

Weak Models For Distributed Computing

Gadi Taubenfeld

IDC, Israel



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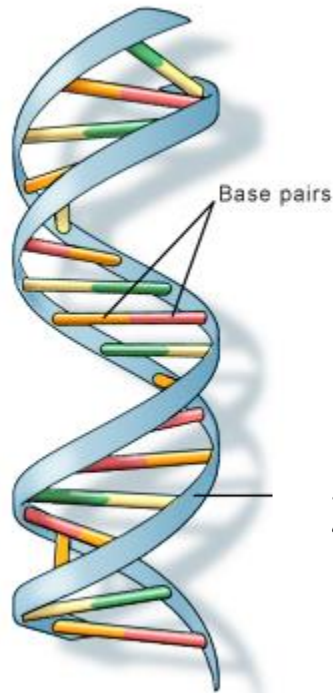
IDC, Israel

Gadi: I am not an implementor of tools, programming languages, or platforms!

Annie: ... pls mention where computers can help you except from text editor and a slides editor ...

Part I

Genome-Wide Epigenetic Modifications as a Shared Memory Consensus Problem



U.S. National Library of Medicine



Ziv Bar-Joseph
CMU



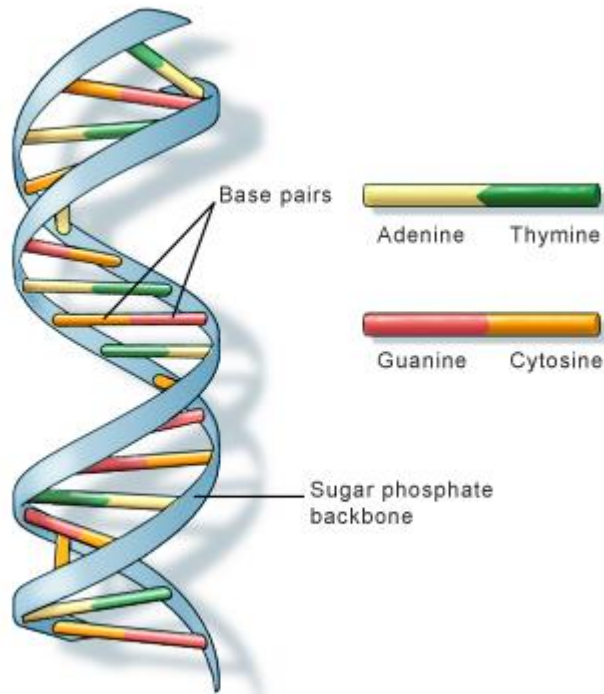
Sabrina Rashid
CMU



Gadi Taubenfeld
IDC

The human genome

The entire DNA of a single human cell



U.S. National Library of Medicine

- Two meters long
- 3 billion base pairs
- About 25,000 genes

(Only about 1 percent of DNA is made up of protein-coding genes)

Chromatin

Package DNA into a small volume to fit into the nucleus of a cell

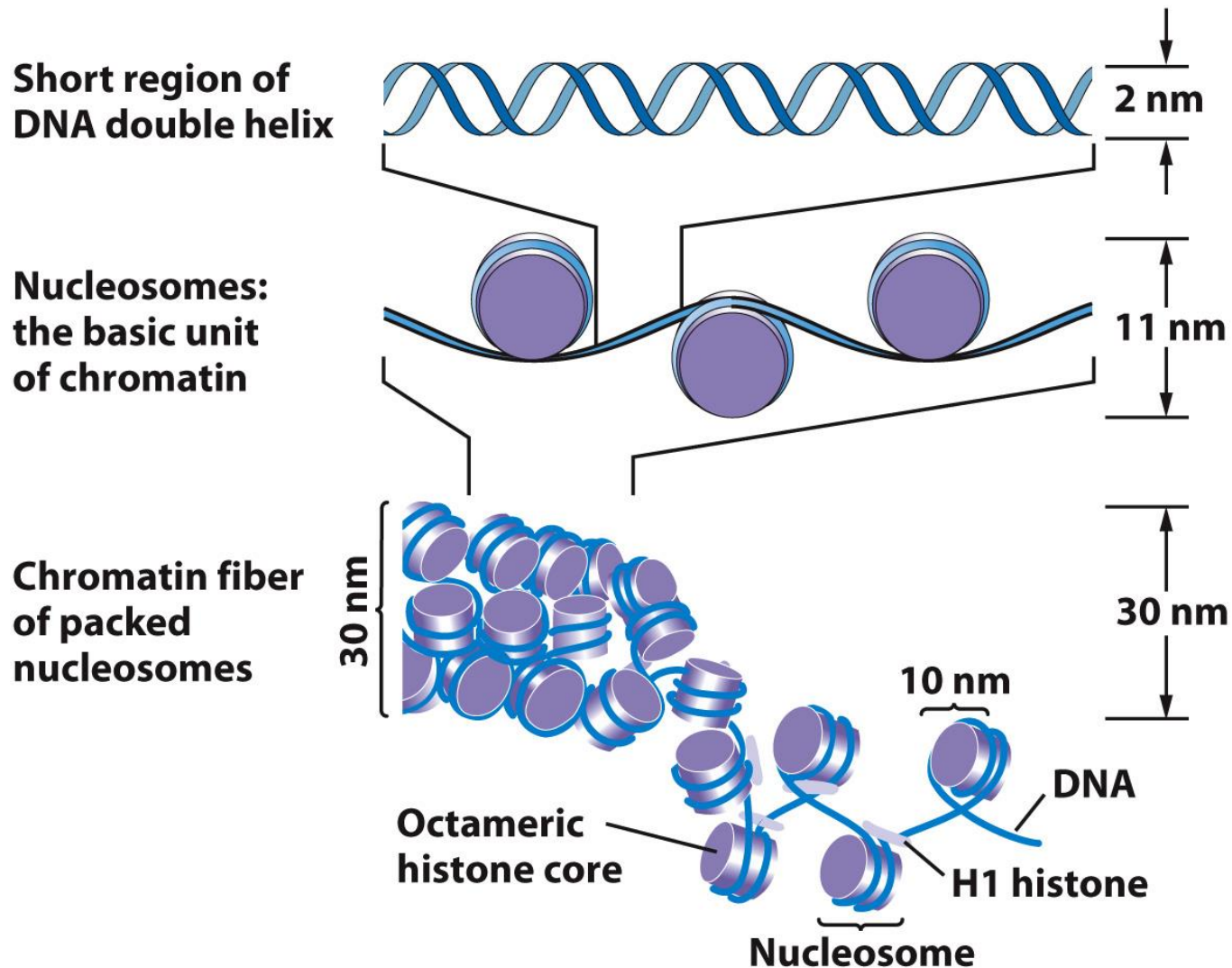


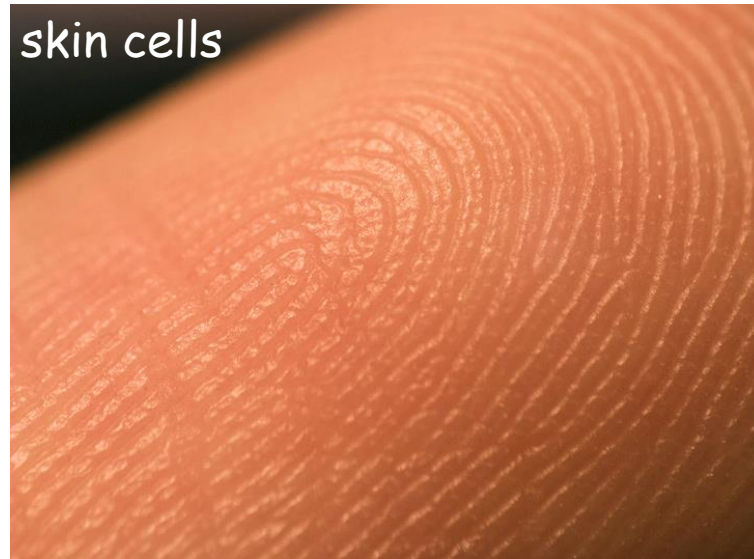
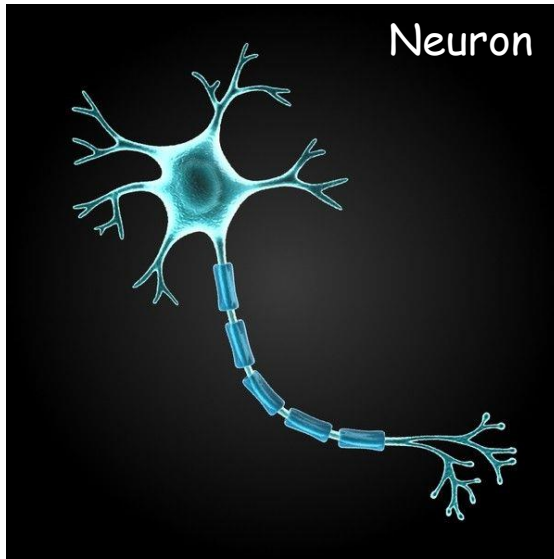
Figure 12-11a

Introduction to Genetic Analysis, Eleventh Edition

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Cells types & DNA

Q: How can an organism have different cell types yet one genome?

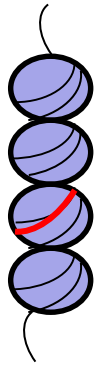


A: Each cell expresses, or turns on, only a fraction of its genes. The rest of the genes are repressed, or turned off.

Regulation of gene expression

Turning genes on and off

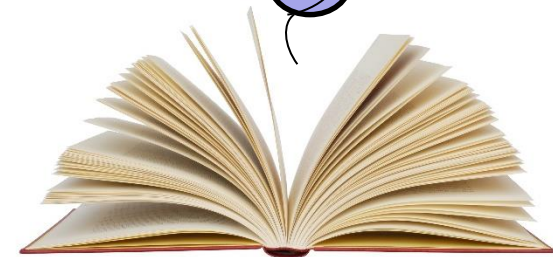
Condensed chromatin



Off



Open chromatin



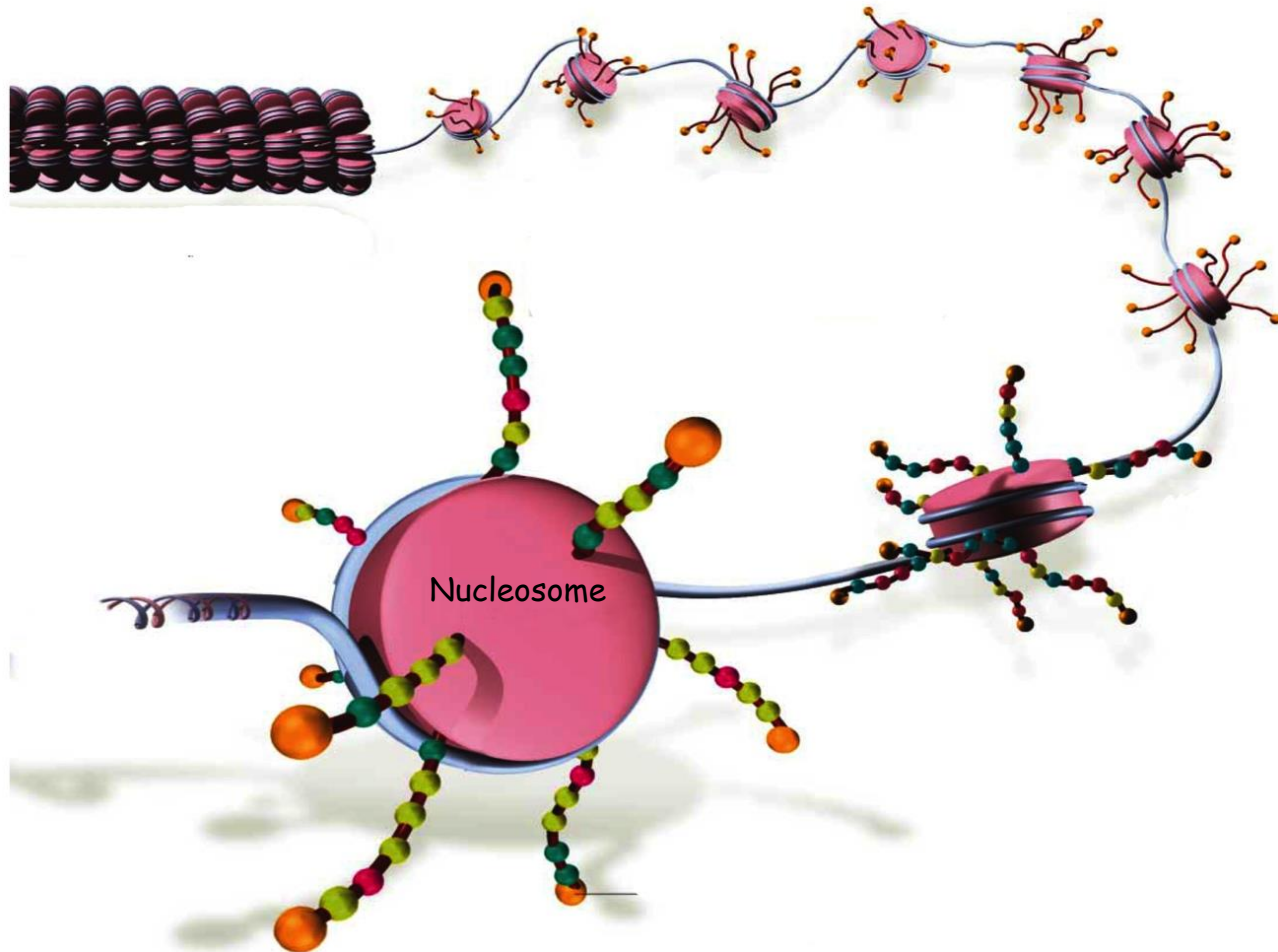
On

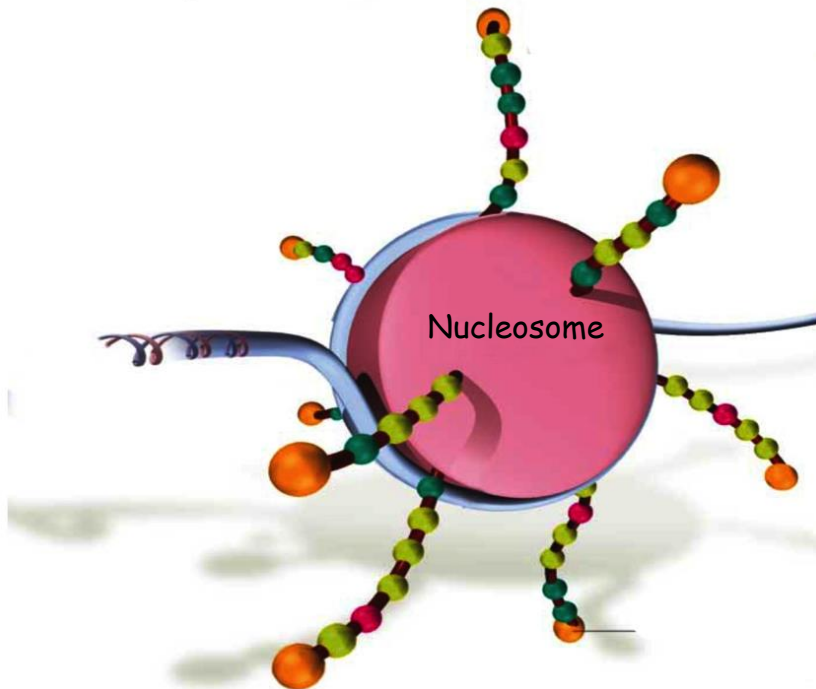
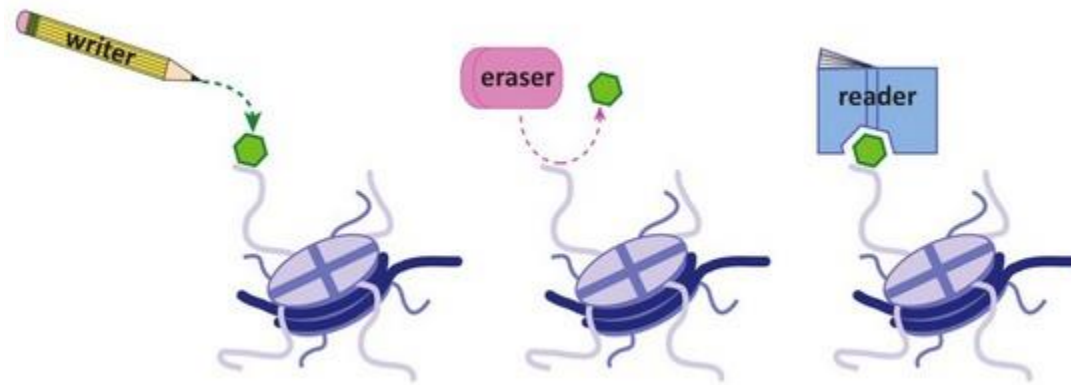
Environmental influences, such as a person's diet, stress and exposure to pollutants, impact gene expression.



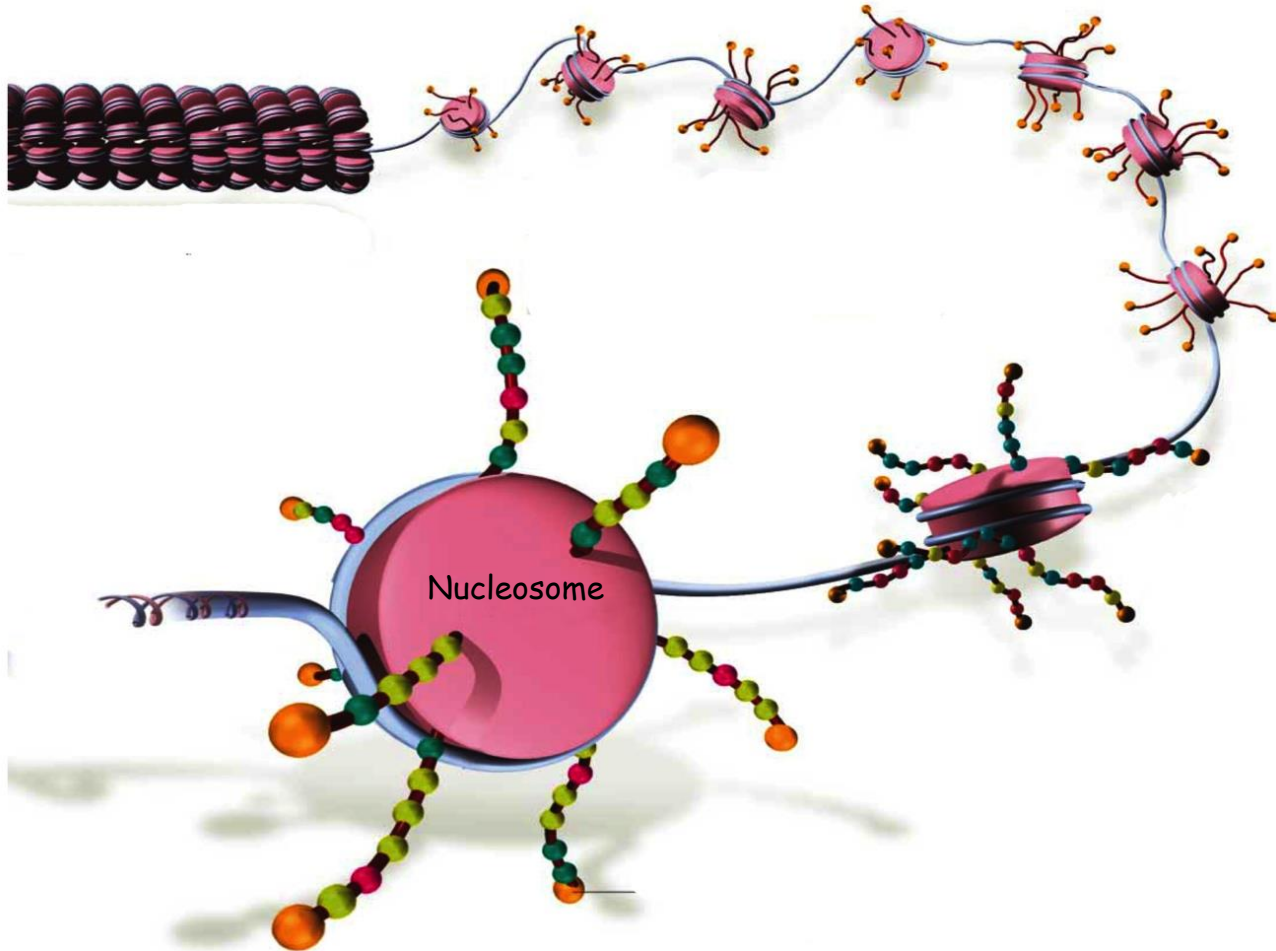
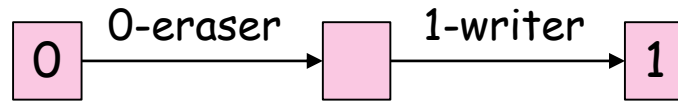
Epigenetics

Modifications that do not change the DNA and affect gene activity

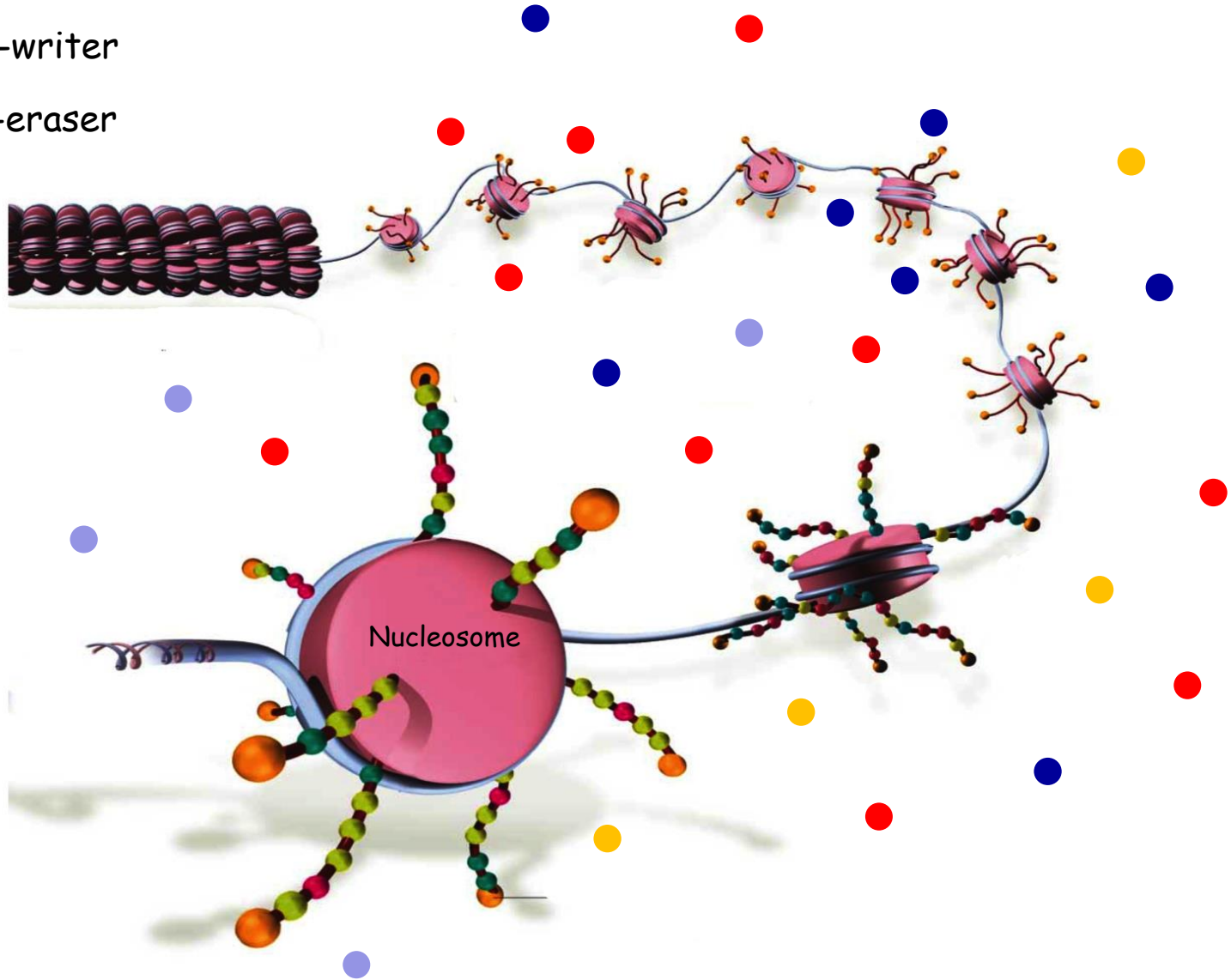




- 1-writer
- 0-eraser
- 0-writer
- 1-eraser

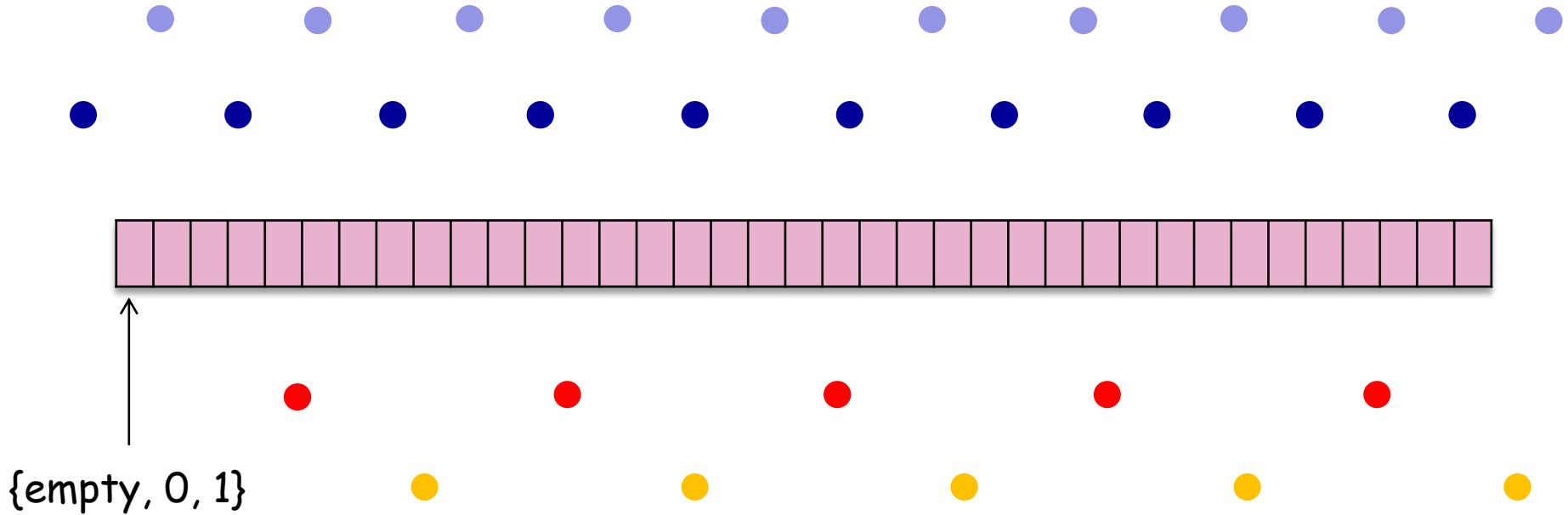


- 1-writer
- 0-eraser
- 0-writer
- 1-eraser



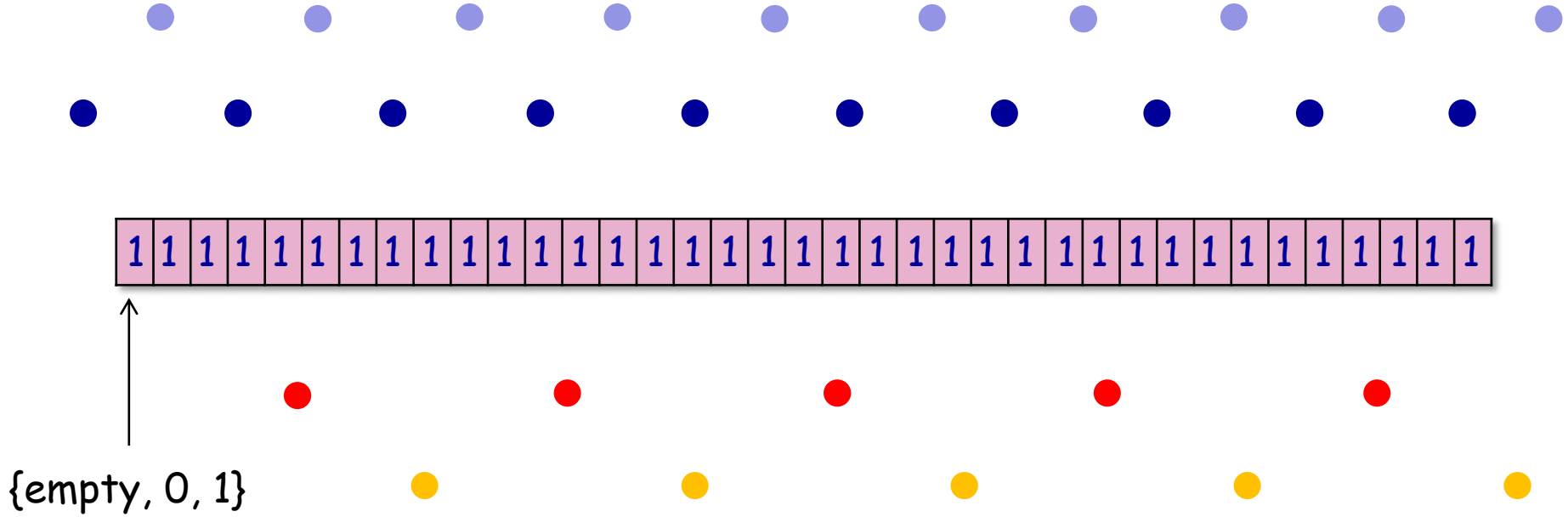
The epigenetic consensus problem

- 1-writer
- 0-eraser
- 0-writer
- 1-eraser



The epigenetic consensus problem

- 1-writer
- 0-eraser
- 0-writer
- 1-eraser



Very weak model

- Randomization
- Anonymous processes (no identifiers)
- **Anonymous shared memory**
- Memory-less processes (well may 1-2 bits)
- A transition from 0 to 1 cannot occur directly
- No sense of direction
- Self-stabilization

*We've got FSMs...
what else do we need?*



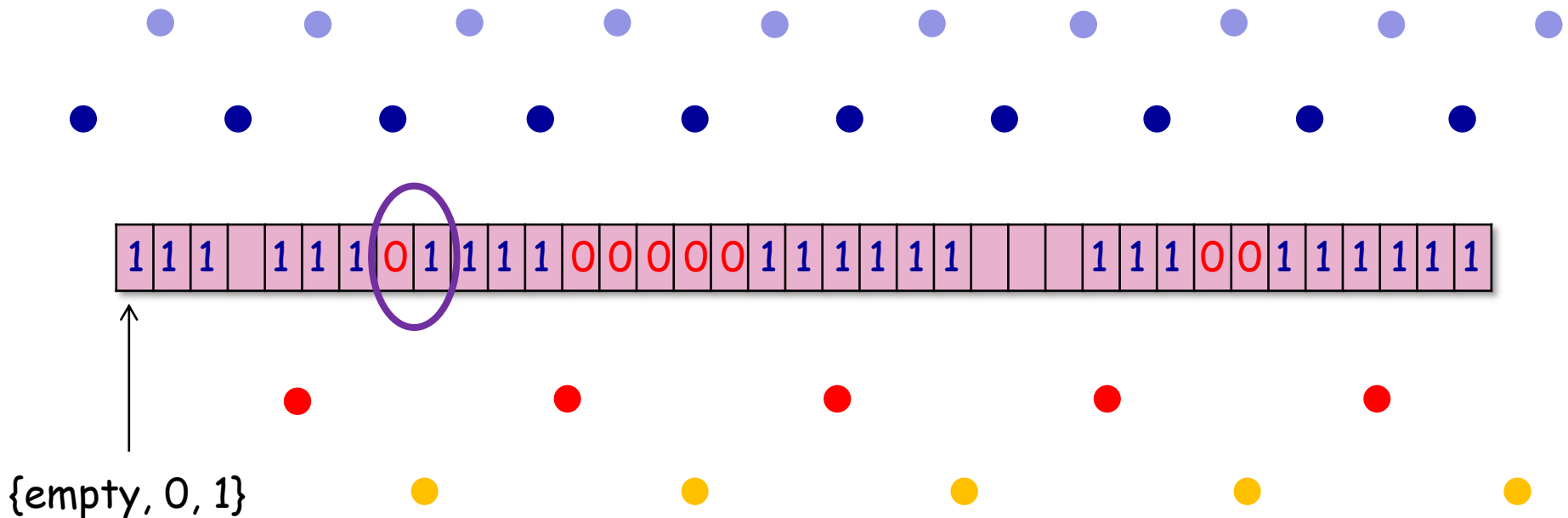
We present an algorithm that matches the biological assumptions, prove its correctness and derive bounds on its expected run time both theoretically and in simulations.

The epigenetic consensus problem

Corollary 6.2 Assume $W_1/W_0 \geq 3$. The probability that the final decision value is 1 is more than

$$(1 - (1/3)^N) \times (1 - e^{-N/12})$$

Corollary 7.2 Assuming $W_1/W_0 \geq 3$, $E[T] \leq 6.4N^2$



Annie's question: Where computers can help you except from text editor and a slides editor ?



Simulations

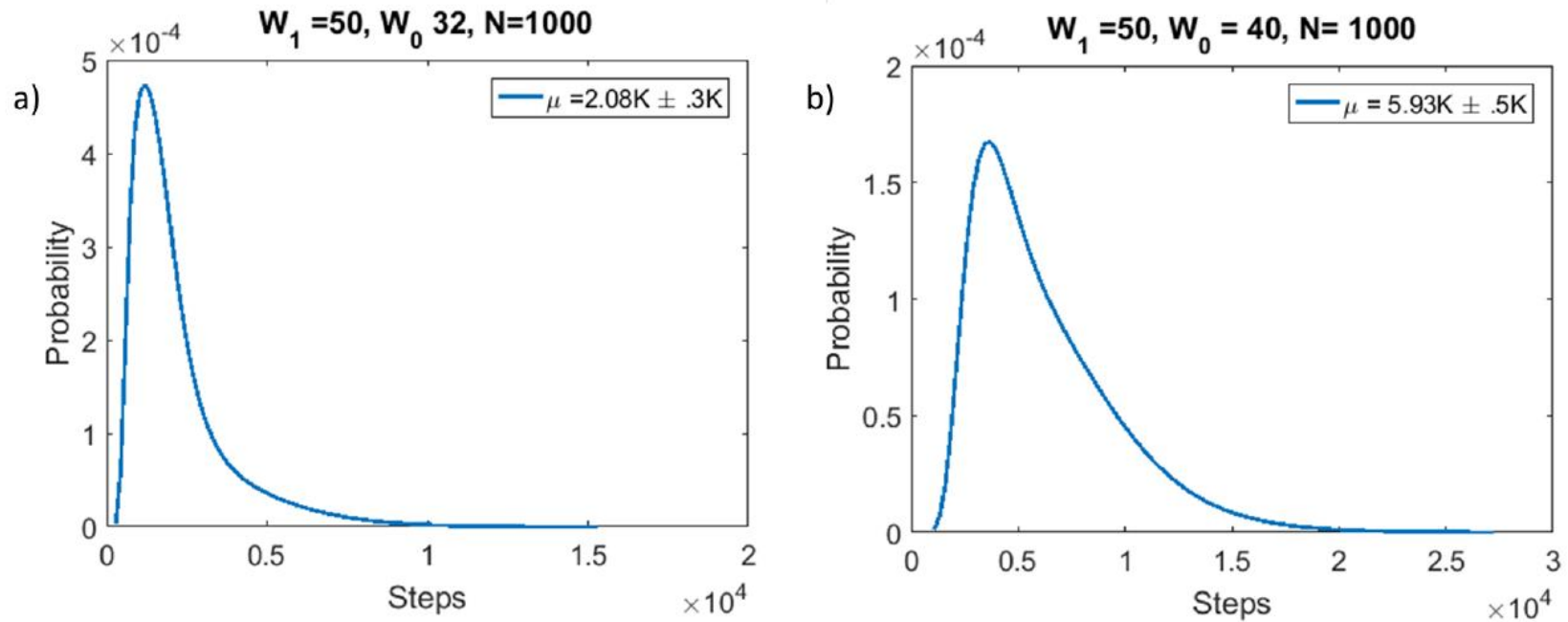


Figure 1: Distribution of number of steps to reach consensus. Plots summarize 300 random runs of the algorithm. a) low and b) high level of competitions between 1-writers and 0-writers. μ denotes the average time to reach consensus.

Simulations

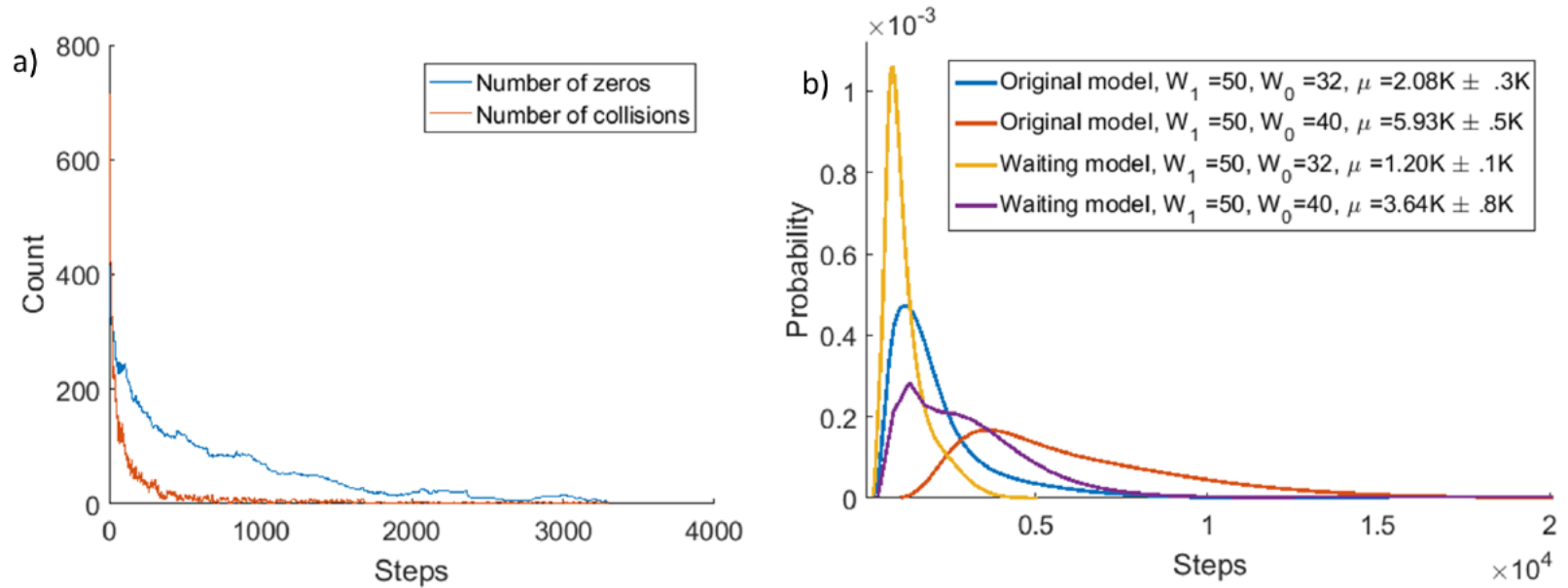


Figure 2: a) Number of zeros and collisions vs steps in the algorithm. While the initial number of collisions is a linear function of the number of 0's, we observe that towards the end of the algorithm there are very few collisions while the number of 0's remains relatively high. b) Comparison between the proposed original model and a revised model that allows writers to attach themselves with the erasers and the erasers wait until a collision is resolved. Here we can see that the waiting version is faster at both competition level compared to the original model.

Conclusion #1

Weak models are interesting!



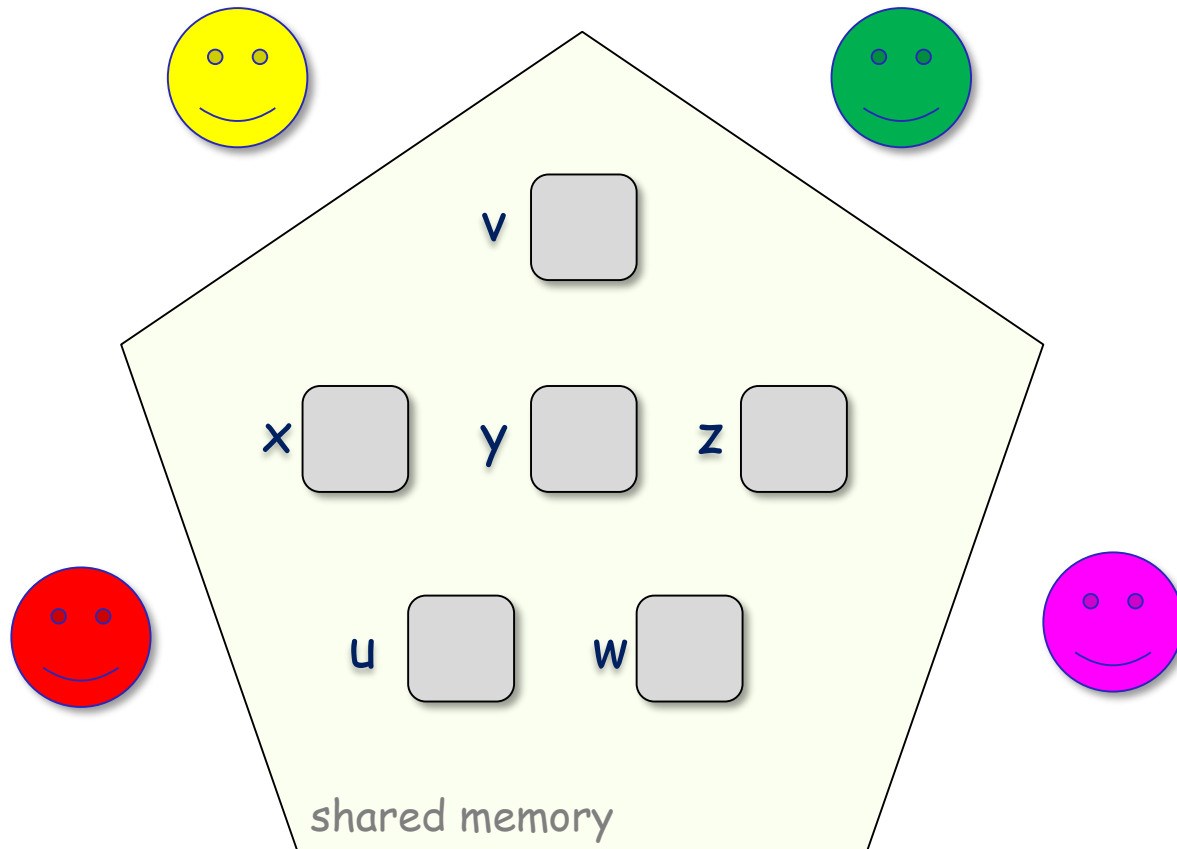
Part II

Anonymous Shared Memory



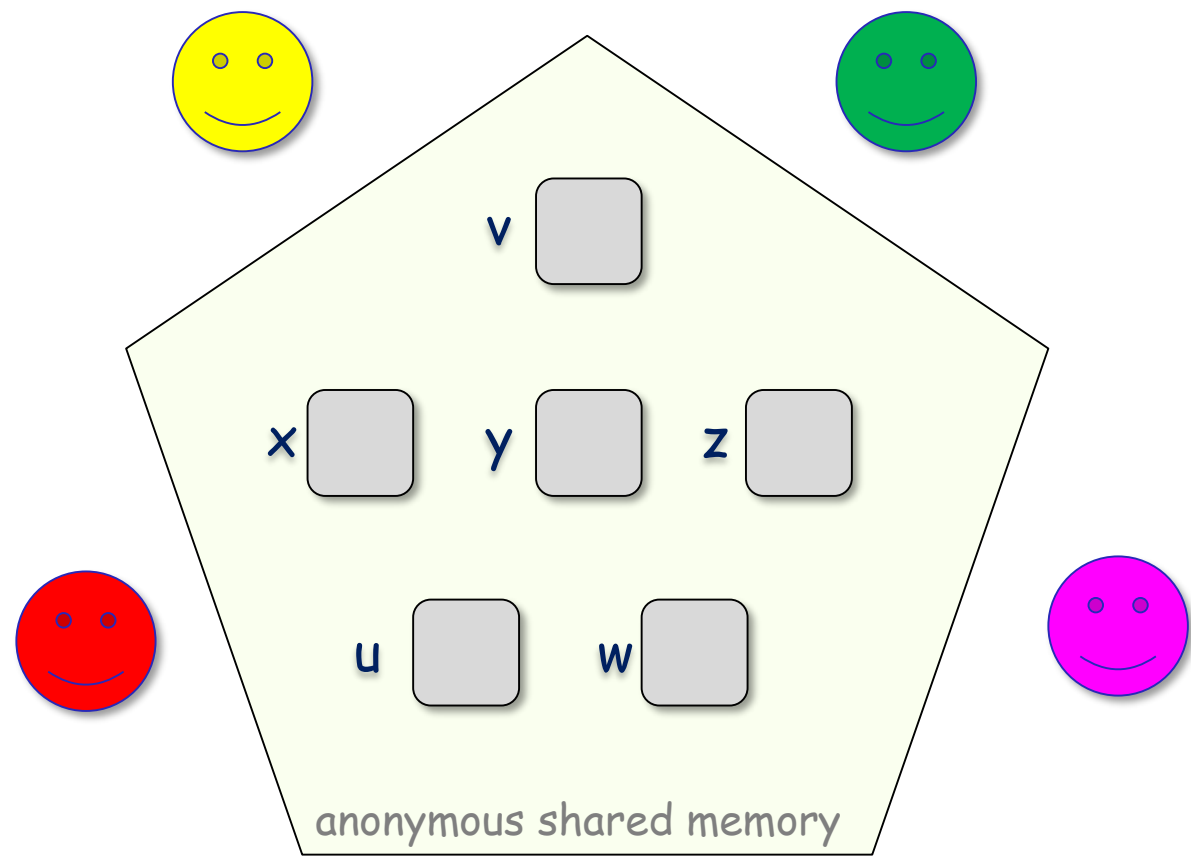
Classical view of SM

Objects have names



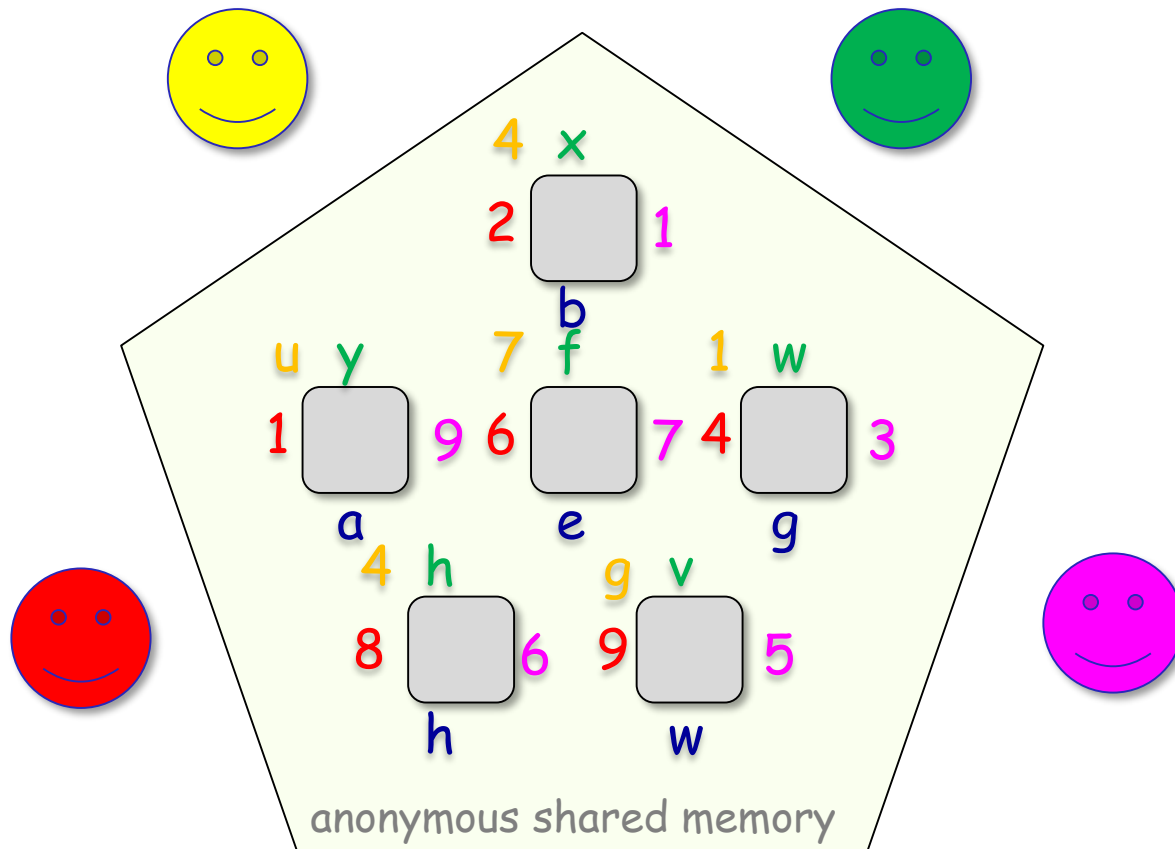
Anonymous shared memory

NO prior agreement on the names of the objects!



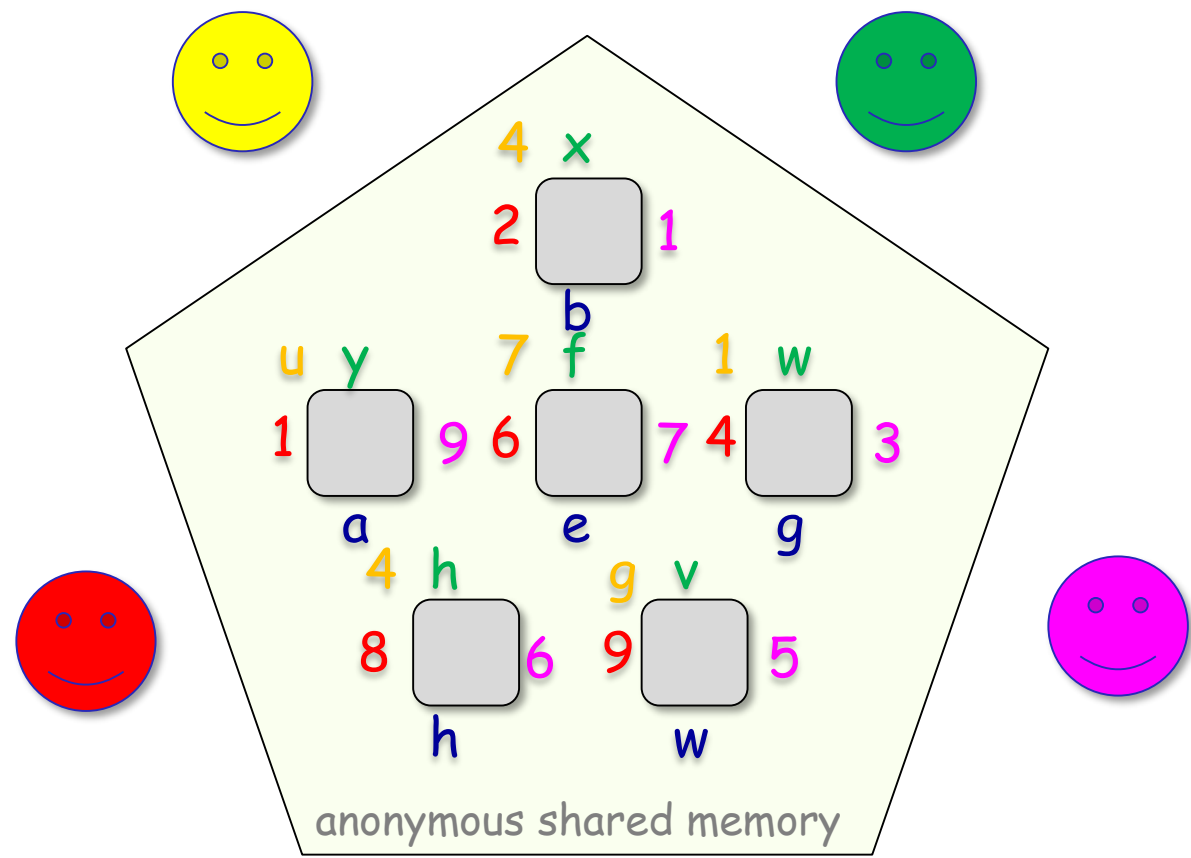
Anonymous shared memory

New Model



Anonymous shared memory

Coordination without prior agreement
 by *Gadi Taubenfeld*



Algorithms & space bounds

Algorithms	Can do ✓	Cannot do ✗
Deadlock-free symmetric mutual exclusion for <u>two</u> processes	odd # of registers	even # of registers
Obstruction-free consensus for $n \geq 2$ processes	$2n-1$ or more	n or less
Obstruction-free adaptive perfect renaming for $n \geq 2$ processes	$2n-1$ or more	n or less

(The # of registers is not 1)

Optimal Memory-Anonymous Symmetric Deadlock-Free Mutual Exclusion

- **Theorem.** For every $n \geq 1$, there is a symmetric deadlock-free mutual exclusion algorithm for n processes using $m \geq 1$ anonymous **R/W registers** if and only if for every positive integer $1 < k \leq n$, m and k are relatively prime.
- The same result holds also for RMW registers! *



Zahra
Aghazadeh



Damien
Imbs



Michel
Raynal



Philipp
Woelfel



Gadi
Taubenfeld

* It is trivial to do also with one RMW register.

Resolving two open problems

For a universe which includes (also) anonymous objects,

- Are atomic read/write registers the weakest objects ?
- Are deterministic (oblivious) objects with the same set-consensus number have the same computational power ?

NO

Conclusion #2

Weak models are interesting!



Part III

Fractions in Distributed Computing

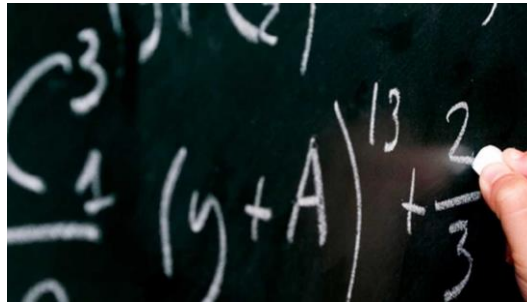
Egypt
1600 B.C.



Fractions were studied by Egyptians mathematicians around 1600 B.C.

However, fractions, as we use them today, didn't exist in Europe until the 17th century.

Europe
17th century



Dist. Comp.
???

Part III

Fractions in Distributed Computing

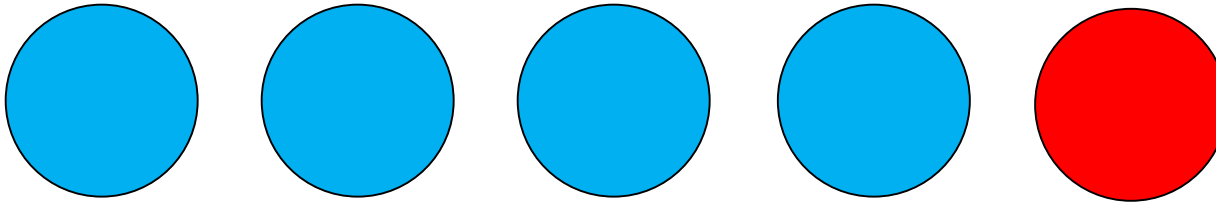
- We understand what it means to tolerate **one** process failure.
- But what does it mean to tolerate **0.8** process failure ?



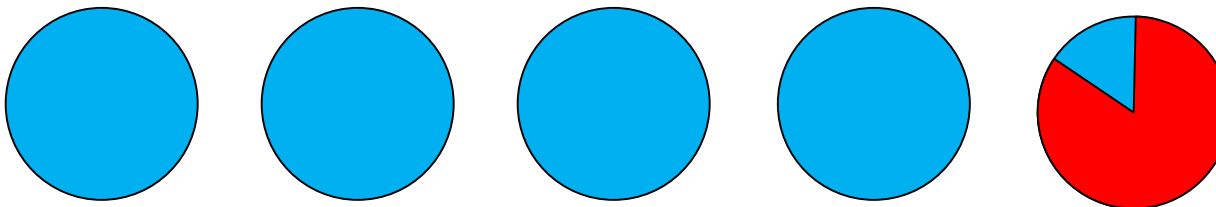
Motivation

Something is better than nothing

❖ FLP: Impossibility of consensus in the presence of a single failure.

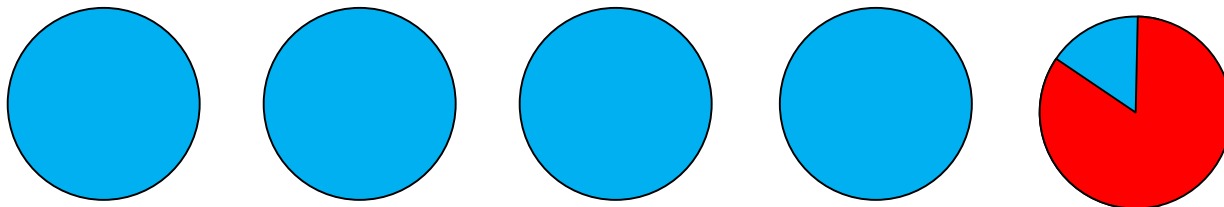


❖ Is consensus possible in the presence of a single weak failure?



Weak Failures: Definitions, Algorithms and Impossibility Results by *Gadi Taubenfeld*

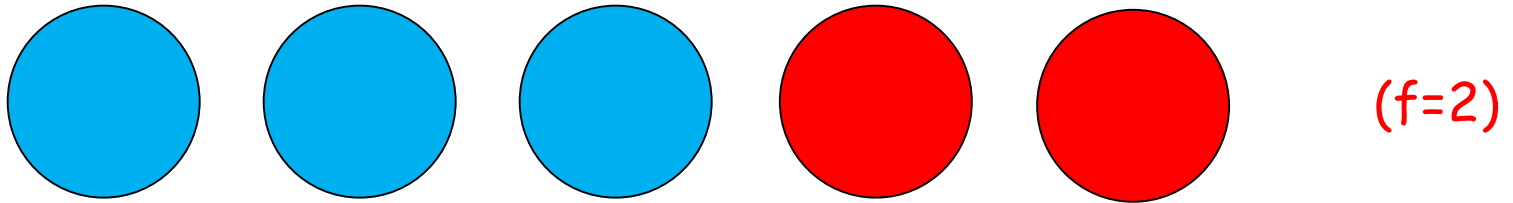
❖ Is consensus possible in the presence of a single weak failure?



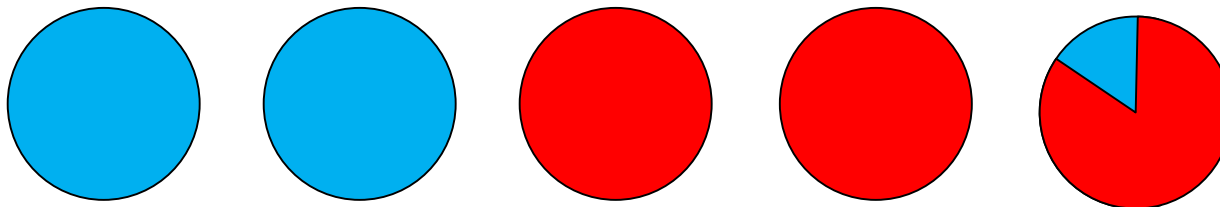
Motivation

Generalizing from the previous example

- ❖ Suppose you can solve a problem in the presence of f traditional failures, but not in the presence of $f+1$ such failures.

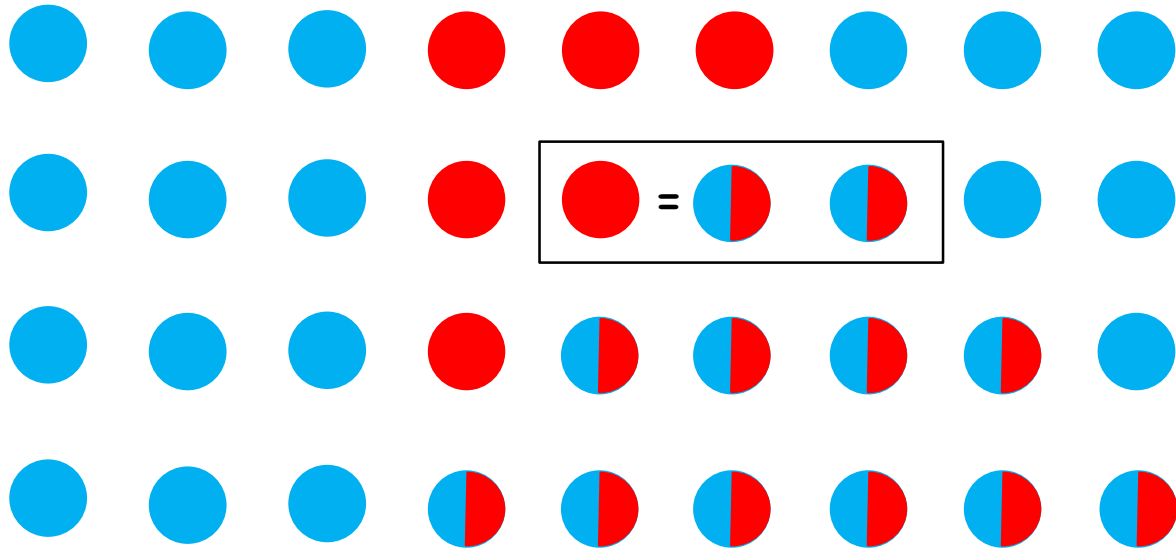


- ❖ Maybe it is possible to solve the problem in the presence of f traditional failures **plus** several weak failures.



Set agreement and renaming in the presence of contention-related crash failures

SSS 2018



Anais Durand



Michel Raynal



Gadi Taubenfeld

Conclusion #3

Weak models are interesting!



