Formal Methods for Software Development

Model Checking with Temporal Logic

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21st September 2018

Model Checking

Check whether a formula is valid in all runs of a transition system.

Given a transition system $\mathcal{T}$ (e.g., derived from a PROMELA program).

Verification task: is the LTL formula ϕ satisfied in all traces of $\mathcal{T}$, i.e.,

$$\mathcal{T} \models \phi \quad ?$$

LTL Model Checking—Overview

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When What?

this lecture

3.–5. product of transition system and Büchi automaton (construction and analysis)

When What?

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3.–5. product of transition system and Büchi automaton (construction and analysis)

next lecture

1. translating LTL into generalised Büchi automata
2. generalised Büchi automata and their normalisation

Product of Transition System and Büchi Automaton

A model checking graph is a directed graph with initial and accepting nodes.

Definition (Model Checking Graph)

A model checking graph $(N, \rightarrow, N_0, N_a)$ is composed of:

- finite, non-empty set of nodes N
- an 'arrow' relation $\rightarrow \subseteq N \times N$
- a non-empty set of initial nodes $N_0 \subseteq N$
- a set of accepting nodes $N_a \subseteq N$

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1. transition systems without terminal states:
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Can always be achieved by adding 'trap states' or 'trap locations', resp.

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We assume a set of atomic propostions AP.

Definition (Product of Transition System and Büchi Automaton)

Let $\mathcal{T} = (S, \rightarrow, S_o, L)$ be a transition system over AP and $\mathcal{B} = (Q, \delta, Q_0, F)$ be a Büchi automaton over the alphabet 2^{AP}.
Then, $\mathcal{T} \otimes \mathcal{B}$ is the following model checking graph:

$$\mathcal{T} \otimes \mathcal{B} = (S \times Q, \rightarrow', N_0, N_a)$$

where:

- $\langle s, q \rangle \rightarrow' \langle s', q' \rangle$ iff $s \rightarrow s'$ and $(q, L(s'), q') \in \delta$

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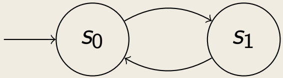
Model Checking Example

Assume $AP = \{red, green\}$

$\mathcal{T}$:



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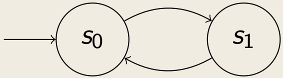

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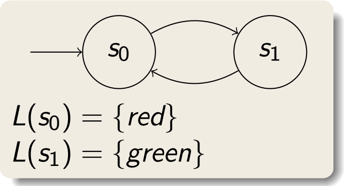

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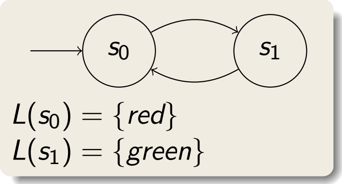

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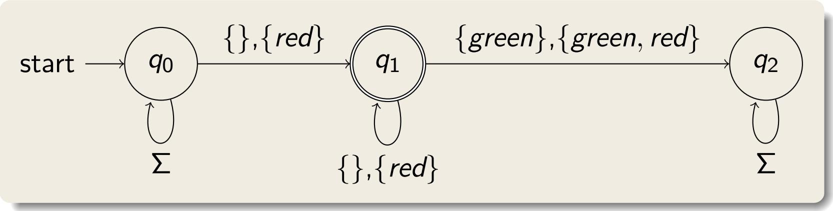

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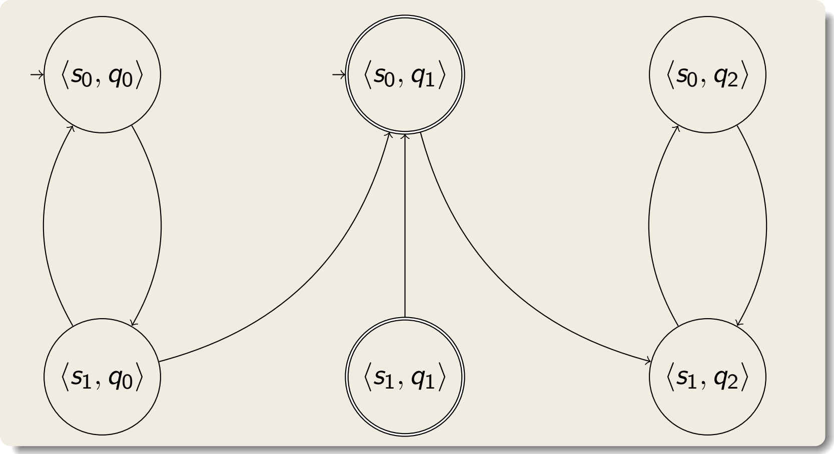

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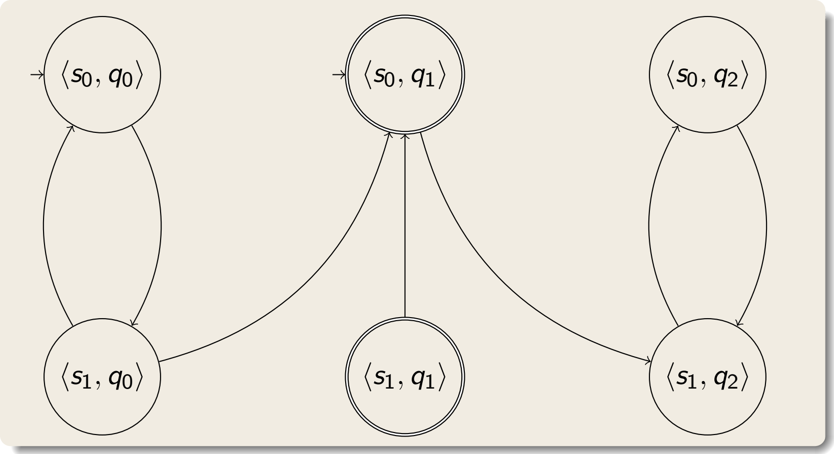

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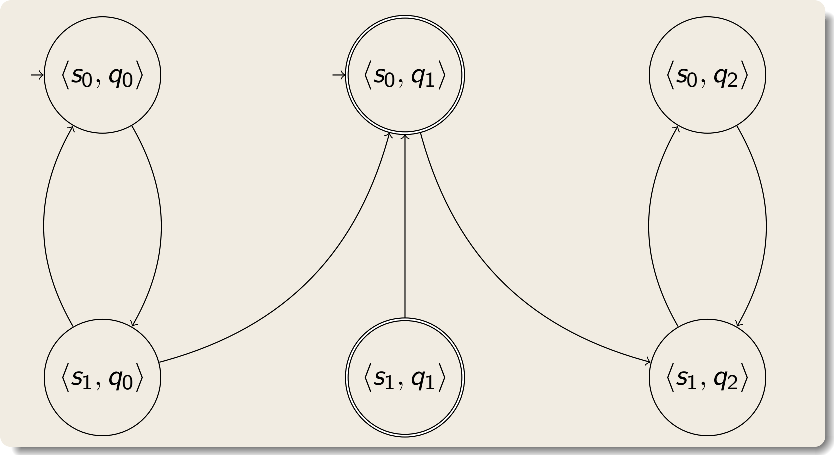

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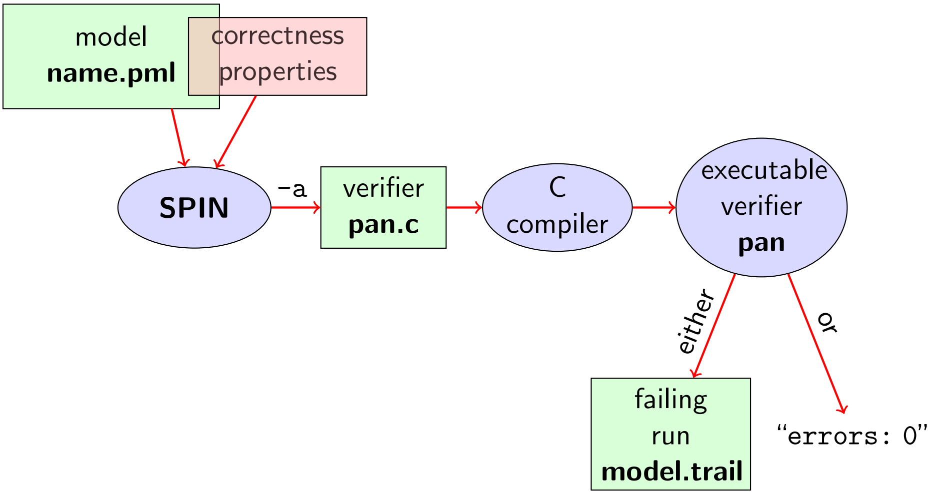
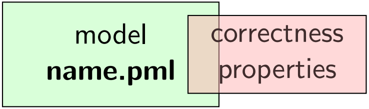
Model Checking with Spin



Stating Correctness Properties



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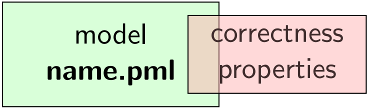

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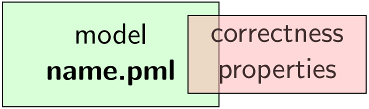

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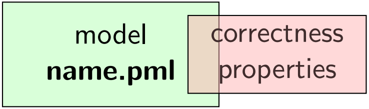

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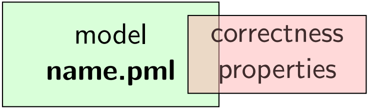

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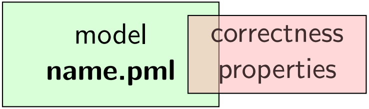

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Preliminaries

1. Accept labels in PROMELA $\leftrightarrow$ Büchi automata
2. Fairness

Preliminaries 1: Acceptance Cycles

Definition (Accept Location)

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Accept locations can be used to specify cyclic behavior

Definition (Acceptance Cycle)

A run which infinitely often passes through an accept location is called an acceptance cycle.

Acceptance cycles are mainly used in never claims (see below), to define (undesired) infinite behavior

Preliminaries 2: Fairness

Does this model terminate in each run?

Simulate: `start/fair.pml`

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byte n = 0;
bool flag = false;

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Definition (Weak Fairness)

A run is called weakly fair iff the following holds:
each continuously executable statement is executed eventually.

Model Checking of Temporal Properties

Many correctness properties not expressible by assertions

- All properties that involve state changes
- Temporal logic expressive enough to characterize many (but not all) Linear Time properties

In this course: "temporal logic" synonymous with "linear temporal logic"

Today: model checking of properties formulated in temporal logic

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these are temporal properties $\Rightarrow$ use temporal logic

Boolean Temporal Logic

Numerical variables in expressions

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- Propositional LTL as introduced so far only knows propositions
- Slight generalisation of LTL required

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In Boolean Temporal Logic, atomic building blocks are *Boolean expressions* over PROMELA variables

Boolean Temporal Logic over PROMELA

Set For_{BTL} of Boolean Temporal Formulas (simplified)

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Boolean Temporal Logic over Promela

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$$!\,\phi,\quad \phi\ \&\&\ \psi,\quad \phi\ ||\ \psi,\quad \phi\ {->}\ \psi,\quad \phi\ {<->}\ \psi$$

$$[\,]\phi,\quad <>\phi,\quad \phi\,\mathtt{U}\,\psi$$

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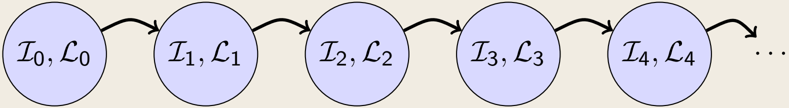
Semantics of Boolean Temporal Logic

A trace τ through a PROMELA model M



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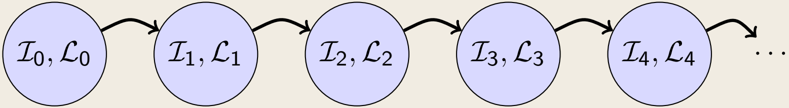

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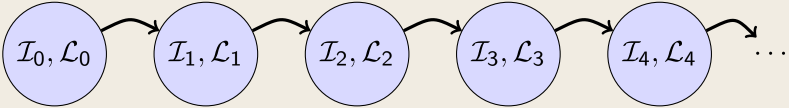

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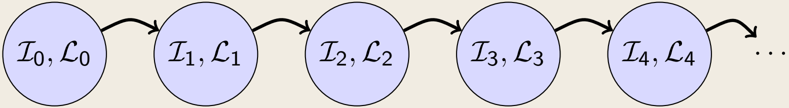

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Evaluating other formulas $\in For_{BTL}$ in traces τ: see previous lecture

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Example

TL formula `[](critical <= 1)`

"It is guaranteed throughout each run that at most one process visits its critical section at any time."

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- state that something 'good' is guaranteed throughout each run
- each violating run violates the property after *finitely* many steps

Example

TL formula `[](critical <= 1)`

"It is guaranteed throughout each run that at most one process visits its critical section at any time."

or, equivalently:
"It will never happen that more than one process visits its critical section."

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Example

TL formula `[](critical <= 1)`

"It is guaranteed throughout each run that at most one process visits its critical section at any time."

or, equivalently:
"It will never happen that more than one process visits its critical section."

Any violating run would have `(critical <= 1)` after *finite* time

Applying Temporal Logic to Critical Section Problem

We want to verify `[](critical<=1)` as a correctness property of:

```
active proctype P() {
  do :: /* non-critical activity */
        atomic {
          !inCriticalQ;
          inCriticalP = true
        }
        critical++;
        /* critical activity */
        critical--;
        inCriticalP = false
  od
}

/* similarly for process Q */
```

Model Checking a Safety Property with Spin

Command Line Execution

Add definition of TL formula to Promela file

Example `ltl atMostOne { [](critical <= 1) }`

General `ltl` *name* `{` *TL-formula* `}`

can define more than one formula

```
> spin -a file.pml
> gcc -DSAFETY -o pan pan.c
> ./pan -N name
```

Demo: `target/safety1.pml`

Model Checking a Safety Property with Spin

Command Line Execution

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General `ltl` *name* `{` *TL-formula* `}`

can define more than one formula

```
> spin -a file.pml
> gcc -DSAFETY -o pan pan.c
> ./pan -N name
```

Demo: `target/safety1.pml`

- The '`ltl` *name* `{` *TL-formula* `}`' construct must be part of your lab submission!

Model Checking a Safety Property with SPIN

Command Line Execution

Add definition of TL formula to PROMELA file

Example `ltl atMostOne { [](critical <= 1) }`

General `ltl` *name* `{` *TL-formula* `}`

can define more than one formula

```
> spin -a file.pml
> gcc -DSAFETY -o pan pan.c
> ./pan -N name
```

Demo: `target/safety1.pml`

- The '`ltl` *name* `{` *TL-formula* `}`' construct must be part of your lab submission!

`ltl` definitions not part of Ben Ari's book (SPIN$\leq$ 6): ignore 5.3.2, etc.

Model Checking a Safety Property using Web Interface

1. add definition of TL formula to PROMELA file

   Example `ltl atMostOne { [](critical <= 1) }`

   General `ltl` *name* `{` *TL-formula* `}`

   can define more than one formula
2. load PROMELA file into web interface
3. ensure `Safety` is selected
4. enter name of LTL formula in according field
5. select `Verify`

Demo: `safety1.pml`

Model Checking a Safety Property using jSpin

1. add definition of TL formula to Promela file

   Example `ltl atMostOne { [](critical <= 1) }`

   General `ltl` *name* { *TL-formula* }

   can define more than one formula
2. load Promela file into jSpin
3. write *name* in 'LTL formula' field
4. ensure `Safety` is selected
5. select `Verify`
   - (corresponds to command line `./pan -N` *name* `...`)
6. (if necessary) select `Stop` to terminate too long verification

Demo: `safety1.pml`

Temporal Model Checking without Ghost Variables

We want to verify mutual exclusion without using ghost variables.

```
bool inCriticalP = false, inCriticalQ = false;

active proctype P() {
  do :: atomic {
          !inCriticalQ;
          inCriticalP = true
        }
cs:     /* critical activity */
        inCriticalP = false
  od
}

/* similar for process Q with same label cs: */

ltl mutualExcl { []!(P@cs && Q@cs) }
```

Demo: `start/noGhost.pml`

Never Claims: Processes trying to show user wrong

Büchi automaton, as PROMELA process, for negated property

1. Negated TL formula translated to '**never**' process
2. Accepting locations in Büchi automaton represented with help of accept labels ("`accept`*xxx*`:`")
3. If one of these reached infinitely often, the orig. property is violated

Never Claims: Processes trying to show user wrong

Büchi automaton, as PROMELA process, for negated property

1. Negated TL formula translated to '**never**' process
2. Accepting locations in Büchi automaton represented with help of accept labels ("accept*xxx*:")
3. If one of these reached infinitely often, the orig. property is violated

Example (Never claim for <>p, simplified for readability)

```
never {  /* !<>p */
  accept_xyz: /* passed ∞ often iff !<>p holds */
  do
  :: !p
  od
}
```

Liveness Properties

Liveness Properties

- state that something good (ϕ) eventually happens in each run
- each violating requires *infinitely* many steps

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- state that something good (ϕ) eventually happens in each run
- each violating requires *infinitely* many steps

Example

`<>csp`

(with `csp` a variable only true in the critical section of `P`)

"in each run, process `P` visits its critical section eventually"

Applying Temporal Logic to Starvation Problem

We want to verify `<>csp` as a correctness property of:

```
active proctype P() {
  do :: /* non-critical activity */
        atomic {
          !inCriticalQ;
          inCriticalP = true
        }
        csp = true;
        /* critical activity */
        csp = false;
        inCriticalP = false
  od
}

/* similarly for process Q */
/* there, using csq        */
```

Model Checking a Liveness Property using jSpin

1. open Promela file `liveness1.pml`
2. write `ltl pWillEnterC { <>csp }` in Promela file
   (as first `ltl` formula)
3. ensure that `Acceptance` is selected
   (Spin will search for *accepting* cycles through the never claim)
4. *for the moment* uncheck `Weak Fairness` (see discussion below)
5. select `Verify`

Verification Fails

Demo: `start/liveness1.pml`

Verification fails!

Why?

Verification Fails

Demo: `start/liveness1.pml`

Verification fails!

Why?

The liveness property on one process "had no chance".
Not even weak fairness was switched on!

Model Checking Liveness with Weak Fairness using jSpin

Always check `Weak fairness` when verifying liveness

1. open Promela file
2. write `ltl pWillEnterC { <>csp }` in Promela file (as first `ltl` formula)
3. ensure that `Acceptance` is selected (Spin will search for *accepting* cycles through the never claim)
4. ensure `Weak fairness` is checked
5. select `Verify`

Model Checking Liveness using Web Interface

1. add definition of TL formula to PROMELA file

   Example `ltl pWillEnterC { <>csp }`

   General `ltl` *name* `{` *TL-formula* `}`

   can define more than one formula
2. load PROMELA file into web interface
3. ensure `Acceptance` is selected
4. enter name of LTL formula in according field
5. ensure `Weak fairness` is checked
6. select `Verify`

Demo: `liveness1.pml`

Model Checking Liveness using SPIN directly

Command Line Execution

Make sure `ltl` *name* { *TL-formula* } is in *file*.pml

```
> spin -a file.pml
> gcc -o pan pan.c
> ./pan -a -f [-N name]
```

`-a` acceptance cycles, `-f` weak fairness

Demo: `start/liveness1.pml`

Limitation of Weak Fairness

Verification fails again!

Why?

Limitation of Weak Fairness

Verification fails again!

Why?

Weak fairness is too weak . . .

Definition (Weak Fairness)

A run is called weakly fair iff the following holds:
each continuously executable statement is executed eventually.

Limitation of Weak Fairness

Verification fails again!

Why?

Weak fairness is too weak . . .

Definition (Weak Fairness)

A run is called weakly fair iff the following holds:
each continuously executable statement is executed eventually.

Note that `!inCriticalQ` is not continuously executable!

Limitation of Weak Fairness

Verification fails again!

Why?

Weak fairness is too weak . . .

Definition (Weak Fairness)

A run is called weakly fair iff the following holds:
each continuously executable statement is executed eventually.

Note that `!inCriticalQ` is not continuously executable!

Restriction to weak fairness is principal limitation of Spin

Here, liveness needs strong fairness, which is not supported by Spin.

Revisit `fair.pml`

- Specify liveness of `fair.pml` using labels

Revisit `fair.pml`

- Specify liveness of `fair.pml` using labels
- Prove termination

Demo: `target/fair.pml`

Revisit `fair.pml`

- Specify liveness of `fair.pml` using labels
- Prove termination
- Here, weak fairness is needed, *and sufficient*

Demo: `target/fair.pml`

Literature for this Lecture

Ben-Ari Chapter 5
except Sections 5.3.2, 5.3.3, 5.4.2
(`ltl` construct replaces **#define** and `-f` option of SPIN)