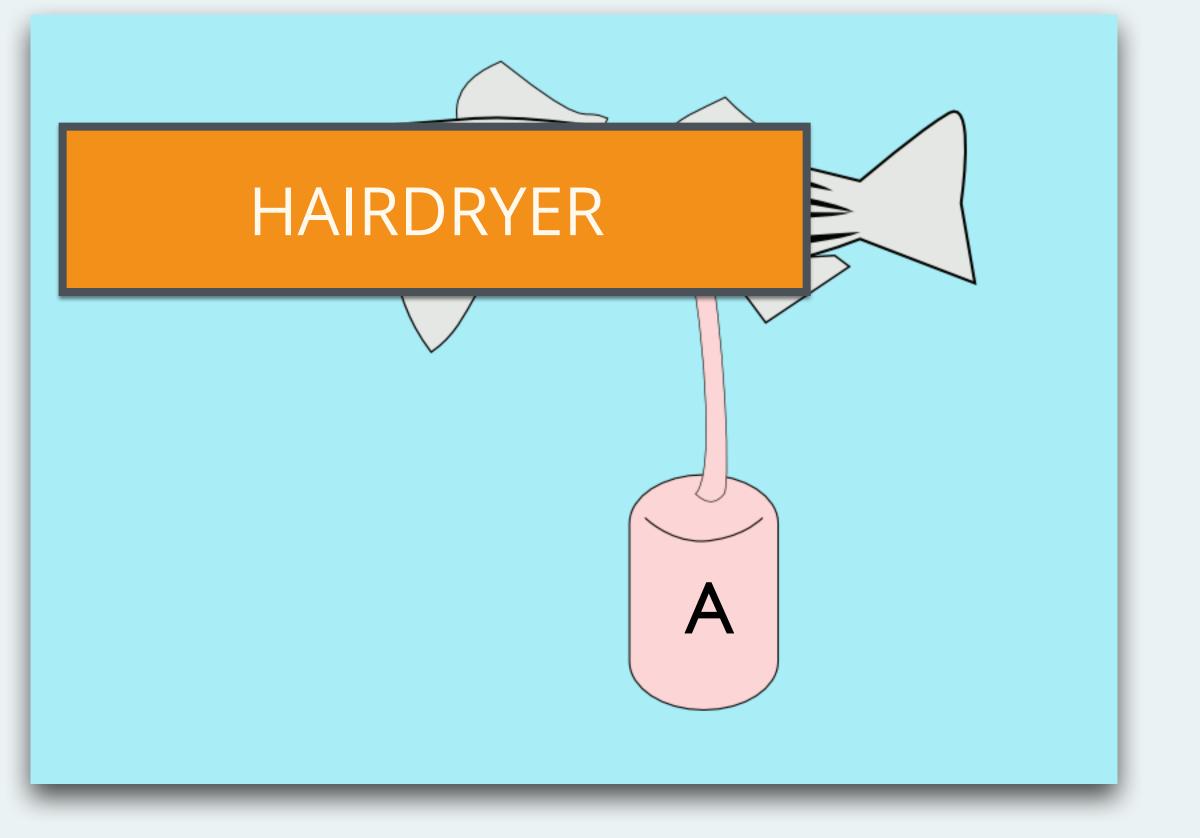
# Dynamo The Shopping Cart

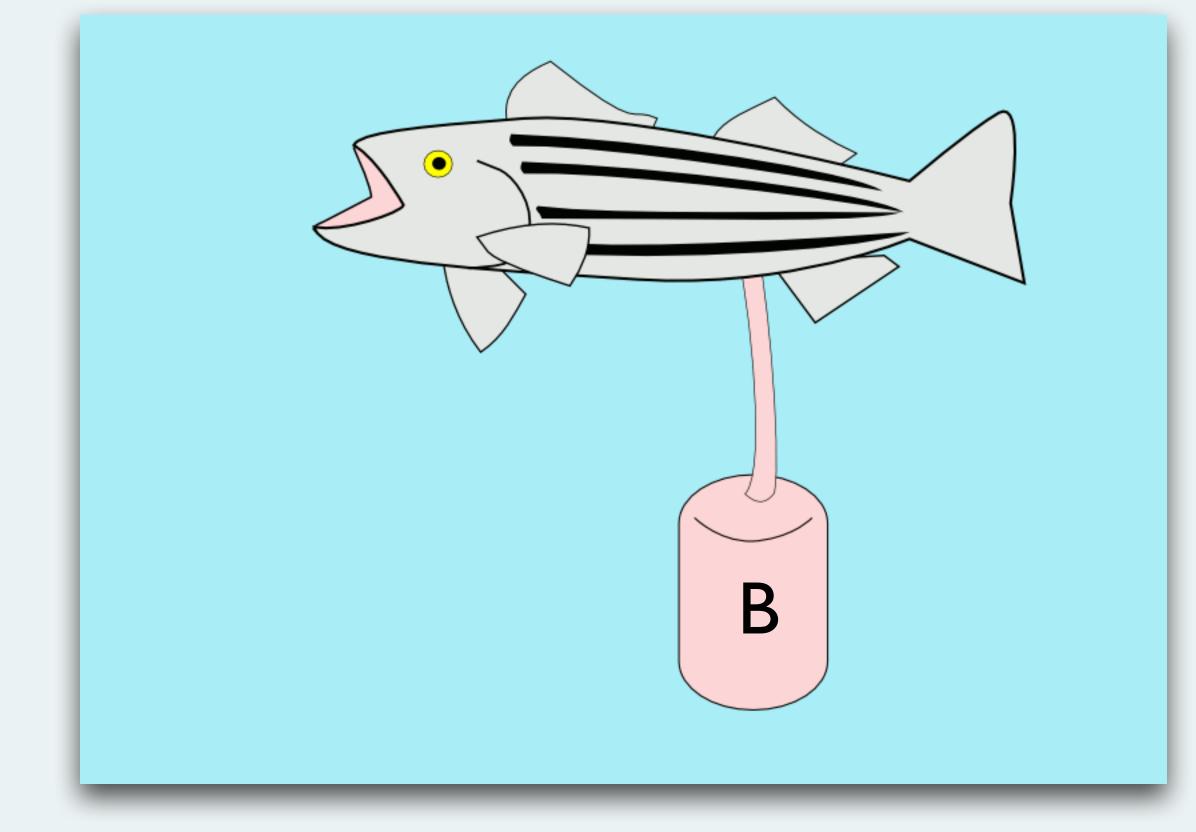




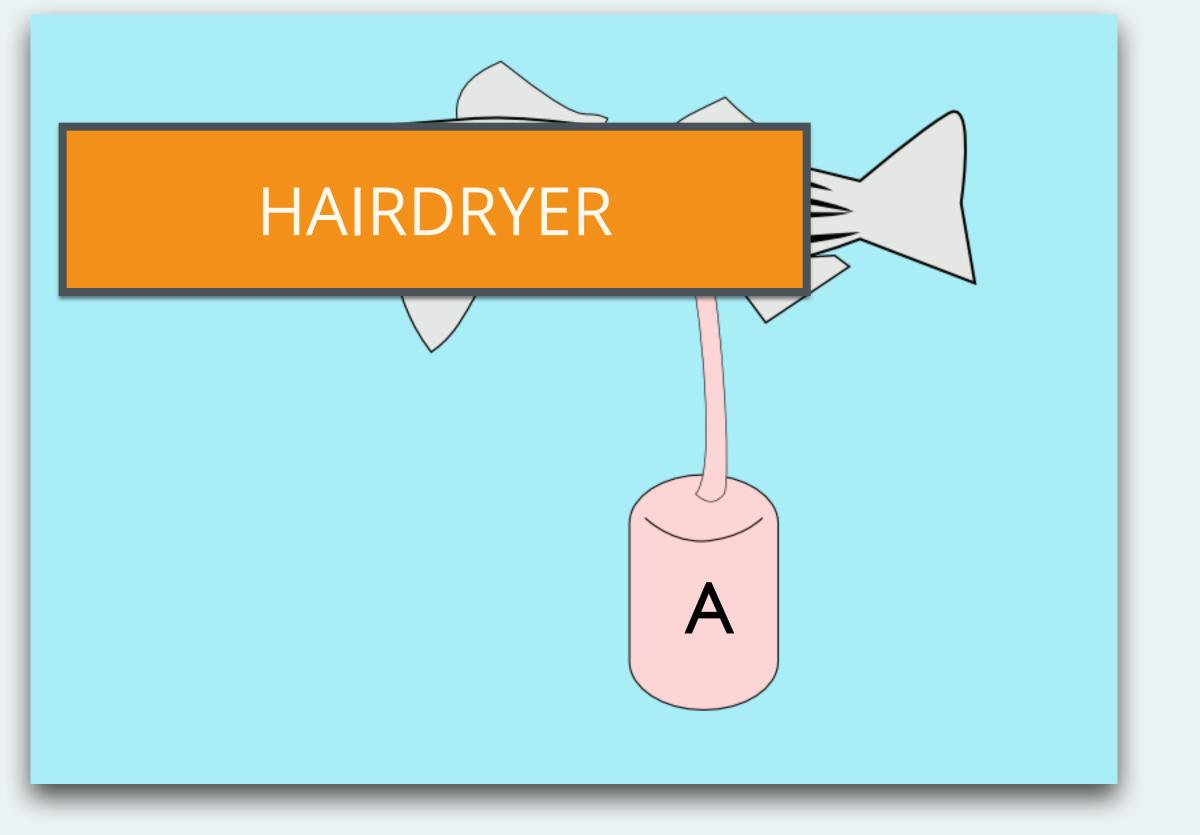
HAIRDRYER

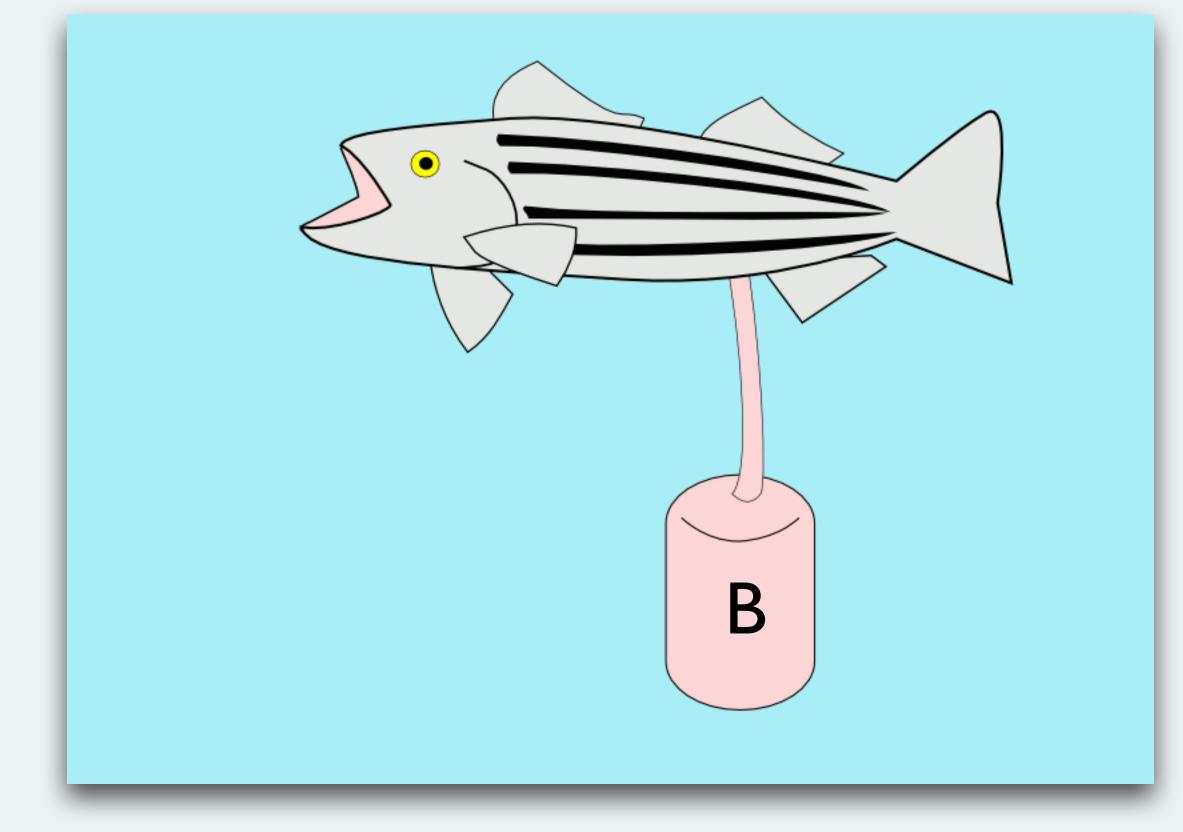






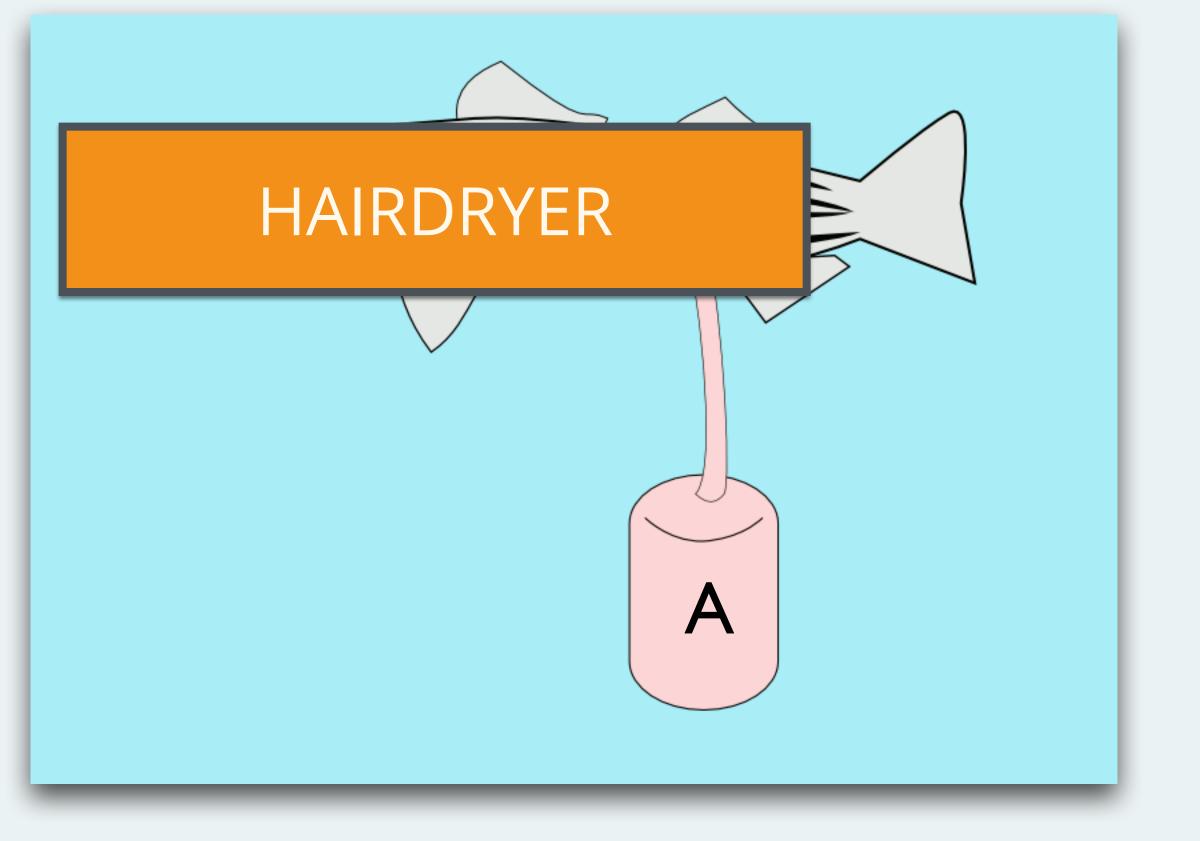


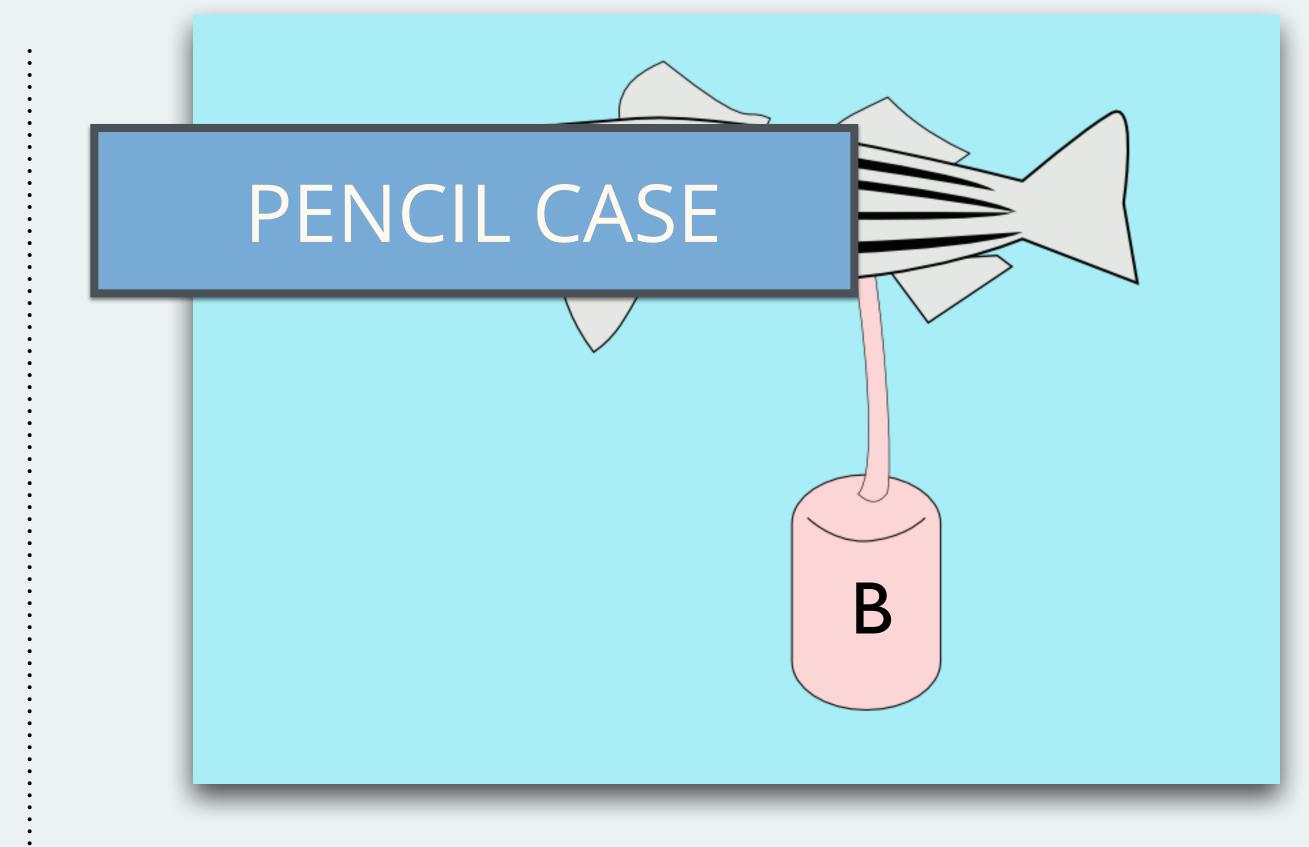




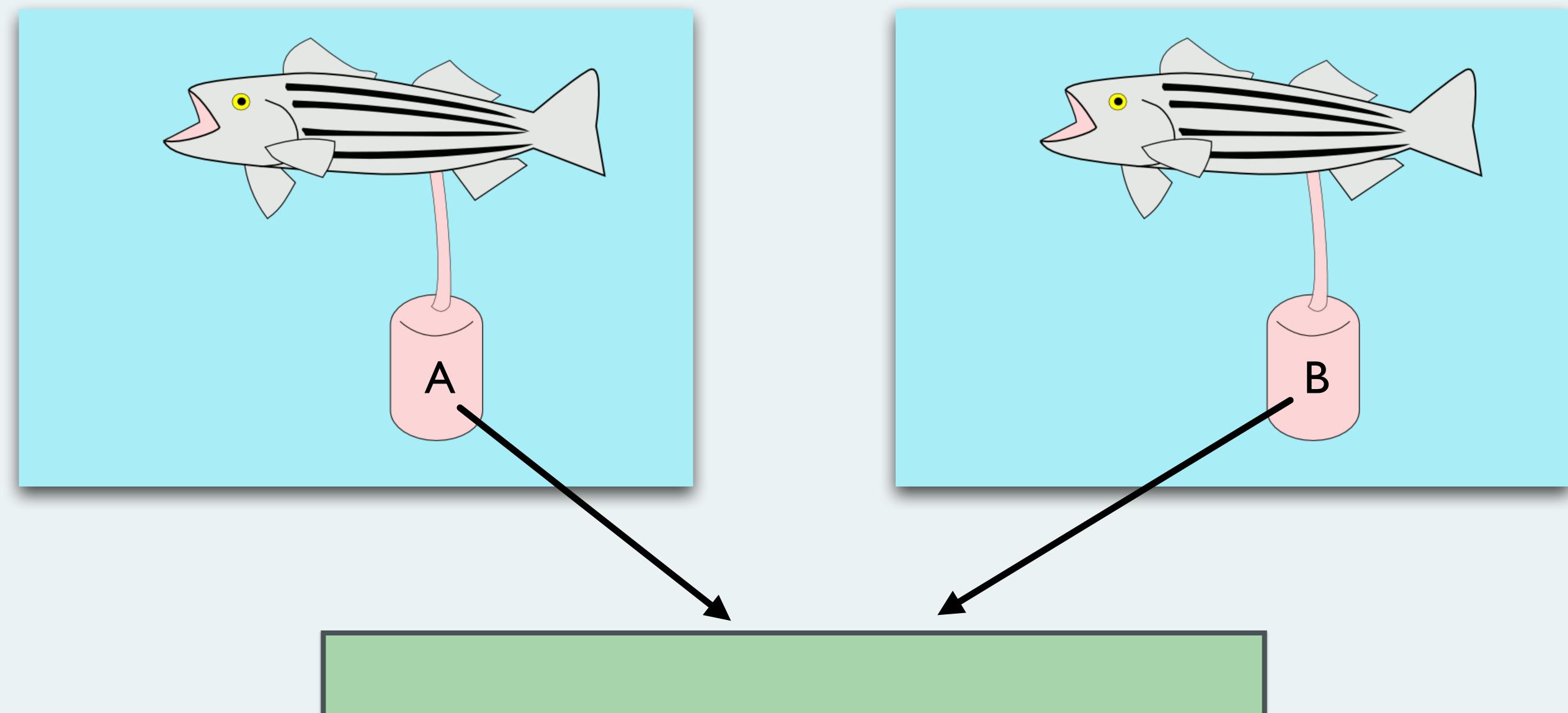
PENCIL CASE











[HAIRDRYER], [PENCIL CASE]



## Set Union of Values Simples, right?

# IVA Erge Deterministic

## Deterministic Idempotent

### Deterministic Idempotent Associative

## Merge

Deterministic Idempotent Associative Commutative

## Removes?

Set Union?
"Anomaly"
Reappear

#### Absence

How can you tell if X is missing from A but present in B because A hasn't yet seen the addition, or if A has removed it already?



# Complexity

## AG HOC

# 

# CRDTS Convergent Replicated Data Types

# CRDIS Commutative Replicated Data Types

## CRDIS

Conflict Free Data Structures

# Theory





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#### A comprehensive study of Convergent and Commutative Replicated Data Types

Marc Shapiro, INRIA & LIP6, Paris, France
Nuno Preguiça, CITI, Universidade Nova de Lisboa, Portugal
Carlos Baquero, Universidade do Minho, Portugal

Marek Zawirski, INRIA & UPMC, Paris, France







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#### An Optimized Conflict-free Replicated Set

Annette Bieniusa, INRIA & UPMC, Paris, France
Marek Zawirski, INRIA & UPMC, Paris, France
Nuno Preguiça, CITI, Universidade Nova de Lisboa, Portugal
Marc Shapiro, INRIA & LIP6, Paris, France
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Valter Balegas, CITI, Universidade Nova de Lisboa, Portugal

Sérgio Duarte CITI, Universidade Nova de Lisboa, Portugal



#### **Dotted Version Vectors: Logical Clocks for Optimistic Replication**

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#### **Abstract**

In cloud computing environments, a large number of users access data stored in highly available storage systems. To provide good performance to geographically disperse users and allow operation even in the presence of failures or network partitions, these systems often rely on optimistic replication solutions that guarantee only eventual consistency. In this scenario, it is important to be able to accurately and efficiently

The mentioned systems follow a design where the data store is always writable. A consequence is that replicas of the same data item are allowed to diverge, and this divergence should later be repaired. Accurate tracking of concurrent data updates can be achieved by a careful use of well established causality tracking mechanisms [5], [6], [7], [8]. In particular, for data storage systems, version vectors [6] enables the system to compare any pair of replica versions and detect if





This project is funded by the European Union, 7th Research Framework Programme, ICT call 10, grant agreement n°609551.



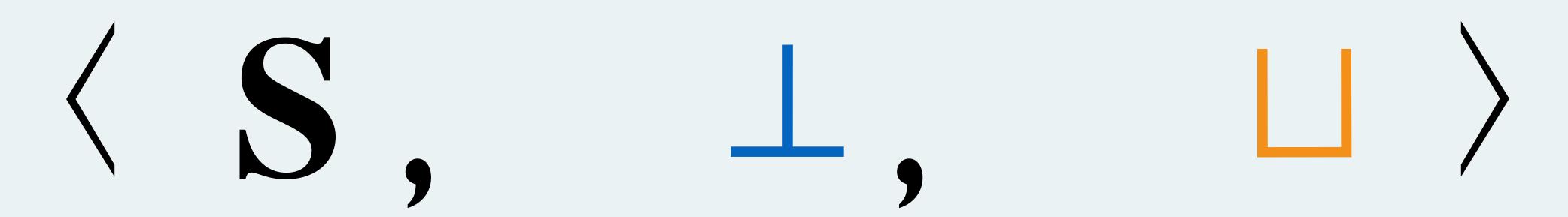


### Join Semi-lattice



#### Join Semi-lattice

Partially ordered set; Bottom; least upper bound





#### Join Semi-lattice

Associativity:  $(X \cup Y) \cup Z = X \cup (Y \cup Z)$ 



**Commutativity:** X⊔Y = Y⊔X



Idempotent: X\(\triangle X\)

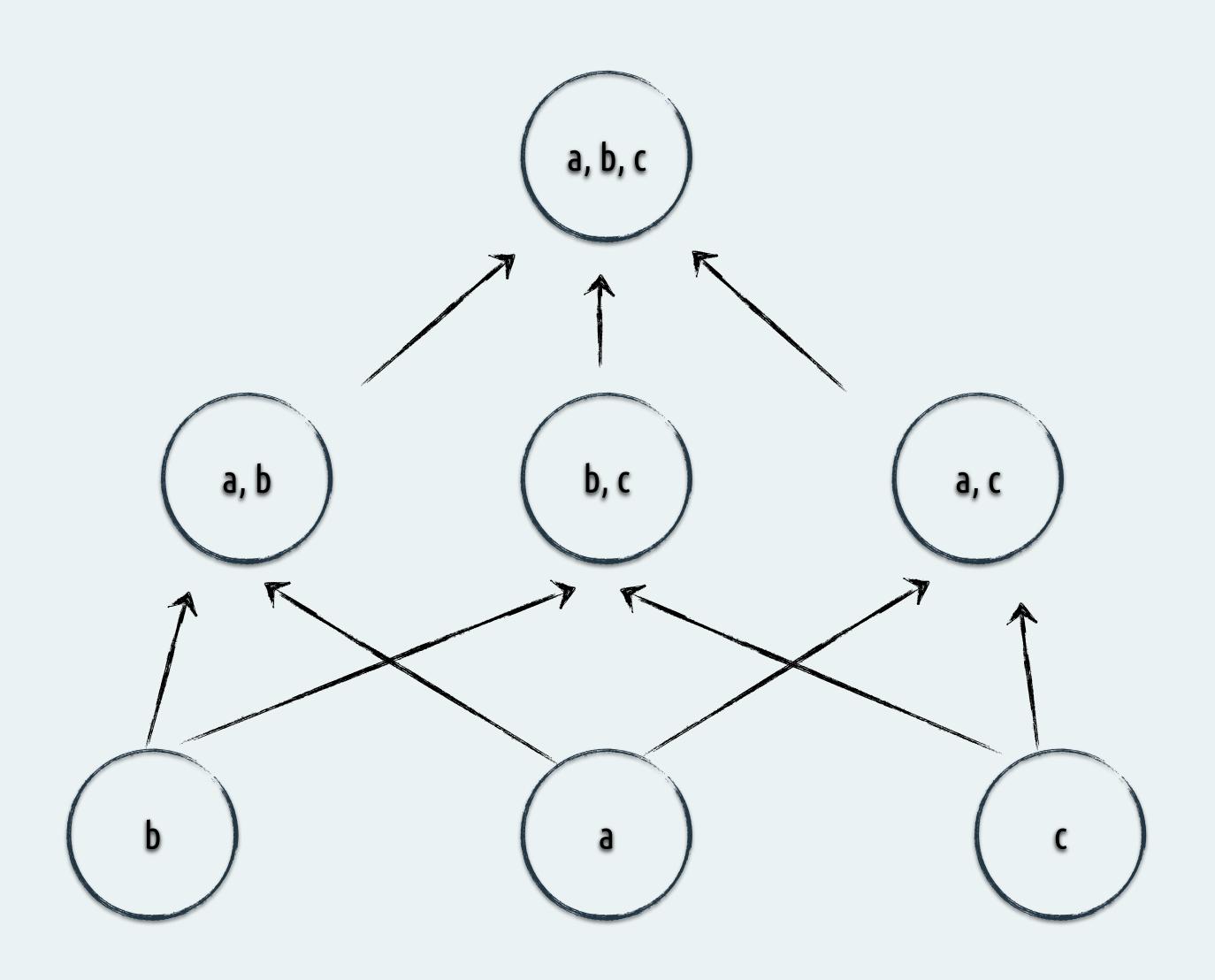


Objects grow over time; merge computes LUB

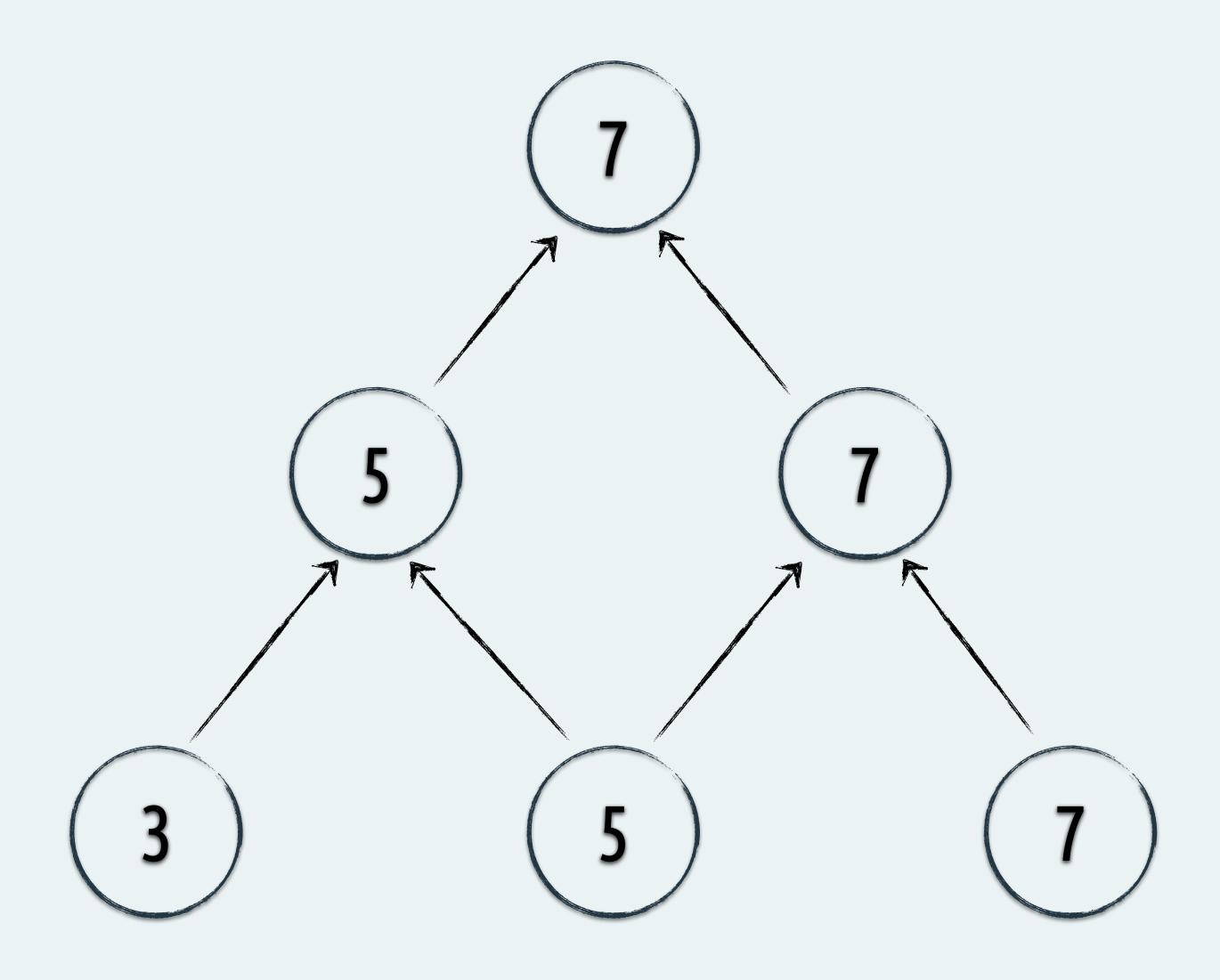


Examples

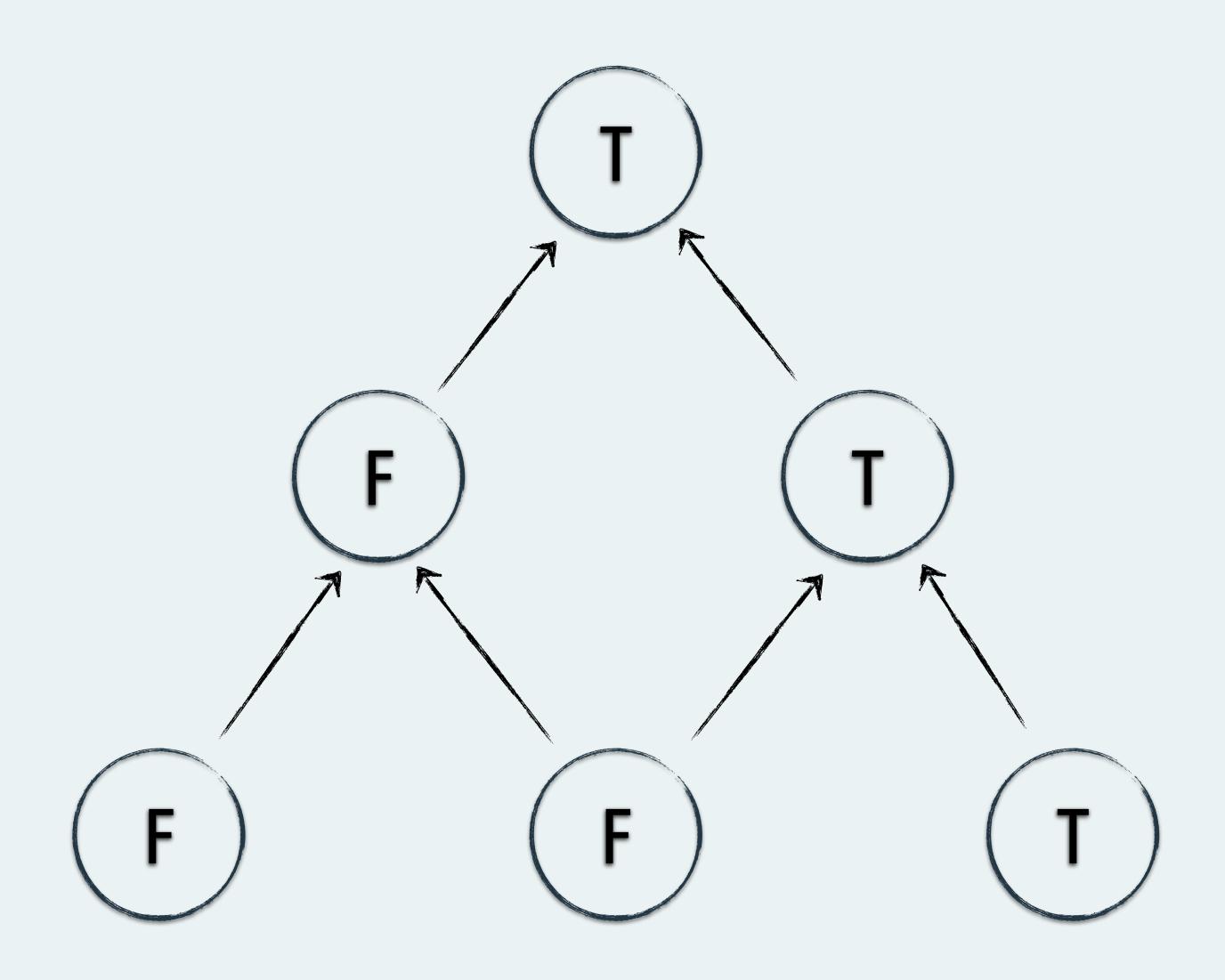




Set; merge function: union.



Increasing natural; merge function: max.



Booleans; merge function: or.

### Deterministic Idempotent Associative Commutative

## 

https://www.cs.indiana.edu/~lkuper/papers/lvars-fhpc13.pdf

#### Pick your semantic

Add wins
Remove wins
Keep both



Replicated Data Types: Specification, Verification, Optimality

Sebastian Burckhardt, Alexey Gotsman, Hongseok Yang, Marek Zawirski

# Trade Off More metadata == bigger objects

## Actors?

Version Vectors
Entry Per Actor

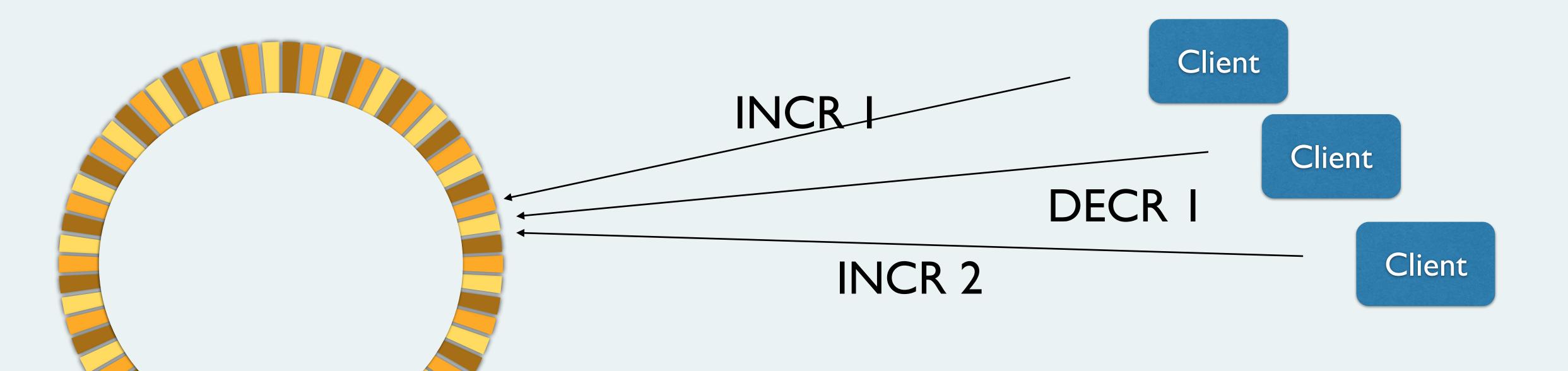
### PN-Counter

Counters: O(Actors)

 $A = [{a, 1,0}]$ 

 $B=[\{b, 0, 1\}]$ 

 $C = [\{c, 2, 0\}]$ 



 $A = [{a, 1,0}]$ 

 $B=[\{b, 0, 1\}]$   $C=[\{c, 2, 0\}]$ 

 $A = [{a, 1,0}]$ 

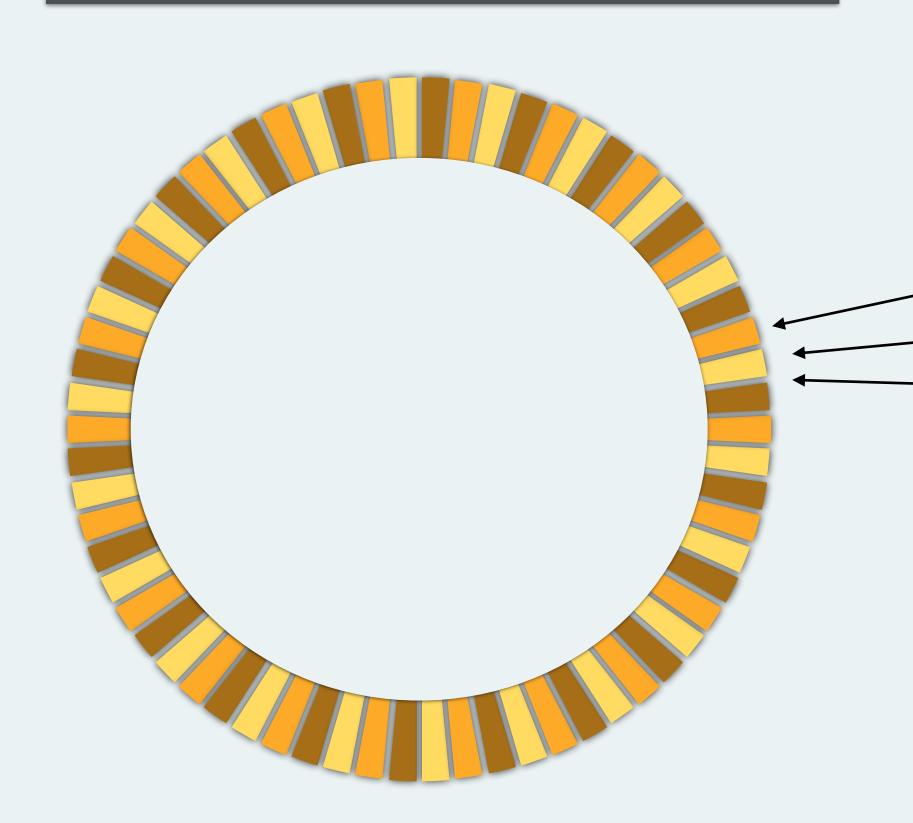
 $B=[\{a, 1, 0\}, \{b, 0, 1\}]$ 

 $C = [\{b, 0, 1\}, \{c, 2, 0\}]$ 

 $A = [{a, 1, 3}, {b, 2, 1}]$ 

B=[{a, 0, 1}, {b, 4, 2}, {c, 1, 0}]

 $C = [\{b, 2, 1\}, \{c, 3, 1\}]$ 



INCR, INCR, DECR,...

DECR, INCR,...

INCR, DECR,...

Client

Client

Client

```
A = [{a, 1, 3}, {b, 2, 1}]
```

```
B=[{a, 0, 1}, {b, 4, 2}, {c, 1, 0}]
```

```
C = [\{b, 2, 1\}, \{c, 3, 1\}]
```

```
AB=[{a, 1, 3}, {b, 4, 2}, {c, 1, 0}]
```

```
AB=[{a, 1, 3}, {b, 4, 2}, {c, 1, 0}]
```

$$P=1+4+3$$
 $N=3+2+1$ 
 $8-6=2$ 

Sets: Add, Remove, Membership.

Sets: Add wins O(Actors + Elements)

# Maps: Recursive; Associative Array; Nestable

Maps: Update wins; O(Actors + Elements)

## Composition

Maps: LWW-Register, Booleans, Sets and Maps

#### Use Case

- Mobile game progress data
  - Game State
- Non trivial merge

#### Sample JSON

```
{"gold": 500,
 "wood": 1250,
 "stone": 100,
 "buildings": [
   "house",
   "forge",
   "farm"
```

#### Desired Outcome

- Express updates as operations
- Apply related updates together
- Avoid "hand coded" resolution

#### Conceptual

#### Build Tower ::

- subtract 250 gold
- subtract 500 wood
- subtract 100 stone
- add "tower" to buildings

#### JSON Equivalent

```
{"gold_counter":
          {"decrement": 250},

"wood_counter":
          {"decrement": 500},

"stone_couner":
          {"decrement": 100},

"buildings_set":
          {"add": "tower"}}
```

# riak ot

git clone git@github.com:basho/riak\_dt.git

# Causality Version Vectors

[{a, 1}, {b, 3}, {c, 2}]



#### Causality Version Vectors

[{a, 2}, {b, 3}, {c, 2}]

[{a, 1}, {b, 3}, {c, 2}]



# Causality Version Vectors

[{a, 2}, {b, 3}, {c, 2}]

[{a, 2}, {d, 1}, {c, 2}]

[{a, 1}, {b, 4}, {c, 2}]

[{a, 2}, {b, 4}, {c, 2}]



# Causality Version Vectors

'Dots' are Events



### Causality

#### 'Dots' are Events

[{a, 2}, {b, 3}, {c, 2}]

{b, 1} {b, 2} {b, 3}



GSET

212-515

# Evolution of a Set U-SET

# Evolution of a Set

Adds

a1 Shelly

a2 Bob

a3 Pete

a4 Anna

b2 Shelly

Removes

a1 Shelly

a2 Bob

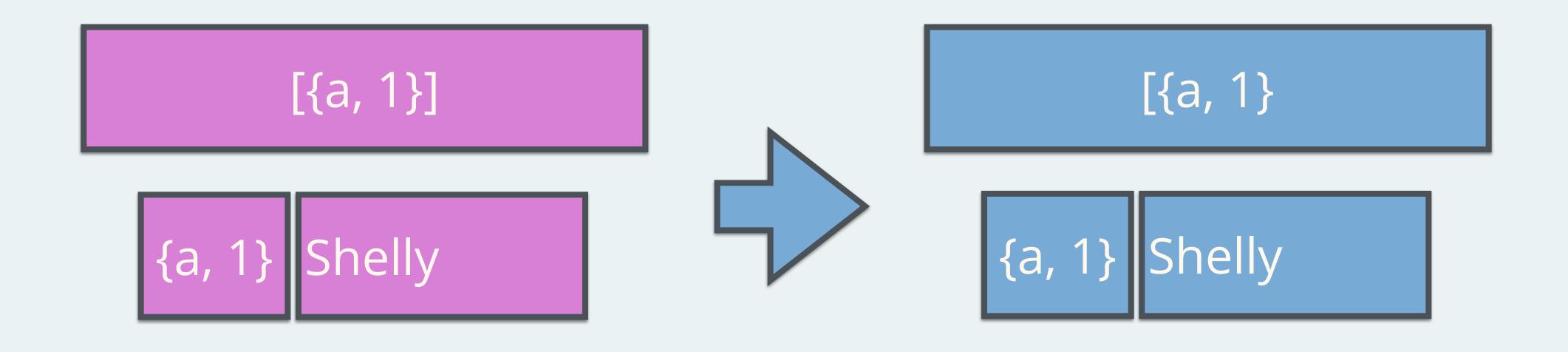
a3 Pete

# Evolution of a Set

### Evolution of a Set

[{a, 1}]

{a, 1} | Shelly



[{a, 1}]

{a, 1} Shelly

[{a, 1}, {b, 3}]

{a, 1} Shelly

{b, 1} Bob

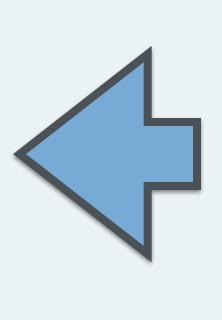
{b, 2} Phil

{b, 3} Pete

[{a, 1}, {b,3}]

[{a, 1}, {b, 3}]

{a, 1} Shelly {b, 1} Bob {b, 2} Phil Pete



{a, 1} Shelly {b, 1} Bob {b, 2} Phil {b, 3} Pete

[{a, 2}, {b, 3}]

{a, 1} Shelly

{b, 1}

Bob

{b, 3} Pete

{a, 2} Anr

[{a, 1}, {b, 3}]

{a, 1} Shelly

{b, 1} Bob

{b, 2} Phil

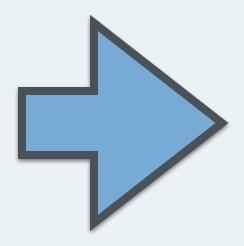
{b, 3} Pete

[{a, 2}, {b, 3}]

[{a, 2}, {b, 3}]

{a, 1} Shelly

Bob



{a, 1} | Shelly

{b, 1}

Bob

Pete

{b, 3} Pete

{a, 2} | Anna

#### Quickchecking Our Work

- OR-Set (Inefficient, Simple)
- ORSWOT (Complex)
- EQC Statem (OR-Set IS the Model)



#### Quickchecking Our Work

- Single Key
- 2-20 "Replicas"
- Riak as a list of 3-tuples
  - [{actor(), orset(), orswot()}]



#### Quickchecking Our Work

- Generate Commands
  - Add, Remove, Merge
- Test for equivalence
  - Per replica per command (post condition)
  - Merge all replicas



### Riak 2.0 Beta

http://docs.basho.com/riak/ 2.0.0beta2/downloads/

### And Then?

- Ad Hoc Composability
- More compact CRDTs
- Delta-Mutators (Baquero et al)
- Actor Garbage Collection



## Questions?

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