

Programming and Program Verification with Dependent Types in Agda

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Type Systems

- Basic types `Int`, `Word` from machine (CPU)
- Data types (record, objects, variants) describe structures
- Higher types (function, polymorphism, monads) describe computations
- Strong typing: check at compile-time
 - 1 Enable code optimizations
 - 2 Catch errors early
 - 3 Document code and libraries
 - 4 Allow (partial) **verification**

Types for verification

- Types for physical units (F#)
- Polymorphic types: A more general type has less inhabitants
- Security types: Control access to resources
- Resource types: Control complexity (time/space)
- **Dependent types**
 - Express arbitrary (logical) properties of a function in its type
 - Pre- and postconditions
- Vision: dependent types as universal type system
- Lightweight incorporation of security, resource, . . . types

Dependent types

- Dependent function type Π :

$$(x : A) \rightarrow B \ x$$

Codomain B can depend on argument value x .

- Dependent record type Σ :

record Seq (A : Set) : Set where
field

length : \mathbb{N}

elements : Vec A length

The type of a field can depend on the value of a previous field.

- **Full** dependent types:

Sum : $\mathbb{N} \rightarrow \text{Set}$

Sum 0 = \mathbb{N}

Sum (n + 1) = $\mathbb{N} \rightarrow \text{Sum } n$

Shape of type can depend on value.

Curry-Howard-Correspondence

- Constructive propositional logic corresponds to simply-typed lambda-calculus

Logic		Programming	
Proposition		Type	
Proof		Program	
Implication	\Rightarrow	Function type	\rightarrow
Conjunction	\wedge	Cartesian product	\times
Disjunction	\vee	Disjoint sum	$+$
Truth	\top	Biggest type (unit)	\top
Absurdity	\perp	Least type (empty)	\perp

- One language for types and specifications
- One language for programming and proving

Curry-Howard for Predicate Logic

- Constructive predicate logic corresponds to dependently-typed lambda-calculus (Martin-L" of Type Theory)

Logic		Programming	
Proposition		Type	
Proof		Program	
Universal q.	$\forall x : \tau. P x$	Depend. function t.	$(x : A) \rightarrow B x$
Existential q.	$\exists x : \tau. P x$	Depend. pair t.	$(x : A) \times B x$
Predicates	$P t$	Indexed type	$A t$
Type	τ	Type	A
Term	t	Program	t

- Unification of concepts \Rightarrow lean language.
- Program extraction: discard proofs

Some Dependently Typed Languages

- NuPrl: Extensional Type Theory (ETT)
Cornell, US (Bob Constable, ...) 1980s–
- Coq: Impredicative Intensional Type Theory (ITT)
INRIA, France (Huet, Coquand, ...) 1980s–
- Epigram: Predicative ITT
McKinna, McBride 2000s–
- Agda: Predicative ITT
Chalmers, Sweden 1990s–: Alf, Alfa, Agda, AgdaLight
Agda 2 (Norell, ...) 2006–

Agda 2

- Haskell-like syntax +
 - mixfix identifiers
 - Parametrized modules
 - People
 - Implemented by Ulf Norell
 - Library by Nils Anders Danielsson (Nisse)
 - Maintained by Ulf, Nisse, and me
 - Contributions by 20 more...
 - Compiles to
 - Haskell (M. Takeyama)
 - C (via Epic) (Gustafsson)
 - JavaScript (A. Jeffrey)
 - Agda is implemented in Haskell
- ```
cabal install Agda
```

# Agda tour (interactive)

- Numbers
- Vectors
- Logic
- Sorted lists

# Theory of Agda

- Type checking = proof checking
- Termination
- Automation
  - Higher-order unification
  - Program synthesis = proof search
  - Overloading via instance search (type classes)
- Program extraction
  - Identify computationally irrelevant terms (proofs)
  - Eliminate uniquely determined terms (singletons)
  - Exploit rich type information for optimizations

# Termination

- Looping programs prove anything

`loop` :  $\perp$

`loop` = `loop`

- Agda accepts only terminating programs
- Current termination checker searches for lexicographic structural descent
- Plan: certify termination by sized types
  - Type-based termination (Pareto 1996, Giménez 1998, ...)
  - My PhD thesis: non-dependent language + sized types
  - MiniAgda: study sized types and pattern matching
- Coinduction: non-terminating, but **productive** programs
  - Induction: consume input in each iteration
  - Coinduction: produce output in each iteration

# Consistency

- Logical consistency proved by model
- Term model: Model types by sets of terminating programs
- Proves also decidability of type checking
- No model for all of Agda published
  - Dependently-typed lambda-calculus
  - Inductive types
  - Indexed types
  - Recursive functions and pattern matching
  - Universes
- Current status: extrapolation from models that consider a subset of features
- Plan: simpler model through sized types

# Unification

- Type-checking decidable, but infeasible
- Excessive amount of (repetitive) type arguments

$$\_ \circ \_ : \{A B C : \text{Set}\} \rightarrow (B \rightarrow C) \rightarrow (A \rightarrow B) \rightarrow A \rightarrow C$$

$$g \circ f = \lambda x \rightarrow g (f x)$$

$$\_ \circ \_ \{A\} \{B\} \{C\} g f = \lambda x \rightarrow g (f x)$$

- Reconstruction: Omit arguments, solve by unification
- Higher-order unification = solve equations over lambda-terms
- Undecidable (Hilbert's 10th problem)  $\Rightarrow$  constraints
  - Agda: Sound and complete constraint simplification
  - Only finds unique solutions
  - Coq: more solving by first-order unification, but also wrong solutions

# Higher-Order Unification

- Some constraints have no unique solutions

$$\lambda x \text{ true} = \text{true}$$

$$\lambda x \ x = x + 1$$

- Postpone constraints, try to gather more information

- Unification for records:

*Abel, Pientka, TLCA 2011*

- Beware of type checking modulo constraints:

If constraints are unsolvable, term might be ill-typed/looping

- Satisfying theoretical account lacking!

## Computational irrelevance

- Declare some programs as proofs by typing
$$f : \cdot (x : A) \rightarrow B \ x$$
 $f$  and  $B$  do not computationally depend on  $x$
- $\cdot A$  is “squash  $A$ ”.
- Irrelevant arguments can be erased during program extraction.

## Conclusion

- Agda: Unified proof/programming language via Curry-Howard
- Lean syntax, but complicated theory
- Vision: universal type system
- Subsume the next 700 type systems
- Need: more automation for boilerplate!
- Tactics!

## Get involved!

- Try Agda!
- Formalize your papers with Agda!
- Write libraries!
- Write documentation (Agda wiki)!
- Fix bugs [code.google.com/agda/issues](https://code.google.com/agda/issues/) | !
- Contribute features!
- Contribute theory!