

Forcing and non principal ultrafilter

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Abstract

The goal of this note is to present a simple proof of the fact that analysis extended with the existence of a non principal ultrafilter of natural number is conservative over analysis with dependent choice. The proof is purely syntactical and is a variation of an argument presented by Levin [1].

A.M. Levin “One conservative extension of formal mathematical analysis with a scheme of dependent choice” (1977)

Forcing over the system $\text{HA}^\omega + \text{EM} + \text{DC}$ (for well-ordering of the reals)

Theorem: *If $\text{HA}^\omega + \text{EM} + \text{DC} + \text{SUF} \vdash A$ then $\text{HA}^\omega + \text{EM} + \text{DC} \vdash A$*

The terms of the language are simply typed lambda terms. We have two basic types N (natural numbers) and N_2 (booleans). The atomic formulae are simply the terms of type N_2 . There are two terms $0, 1$ of type N_2 and we identify 1 with the true formula \top and 0 with the false formula \perp .

The formulae are

$$\varphi ::= \varphi \rightarrow \varphi \mid t \mid \forall x.\varphi$$

where t is a term of type N_2 (decidable atomic formula)

We use n, m, \dots for variables over the type N . Example: $\forall n.\exists^c m.n < m$.

$\neg\varphi$ to be $\varphi \rightarrow \perp$

$\exists^c x.\varphi$ is $\neg\forall x.\neg\varphi$

The system HA^ω is intuitionistic with the usual rules of natural deduction and induction over natural numbers and boolean. The rule EM is $(\neg\neg\varphi) \rightarrow \varphi$ which is equivalent to $\varphi \vee \neg\varphi$. The rule DC is

$$\forall n.\forall x.\exists y.\varphi(n, x, y) \rightarrow \forall u.\exists f.\varphi(0, u, f(0)) \wedge \forall n.\varphi(n, f(n), f(n+1))$$

The rule CC is

$$\forall n.\exists y.\varphi(n, y) \rightarrow \exists f.\forall n.\varphi(n, f(n))$$

We add a new symbol μ and new atomic formula $\mu(f)$ for f of type $N \rightarrow N_2$

We consider now the extension of the theory HA^ω with the axioms (we could add the selectivity axiom)

$$\begin{aligned} \mu(1) \quad \mu(fg) &\leftrightarrow (\mu(f) \wedge \mu(g)) \\ \mu(f) \vee^c \mu(1-f) \quad \mu(f) &\rightarrow \forall m.\exists^c n > m.f(m) \end{aligned}$$

We use letters p, q, r, \dots to denote *forcing conditions*, here simply terms of type $N \rightarrow N_2$. One can think of forcing conditions as *decidable subsets* of \mathbb{N} .

We define a formula $p \Vdash \varphi$ by induction on φ where φ is an extended formula (which may contain the new symbol μ) and p is of type $N \rightarrow N_2$.

$$I(p) \text{ is } \forall n.\exists m > n.p(m) \quad F(p) \text{ is } \exists n.\forall m > n.\neg p(m)$$

$\mu(f) \rightarrow I(f)$
 $p \leq q$ is $F(p(1 - q))$
 $p \Vdash \mu(f)$ is $p \leq f$
 $p \Vdash \varphi$ is $I(p) \rightarrow \varphi$ if φ is a boolean
 $p \Vdash \varphi_0 \rightarrow \varphi_1$ is $\forall q \leq p. (q \Vdash \varphi_0) \rightarrow (q \Vdash \varphi_1)$
 $p \Vdash \forall x. \varphi$ is $\forall x. (p \Vdash \varphi)$

We can add other connectives and existential quantification

Not needed if we are only interested in classical logic

Proposition: *If $\varphi_1, \dots, \varphi_n \vdash \varphi$ and $p \Vdash \varphi_1, \dots, p \Vdash \varphi_n$ then $p \Vdash \varphi$*

Using EM

Proposition: *We have $p \Vdash \varphi_0 \vee^c \varphi_1$ iff*

$$\forall q \leq p. \exists r \leq q. (r \Vdash \varphi_0) \vee^c (r \Vdash \varphi_1)$$

and $p \Vdash \exists^c x. \varphi$ iff

$$\forall q \leq p. \exists r \leq q. \exists^c x. r \Vdash \varphi$$

Proposition: *We have (classical version of the comprehension axiom)*

$$p \Vdash (\forall n. \varphi(n, 0) \vee^c \varphi(n, 1)) \rightarrow \exists^c f. \forall n. \varphi(n, f(n))$$

This expresses that there are no more decidable functions in the extension than in the ground model

Proposition: *We have (countable choice)*

$$p \Vdash (\forall n. \exists^c x. \varphi(n, x)) \rightarrow \exists^c f. \forall n. \varphi(n, f(n))$$

All the axioms of non principal ultrafilters are forced

We have $\text{HA}^\omega \vdash (I(p) \rightarrow \varphi) \leftrightarrow (p \Vdash \varphi)$ if φ does not mention μ

$\text{HA}^\omega + \text{EM} + \text{DC} + \text{SUF} \vdash \varphi$ implies $\text{HA}^\omega + \text{EM} + \text{DC} \vdash (\Vdash \varphi)$ and hence $\text{HA}^\omega + \text{EM} + \text{DC} \vdash \varphi$

So we have a computational interpretation of non principal ultrafilters

Levin (1977) does the same with a well-ordering of the reals, which justifies also the continuum hypothesis

References

- [1] A.M. Levin. One conservative extension of formal mathematical analysis with a scheme of dependent choice 1977