Simplicial set semantics of higher inductive types

Introduction

The goal of this note is to show that, by combining some results in the references[1], [2] and [3], we do get a model of higher inductive types in simplicial sets, extending and simplifying Voevodsky's model of univalent type theory. This follows essentially what was announced by Andrew Swan in some discussions¹ and this note tries to record the details of this argument.

To simplify the presentation, we limit ourselves to explain how to represent the suspension operation as an operation in $U \to U$ where U is some univalent universe in the simplicial set model.

Another goal of this note is to record the fact that the ideas used for the cubical set model, presented axiomatically in [3], can also be combined with [4] to give a quite simple way to build the standard Quillen model structure on *simplicial* sets. The use of classical logic is limited to one point (for establishing the logical equivalence of the "uniform" notion of Kan fibration with the usual definition).

1 A reformulation of Kan composition

Like in [3], we use the internal language of presheaf models. We write I, J, K, \ldots the objects of the base category (nonempty finite linear posets). We write \mathbb{F} the presheaf such that $\mathbb{F}(I)$ is the set of decidable sieves on I. (If the metalanguage is classical then $\mathbb{F} = \Omega$ is the subobject classifier, but we want here to limit the use of classical logic to one place.) We let \mathbb{I} be the present Δ^1 . Internally, \mathbb{I} has a (bounded) distributive lattice structure and the can follow the setting of [1, 3]. We have a dependent type $[\psi]$ for $\psi : \mathbb{F}$ where $[\psi]$ is given by $[\psi](I, S) = \{0 \mid 1_I \in S\}$. A filling operation for a dependent type A over a type Γ is an operation which given $\gamma : \Gamma^{\mathbb{I}}$ and $\psi : \mathbb{F}$ and a partial section in $\Pi(i : \mathbb{I})[\psi \lor i = b] \to A\gamma(i)$ with b = 0 or 1, extends it to a total section in $\Pi(i : \mathbb{I})A\gamma(i)$. We write Fill(Γ, A) the type of such operations.

If Γ is the terminal object we get the notion of *fibrancy* structure. A fibrancy structure on a presheaf X is an operation which, given $\psi : \mathbb{F}$ and a partial section in $\Pi(i : \mathbb{I})[\psi \lor i = b] \to X$ extends it to a total section in $X^{\mathbb{I}}$. We write $\mathsf{Fib}(X)$ the type of such operations.

In general to give only $\Pi(\rho: \Gamma)\mathsf{Fib}(A\rho)$ is weaker than to give an element in $\mathsf{Fib}(\Gamma, A)$.

In [1] we present a type of transport structure, which together with an element in $\Pi(\rho: \Gamma)\mathsf{Fib}(A\rho)$ produces an element in $\mathsf{Fib}(\Gamma, A)$. It is given by an operation which give ψ and $\gamma: \Gamma^{\mathbb{I}}$ which is constant on ψ and a partial section in $\Pi(i:\mathbb{I})[\psi \lor i = b] \to A\gamma(i)$ which is constant on ψ , extends it to a total section in $\Pi(i:\mathbb{I})A\gamma(i)$.

2 Suspension operation

Given X, we define a Susp X-algebra to be a type A with a fibrancy structure h_A and two points n_A, s_A and a family of paths $l_A : X \to A^{\mathbb{I}}$ connecting n_A to s_A . There is a natural notion of Susp X-algebra and we can show [1] externally² that for any X there exists an initial Susp X-algebra denoted simply by Susp X. It has three constructors N, S: Susp X and merid x i: Susp X for x: X and i: I.

¹This discussion can be found at https://groups.google.com/d/msg/homotopytypetheory/bNHRnGiF5R4/3RYz1YFmBQAJ. ²Thanks to one referee for pointing out to us that we don't need any special property of the interval for this operation,

Theorem 2.1 suspX satisfies the dependent elimination rule for the suspension: given a family of type P over suspX with a composition structure, and n in P N and s in P S and l x i in P (merid x i) such that l x 0 = n and l x 1 = s there exists a map elim : $\Pi(x : \text{susp}X)P x$ such that elim N = n and elim S = s and elim (merid x i) = l x i.

If A is a dependent type over Γ then we define Susp A by $(\text{Susp } A)\rho = \text{Susp } (A\rho)$. It is then possible to show [1] that from any element in Fill(Γ , A) we can build an element in Fill(Γ , Susp A).

Let \mathcal{U} be a Grothendieck universe. If A is \mathcal{U} -valued presheaf on the category of elements of Γ , then so is Susp A.

3 Kan fibration and simplicial set model

We can relate this internal notion of filling structure to the usual notion of Kan fibration.

Theorem 3.1 If Γ is a presheaf and A a presheaf on the category of elements of Γ then the following conditions are equivalent

- 1. $\Gamma.A \rightarrow \Gamma$ is a Kan fibration
- 2. $\Gamma.A \to \Gamma$ has the right lifting property w.r.t. any pushout product of a monomorphism and an endpoint inclusion in Δ^1
- 3. there exists an element in $Fill(\Gamma, A)$.

Proof. The equivalence between the two first points is a classic result in the theory of simplicial sets (e.g. Goerss-Jardine, Proposition 4.2; this is the only place where one uses classical logic, and more precisely decidability of degeneracy and axiom of choice). The equivalence of the second and third conditions is proved elegantly in [2], by using the notion of Leibnitz product and exponential.

We introduce the following notation $\mathsf{Type}_0(\Gamma)$ is the set of \mathcal{U} -valued presheaves on the category of elements of Γ , and $\mathsf{FType}_0(\Gamma)$ is the set of \mathcal{U} -valued presheaves on the category of elements of Γ together with a filling structure and $\mathsf{KType}_0(\Gamma)$ is the subpresheaf of $\mathsf{Type}_0(\Gamma)$ of \mathcal{U} -valued presheaves on the category of elements of Γ for which there exists a filling structure. By Theorem 3.1, $\mathsf{KType}_0(\Gamma)$ is equivalently the set of \mathcal{U} -valued presheaves A on the presheaf Γ such that $\Gamma.A \to \Gamma$ is a Kan fibration.

All of these define presheaves on the category of presheaves in a canonical way.

We define U(I) to be the set $\mathsf{KType}_0(Yon(I))$. We define a presheaf El on the category of elements of U by $El(I, X) = X(I, 1_I)$, so that El is \mathcal{U} -valued. If A is a \mathcal{U} -valued presheaf on the category of elements of Γ with a filling structure, there exists a unique map $|A| : \Gamma \to U$ such that El|A| = A.

It can be shown that $U.El \to U$ is a Kan fibration. By Theorem 3.1, we have a global element in Fill(U, El) (this is the only place where classical logic is used).

We expand the difference between this model (where filling is a property) and the "cubical" set models (where filling is a structure). If we define V(I) to be the set $\mathsf{FType}(Yon(I))$, then this defines a presheaf, but there will not be a natural bijection between $\Gamma \to V$ and $\mathsf{FType}_0(\Gamma)$. On the contrary, when we define U(I) as above, then there is a natural bijection between the set $\Gamma \to U$ and $\mathsf{KType}_0(\Gamma)$. We have a map $\mathsf{susp}: U \to U$ such that $El(\mathsf{susp} X) = \mathsf{Susp}(ElX)$ for X: U.

To show that we have an elimination rule, we proceed as for showing the existence of the elimination rule for the identity type in the simplicial set model. We consider the context (using extension types notation)

 $X: U, \ P: El(\mathsf{susp}\ X) \to U, \ n: P \mathsf{N}, \ s: P \mathsf{S}, \ l: \Pi(x: ElX) \Pi(i:\mathbb{I}) P \text{ (merid } x \text{ } i)[i=0 \mapsto n, \ i=1 \mapsto s]$

and in this context (like in [1]) we build an element in elim : $\Pi(x : El(susp X))P x$ such that elim N = nand elim S = s and elim (merid x i) = l x i.

4 Quillen model structure on simplicial sets and one conjecture

Using [3], it can be shown that U is fibrant (i.e. has a fibrant structure). We can follow [4] and build a model structure on simplicial sets. The only use of classical logic is in the existence of a filling structure in Fill(U, El) which is a consequence of Theorem 3.1.

It follows also from what we presented that it is possible to interpret in simplicial sets the version of cubical type theory (based on distributive lattice) where the composition operation are new *constants* (they satisfy only the substitution laws but no computation rules for composition of dependent products and sums and paths and universes). This system has also models in cubical sets (where we can compute). A conjecture is that it should be possible to use the glueing technique (as in [5]) to show that this formal system satisfies Voevodsky's conjecture: any closed term of type natural numbers should be path equal to a numeral. This would be one way to show that various versions of cubical type theory (that are extensions of this "constant" version by new computation rules) give the same values for a closed term of type natural numbers in ordinary dependent type theory extended with univalence.

References

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