

# Functional Datastructures

# Efficiency

Consider a naive reverse definition

```
reverse :: [a] -> [a]
reverse []      = []
reverse (x:xs) = reverse xs ++ [x]
```

```
(++) :: [a] -> [a] -> [a]
[] ++ ys      = ys
(x:xs) ++ ys = x:(xs ++ ys)
```

Note: reverse and (++)  
are part of the Prelude

How many (++) calls  
needed to produce all  
elements of xs ++ ys?

$O(\text{length } xs)$

# Efficiency

- Reversing a list takes (length xs) calls to reverse
- Each call to reverse costs  $O(\text{length (reverse xs)}) = O(\text{length xs})$
- So reversing a list of length n requires approx  $(n-1) + (n-2) + \dots + 1 = O(n*n)$  steps

```
reverse :: [a] -> [a]
reverse []      = []
reverse (x:xs) = reverse xs ++ [x]
```

# Fast Reverse

- Quicker reverse avoids using append.  
Idea: use an accumulating parameter

```
reverse :: [a] -> [a]
reverse xs = revInto [] xs
  where revInto ys []      = ys
        revInto ys (x:xs) = revInto (x:ys) xs
```

A helper function

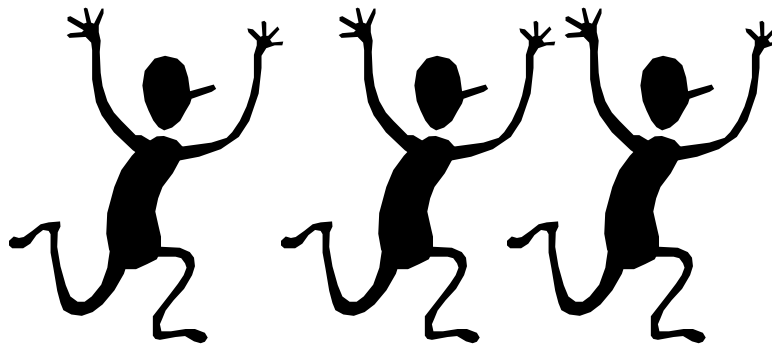
accumulating  
parameter – it  
accumulates the  
answer

# Data Structures

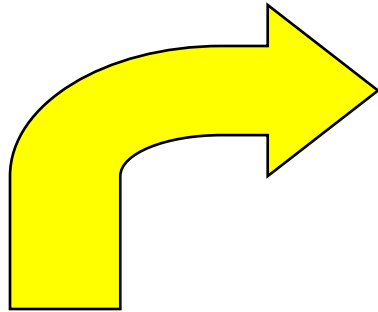
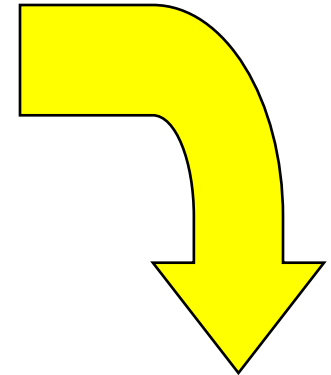
- Datatype
  - A model of something that we want to represent in our program
- Data structure
  - A particular way of *storing* data
  - How? Depending on what we want to do with the data
- Today: one example
  - Queue

# What is a Queue?

Join at the back



Leave at front



## Examples

- Files to print
- Processes to run
- Tasks to perform

# What is a Queue?

A *queue* contains a sequence of values. We can add elements at the back, and remove elements from the front.

We'll implement the following operations:

```
empty    :: Q a           -- an empty queue
add      :: a -> Q a -> Q a -- add element at back
remove   :: Q a -> Q a    -- remove an element from front
front    :: Q a -> a      -- inspect the front element
isEmpty  :: Q a -> Bool   -- check if the queue is empty
```

# First Try

**data** Q a = Q [a] **deriving** (Eq, Show)

empty = Q []

add x (Q xs) = Q (xs++[x])

remove (Q (x:xs)) = Q xs

front (Q (x:xs)) = x

isEmpty (Q xs) = null xs



# Works, but slow

$\text{add } x \text{ (Q } xs) = \text{Q } (xs++[x])$

$[] \quad ++ \text{ } ys = ys$

$(x:xs) ++ \text{ } ys = x : (xs++ys)$

As many recursive calls as there are elements in  $xs$

Add 1, add 2, add 3, add 4, add 5...

Time is the *square* of the number of additions

# A Module

- Implement the result in a *module*
- Use as specification
- Hides the internals (representation)
- Allows the re-use
  - By other programmers
  - Of the same names

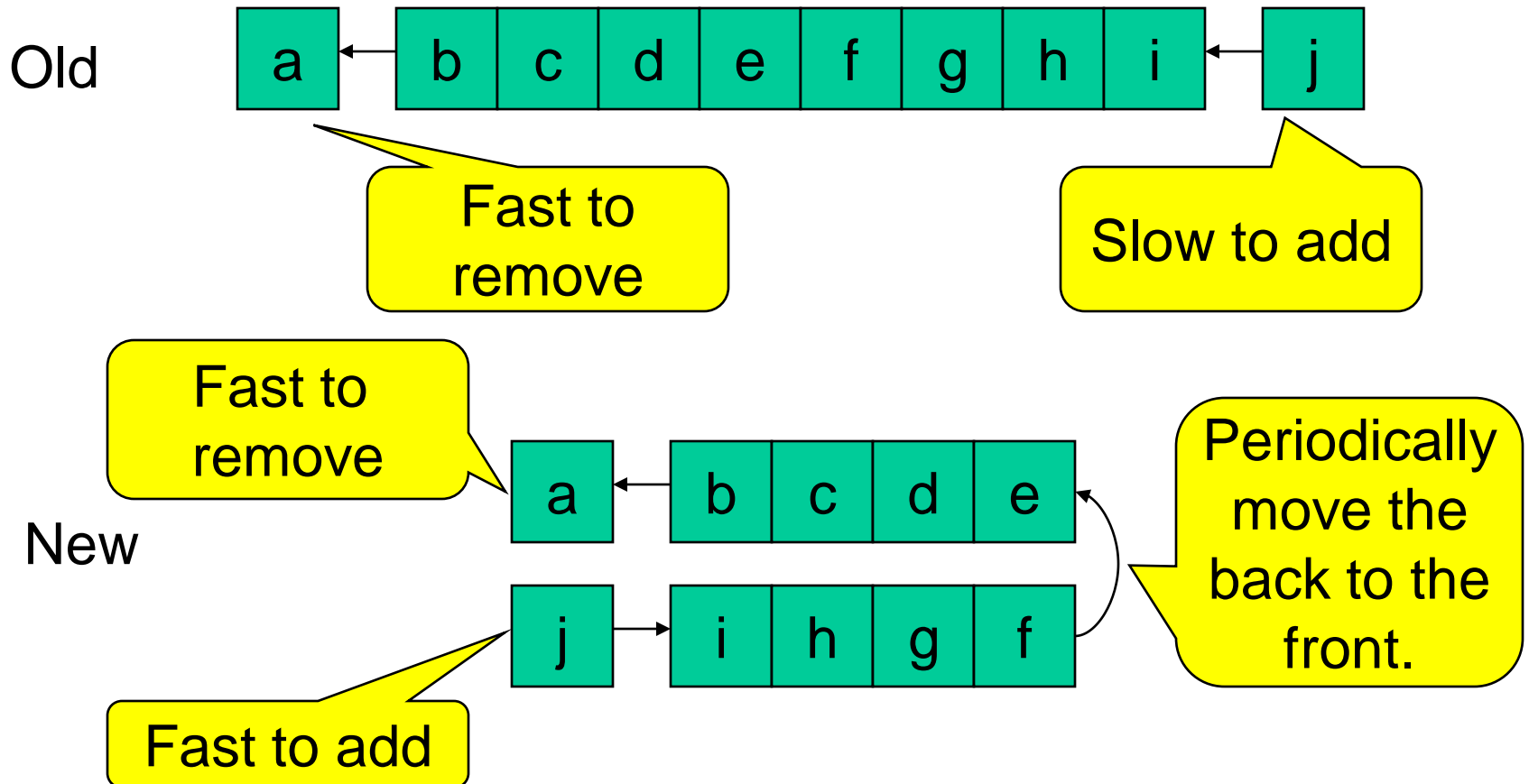
# SlowQueue Module

```
module SlowQueue where
```

```
data Q a = Q [a] deriving (Eq, Show)
```

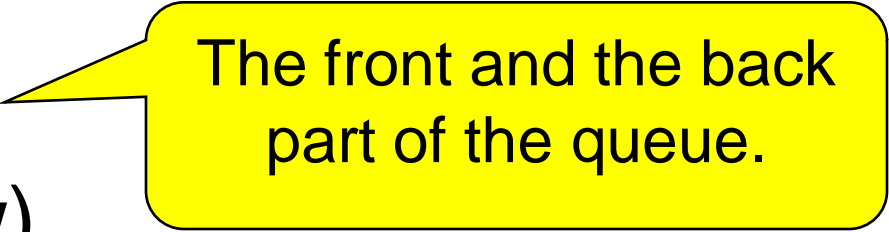
```
empty           = Q []  
add x (Q xs)    = Q (xs++[x])  
remove (Q (x:xs)) = Q xs  
front (Q (x:xs)) = x  
isEmpty (Q xs)  = null xs
```

# New Idea: Store the Front and Back Separately



# Smart Datatype

**data** Q a = Q [a] [a]  
**deriving** (Eq, Show)

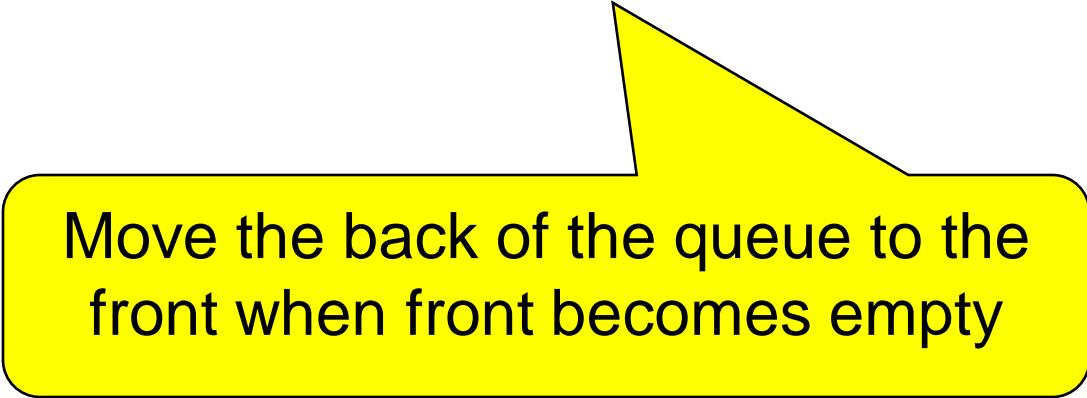


The front and the back  
part of the queue.

Invariant: front is empty only when the  
back is also empty

# Smart Operations

empty = Q [] []  
isEmpty q = q == empty  
add x (Q front back) = **fixQ** (Q front (x:back))  
front (Q (x:front) back) = x  
remove (Q (x:front) back) = **fixQ** (Q front back)



Move the back of the queue to the front when front becomes empty

# Flipping

$\text{fixQ } (Q \ [] \ \text{back}) = Q \ (\text{reverse back}) \ []$   
 $\text{fixQ } q = q$

- `fixQ` takes one call per element
- Each element is flipped exactly once, so
  - $O(1)$  to add,  $O(1)$  to `fixQ`,  $O(1)$  to remove.

# Wrapping it up

```
module Queue (Q,  
              empty, add, remove,  
              front, isEmpty  
              ) where
```

Exports type  
Q but not the  
constructor

```
*Main> :i Q  
data Q a    -- Defined at Queue.hs:11:5  
*Main> front (Q [1,2] [3])  
<interactive>:1:0: Not in scope: data constructor `Q'
```



# Exported Constructors

```
module Queue (Q(Q)  
  empty  
  from  
  ) where
```

Exports type  
Q and the  
constructor Q

Not a good idea here: allows  
client to

- become dependent on  
internal implementation  
details
- break datatype invariants

```
*Main> :i Q  
data Q a = Q [a] [a] -- Defined at Queue.hs:11:5  
*Main> Q [] [3]  
Q [] [3]
```

# How can we test the smart functions?

- By using the original implementation as a *reference*
- The behaviour should be "the same"
  - Check results
- First version is an *abstract model* that is "obviously correct"

# Later we will see:

- How to make QuickCheck work for our own datatypes
  - We need to tell it how to generate random values
- How to test the equivalence of the reference and efficient implementations
  - we need to add conversion functions
- How to test the intended invariants