

Exam

Data structures DIT960

Time	Monday 30 th May 2016, 14:00–18:00
Place	Hörsalsvägen
Course responsible	Nick Smallbone, tel. 0707 183062

The exam consists of **six questions**. For each question you can get a **G** or a **VG**.

To get a **G** on the exam, you need to answer **three** questions to **G** standard.

To get a **VG** on the exam, you need to answer **five** questions to **VG** standard.

A fully correct answer for a question will get a **VG**. An answer with small mistakes will get a **G**. An answer with large mistakes will get a **U**.

When a question asks for **pseudocode**, you can use a mixture of English and programming notation to describe your solution, and should give enough detail that a competent programmer could easily implement your solution.

Allowed aids One A4 piece of paper of notes, which should be handed in after the exam. You may use both sides.

You may also bring a dictionary.

Note Begin each question on a new page.

Write your anonymous code (*not* your name) on every page.

Good luck!

1. The following algorithm takes as input an array, and returns the array with all the duplicate elements removed. For example, if the input array is {1, 3, 3, 2, 4, 2}, the algorithm returns {1, 3, 2, 4}.

```
S = new empty set
A = new empty dynamic array
for every element x in input array
    if not S.member(x) then
        S.insert(x)
        A.append(x)

return A
```

What is the big-O complexity of this algorithm, if the set is implemented as:

- a) an AVL tree?
- b) a hash table?

For G: write the complexity in terms of n , the size of the input array.

For VG: write the complexity in terms of n and m , where n is the size of the input array and m is the number of distinct elements in the array (i.e. the number of elements ignoring duplicates).

2. Suppose you have the following hash table, implemented using *linear probing*. The hash function we are using is the identity function, $h(x) = x$.

0	1	2	3	4	5	6	7	8
9	18		12	3	14	4	21	

- a) In which order could the elements have been added to the hash table?
There are several correct answers, and you should give all of them.
Assume that the hash table has never been resized, and no elements have been deleted yet.
- A 9, 14, 4, 18, 12, 3, 21
 - B 12, 3, 14, 18, 4, 9, 21
 - C 12, 14, 3, 9, 4, 18, 21
 - D 9, 12, 14, 3, 4, 21, 18
 - E 12, 9, 18, 3, 14, 21, 4
- b) Remove 3 from the hash table, and write down how it looks afterwards.
- c) **For VG only:** if we want a hash table that stores a set of *strings*, one possible hash function is the string's length, $h(x) = x.length$.

Is this a good hash function? Explain your answer.

3. Design an algorithm that takes two arrays, and returns *true* if the arrays are disjoint, i.e. have no elements in common.

You may freely use standard data structures and algorithms from the course in your solution, without explaining how they are implemented.

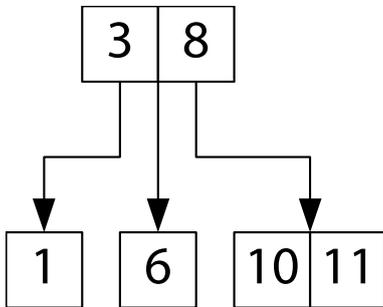
Write down your algorithm as **pseudocode**. You don't need to write Java code, but be precise – a competent programmer should be able to take your description and easily implement it.

For a G: your algorithm should take $O(n \log n)$ time.

For a VG: your algorithm should take $O(n \log m)$ time, where n is the size of the *larger* array and m is the size of the *smaller* array.

Hint for VG: since $n \geq m$, this is the same as $O(n \log m + m \log m)$.

4. Look at the 2-3 tree below.



- a) Draw the 2-3 tree as an AA tree. Mark each node with its level.¹
- b) Insert 9 into the AA tree using the AA tree insertion algorithm. Write down the final tree.

¹ If you have studied the version of AA trees that uses colours, you may write down the colours instead of the levels.

5. A *bidirectional map* is a map which supports bidirectional lookup: given a key, you can find the corresponding value, and given a value, you can find the corresponding key.

In a bidirectional map there is always a one-to-one relationship between keys and values. In other words, each key has exactly one value, and each value is found under exactly one key.

It supports the usual map operations:

- `new()`: create a new, empty map
- `insert(k, v)`: add the mapping $k \rightarrow v$ to the map; **any existing mapping with key k or value v is removed.**
- `lookup(k)`: if the map contains a mapping $k \rightarrow v$, return v
- `delete(k)`: if the map contains a mapping $k \rightarrow v$, delete it

plus the following *reverse lookup* operation:

- `rlookup(v)`: if the map contains a mapping $k \rightarrow v$, return k

The following example shows what the various operations do.

Operation	Result
<code>new()</code>	Map is {}
<code>insert(1, 2)</code>	Map is {1 → 2}
<code>insert(3, 4)</code>	Map is {1 → 2, 3 → 4}
<code>lookup(1)</code>	Returns 2
<code>insert(4, 2)</code>	Map is {3 → 4, 4 → 2}. Notice that the mapping 1 → 2 is replaced by 4 → 2.
<code>rlookup(2)</code>	Returns 4
<code>delete(4)</code>	Map is {3 → 4}

We can implement a bidirectional map using *two* maps, each implemented as e.g. an AA tree:

- forward is a map from keys to values.
In the example above, it contains $3 \rightarrow 4$ and $4 \rightarrow 2$.
- back is a map from values to keys.
In the example above, it contains $4 \rightarrow 3$ and $2 \rightarrow 4$.

The invariant is that the two maps always contain the same data: forward contains the mapping $k \rightarrow v$, if and only if back contains the mapping $v \rightarrow k$.

We can then implement `new`, `lookup` and `rlookup` as follows:

- `new`: set forward and back to be empty AA trees
- `lookup(k)`: look up k in forward using the BST lookup algorithm
- `rlookup(v)`: look up v in back using the BST lookup algorithm

Your task is to implement the remaining operations, `insert` and `delete`, with $O(\log n)$ complexity.

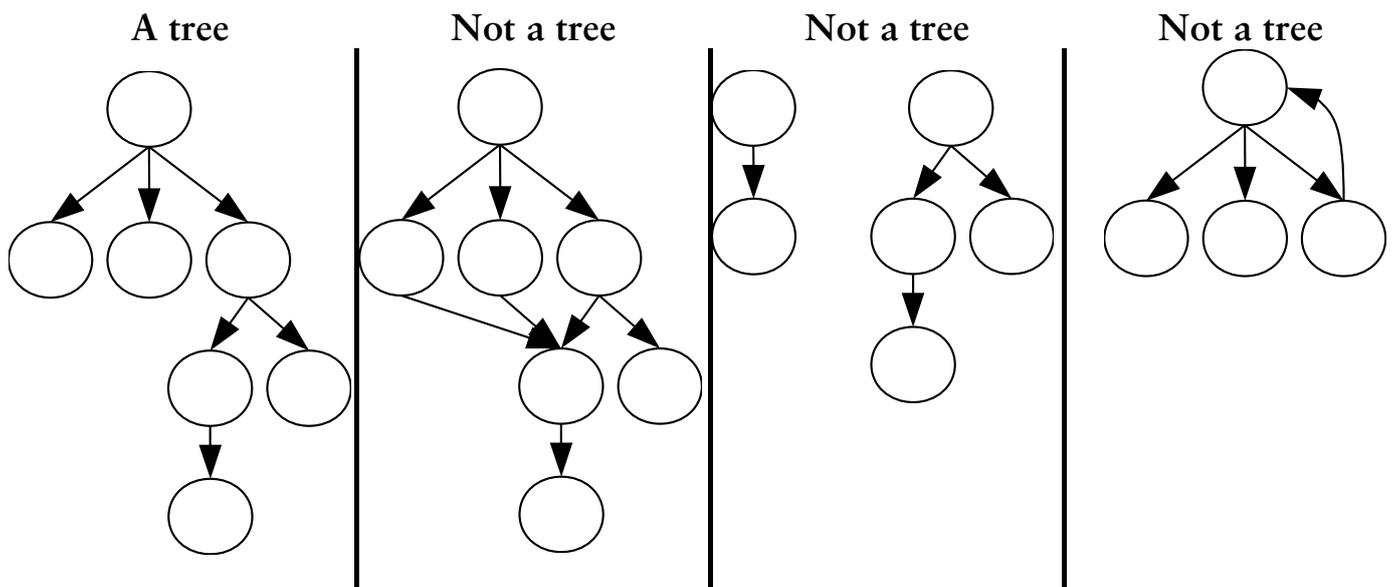
Be careful: the algorithm is more complicated than it seems. Be extra careful that you maintain the data structure invariant. It's also a good idea to test your solution on the previous page's example.

Give **pseudocode** for each of the operations. You don't need to write Java code, but be precise – a competent programmer should be able to take your description and easily implement it..

You may freely use standard data structures and algorithms from the course in your solution, including insertion/lookup/deletion in a map, without explaining how they are implemented.

6. We can think of a tree as being a special kind of directed graph. To model a tree as a graph, we make the nodes of the tree become nodes in the graph, and draw an edge from a parent node to each of its children.

The drawing on the left shows a tree as a graph; the other three directed graphs do *not* correspond to a tree.



Suppose we want to check if a given directed graph corresponds to a tree. What properties should we check that the graph has? Write down a list of properties such that, if a directed graph has those properties, it must be a tree. You can refer to standard properties of graphs in your answer without explaining them.

ALGORITHMS BY COMPLEXITY

MORE COMPLEX →

LEFTPAD QUICKSORT

GIT
MERGE

SELF-
DRIVING
CAR

GOOGLE
SEARCH
BACKEND

SPRAWLING EXCEL SPREADSHEET
BUILT UP OVER 20 YEARS BY A
CHURCH GROUP IN NEBRASKA TO
COORDINATE THEIR SCHEDULING

<http://www.xkcd.com/1667>