Parallelization and lock-free

programming

### AtomicInteger x = new AtomicInteger(0);

### thread t

```
int v;

1  do {
2    v = x.get();
3    v = v + 1;
4  } while (!x.compareAndSet(v - 1, v));
```

- 1. it is starvation free
- 2. it is lock free
- 3. it is lock free, and hence also wait free
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int v;

1  for (int i = 0; i < 10_000; i++) {
2    v = x.get();
3    v = v + 1;
4    if (x.compareAndSet(v - 1, v))
5       break;
6  }</pre>
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Note that the increment may fail after trying 10'000 times.

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- 1. the sum of elements in L
- 2. the maximum of elements in L
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How many tasks may execute in parallel when computing the factorial of n?

```
class Factorial extends RecursiveTask<Integer> {
  int n; // number to compute factorial of
  protected Integer compute() {
    if (n <= 1) return 1;
    Factorial f = new Factorial(n - 1);
    f.fork():
    return n * f.join();
 1. n! (the factorial of n)
 2. n
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- 3. it depends on the number of available cores
- 4. there is practically no parallelism

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How many tasks may execute in parallel when computing the sum of integers from 1 to n?

```
class Sum extends RecursiveTask<Integer> {
  int m, n; // sum integers from m to n
  protected Integer compute() {
    if (m > n) return 0;
    if (m == n) return m:
    int mid = m + (n-m)/2; // mid point
    Sum\ lower = new\ Sum(m, mid);
    Sum upper = new Sum(mid+1, n);
    lower.fork(); upper.fork();
    return lower.join() + upper.join();
 1. 2<sup>n</sup> (2 to the power of n)
 2. n<sup>2</sup> (the square of n)
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