Secure Programming via Libraries

A library for information-flow in Haskell

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CHALMERS

Encoding information-flow in Haskell [Li, Zdancewic 06]

- Show that it is possible to guarantee IFC by a library
- Implementation in Haskell using Arrows [Hughes 98]
- Arrows? A generalization of Monads [Wadler 01]
- Pure values only
 - No side-effects
- One security label for data
 - All secret or all public!



CHALMERS

Encoding information-flow in Haskell [Tsai, Russo, Hughes 07]

- Extend the library by Li and Zdancewic
 - More than one security label for data
 - Concurrency
- Major changes in the library
 - New arrows
 - Lack of arrow notation
- Why arrows?
 - Li and Zdancewic argue that *monads are not suitable for the design of such a library*

A lightweight library for Information-flow in Haskell [Russo, Claessen, Hughes 08]

- Lightweight
 - Approximately 325 lines of code
 - Static type-system of Haskell to enforce non-interference
 - Dynamic checks when declassification occurs
- Use Monads (not Arrows!)
 - Programmers are more familiar with Monads than Arrows

A lightweight library for Information-flow in Haskell [Russo, Claessen, Hughes 08]

- The library relies on Haskell
 - Capabilities to maintain abstraction of data types
 - Haskell module system
 - Haskell is strongly typed
 - We cannot cheat!

unsafePerformIO :: IO a -> a unsafeCoerce :: a -> b

- There are extensions of Haskell that break these two requirements!
- For a full list, please visit the proposal of SafeHaskell
 - An extension of Haskell to disallow those dangerous features than can jeopardize security
 - Join work with David Mazieres et al. at Stanford university.

Why Haskell?

- Clear separation of pure computations with those with side-effects
- Every computation with side-effects is encapsulated into the IO monad
- Side-effects can encode information about secret data
- It is necessary to control them
 - It is known where they occur! Just look at the type!

Side-effects and IO

• Just look at the type!

```
f1 :: Eq a => a -> [a] -> ([a], Bool)
```

f2 :: (Show a, Eq a) => Int -> a -> ([a], IO Bool)

- All bets are off if an IO computation comes from untrustworthy code
 - It is not known the side-effects that it will produce f1 x xs = (take 10 (cycle xs), elem x xs)

```
f2 n x = (take n (iterate id x),

do putStrLn "Hi!"

putStrLn "The arguments of the function are"

putStrLn $ "x = " ++ show x

putStrLn $ "n = " ++ show n

return True )
```

Secure Pure Computations f :: (Char {- secret -}, Int) -> (Char {- secret -}, Int) YES f(c, i) = (chr(ord c + i), i)f (c, i) = (chr(ord c + i), ord c) NO f(c, i) = (chr(ord c + 1), i+1)YES f (c, i) | c > 65 = (c, 42) | otherwise = (c, i) NO

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A Security Monad for Pure Computations

- Security monad
 - It assigns a security level to data
 - Once inside the monad, it is not possible to escape!

data Sec s a -- abstract
instance Monad (Sec s)

• We represent security levels by singleton types

secret :: Sec H Int H secret = ...

known :: Sec L Int known = ...

Using Sec

Using Sec

Security Guarantee

Type checks!

Non-interferece



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A Security Monad for Pure Computations

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data Sec s a -- abstract
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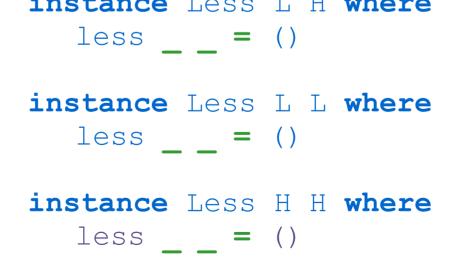
- We represent security levels by singleton types
- What about the security lattice?

Security Lattice

- We model it using type classes in Haskell
 - Constrains to polymorphic types

class Less s s' where less :: $s \rightarrow s' \rightarrow ()$

 Encoding two-point lattice is just provide instances for that type class
 instance Less L H where



Security Monad and the Security Lattice

Push up information in the security lattice

up :: Less s s' => Sec s a -> Sec s' a

- It allows to convert public values into secrets
 fup :: Sec L Int -> Sec H Char
 fup sec_i = do i <- up (sec_i) return (chr i)
- What if it is possible to make the following instance?
 instance Less H L where
 less _ = ()

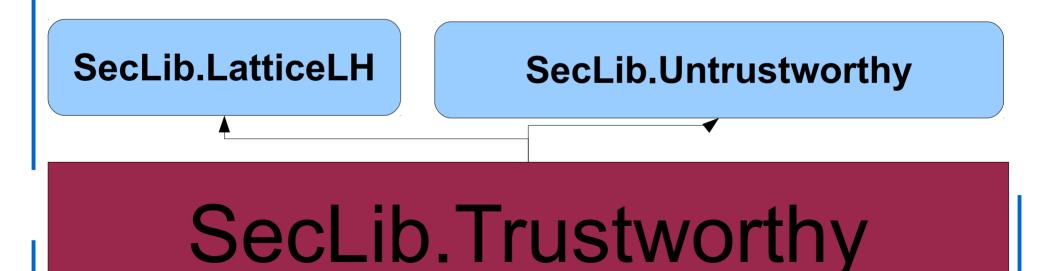
Security Monad and the Security Lattice

- The library works as long as
 - Attackers cannot define method less for arbitrary instances of the type class Less
- How to ensure that?
 - Mainly by the abstraction power of Haskell's module system

Arquitecture

module X where

import SecLib.Untrustworthy
import SecLib.LatticeLH





Importing SecLib.Trustworthy

- SecLib.Trustworthy must not be imported by untrustworthy code
 - Otherwise, no security guarantees are possible

instance Less H L where
 less _ = ()

Other Assumptions

- The monad Sec s must remain abstract
 - Guarantee by the installation of the library
 - Sec.hs is not an exposed module
- Use of unsafe Haskell extensions
 - StandaloneDeriving
 - System.IO.Unsafe
 - unsafePerformIO, unsafeIterleaveIO, etc.
 - OverlappingInstances
- Check SafeHaskell (work-in-progress)
 - A Haskell extension to safely execute untrusted Haskell code

Security API for Pure Computations

data Sec s a -- abstract
instance Monad (Sec s)

up :: Less s s' => Sec s a -> Sec s' a

module X where

import SecLib.Untrustworthy
import SecLib.LatticeLH