

Formal Methods for Software Development

Reasoning about Programs with Dynamic Logic

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(JAVA) Dynamic Logic

Typed FOL

- ▶ + (JAVA) programs p
- ▶ + modalities $\langle p \rangle \phi$, $[p] \phi$ (p program, ϕ DL formula)
- ▶ + ... (later)

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Remark on Hoare Logic and DL

In Hoare logic $\{Pre\} p \{Post\}$

(Pre, Post must be FOL)

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In Hoare logic $\{Pre\} p \{Post\}$

(Pre, Post must be FOL)

In DL $Pre \rightarrow [p]Post$

(Pre, Post any DL formula)

Proving DL Formulas

An Example

```
∀ int x;  
(x ≥ 0 ∧ n = x →  
  [ i = 0; r = 0;  
    while(i < n){i = i + 1; r = r + i;}  
    r = r + r - n;  
  ]r = x * x)
```

How can we prove that the above formula is valid
(i.e. satisfied in all states)?

Semantics of DL Sequents

$\Gamma = \{\phi_1, \dots, \phi_n\}$ and $\Delta = \{\psi_1, \dots, \psi_m\}$ sets of DL formulas where all logical variables occur bound.

Recall: $\mathcal{S} \models (\Gamma \Rightarrow \Delta)$ iff $\mathcal{S} \models (\phi_1 \wedge \dots \wedge \phi_n) \rightarrow (\psi_1 \vee \dots \vee \psi_m)$

Define semantics of DL sequents identical to semantics of FOL sequents

Definition (Validity of Sequents over DL Formulas)

A sequent $\Gamma \Rightarrow \Delta$ over DL formulas is **valid** iff

$$\mathcal{S} \models (\Gamma \Rightarrow \Delta) \text{ in all states } \mathcal{S}$$

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Consequence for program variables

Initial value of program variables implicitly “**universally** quantified”

Symbolic Execution of Programs

Sequent calculus decomposes top-level operator in formula.
What is “top-level” in a sequential program $p; q; r; ?$

Symbolic Execution

- ▶ Follow the **natural control flow** when analysing a program
- ▶ Values of some variables unknown: **symbolic state representation**

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Example

Compute the final state after termination of

$x = x + y; y = x - y; x = x - y;$

Symbolic Execution of Programs Cont'd

Typical form of DL formulas in symbolic execution

$$\langle \text{stmt}; \text{rest} \rangle \phi \quad [\text{stmt}; \text{rest}] \phi$$

- ▶ Rules symbolically execute *first* statement (“**active statement**”)
- ▶ Repeated application of such rules corresponds to **symbolic program execution**

Symbolic Execution of Programs Cont'd

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Example (`symbolicExecution/simpleIf.key`,
Demo, active statement only)

```
\programVariables {  
  int x; int y; boolean b;  
}  
\problem {  
  \<{ if (b) { x = 1; } else { x = 2; } y = 3; }\> y > x  
}
```

Symbolic Execution of Programs Cont'd

Symbolic execution of conditional

$$\text{if } \frac{\Gamma, b = \text{TRUE} \Rightarrow \langle p; \text{rest} \rangle \phi, \Delta \quad \Gamma, b = \text{FALSE} \Rightarrow \langle q; \text{rest} \rangle \phi, \Delta}{\Gamma \Rightarrow \langle \text{if } (b) \{ p \} \text{ else } \{ q \} ; \text{rest} \rangle \phi, \Delta}$$

Symbolic execution must consider all possible execution branches

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Symbolic execution must consider all possible execution branches

Symbolic execution of loops: unwind

$$\text{unwindLoop} \frac{\Gamma \Rightarrow \langle \text{if } (b) \{ p; \text{while } (b) p \}; \text{rest} \rangle \phi, \Delta}{\Gamma \Rightarrow \langle \text{while } (b) \{ p \}; \text{rest} \rangle \phi, \Delta}$$

Updates for KeY-Style Symbolic Execution

Needed: a Notation for Symbolic State Changes

- ▶ Symbolic execution should “walk” through program in natural **forward** direction
- ▶ Need **succinct representation** of state changes, effected by each symbolic execution step
- ▶ Want to **simplify** effects of program execution **early**
- ▶ Want to **apply** state changes **late**
(to branching conditions and post condition)

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We use dedicated notation for state changes: **updates**

Explicit State Updates

Definition (Syntax of Updates, Updated Terms/Formulas)

If v is program variable, t FOL term type-conformant to v , t' any FOL term, and ϕ any DL formula, then

- ▶ $\{v := t\}$ is an update
- ▶ $\{v := t\}t'$ is DL term
- ▶ $\{v := t\}\phi$ is DL formula

Definition (Semantics of Updates)

State \mathcal{S} interprets program variables v with $\mathcal{I}_{\mathcal{S}}(v)$, β variable assignment for logical variables in t , define semantics ρ as:

$$\rho_{\beta}(\{v := t\})(\mathcal{S}) = \mathcal{S}' \text{ where } \mathcal{S}' \text{ identical to } \mathcal{S} \text{ except } \mathcal{I}_{\mathcal{S}'}(v) = \text{val}_{\mathcal{S},\beta}(t)$$

Explicit State Updates Cont'd

Facts about updates $\{v := t\}$

- ▶ Update semantics similar to that of assignment
- ▶ Value of update also depends on \mathcal{S} and **logical** variables in t , i.e., β
- ▶ Updates are **not assignments**: right-hand side is FOL term
 - $\{x := n\}\phi$ cannot be turned into assignment if n is a logical variable
 - $\{x=i++;\}\phi$ cannot (immediately) be turned into update
- ▶ Updates are **not equations**: they **change** value of v

Computing Effect of Updates (Automated)

Rewrite rules for update followed by ...

program variable $\left\{ \begin{array}{l} \{x := t\}x \rightsquigarrow t \\ \{x := t\}y \rightsquigarrow y \end{array} \right.$

logical variable $\{x := t\}w \rightsquigarrow w$

complex term $\{x := t\}f(t_1, \dots, t_n) \rightsquigarrow f(\{x := t\}t_1, \dots, \{x := t\}t_n)$

atomic formula $\{x := t\}p(t_1, \dots, t_n) \rightsquigarrow p(\{x := t\}t_1, \dots, \{x := t\}t_n)$

FOL formula $\left\{ \begin{array}{l} \{x := t\}(\phi \ \& \ \psi) \rightsquigarrow \{x := t\}\phi \ \& \ \{x := t\}\psi \\ \dots \\ \{x := t\}(\forall \tau y; \phi) \rightsquigarrow \forall \tau y; (\{x := t\}\phi) \end{array} \right.$

program formula No rewrite rule for $\{x := t\}\langle prog \rangle\phi$

Substitution delayed until *prog* symbolically executed

Assignment Rule Using Updates

Symbolic execution of assignment using updates

$$\text{assign} \frac{\Gamma \Rightarrow \{x := t\} \langle rest \rangle \phi, \Delta}{\Gamma \Rightarrow \langle x = t; rest \rangle \phi, \Delta}$$

- ▶ Works as long as t is 'simple' (has no side effects)
- ▶ For every built-in Java operation, we need a separate rule (for $x = t_1 + t_2$ and $x = t_1 - t_2$ etc.)

Demo

updates/assignmentToUpdate.key

How to apply updates on updates?

Example

Symbolic execution of

```
t=x; x=y; y=t;
```

yields:

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{t := x}{x := y}{y := t}
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Need to compose three sequential state changes into a single one:

Parallel Updates

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parallel updates

Parallel Updates Cont'd

Definition (Parallel Update)

A **parallel update** has the form $\{v_1 := r_1 \parallel \dots \parallel v_n := r_n\}$, where each $\{v_i := r_i\}$ is simple update

- ▶ All r_i computed in **old state** before update is applied
- ▶ Updates of all program variables v_i executed **simultaneously**
- ▶ Upon **conflict** $v_i = v_j, r_i \neq r_j$ later update ($\max\{i, j\}$) wins

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$$\{v_1 := r_1\}\{v_2 := r_2\} = \{v_1 := r_1 \parallel v_2 := \{v_1 := r_1\}r_2\}$$

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$$\{v_1 := r_1 \parallel \dots \parallel v_n := r_n\}x = \begin{cases} x & \text{if } x \notin \{v_1, \dots, v_n\} \\ r_k & \text{if } x = v_k, x \notin \{v_{k+1}, \dots, v_n\} \end{cases}$$

Symbolic Execution with Updates (by Example)

$\Rightarrow x < y \rightarrow \langle t=x; x=y; y=t; \rangle y < x$

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Parallel Updates Cont'd

Demo

updates/swap1.key

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Example

symbolic execution of `x=x+y; y=x-y; x=x-y;` gives

$$(\{x := x+y\}\{y := x-y\})\{x := x-y\}$$

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Parallel updates store intermediate state of symbolic computation

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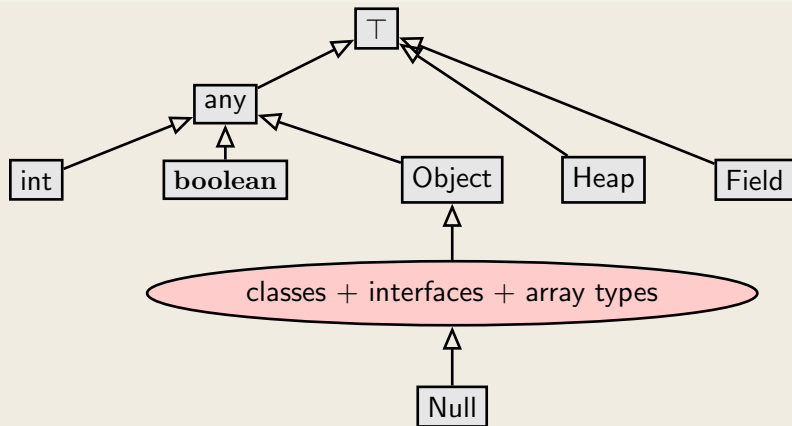
Instead

Quantify over **value**, and **assign** it to program variable:

$\forall \tau x; \{i := x\} \langle \dots i \dots \rangle \phi$

Modelling Java in FOL: Fixing a Type Hierarchy

Signature based on Java's type hierarchy



Each interface and class in API and in target program becomes type with appropriate subtype relation

Modelling the Heap in FOL

The Java Heap

Objects are stored on (i.e., in) the **heap**.

- ▶ Status of heap changes during execution
- ▶ Each heap associates values to object/field pairs

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Each element of data type Heap represents a certain heap status.

Two functions involving heaps:

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`store(h, o, f, v)` returns heap like h , but with v associated to $o.f$
- ▶ in F_Σ : any `select(Heap, Object, Field)`;
`select(h, o, f)` returns value associated to $o.f$ in h

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Modelling instance fields

Person
<code>int age</code> <code>int id</code>
<code>int setAge(int newAge)</code> <code>int getId()</code>

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Reading Field `id` of `Person p`

FOL notation `select(h, p, id)`

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Key notation `p.id@h` (abbreviating `select(h, p, id)`)
`p.id` (abbreviating `select(heap, p, id)`)^a

^a`heap` is special program variable for “current” heap; mostly implicit in `o.f`

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KeY notation $h[p.\text{id} := 6238]$ (notation for store, not update)

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We formalise the *behaviour*, with an algebra of heap operations:

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$$(o \neq o' \vee f \neq f') \rightarrow \text{select}(\text{store}(h, o, f, x), o', f') = \text{select}(h, o', f')$$

Example

$$\text{select}(\text{store}(h, o, f, 15), o, f) \rightsquigarrow 15$$

$$\text{select}(\text{store}(h, o, f, 15), o, g) \rightsquigarrow \text{select}(h, o, g)$$

$$\text{select}(\text{store}(h, o, f, 15), u, f) \rightsquigarrow$$

$$\text{if } (o = u) \text{ then } (15) \text{ else } (\text{select}(h, u, f))$$

Shorthand Notations for Heap Operations

<code>o.f@h</code>	is	<code>select(h, o, f)</code>
<code>h[o.f := v]</code>	is	<code>store(h, o, f, v)</code>

Shorthand Notations for Heap Operations

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Pretty Printing

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$h[o.f := v][o'.f' := v']$ is `store(store(h, o, f, v), o', f', v')`

Very-Shorthand Notations for **Current** Heap

Current heap always in special variable `heap`.

$o.f$ is `select(heap, o, f)`

$\{o.f := v\}$ is `update {heap := heap[o.f := v]}`

Modelling the Heap in FOL—The Full Story

Is formula $\text{select}(h, p, \text{id}) \geq 0$ type-safe?

Modelling the Heap in FOL—The Full Story

Is formula `select(h, p, id) >= 0` type-safe?

1. Return type is any—need to 'cast' to `int`
2. There can be many fields with name `id`

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Real Field Access

`int::select(h, p, Person::$id) >= 0` is type-safe

- ▶ `int::select` is a function name, not a cast
- ▶ can be understood *intuitively* as `(int)select`

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General

For each `T` typed field `f` of class `C`, F_{Σ} contains

- ▶ a constant declared as `Field C::$f`
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Everything **blue** is a function name

Field Update Assignment Rule

Changing the value of fields

How to (symbolically) execute assignment to field, e.g., `p.age=18;` ?

$$\text{assign} \frac{\Gamma \Rightarrow \{o.f := t\} \langle rest \rangle \phi, \Delta}{\Gamma \Rightarrow \langle o.f = t; rest \rangle \phi, \Delta}$$

Admit on left-hand side of update **JAVA location expressions**

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But is this rule correct? See below.

Dynamic Logic: KeY input file

```
\javaSource "path to source code referenced in problem";  
  
\programVariables { Person p; }  
  
\problem {  
    \<{    p.age = 18;    }\> p.age = 18  
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Demo

updates/firstAttributeExample.key

Refined Semantics of Program Modalities

Does abrupt termination count as normal termination?

No! Need to distinguish **normal** and **exceptional** termination

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Refined Semantics of Program Modalities

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- ▶ $[p] \phi$: If p terminates **normally** then formula ϕ holds in final state (partial correctness)

Abrupt termination on top-level counts as non-termination!

Example Reconsidered: Exception Handling

```
\javaSource "path to source code";  
  
\programVariables {  
  ...  
}  
  
\problem {  
  p != null -> \<{ p.age = 18; }\> p.age = 18  
}
```

Only provable when no top-level exception thrown

Demo

updates/secondAttributeExample.key

The Self Reference

Modeling reference `this`

Special name for the object whose JAVA code is currently executed:

in JML: Object `this`;

in Java: Object `this`;

in KeY: Object `self`;

Default assumption in JML-KeY translation: `self != null`

Which Objects do Exist?

How to model **object creation** with **new** ?

Which Objects do Exist?

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Constant Domain Assumption

Assume that domain \mathcal{D} is the same in all states $(\mathcal{D}, \delta, \mathcal{I}) \in States$

Consequence:

Quantifiers and modalities commute:

$$\models (\forall T x; [p]\phi) \leftrightarrow [p](\forall T x; \phi)$$

Object Creation (background; no need to remember this)

Realizing Constant Domain Assumption

- ▶ Implicitly declared field `boolean <created>` in class `Object`
- ▶ `<created>` has value `true` iff argument object has been created
- ▶ Object creation modeled as `{heap := create(heap, ob)}` for not (yet) created `ob` (essentially sets `<created>` field of `ob` to `true`)

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$$\frac{\Gamma, \text{ob. } \langle \text{created} \rangle = \text{FALSE} \Rightarrow \{ \text{heap} := \text{create}(\text{heap}, \text{ob}) \} \{ \text{o} := \text{ob} \} \langle \text{o.} \langle \text{init} \rangle (\text{param}); \omega \rangle \phi, \Delta}{\Gamma \Rightarrow \langle \text{o} = \text{new } T(\text{param}); \omega \rangle \phi, \Delta}$$

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`ob` is a fresh program variable

Alternatives exist in the literature. E.g.:

[Ahrendt, de Boer, Grabe, *Abstract Object Creation in Dynamic Logic – To Be or Not To Be Created*, Springer, LNCS 5850]

Titlepage

Symbolic Execution

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Parallel Updates

Modeling OO Programs

Self

Object Creation

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Abrupt Termination

Null Pointers

Aliasing

Summary

Literature

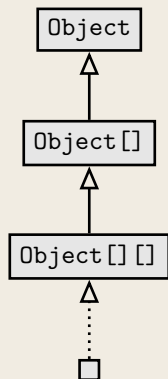
Dynamic Logic to (almost) full Java

KeY supports full **sequential** Java, with some limitations:

- ▶ Limited concurrency
- ▶ No generics
- ▶ No I/O
- ▶ No dynamic class loading or reflection
- ▶ Ongoing work to support floating point arithmetic
- ▶ API method calls: need either JML contract or implementation

Java Features in Dynamic Logic: Arrays

Arrays



- ▶ JAVA type hierarchy includes array types
- ▶ Types ordered according to JAVA subtyping rules
- ▶ Function $\text{arr} : \text{int} \rightarrow \text{Field}$ turns integer index into type Field (required in store).
- ▶ Store array elements on heap
- ▶ Value of $a[i]$ in heap $\text{store}(\text{heap}, a, \text{arr}(i), 8)$ is 8
- ▶ Arrays a and b can refer to same object (aliasing)

Java Features in Dynamic Logic: Complex Expressions

Complex expressions with side effects

- ▶ JAVA expressions may have **side effects**, due to method calls, increment/decrement operators, nested assignments
- ▶ FOL terms have **no** side effect on the state

Example (Complex expression with side effects in Java)

```
int i = 0; if ((i=2)>= 2) i++;   value of i ?
```


Complex Expressions Cont'd

Decomposition of complex terms by symbolic execution

Follow the rules laid down in JAVA Language Specification

Local code transformations

$$\text{evalOrderIteratedAssgnmt} \frac{\Gamma \Rightarrow \langle y = t; x = y; \omega \rangle \phi, \Delta}{\Gamma \Rightarrow \langle x = y = t; \omega \rangle \phi, \Delta} \quad t \text{ simple}$$

Temporary variables store result of evaluating subexpression

$$\text{ifEval} \frac{\Gamma \Rightarrow \langle \text{boolean } v0; v0 = b; \text{if } (v0) \text{ p}; \omega \rangle \phi, \Delta}{\Gamma \Rightarrow \langle \text{if } (b) \text{ p}; \omega \rangle \phi, \Delta} \quad b \text{ complex}$$

Java Features in Dynamic Logic: Abrupt Termination

Abrupt Termination: Exceptions and Jumps

Redirection of control flow via return, break, continue, exceptions

$$\langle \text{try } \{p\} \text{ catch}(T \ e) \{q\} \text{ finally } \{r\} \omega \rangle \phi$$

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Rule tryThrow matches **try-catch** in pre-/postfix and active throw

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Demo

exceptions/try-catch.key

Null pointer exceptions

There are no “exceptions” in FOL: \mathcal{I} total on FSym

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^aCan be changed with Taclet Option runtimeExceptions

Field Update Assignment Rule Revisited (A)

Changing the value of fields

How to (symbolically) execute assignment to field?

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π is the “inactive prefix”, any number of opening try blocks: $(\text{try}\{\})^*$

Field Update Assignment Rule Revisited (B)

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Revisit: Field Assignment Demo

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Demo

updates/firstAttributeExample.key

Java Features in Dynamic Logic: Aliasing

Demo

aliasing/attributeAlias1.key

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Demo

aliasing/attributeAlias1.key

Reference Aliasing

Alias resolution causes **proof split**

Summary

- ▶ Most JAVA features covered in KeY
- ▶ Degree of automation for loop-free programs is very high
- ▶ Handling of loops: last lecture

Literature for this Lecture

KeYbook *W. Ahrendt, B. Beckert, R. Bubel, R. Hähnle, P. Schmitt, M. Ulbrich, editors.*

Deductive Software Verification - The KeY Book

Vol 10001 of *LNCS*, Springer, 2016

(E-book at link.springer.com)

- ▶ B. Beckert, V. Klebanov, B. Weiß, **Dynamic Logic for Java**
Chapter 3 in [KeYbook]
on the surface only: Sections 3.1, 3.2, 3.4, 3.5.5, 3.5.6, 3.5.7, 3.6
- ▶ *W. Ahrendt, S. Grebing, Using the KeY Prover*
Chapter 15 in [KeYbook]