

A Poor Man's Concurrency Monad

Koen Claessen

without adding primitives,
we construct a way to lift
any monad into a limited,
but useful concurrent
setting.

Monads

- abstraction from computation

class Monad m where

$(\gg=) :: m a \rightarrow (a \rightarrow m b) \rightarrow m b$

return :: $a \rightarrow m a$

- we use special notation

do	$a \leftarrow \text{expr}_1$		$\text{expr}_1 \gg= \lambda a \rightarrow$
	expr_2		$\text{expr}_2 \gg= \lambda _ \rightarrow$
	$b \leftarrow \text{expr}_3$		$\text{expr}_3 \gg= \lambda b \rightarrow$
	expr_4		expr_4

Writer Monad

- can produce some output during computation

class Monad m \Rightarrow Writer m

where

write :: String \rightarrow m ()

- An implementation could be:

- type W a = (a, String)

- instance Monad W where

m \gg k = let (a, s) = m

(b, s') = k a

in (b, s++s')

return a = (a, "")

- instance Writer W where

write s = ((), s)

- output :: W a \rightarrow String

output (a, s) = s

Monad Transformer

- adds a feature to an existing monad

```
class MonadT t where
  lift :: Monad m
      => m a -> (t m) a
```

- examples:
 - state
 - exception
 - non determinism
- "compose your own monad" - LEGO

Concurrency

- * interleaving actions
- * atomic actions are actions in some monad
- * round robin scheduler

* process has to consist of initial action + future.

Actions

We build actions from three different constructions:

atomic actions, forked actions and no-action.

data Action m
= Atom (m (Action m))
| Fork (Action m)
 (Action m)
| Stop

We use constructors:

- general & simple
- expressive

See also Scholz [2].

Continuation

specifies what to do with result.

type $C\ a =$
 $(a \rightarrow \text{Action}) \rightarrow \text{Action}$

parametrize over a monad:

type $C\ m\ a =$
 $(a \rightarrow \text{Action}\ m) \rightarrow \text{Action}\ m$

for some type Action that stands for a process.

It is a monad:

instance $\text{Monad}\ (C\ m)$ where

$m \gg= k = \backslash \text{cont} \rightarrow m$
 $(\backslash a \rightarrow k\ a\ \text{cont})$

$\text{return}\ a = \backslash \text{cont} \rightarrow \text{cont}\ a$

Useful Operations

Some functions that make life easier.

- Turn a $C\ m\ a$ into an Action:

$action :: C\ m\ a \rightarrow Action\ m$

$action\ c = c\ (\backslash a \rightarrow Stop)$

- Turn an $m\ a$ into an (atomic) $C\ m\ a$:

$atom :: m\ a \rightarrow C\ m\ a$

$atom\ m = \backslash cont \rightarrow$

$Atom\ (do\ a \leftarrow m$
 $\quad\quad\quad return\ (cont\ a))$

- End a process (the empty process):

$stop :: C\ m\ a$

$stop = \backslash cont \rightarrow Stop$

Fork

Some operations on fork:

- 'Imperative' fork:

$\text{fork} :: C\ m\ a \rightarrow C\ m\ ()$
 $\text{fork } c = \backslash \text{cont} \rightarrow \text{Fork}$
 $(\text{action } c) (\text{cont } ())$

- 'Algebraic' or symmetrical fork:

$\text{par} :: C\ m\ a \rightarrow C\ m\ a \rightarrow C\ m\ a$
 $\text{par } c_1\ c_2 = \backslash \text{cont} \rightarrow$
 $\text{Fork } (c_1\ \text{cont})\ (c_2\ \text{cont})$

Running a C

Ideally, we would like
a function

$\text{run} :: C\ m\ a \rightarrow m\ a$

this is "not" possible, due
to typing problems.

We will define a function

$\text{run} :: C\ m\ a \rightarrow m\ ()$

This means we'll only get
the side-effects of the
computation.

Round Robin

simple scheduler.

round :: [Action m] → m ()

round [] = return ()

round (p:ps) =

case p of

- Atom ma →

do p' ← ma

round (ps++ [p'])

- Fork p1 p2 →

round (ps++ [p1, p2])

- Stop →

round ps

Using C

- We can use the scheduler to define:

run :: C m a \rightarrow m ()

run c = round [action c]

- We can construct C's with atom, fork, stop, and can run them using run .

C is a Monad Transformer

C can be made an instance of MonadTrans.

instance MonadTrans C
where
lift = atom

All lifted actions become atomic actions in the new setting.

Example 1: Writer

We lift every writer monad:

instance Writer m =>
 Writer (C m) where
 write s = lift (write s)

Every write action is now atomic.

example :: C W ()

example = do write "hej!"
 fork (loop "apa")
 fork (loop "hund")

where

loop s = do write s
 loop s

will result in:

~~hej! apa.hund apa.hund apa....~~

hej! apa.hund apa.hund apa....

Example 2: Another lifting

We can lift writers in a different way:

```
instance Writer m =>
    Writer (C m) where
    write "" = return ()
    write (c:s) = do lift (write [c])
                    write s
```

a write action is now split up in atomic actions for each character.

hej! ahpuanadphaupn....