Complexity

(Weiss chapter 5)

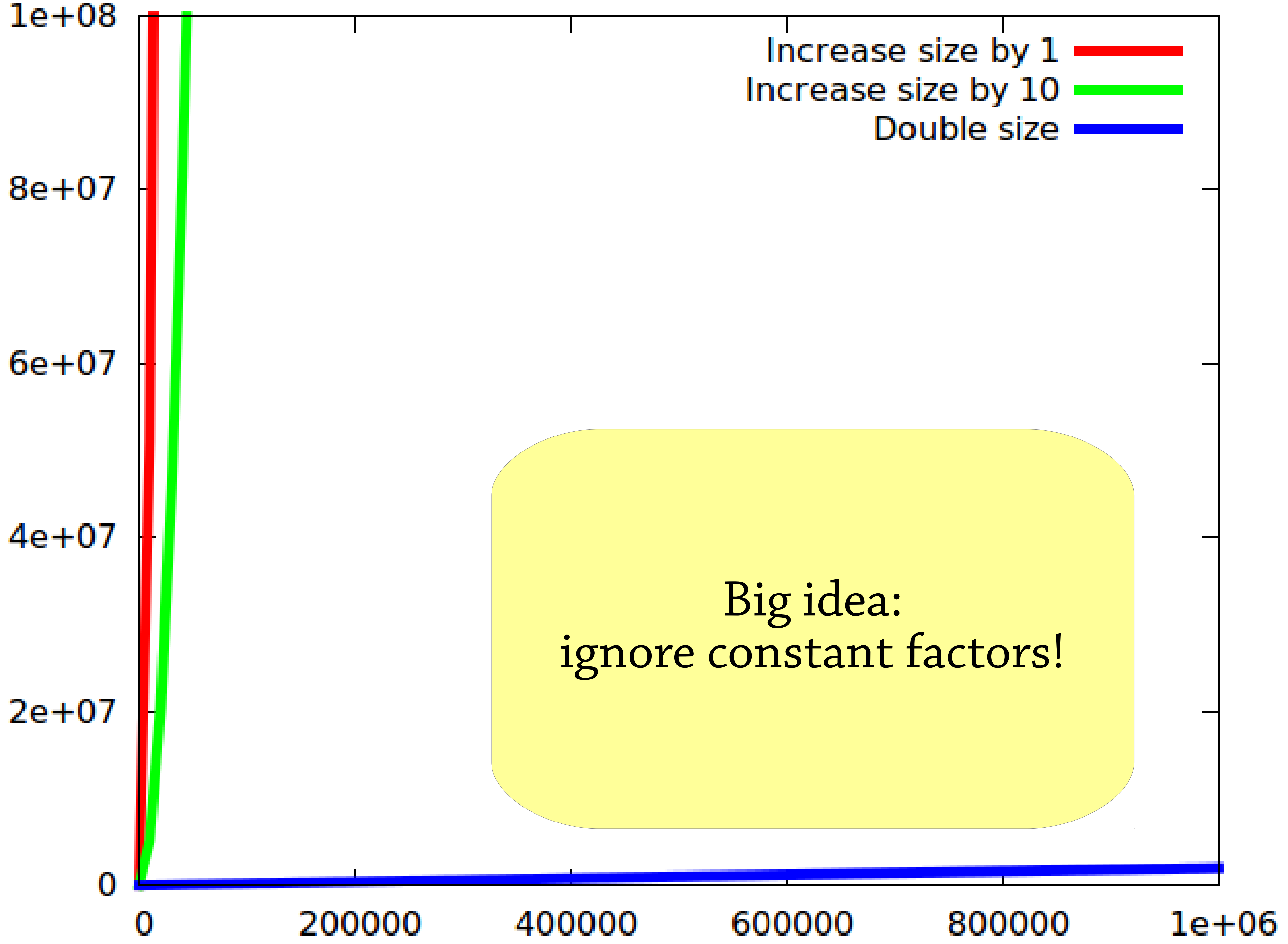
Complexity

This lecture is all about *how to describe the performance of an algorithm*

Last time we had three versions of the file-reading program. For a file of size *n*:

- The first one needed to copy $n^2/2$ characters
- The second one needed to copy $n^2/200$ characters
- The third needed to copy 2n characters

We worked out these formulas, but it was a bit of work – now we'll see an easier way


1e+08
8e+07
6e+07
4e+07
2e+07
0
0
200000
400000
600000
800000
1e+06
Increase size by 1
Increase size by 10
Double size
Big idea:
ignore constant factors!

Why do we ignore constant factors?

Well, when n is 1000000…

- $\log_2 n$ is 20
- n is 1000000
- n^2 is 1000000000000
- 2^n is a number with 300,000 digits…

Given two algorithms:

- The first takes $1000000 \log_2 n$ steps to run
- The second takes 0.00000001×2^n

The first is miles better!
Constant factors *normally* don't matter

Big O notation

Instead of saying…

- The first implementation copies $n^2/2$ characters
- The second copies $n^2/200$ characters
- The third copies $2n$ characters

We will just say…

- The first implementation copies **$O(n^2)$** characters
- The second copies **$O(n^2)$** characters
- The third copies **$O(n)$** characters

$O(n^2)$ means "proportional to n^2" (almost)

Time complexity

With big-O notation, it doesn't matter whether we count steps or time!

Suppose an algorithm takes $n^2/2$ steps, which is $O(n^2)$

And suppose each step takes 100ns to run

Then the algorithm takes $50n^2$ ns, which is also $O(n^2)$!

We say that the algorithm has $O(n^2)$ *time complexity* or simply *complexity*

Do we really need
big O notation?

What happens without big O?

How many steps does this function take on an array of length *n* (in the worst case)?

```
Object search(Object[] a, Object x) {
  for(int i = 0; i < a.length; i++) {
    if (a[i].equals(target))
      return a[i];
  }
  return null;
}
```

Assume that loop body takes 1 step

What happens without big O?

How many steps does this function take on an array of length n (in the worst case)?

```
Object search(Object[] a      ject x) {
  for(int i = 0; i <          ) {
    if (a[i].equal
      return a[i
  }
  return null;
}
```

Answer:
n

What about this one?

```
boolean unique(Object[] a) {
  for(int i = 0; i < a.length; i++)
    for (int j = 0; j < a.length; j++)
      if (a[i].equals(a[j]) && i != j)
        return false;
  return true;
}
```

What about this one?

```
boolean unique(Object[  a)
  for(int i = 0;  <          ++)
    for (int j =              ++)
      if (a[i                 != j)
        ret
  return tru
}
```

Outer loop runs n times
Each time, inner loop runs n times

Total: $n \times n = n^2$

What about this one?

```
boolean unique(Object[] a) {
  for(int i = 0; i < a.length; i++)
    for (int j = 0; j < i; j++)
      if (a[i].equals(a[j]))
        return false;
  return true;
}
```

Loop runs to *i* instead of *n*

Some hard sums

When $i = 0$, inner loop runs 0 times

When $i = 1$, inner loop runs 1 time

…

When $i = n$-1, inner loop runs n-1 times

Total:

- $\sum_{i=0}^{n-1} i = 0 + 1 + 2 + \ldots + n\text{-}1$

which is $n(n\text{-}1)/2$

What about this one?

```
boolean unique(Object[] a) {
  for(int i = 0; i < a.length; i++)
    for (int j = 0; j <      ++)
      if (a[i].equal
        return fal
  return true;
}
```

Answer:
$n(n\text{-}1)/2$

What about this one?

```
void something(Object[] a) {
  for(int i = 0; i < a.length; i++)
    for (int j = 0; j < i; j++)
      for (int k = 0; k < j; k++)
        “something that takes 1 step”
}
```

More hard sums

$$\sum_{i=0}^{n-1} \sum_{j=0}^{i-1} \sum_{k=0}^{j-1} 1$$

Outer loop:
i goes from 0 to *n-1*

Middle loop:
j goes from 0 to i-*1*

Inner loop:
k goes from 0 to j-*1*

Counts: how many values i, j, k where
$0 \le i < n$, $0 \le j < i$, $0 \le k \le j$

More hard sums

$$\sum_{i=0}^{n-1} \sum_{j=0}^{i-1} \sum_{k=}^{j-1} 1$$

I have no idea
how to solve this!
Wolfram Alpha says it's
n(*n*-1)(*n*-2)/6

Counts: how many values i, j, k where
$0 \le i < n$, $0 \le j < i$, $0 \le k \le j$

What about this one?

```
void something(Object[] a) {
  for(int i = 0; i < a.lengt ; i++)
    for (int j = 0; j <         ++)
      for (int k =
        "something                step"
}
```

Answer:
$n(n-1)(n-2)/6$,
apparently

This is just horrible!
Isn't there a better way?

Using big O complexity

```
void something(Object[] a) {
  for(int i = 0; i < a.length; i++)
    for (int j = 0; j < i; j++)
      for (int k = 0;   < j; k++)
        “something tha         e 1 step”
}
```

Three nested loops,
all running from
0 to at most n…
Answer: **O(n^3).**

Why ignore constant factors?

We lose some precision by throwing away constant factors

- ...you probably *do* care about a factor of 100 performance improvement

On the other hand, life gets much simpler:

- A small phrase like $O(n^2)$ tells you a lot about how the performance *scales* when the input gets big
- It turns out to be much easier to calculate big-O complexity than a precise formula

Big O is normally a good compromise!

Big O, formally

Big O measures the growth of a *mathematical function*

- Typically a function T(n) giving the number of steps taken by an algorithm on input of size n
- But can also be used to measure *space complexity* (memory usage) or anything else

So for the file-copying program:

- $T(n) = n^2/2$
- $T(n)$ is in $O(n^2)$

Big O, formally

What does it mean to say "$T(n)$ is $O(n^2)$"?

We could say it means $T(n)$ is proportional to n^2

- i.e. $T(n) = kn^2$ for some k

But this is too restrictive!

- What if e.g. $T(n) = kn(n-1)$, or $T(n) = kn^2 + 1$? We still want it to be $O(n^2)$!

Big O, formally

Instead, we say that $T(n)$ is $O(n^2)$ if:

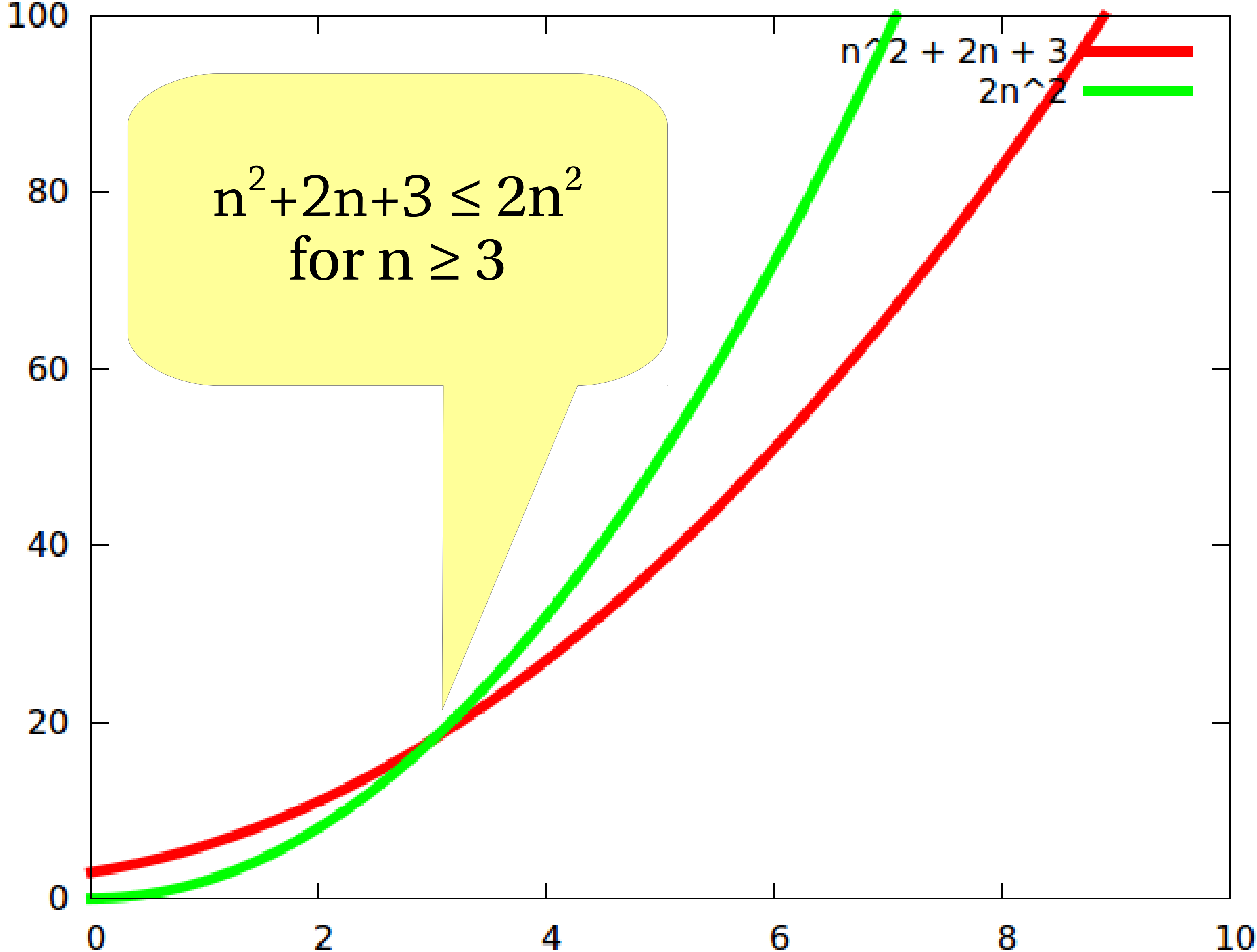
- $T(n) \leq kn^2$ for some k,
  i.e. $T(n)$ is proportional to n^2 *or lower*!
- This only has to hold for *big enough* n:
  i.e. for all n above some threshold n_0

If you draw the graphs of $T(n)$ and kn^2, at some point the graph of kn^2 must permanently overtake the graph of $T(n)$

- In other words, $T(n)$ grows more slowly than kn^2

Note that big-O notation is allowed to *overestimate* the complexity!

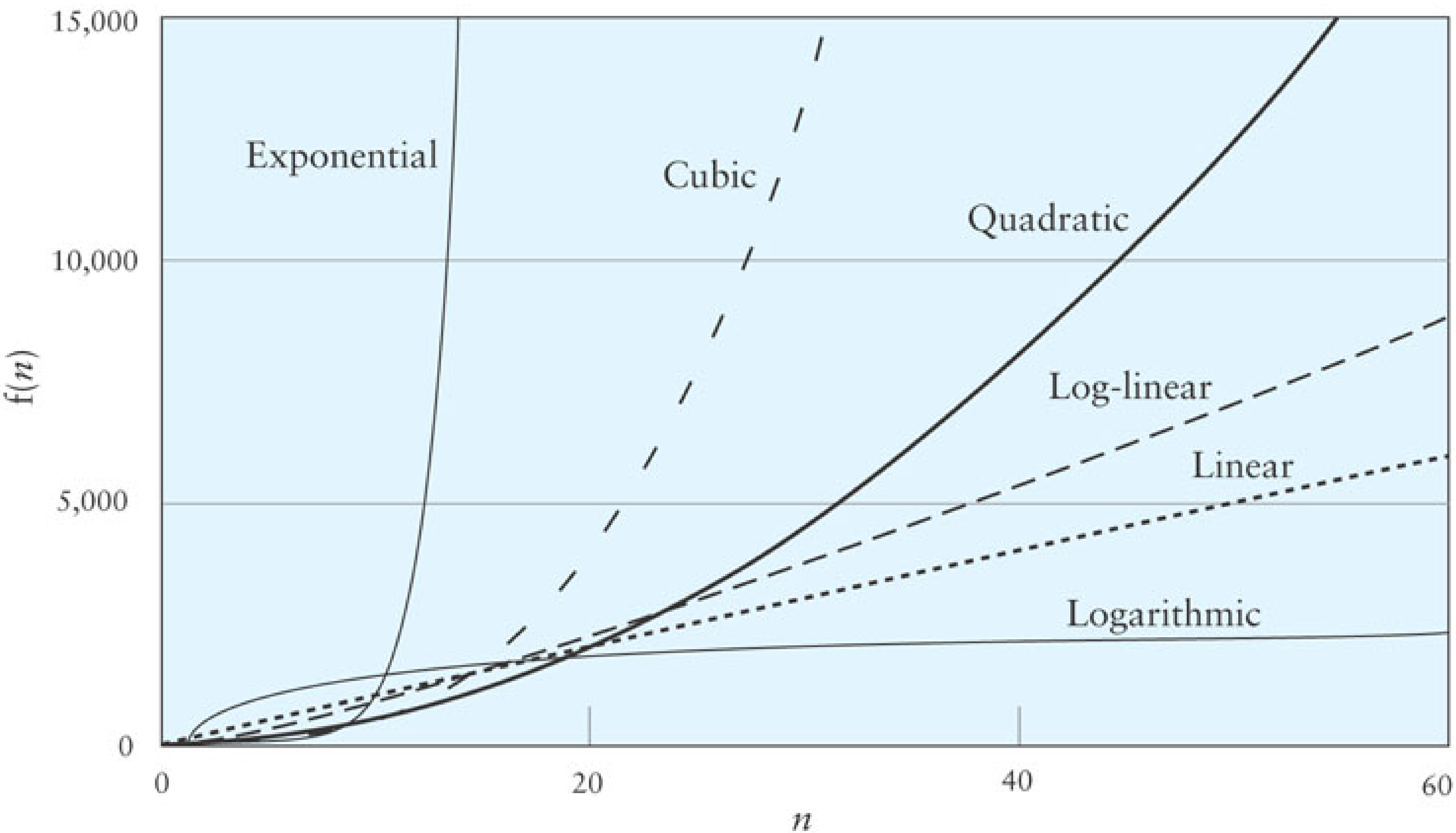
An example: $n^2 + 2n + 3$ is $O(n^2)$



Exercises

- Is $n^2 + 2n + 3$ in $O(n^3)$?
- Is it in $O(n)$?
- Is $3n + 5$ in $O(n)$?
- Why do we need the threshold?

Big-O	Name
$O(1)$	Constant
$O(\log n)$	Logarithmic
$O(n)$	Linear
$O(n \log n)$	Log-linear
$O(n^2)$	Quadratic
$O(n^3)$	Cubic
$O(2^n)$	Exponential
$O(n!)$	Factorial


15,000
10,000
5,000
0
f(n)
Exponential
Cubic
Quadratic
Log-linear
Linear
Logarithmic
0
20
40
60
n

Growth rates

Imagine that we double the input size from n to 2n.

If an algorithm is...

- $O(1)$, then it takes the same time as before
- $O(\log n)$, then it takes a constant amount more
- $O(n)$, then it takes twice as long
- $O(n \log n)$, then it takes twice as long plus a little bit more
- $O(n^2)$, then it takes four times as long

If an algorithm is $O(2^n)$, then adding *one element* makes it take twice as long

Big O tells you *how the performance of an algorithm is affected by the input size*

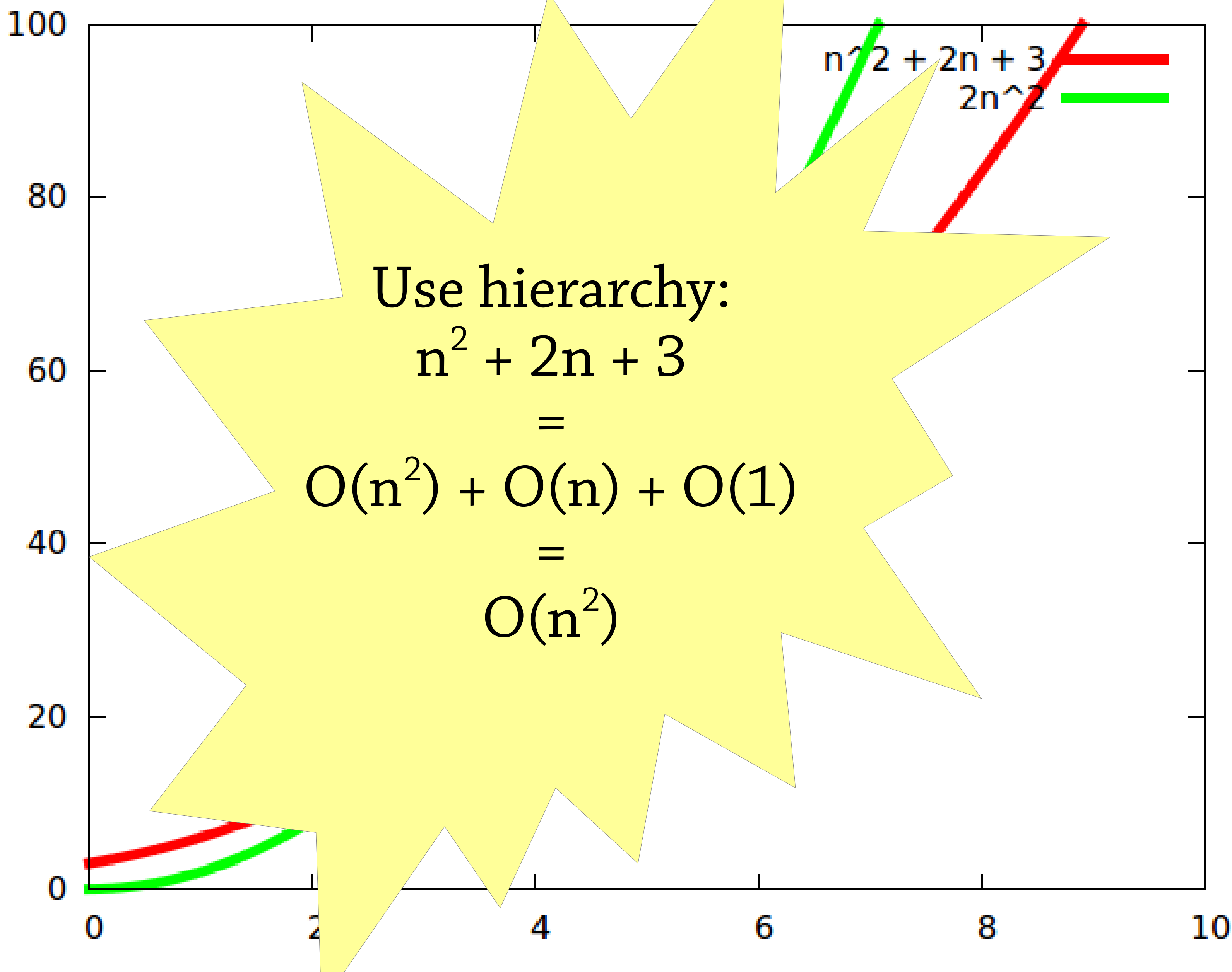
Adding big O (a hierarchy)

$O(1) < O(\log n) < O(n) < O(n \log n) < O(n^2) < O(n^3) < O(2^n)$

When adding a term lower in the hierarchy to one higher in the hierarchy, the lower-complexity term disappears:

- $O(1) + O(\log n) = O(\log n)$
- $O(\log n) + O(n^k) = O(n^k)$ (if $k \geq 0$)
- $O(n^j) + O(n^k) = O(n^k)$, if $j \leq k$
- $O(n^k) + O(2^n) = O(2^n)$

An example: $n^2 + 2n + 3$ is $O(n^2)$



Use hierarchy:

$$n^2 + 2n + 3$$
$$=$$
$$O(n^2) + O(n) + O(1)$$
$$=$$
$$O(n^2)$$

Quiz

What are these in Big O notation?

- $n^2 + 11$
- $2n^3 + 3n + 1$
- $n^4 + 2^n$

Just use hierarchy!

$n^2 + 11 = O(n^2) + O(1) = O(n^2)$

$2n^3 + 3n + 1 = O(n^3) + O(n) + O(1) = O(n^3)$

$n^4 + 2^n = O(n^4) + O(2^n) = O(2^n)$

Multiplying big O

$O(\text{this}) \times O(\text{that}) = O(\text{this} \times \text{that})$

- e.g., $O(n^2) \times O(\log n) = O(n^2 \log n)$

You can drop constant factors:

- $k \times O(f(n)) = O(f(n))$, if k is constant
- e.g. $2 \times O(n) = O(n)$

(Exercise: show that these are true)

Quiz

What is $(n^2 + 3)(2^n \times n) + \log_{10} n$ in Big O notation?

Answer

$(n^2 + 3)(2^n \times n) + \log_{10} n$
$= O(n^2) \times O(2^n \times n) + O(\log n)$
$= O(2^n \times n^3) + O(\log n)$ (multiplication)
$= O(2^n \times n^3)$ (hierarchy)

$\log_{10} n = \log n / \log 10$
i.e. log n times a constant factor

Complexity of a program

```
boolean unique(Object[] a) {
  for(int i = 0; i < a.length; i++)
    for (int j = 0; j < a.length; j++)
      if (a[i].equals(a[j]) && i != j)
        return false;
  return true;
}
```

Complexity of a program

```
boolean unique(Object[] a) {
  for(int i = 0; i < a.le
    for (int j = 0; j < a.len
      i   (a[i] equals(a[j]) && i != j)
                 e;
  ret
}
```

Outer loop runs
n times:
$O(n) \times O(n) = O(n^2)$

Inner loop runs
n times:
$O(n) \times O(1) = O(n)$

Loop body:
$O(1)$

Complexity of loops

The complexity of a loop is:
the number of times it runs
times the complexity of the body

What about this one?

```
void function(int n) {
  for(int i = 0; i < n*n; i++)
    for (int j = 0; j < n/2; j++)
      “something taking O(1) time”
}
```

What about this one?

```
void function(int n) {
  for(int i = 0; i < n*n,
    for (int j = 0; j < n/2; j
      “s ething taking O(1) time”
}
```

Outer loop runs
n^2 times:
$\mathbf{O(n^2)} \times O(n) = O(n^3)$

Inner loop runs
$n/2 = \mathbf{O(n)}$ times:
$O(n) \times O(1) = O(n)$

Loop body:
$O(1)$

Here's a new one

```
boolean unique(Object[] a) {
  for(int i = 0; i < a.length; i++)
    for (int j = 0; j < i; j++)
      if (a[i].equals(a[j]))
        return false;
  return true;
}
```

Here's a new one

```
boolean unique(Object[] a) {
  for(int i = 0; i < a.length; i++)
    for (int j = 0; j < i; j++)
      if (a[i].equals(a[j]))
            e;
  ret
}
```

Inner loop is
$i \times O(1) = O(i)$??
But it should be
in terms of n?

Body is $O(1)$

Here's a new one

```
boolean unique(Object[] a) {
  for(int i = 0; i < a.length; i++)
    for (int j = 0; j < i; j++)
      if (a[i].equals(a[j]))
        e;
  ret
}
```

i < n, so **i is O(n)**
So loop runs **O(n)**
times, complexity:
O(n) × O(1) = O(n)

Body is O(1)

Here's a new one

```
boolean unique(Object[] a) {
  for(int i = 0; i < a.le
    for (int j = 0; j < i; j+
      i  (a[i].equals(a[j]))
              e;
  ret
}
```

Outer loop runs n times: $O(n) \times O(n) = O(n^2)$

$i < n$, so **i is $O(n)$**
So loop runs **$O(n)$** times, complexity:
$O(n) \times O(1) = O(n)$

Body is $O(1)$

The example from earlier

```
void something(Object[] a) {
  for(int i = 0; i < a.length; i++)
    for (int j = 0; j < i; j++)
      for (int k = 0; k < j; k++)
        “something that takes 1 step”
}
```

$i < n$, $j < n$, $k < n$,
so all three loops run **O(n)** times
Total runtime is
$O(n) \times O(n) \times O(n) \times O(1) = \mathbf{O(n^3)}$

What's the complexity?

```
void something(Object[] a) {
  for(int i = 0; i < a.length; i++)
    for (int j = 1; j < a.length; j *= 2)
      … // something taking O(1) time
}
```

Outer loop is
O(n log n)

What's the complexity?

Inner loop is
O(log n)

```
void something(Object[] a) {
  for(int i = 0; i < a.length; i+
    for (int j = 1; j < a.length; j *= 2)
      … // something taking O(1) time
}
```

A loop running through i = 1, 2, 4, …, n
runs **O(log n)** times!

While loops

```
long squareRoot(long n) {
    long i = 0;
    long j = n+1;
    while (i + 1 != j) {
        long k = (i + j) / 2;
        if (k*k <= n) i = k;
        else j = k;
    }
    return i;
}
```

Each iteration takes O(1) time…
but how many times does the loop run?

While loops

```
long squareRoot(long n) {
    long i = 0;
    long j = n+1;
    while (i + 1 != j) {
        long k = (i + j) / 2;
        if (k*k <= n) i = k;
        else j = k;
    }
    return i;
}
```

Each iteration takes O(1) time

...and halves `j-i`, so **O(log n)** iterations

Summary: loops

Basic rule for complexity of loops:

- Number of iterations times complexity of body
- for (int i = 0; i < n; i++) ...: n iterations
- for (int i = 1; i ≤ n; i *= 2): O(log n) iterations
- While loops: same rule, but can be trickier to count number of iterations

If the complexity of the body depends on the value of the loop counter:

- e.g. O(i), where 0 ≤ i < n
- round it up to O(n)!

Sequences of statements

What's the complexity here?
(Assume that the loop bodies are O(1))

```
for (int i = 0; i < n; i++) …
for (int i = 1; i < n; i *= 2) …
```

Sequences of statements

What's the complexity here?
(Assume that the loop bodies are O(1))

```
for (int i = 0; i < n; i++) …
for (int i = 1; i < n; i *= 2) …
```

First loop: **O(n)**
Second loop: **O(log n)**
Total: O(n) + O(log n) = **O(n)**

For sequences, add the complexities!

A familiar scene

```
int[] array = {};
for (int i = 0; i < n; i++) {
  int[] newArray =
    new int[array.length+1];
  for (int j = 0; j < i; j++)
    newArray[j] = array[j];
  newArray = array;
}
```

Assume that each statement takes O(1) time

A familiar scene

```
int[] array = {};
for (int i = 0; i < n; i
  int[] newArray =
    new int[array.length+1];
  for (int j   0; j < i; j++)
    newArray[j   = array[j];
  newArray =
}
```

Rest of loop body **O(1)**,
so loop body
O(1) + O(n) = **O(n)**

Outer loop:
n iterations,
O(n) body,
so **O(n^2)**

Inner loop
O(n)

A familiar scene, take 2

```
int[] array = {};
for (int i = 0; i < n; i+=100) {
  int[] newArray =
    new int[array.length+100];
  for (int j = 0; j < i; j++)
    newArray[j] = array[j];
  newArray = array;
}
```

A familiar scene, take 2

```
int[] array = {};
for (int i = 0; i < n; i+=100) {
  int[] newArray =
    new int[array.length+100];
  for (int j = 0; j < i; j++)
    newArray[j] = array[j];
  newArray =
}
```

Outer loop:
n/100 iterations,
which is O(n)
O(n) body,
so **$O(n^2)$** still

A familiar scene, take 3

```
int[] array = {0};
for (int i = 1; i <= n; i*=2) {
  int[] newArray =
    new int[array.length*2];
  for (int j = 0; j < i; j++)
    newArray[j] = array[j];
  newArray = array;
}
```

A familiar scene, take 3

```
int[] array = {0};
for (int i = 1; i <= n; i*=2) {
  int[] newArray =
    new int[array.length*2];
  for (int j = 0; j < i; j++)
    newArray[j] = array[j];
  newArray =
}
```

Outer loop:
log n iterations,
O(n) body,
so **O(n log n)**??

A familiar scene, take 3

```
int[] array = {0};
for (int i = 1; i <= n; i*=2) {
  int[] newArray =
    new int[array.length*2];
  for (int j = 0; j < i; j++)
    newArray[j] = array[j];
  newArray =
}
```

Here we
“round up”
O(i) to O(n).
This causes an
overestimate!

A complication

Our algorithm has $O(n)$ complexity, but we've calculated $O(n \log n)$

- An overestimate, but not a severe one

  (If $n = 1000000$ then $n \log n = 20n$)
- This can happen but is normally not severe
- To get the right answer: do the maths

Good news: for “normal” loops this doesn't happen

- If all bounds are n, or n^2, or another loop variable, or a loop variable squared, or …

Main exception: loop variable i doubles every time, body complexity depends on i

Doing the sums

In our example:

- The inner loop's complexity is $O(i)$
- In the outer loop, i ranges over $1, 2, 4, 8, \ldots, 2^a$

Instead of rounding up, we will add up the time for all the iterations of the loop:

$$1 + 2 + 4 + 8 + \ldots + 2^a$$

$$= 2 \times 2^a - 1 < 2 \times 2^a$$

Since $2^a \leq n$, the total time is at most $2n$, which is $O(n)$

A last example

```
for (int i = 1; i <= n; i *= 2) {
   for (int j = 0; j < n*n; j++)
      for (int k = 0; k <= j; k++)
         // O(1)
   for (int j = 0; j < n; j++)
      // O(1)
}
```

A last example

The outer loop runs O(log n) times

The j-loop runs n^2 times

```
for (int i = 1; i <= n; i *= 2) {
    for (int j = 0; j < n*n; j++)
        for (int k = 0; k <= j; k++)
            // O(1)
    for (int j = 0; j < n; j++)
        // O(1)
}
```

This loop is O(n)

`k <= j < n*n`
so this loop is $O(n^2)$

Total: $O(\log n) \times (O(n^2) \times O(n^2) + O(n))$
$= O(n^4 \log n)$

A trick: sums are almost integrals

$$\sum_{x=a}^{b} f(x) \approx \int_{a}^{b} f(x)$$

For example:

$$\sum_{i=0}^{n} i = n(n+1)/2 \qquad \int_{0}^{n} x\,dx = n^2/2$$

Not quite the same, but close!

This trick is accurate enough to give you the right complexity class

See: "Finite calculus: a tutorial for solving nasty sums"

Summary

Big O complexity

- Calculate runtime without doing hard sums!
- Lots of "rules of thumb" that work almost all of the time
- Very occasionally, still need to do hard sums :(
- Ignoring constant factors: seems to be a good tradeoff

Weiss chapter 5