## Exercise 3

## SQL queries

Relational algebra

## Example

Courses

SELECT name FROM Courses WHERE nrStudents > 20

| Code | Name | nrstudents |
| :--- | :--- | ---: |
| AB123 | Math | 19 |
| CD234 | Physics | 24 |
| EF345 | Karaoke | 23 |
| GH456 | PERL | 20 |

$>\pi_{\text {name }}\left(\sigma_{\text {nrStudents }>20}(\right.$ Courses $\left.)\right)$

## The most common question

> What the hell is relational algebra good for?

## The DBMS uses relational algebra

A DBMS may have many different ways of implementing the relational algebra operations.

The aim of query optimization is to choose the most efficient one.
$\rangle$ To do this, it uses formulae that estimate the costs for a number of options and selects the one with the lowest cost.

## Projection

Which attributes

Which columns
, SELECT $A, C, E$
$>\pi_{A, C, E}$


## Selection

- Which tuples

Which rows
, WHERE E > 5
$>\sigma_{E}>5$
DONT confuse SELECT (projection) for selection!!

| $A$ | $B$ | $C$ | $D$ | $E$ | $F$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  | 4 |  |
|  |  |  |  | 7 |  |
|  |  |  |  | 5 |  |
|  |  |  |  | 8 |  |
|  |  |  |  | 3 |  |

## SQL vs. Relational Algebra

SELECT
FROM
WHERE
GROUP BY HAVING
ORDER BY

X
T
C
Y

$$
\tau_{Z}\left(\pi_{x}\left(\sigma_{D}\left(/_{y}\left(\sigma_{c}(T)\right)\right)\right)\right)
$$

D
Z

## Combining tables

- Set (actually bag) operations

Cartesian product

Joins

## Set (actually bag) operations



R UNION ALL S
$R \cup S$
R INTERSECT S $R \cap S$

R MINUS S
$R-S$

$$
\begin{array}{|l|l|}
\hline \mathbf{A} & \boldsymbol{B} \\
\hline 1 & 2 \\
\hline 3 & 4 \\
\hline 1 & 2 \\
\hline 5 & 6 \\
\hline
\end{array}
$$

## Joins

- 3 "basic" joins:
- Cartesian product
- Conditional join

Theta join, Inner join, Equi join, Nonequi join, Natural join

- Outer joín


## Cartesian product



## Conditional join = Inner join



R,S
WHERE C
or
R JOIN S ON C R ${ }^{\wedge} S$

If $C$ is "R.A $=S . B^{\prime \prime}$

| $R, A$ | $R, B$ | $S, B$ | $S, C$ |
| :---: | :---: | :---: | :---: |
| 3 | 4 | 3 | 4 |

If $C$ is equality -
Equij join
If C is inequality Nonequij join

Beware of NULL!

## Special case - Natural join

R NATURAL JOIN S

$$
R \bowtie S
$$

| R,A | B | S.C |
| :---: | :---: | :---: |
| 3 | 4 | 5 |

## Outer join


R FULL OUTER JOIN S ON R.B = S.C

| R.A | R.B | S.C | S.D |
| :--- | :--- | :--- | :--- |
| 1 | 2 | NULL | NULL |
| 3 | 4 | 4 | 5 |
| NULL | NULL | 6 | 7 |

All rows in both/left/right table(s) will appear, and the rest will be filled with null if C does not match.

## There is more...

Grouping
Renaming
Sorting
Also note that the terminology regarding joins is confused.
Inner join = Equi join (orafag.com)
Inner join = Conditional join (Wikipedia)
Theta join = Conditional join (Course book)

