

Recursive Data Types



Modelling Arithmetic Expressions

Imagine a program to help school-children learn arithmetic, which presents them with an expression to work out, and checks their answer.

What is $(1+2)*3$? 8
Sorry, wrong answer!

The expression $(1+2)*3$ is *data* as far as this program is concerned (**not** the same as 9!). How shall we represent it?

A string??

What is "1+2"+"3"?

What is "1+hello world**"?

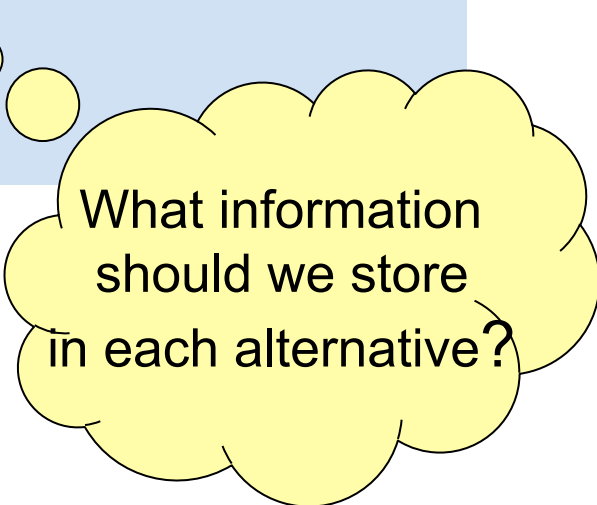
Modelling Expressions

Let's design a datatype to model *arithmetic expressions*
-- not their values, but their structure.

An expression can be:

- a number n
- an addition $a+b$
- a multiplication $a*b$

```
data Expr =  
    Num  
  | Add  
  | Mu1
```



What information should we store in each alternative?

Modelling Expressions

Let's design a datatype to model *arithmetic expressions*
-- not their values, but their structure.

An expression can be:

- a number n
- an addition $a+b$
- a multiplication $a*b$

```
data Expr =  
    Num Integer  
  | Add  Expr Expr  
  | Mul  Expr Expr
```

A recursive data type

Examples

```
data Expr = Num Integer
          | Add Expr Expr
          | Mul Expr Expr
```

The expression: is represented by:

2 Num 2

2+2 Add (Num 2) (Num 2)

(1+2)*3 Mul (Add (Num 1) (Num 2)) (Num 3)

1+2*3 Add (Num 1) (Mul (Num 2) (Num 3))

Example: Evaluation

Define a function

`eval :: Expr -> Integer`

which *evaluates* an expression?

Example: `eval (Add (Num 1) (Mul (Num 2) (Num 3)))`
evaluates to 7

Hint: Recursive types often mean recursive functions!

```
eval :: Expr -> Integer
```

Use pattern matching:
one equation for each case.

```
eval (Num n)      =      n
```

```
eval (Add a b)    =      eval a + eval b
```

```
eval (Mul a b)    =      eval a * eval b
```

a and b are of
type Expr.

Recursive types mean
recursive functions!

Showing Expressions

Expressions will be more readable if we convert them to strings.

```
showExpr :: Expr -> String
```

```
showExpr (Num n) = show n
```

```
showExpr (Add a b) = showExpr a ++ "+" ++ showExpr b
```

```
showExpr (Mul a b) = showExpr a ++ "*" ++ showExpr b
```

```
Main> showExpr (Mul (Num 1) (Add (Num 2) (Num 3)))  
"1*2+3"
```


Quiz

Which brackets are necessary? $1+(2+3)$

$$1+(2*3)$$

$$1*(2+3)$$

What kind of expression *may* need to be bracketed?

When *does* it need to be bracketed?

Quiz

Which brackets are necessary?

$1+(2+3)$

NO!

$1+(2*3)$

NO!

$1*(2+3)$

YES!

Additions

What kind of expression *may* need to be bracketed?

When *does* it need to be bracketed?

Inside multiplications.

Idea

Format *factors* differently:

```
showExpr :: Expr -> String
showExpr (Num n) = show n
showExpr (Add a b) = showExpr a ++ "+" ++ showExpr b
showExpr (Mul a b) = showFactor a ++ "*" ++ showFactor b
```

```
showFactor :: Expr -> String
showFactor (Add a b) = "(" ++ showExpr (Add a b) ++ ")"
showFactor e          = showExpr e
```

Making a Show instance

instance Show Expr where
show = showExpr

```
data Expr = Num Integer | Add Expr Expr | Mul Expr Expr  
deriving (Show, Eq )
```


(Almost) Complete Program

```
questions :: IO ( )
questions =
  do rnd <- newStdGen
     let e = unGen arbitrary rnd 3
         putStr ("What is " ++ show e ++ "? ")
         ans <- getLine
         putStrLn (if read ans == eval e
                    then "Right!" else "Wrong!")
     questions
```

New random seed

An expression generator—needs to be written

What's this?

let: Give name to a result

Opposite of show

Using QuickCheck Generators in Other Programs

- Test.QuickCheck.Gen exports
 - unGen:: Gen a -> StdGen -> Int -> a

QuickCheck generator

The diagram consists of three yellow callout boxes with black outlines. The first box on the left points to 'Gen a' in the signature. The second box in the middle points to 'StdGen'. The third box on the right points to 'Int'.

Random seed

Size parameter for generation

- Size is used, for example, to bound Integers, size of data structures etc.

Generating Arbitrary Expressions

```
instance Arbitrary Expr where
  arbitrary = arbExpr
arbExpr :: Gen Expr
arbExpr =
  oneof [ do n <- arbitrary
          return (Num n)
        , do a <- arbExpr
          b <- arbExpr
          return (Add a b)
        , do a <- arbExpr
          b <- arbExpr
          return (Mul a b) ]
```

Does not work! (why?)

Generates infinite expressions!

Generating Arbitrary Expressions

```
instance Arbitrary Expr where  
  arbitrary = sized arbExpr
```

```
arbExpr :: Int -> Gen Expr
```

```
arbExpr s =  
  frequency [ (1, do n <- arbitrary  
                  return (Num n))  
            , (s, do a <- arbExpr s'  
                  b <- arbExpr s'  
                  return (Add a b))  
            , (s, do a <- arbExpr s'  
                  b <- arbExpr s'  
                  return (Mul a b)) ]
```

```
where s' = s `div` 2
```

Size argument
changes at each
recursive call

Demo

Main> questions

What is $-3*4*-1*-3*-1*-1$? -36

Right!

What is $15*4*(-2+-13+-14+13)$? -640

Wrong!

What is 0? 0

Right!

What is $(-4+13)*-9*13+7+15+12$? dunno

Program error: Prelude.read: no parse

The Program

Crucial line:

failing

```
putStrLn (if read ans==eval e then "Right!"  
           else "Wrong!")
```

```
ans == show (eval e)
```

cannot fail

Reading Expressions

- How about a function
 - `readExpr :: String -> Expr`
- Such that
 - `readExpr "12+173" =`
 - `Add (Num 12) (Num 173)`
 - `readExpr "12+3*4" =`
 - `Add (Num 12) (Mul (Num 3) (Num 4))`

We see how to implement this in the next lecture

Symbolic Expressions

- How about expressions with variables in them?


data Expr = Num Integer

| Add Expr Expr

| Mul Expr Expr

| Var Name

type Name = String



Add **Var** and
change functions
accordingly

Gathering Variables

It is often handy to know exactly which variables occur in a given expression


`vars :: Expr -> [Name]`

`vars (Num n) = []`

`vars (Add a b) = vars a `union` vars b`

`vars (Mul a b) = vars a `union` vars b`

`vars (Var x) = [x]`



From Data.List;
combines two
lists without
duplication

Evaluating Expressions

We would like to evaluate expressions with variables. What is the type?

Table of values for variables

```
eval :: [(Name, Integer)] -> Expr -> Integer
```

```
eval env (Num n) = n
```

```
eval env (Var y) = fromJust (lookup y env)
```

```
eval env (Add a b) = eval env a + eval env b
```

```
eval env (Mul a b) = eval env a * eval env b
```

```
Prelude> :i lookup
```

```
lookup :: (Eq a) => a -> [(a, b)] -> Maybe b
```

Symbolic Differentiation

Differentiating an expression produces a new expression. We implement it as:

```
diff :: Expr -> Name -> Expr
diff (Num n) x           = Num 0
diff (Var y) x | x==y   = Num 1
                  | x/=y = Num 0
diff (Add a b) x        = Add (diff a x) (diff b x)
diff (Mul a b) x        = Add (Mul a (diff b x))
                              (Mul b (diff a x))
```

Variable to differentiate wrt.

Testing differentiate

```
Main> diff (Mul (Num 2) (Var "x")) "x"  
2*1+0*x
```

Not quite what we expected!
-- not *simplified*

What happens?

$$\frac{d}{dx}(2*x) = 2$$

differentiate (Mul (Num 2) (Var "x")) "x"

→ Add (Mul (Num 2) (differentiate (Var "x") "x"))
(Mul (Var "x") (differentiate (Num 2) "x"))

→ Add (Mul (Num 2) (Num 1))

(Mul (Var "x") (Num 0))

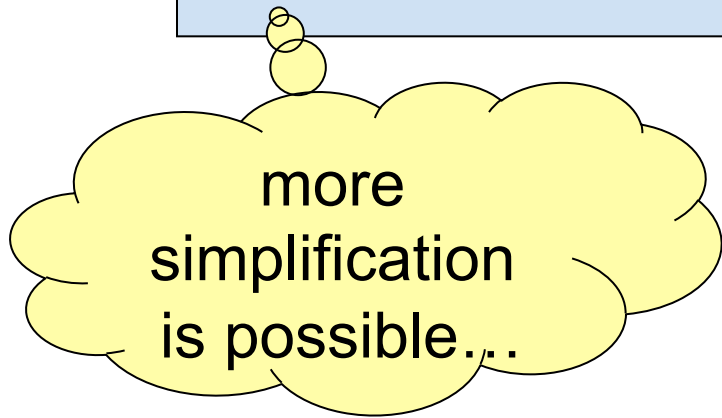

$$2*1 + x*0$$

How can we make differentiate simplify the result?

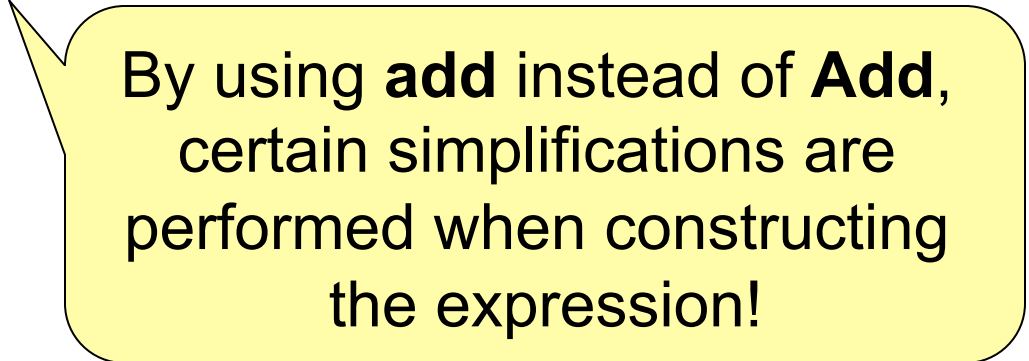
“Smart” Constructors

- Define

```
add :: Expr -> Expr -> Expr
add (Num 0) b           = b
add a      (Num 0)     = a
add (Num x) (Num y)    = Num (x+y)
add a      b           = Add a b
```



more
simplification
is possible...



By using **add** instead of **Add**,
certain simplifications are
performed when constructing
the expression!

Testing add

```
Main> Add (Num 2) (Num 5)
```

```
2+5
```

```
Main> add (Num 2) (Num 5)
```

```
7
```

Symbolic Differentiation

```
diff :: Expr -> Name -> Expr
```

```
diff (Num n) x = Num 0
```

```
diff (Var y) x
```

```
    | x==y = Num 1
```

```
    | x/=y = Num 0
```

```
diff (Add a b) x = add (diff a x) (diff b x)
```

```
diff (Mul a b) x = add (mul a (diff b x))
```

```
                    (mul b (diff a x))
```

“Smart” Constructors -- mul

- How to define mul?

```
mul :: Expr -> Expr -> Expr
mul (Num 0) b           = Num 0
mul a           (Num 0) = Num 0
mul (Num 1) b           = b
mul a           (Num 1) = a
mul (Num x) (Num y)    = Num (x*y)
mul a           b       = Mul a b
```


Expressions

- Expr as a datatype can represent expressions
 - Unsimplified
 - Simplified
 - Results
 - Data presented to the user
- Need to be able to convert between these

An Expression Simplifier

- Simplification function
 - `simplify :: Expr -> Expr`

```
simplify :: Expr -> Expr  
simplify e | null (vars e) = ?  
...
```



Exercises

Testing the Simplifier

```
arbExpr :: Int -> Gen Expr
arbExpr s =
  frequency [ (1, do n <- arbitrary
              return (Num n))
            , (s, do a <- arbExpr s'
              b <- arbExpr s'
              return (Add a b))
            , (s, do a <- arbExpr s'
              b <- arbExpr s'
              return (Mul a b))
            , (1, do x <- elements ["x","y","z"]
              return (Var x))]
```

Cut'n'paste
here should
be refactored

where

$s' = s \text{ `div` } 2$

Testing an Expression Simplifier

- (1) Simplification should not change the value

```
prop_SimplifyCorrect e env =  
  eval env e == eval env (simplify e)
```

```
prop_SimplifyCorrect e (Env env) =  
  eval env e == eval env (simplify e)
```

Generate lists of
values *for variables*

Testing an Expression Simplifier

```
data Env = Env [(Name,Integer)]  
  deriving ( Eq, Show )
```

```
instance Arbitrary Env where  
  arbitrary =  
    do a <- arbitrary  
       b <- arbitrary  
       c <- arbitrary  
       return (Env [("x",a),("y",b),("z",c)])
```

Testing an Expression Simplifier

- (2) Simplification should do a good job

```
prop_SimplifyNoJunk e =  
  noJunk (simplify e)  
where  
  noJunk (Add a b) = not (isNum a && isNum b)  
                    && noJunk a && noJunk b  
  ...
```

You continue at the
exercises!

Exercises

- Build and test an expression simplifier!
- I found *many subtle bugs* in my own simplifier!
 - Often simplifier goes into an infinite loop
 - Use `verboseCheck` instead of `quickCheck` (prints test case *before* every test, so you see them even if the test loops or crashes)

Summary

- Recursive data-types can take many forms other than lists
- Recursive data-types can model *languages* (expressions, natural languages, programming languages)
- Functions working with recursive types are often recursive themselves
- When generating random elements in recursive datatypes, think about the *size*

Next Time

- How to write *parsers*
 - `readExpr :: String -> Expr`