Parsing Expressions Slides by Koen Lindström Claessen & David Sands

Expressions

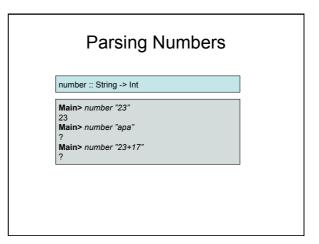
- · Such as
 - 5*2+12
 - 17+3*(4*3+75)
- · Can be modelled as a datatype

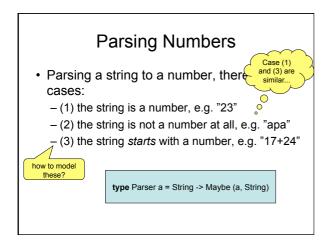
data Expr = Num Int | Add Expr Expr | Mul Expr Expr

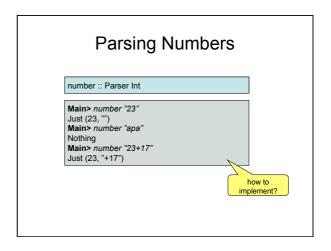
Parsing

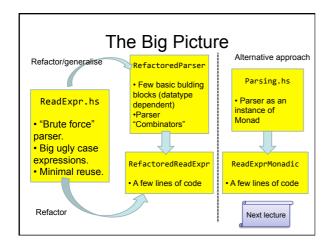
- Transforming a "flat" string into something with a richer structure is called *parsing*
 - expressions
 - programming languages
 - natural language (swedish, english, dutch)
 - _
- Very common problem in computer science
 - Many different solutions

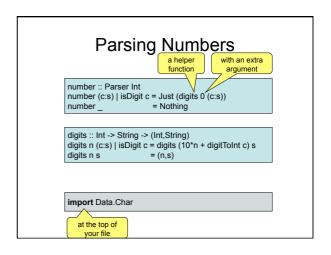
Expressions data Expr = Num Int | Add Expr Expr | Mul Expr Expr | Mul Expr Expr • Let us start with a simpler problem • How to parse data Expr = Num Int but we keep in mind that we want to parse real expressions...

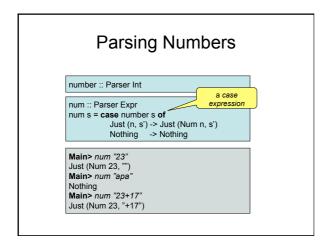


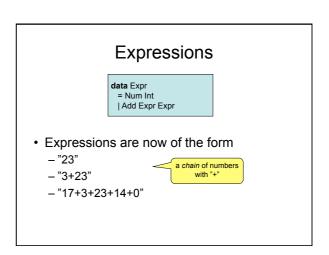


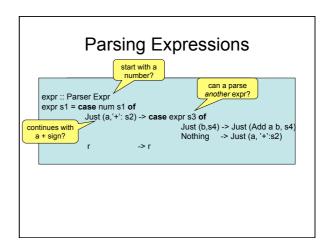










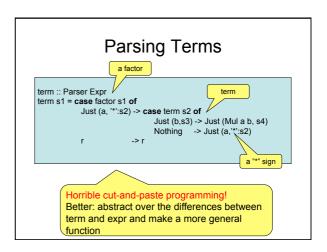


Expressions data Expr = Num Int | Add Expr Expr | Mul Expr Expr | Mul Expr Expr | a chain of terms | a chain of factors with "*" | a chain of factors with "*"

Grammar for Expressions • Parse Expressions according to the following BNF grammar: <expr> ::= <term> | <term> "+" <expr> <term> ::= <factor> | <factor> "*" <term> <factor> ::= "(" <expr> ")" | <number>

```
Parsing Expressions

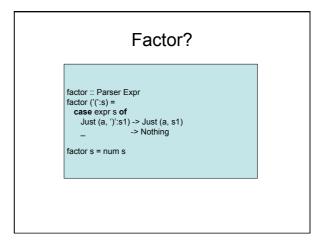
expr:: Parser Expr
expr s1 = case term s1 of
    Just (a,'+':s2) -> case expr s2 of
    Just (b,s3) -> Just (Add a b, s3)
    Nothing -> Just (a, '+':s2)
    r -> r
```



chain p op f s = case p s of Just (n,c:s') | c == op -> case chain p op f s' of Just (m,s") -> Just (f n m,s") Nothing -> Just (n,c:s') r -> r expr, term :: Parser Expr expr = chain term '+' Add term = chain factor '*' Mul

Factor? factor:: Parser Expr factor = num

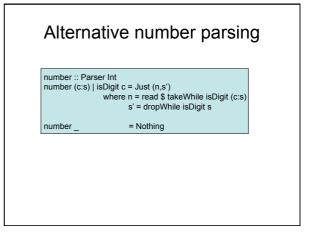
Parentheses • So far no parentheses • Expressions look like - 23 - 23+5*17 - 23+5*(17+23*5+3) a factor can be a parenthesized expression again



```
Reading an Expr

Main> readExpr "23"
Just (Num 23)
Main> readExpr "apa"
Nothing
Main> readExpr "23+17"
Just (Add (Num 23) (Num 17))

readExpr :: String -> Maybe Expr
readExpr s of
Just (a,"") -> Just a
_______-> Nothing
```



Summary

- · Parsing becomes easier when
 - Failing results are explicit
 - A parser also produces the rest of the string
- · Case expressions
 - To look at an intermediate result
- · Higher-order functions
 - Avoid copy-and-paste programming

```
readExpr :: String -> Maybe Expr
readExpr s = case expr s of

Just (a,"") -> Just a

-> Nothing

expr, term :: Parser Expr
expr = chain term '+' Add
term = chain factor '*' Mul

factor :: Parser Expr
factor ('(':s) =
case expr s of
Just (a, ')':s1) -> Just (a, s1)

-> Nothing

factor s = num s
```

The Code (2) chain p op f s = case p s of Just (n,c:s2) | c == op -> case chain p op f s2 of Just (m,s3) -> Just (fn m,s3) Nothing -> Just (n,c:s2) r -> r number :: Parser Int number (c:s) | isDigit c = Just (digits 0 (c:s)) number _ = Nothing digits :: Int -> String -> (Int,String) digits n (c:s) | isDigit c = digits (10*n + digitToInt c) s digits n s = (n,s)

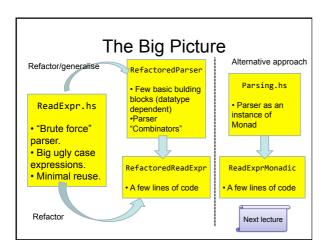
Refactoring the Parser: First Attempt

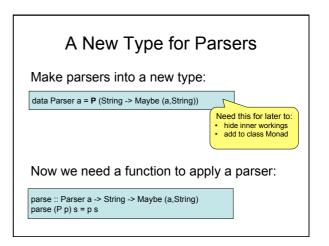
Many operations in our Parser can be made **more general**

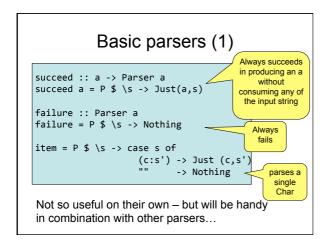
· more reuse, less clutter

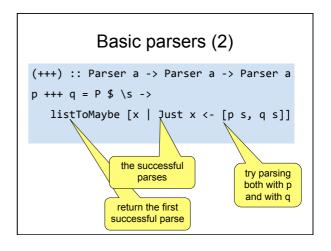
Here we refactor the definition into

- Basic building blocks for parsers (dependent on the type of our Parser)
- Combinators: building blocks for making parsers from other parsers (independent of the type of Parser)









Basic Parsers Lets define some functions to build some basic parsers sat :: (Char -> Bool) -> Parser Char sat prop = P \$ \s -> case s of (c:cs) | prop c -> Just (c,cs) -> Nothing digit = sat isDigit char :: Char -> Parser Char will redefine sat later from more basic parsers

Parse one thing after another

Several ways to parse one thing then another, e.g.

- parse first thing, discard result then parse second thing (function (>->))
- parse first thing, parse and discard a second thing, return result of the first (<-<)
- parse the first thing and then parse a second thing in a way which depends on the value of the first (function (>*>))
- parse a sequence of as many things as possible (functions zeroOrMore, oneOrMore)


```
Derived Parsers

(>->) :: Parser a -> Parser b -> Parser b
p >-> q = p >*> \_ -> q

(as before) throws away the result of first parser

(<-<) :: Parser a -> Parser b -> Parser a
p <-< q = p >*> \a -> q >-> succeed a

throws away the result of second parser

Main> (sat isDigit <-< char '>' ) "2>xxx"

Just ('2', "xxx")
```

```
Parsing sequences to lists

(<:>) :: Parser a -> Parser [a] -> Parser [a] p <:> q = p >*> \a -> pmap (a:) q

zeroOrMore,oneOrMore :: Parser a -> Parser [a]

zeroOrMore p = oneOrMore p +++ succeed []
oneOrMore p = p <:> zeroOrMore p

Main> zeroOrMore (sat isDigit) "1234xxxx"

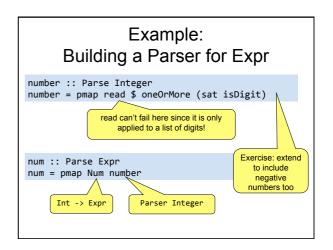
Just ("1234","xxxx")

Main> zeroOrMore (sat isDigit) "x1234xxx"

Just ("","x1234xxx")

Main> (char '@' <:> oneOrMore (char '+')) "@++xxx"

Just ("@++","xxx")
```



Building Parsers with Parsers

Summary (Refactoring)

- By using higher-order programming we can build parser combinators (functions that build parsers from parsers) from which specific parsers can be quickly written.
- Next time: Turning parser combinators into a Monads