

# 1 Graph definitions

## The elements of a graph

**Definition 1** (Graph). A graph  $G = (V, E)$  is a tuple where  $V$  is a set of nodes and  $E$  is a set of edges. Each edge  $e$  is a pair of nodes  $(i, j)$ , with  $i, j \in V$ . If  $(i, j) \in E$  then nodes  $i, j$  are *connected*.

*Remark 1.* If the pairs are ordered, then the graph is directed and edges are depicted with arrows. Source nodes are called *parents* and sink nodes *children*. Otherwise it is undirected, and edges are depicted with lines.

Glossary	
graph	graphe
node	nœud
edge	arête, lien, arc
directed graph	graphe orienté
undirected graph	graphe non-orienté

## Walks, paths and trails

**Definition 2** (Walk (promenade)). A walk from node  $k_1$  to node  $k_n$  is a sequence of nodes  $k_1, \dots, k_n, j$ , where  $(k_t, k_{t+1}) \in E$ .

**Definition 3** (Trail (chemin non-elementaire)). A trail from node  $k_1$  to node  $k_n$  is a walk with no edge  $(k_t, k_{t+1})$  appearing twice.

**Definition 4** (Path (chemin elementaire)). A path from node  $k_1$  to node  $k_n$  is a walk with no node in  $k_2, \dots, k_{n-1}$  appearing twice.

**Definition 5** (Connected graph). A graph is connected when there is a path between any two pairs of nodes.

**Definition 6** (Cycle). A cycle is a path  $k_1, \dots, k_n, k_1$ .

**Definition 7** (Tree). A tree is a connected graph with no cycles.

*Remark 2* (Isomorphism). The same graph can be depicted in many different ways geometrically. This is called isomorphism.

## Assigning meaning to graphs

- By itself, the graph structure may not be very useful.
- We can assign different properties to nodes and edges

## Graph applications

There are many graph applications, as graphs can be used to represent arbitrary relations between things. Examples are

### Example graphs

- Chemical structures
- Biological/Computer/Social networks
- Probabilistic graphical models
- Decision diagrams.
- Preference graphs.
- Decision trees.

### Preferences

**Definition 8** (Preference graph). A preference graph is a graph  $G = (V, E)$  such that  $x \rightarrow y$  if and only if we prefer  $x$  to  $y$ .

If, for any two nodes  $x, y \in V$ , there is a *directed path* from  $x$  to  $y$  or from  $y$  to  $x$ , then we have a total preference order, i.e. we can always say whether we prefer  $x$  to  $y$ .

**Example 1.** Elections where voters vote for a preference list of candidates are of this type. However, combining voters' preferences in a "fair" way is not possible.

### Graph problems

- Shortest path
- Matching
- Largest clique
- Colouring problem
- ...

## 2 Graph characteristics

**Definition 9** (Degree). For undirected graphs, the degree of a node  $i$  is defined as

$$\deg(i) = \sum_{(i,j) \in E} 1,$$

i.e. the number of edges containing  $i$ .

**Definition 10** (In-degree and out-degree). For directed graphs, the indegree of a node  $i$  is defined as

$$\deg^+(i) = \sum_{(j,i) \in E} 1, \quad \deg^-(i) = \sum_{(i,j) \in E} 1,$$

i.e. the number of incoming and outgoing edges to and from  $i$ .

### 3 Random graphs

#### Random graphs

**Definition 11** (Erdős-Rényi Model). A graph  $G$  generated from an  $\mathcal{ER}(n, M)$  model is chosen uniformly from all graphs with  $n$  nodes and  $M$  edges.

**Definition 12** (Gilbert model). A graph  $G$  generated from an  $\mathcal{G}(n, p)$  has  $n$  nodes and the probability of any edge  $(i, j)$  being in the graph is  $p$ .

*Remark 3.* The Gilbert model is equivalent to an ER model where  $M$  is also chosen randomly with probability

$$p^M(1-p)^{\binom{n}{2}-M}.$$

### 4 Graph representations

#### The Laplacian of a graph

**Definition 13** (Graph Laplacian). The Laplacian of graph  $G = (V, E)$  with  $n$  vertices is the  $n \times n$  matrix  $L$ :

$$L_{i,j} \triangleq \begin{cases} \deg(i), & i = j \\ -1, & i \neq j, e(i,j) \in E \\ 0, & \text{otherwise} \end{cases}$$