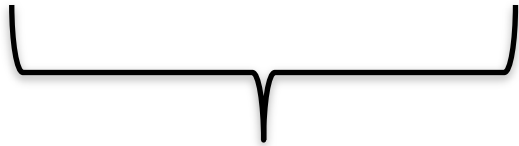


Monads

Monads seen so far:

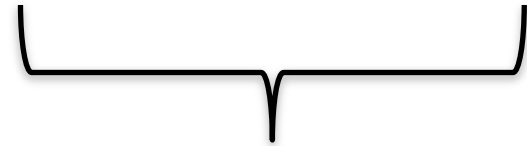
IO vs Gen

IO A



- Instructions to build a value of type A by interacting with the operating system
- Run by the ghc runtime system

Gen A



- Instructions to create a random value of type A
- Run by the QuickCheck library functions to perform random tests

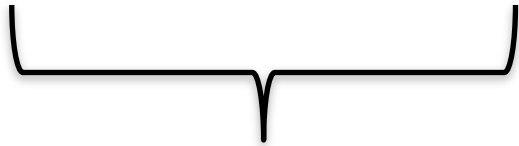
Terminology

- A “*monadic value*” is just an expression whose type is an instance of class Monad
- “*t is a monad*” means t is an instance of the class Monad
- We have often called a monadic value an “*instruction*”. This is not standard terminology – but sometimes they are called “actions”

Monads seen so far:

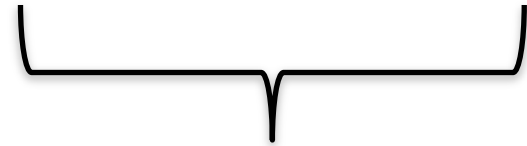
IO vs Gen

IO A



- Instructions to build a value of type A by interacting with the operating system
- Run by the `ghc` runtime system

Gen A



- Instructions to create a random value of type A
- Run by the QuickCheck library functions to perform random tests

Monads = Instructions

- What is the type of doTwice?

```
Main> :i doTwice  
doTwice :: Monad m => m a -> m (a, a)
```

Even the *kind of instructions*
can vary!

Different kinds of instructions,
depending on who obeys
them.

Whatever kind of
result argument
produces, we get a
pair of them

Monads and do notation

- To be an instance of class Monad you need (as a minimal definition) two operations: **>>=** and **return**

```
class Monad m where
  (>>=) :: m a -> (a -> m b) -> m b
  return :: a -> m a
```

```
(>>) :: m a -> m b -> m b
x >> y = x >>= \_ -> y
```

```
fail :: String -> m a
fail msg = error msg
```

Default implementations

The truth about Do

- Do syntax is just a shorthand:

```
do act1
  act2 == act1 >> act2 == act1 >>= \_ -> act2
```

```
do v <- act1
  act2 == act1 >>= \v -> act2
```

The truth about Do

Full translation (I)

```
do act1
  ...
  actn
```

 ==

```
act1 >> do ...
          actn
```

```
do v <- act1
  ...
  actn
```

 ==

```
act1 >>= \v -> do ...
          actn
```

```
do actn
```

 ==

```
actn
```


Example

```
foo :: IO ()  
foo = do  
    filename <- getLine  
    contents <- readFile filename  
    putStrLn contents
```

```
do v <- act1  
...  
actn
```

=

```
act1 >>= \v -> do ...  
                    actn
```

Example

```
foo :: IO ()
foo = do
    filename <- getLine
    contents <- readFile filename
    putStrLn contents
```

```
foo' = getLine >>= \filename ->
    do
        contents <- readFile filename
        putStrLn $ take 100 contents
```

Example

```
foo' = getLine >>= \filename ->
      do
        contents <- readFile filename
        putStrLn $ take 100 contents
```

```
do v <- act1
   ...
   actn
```

=

```
act1 >>= \v -> do ...
                  actn
```

```
do actn
```

=

```
actn
```

Equivalent to

```
foo' = getLine >>= \filename ->  
      do  
        contents <- readFile filename  
        putStrLn $ take 100 contents
```

```
foo'' = getLine >>= \filename ->  
        readFile filename >>= \contents ->  
        putStrLn $ take 100 contents
```

Eta conversion & Monad Law

```
foo'' = getLine >>= \filename ->  
        readFile filename >>= \contents ->  
        putStrLn $ take 100 contents
```

$$f (g x) == (f . g) x$$
$$\backslash v \rightarrow f v == f$$
$$\text{act1} \gg= \backslash v \rightarrow f v \gg= g$$
$$=$$
$$\text{act1} \gg= f \gg= g$$

Eta conversion & Monad Law

```
foo'' = getLine >>= \filename ->  
        readFile filename >>= \contents ->  
        putStrLn $ take 100 contents
```

```
foo''' = getLine >>= readFile >>= putStrLn . take 100
```

The truth about Do

Full Translation (II): Let and pattern matching

```
do let p = e  
  ...  
  actn
```

==

```
let p = e in  
do ...  
  actn
```

```
do pattern <- act1  
  ...  
  actn
```

==

```
let f pattern = do ...  
                actn  
    f _         = fail "Error"  
in  act1 >>= f
```

Pictures from a blog post about functors, applicatives and monads

[http://adit.io/posts/2013-04-17-
functors,_applicatives,_and_monads_in_pictures.html](http://adit.io/posts/2013-04-17-functors,_applicatives,_and_monads_in_pictures.html)

Aditya Y. Bhargava



getline

```
getline :: IO String
```



readFile

```
readFile :: FilePath -> IO String
```



putStrLn

```
putStrLn :: String -> IO ()
```

All three functions take a value (or no value) and produce an IO “wrapped” value

The function `>>=` allows us to join them together

`getLine >>= readFile >>= putStrLn`

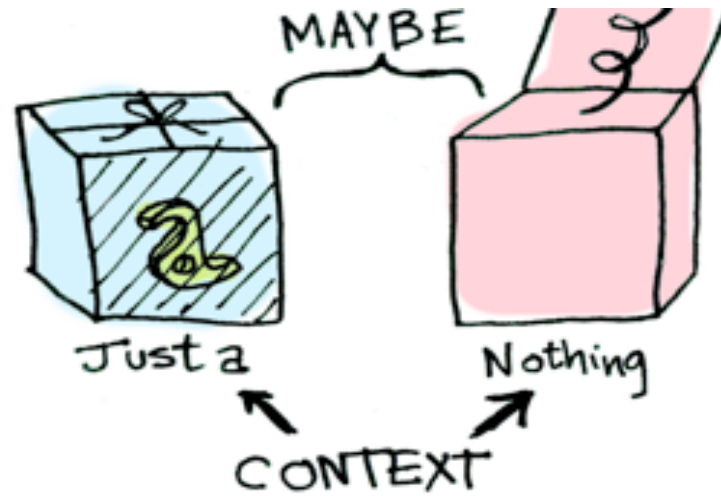


1. GET USER
INPUT

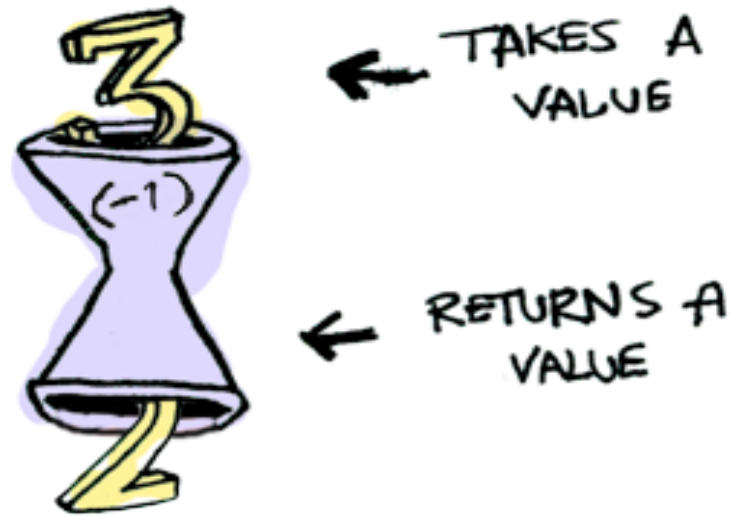


2. USE IT TO
READ A FILE

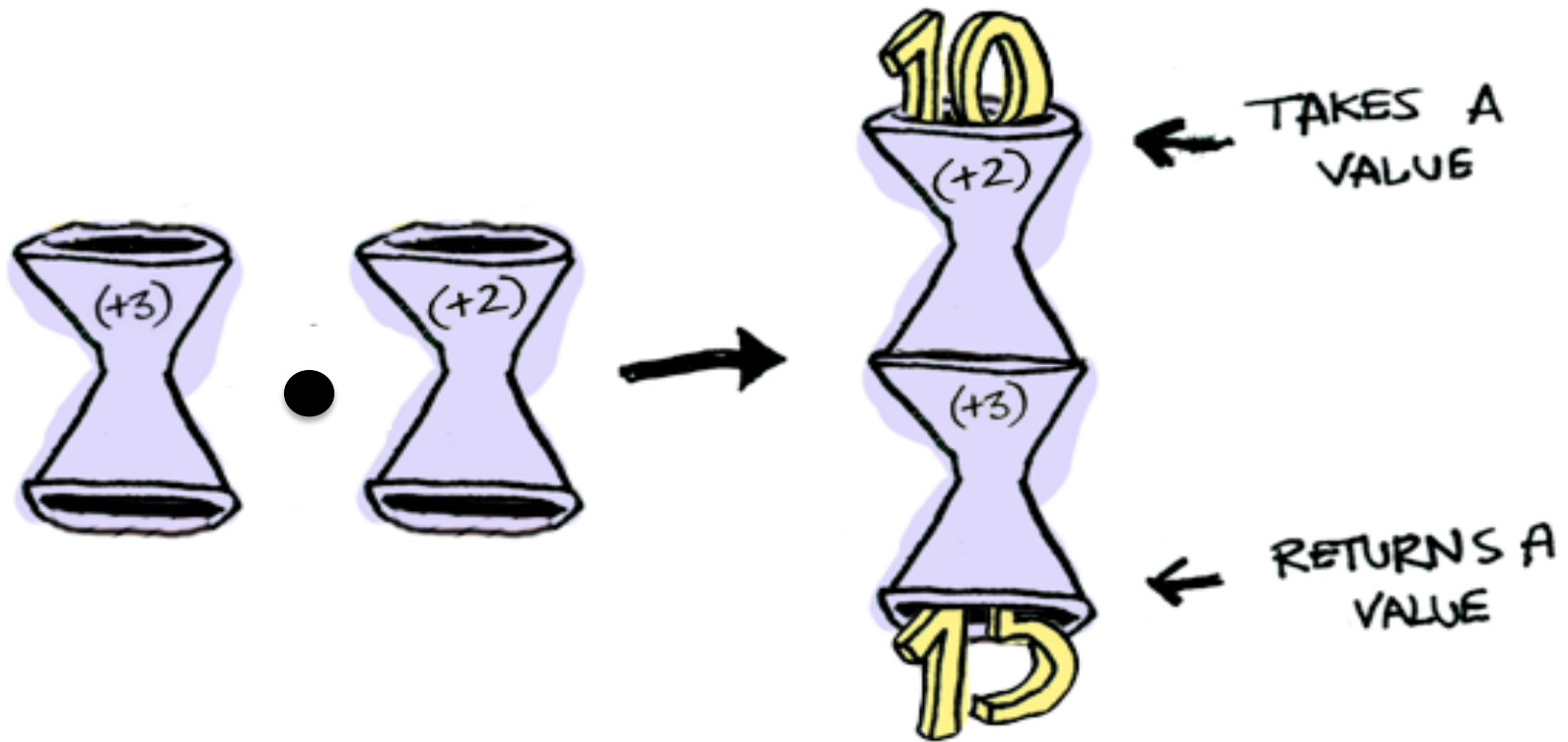
Maybe



Here is a function

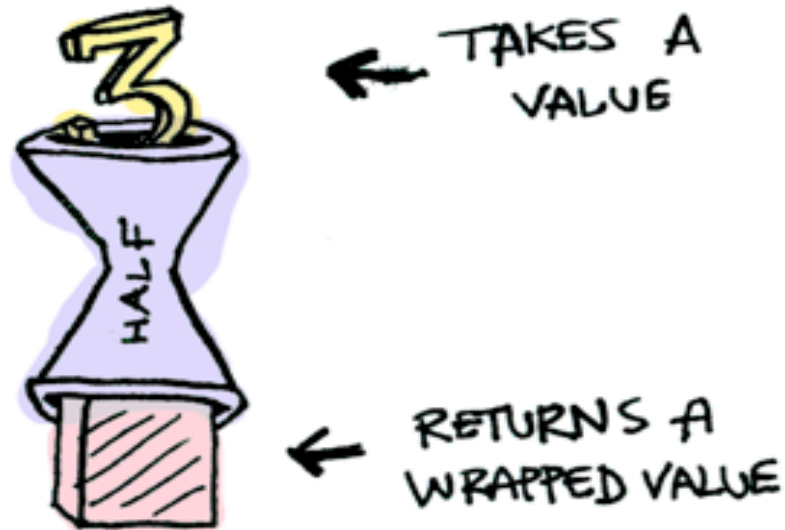


They can be composed



Here is a function

```
half x = if even x  
        then Just (x `div` 2)  
        else Nothing
```



What if we feed it a wrapped value?



We need to use `>>=` to shove our wrapped value into the function

>>=



>>=

Here's how it works:

```
> Just 3 >>= half
Nothing
> Just 4 >>= half
Just 2
> Nothing >>= half
Nothing
```

What's happening inside? `Monad` is another typeclass. Here's a partial definition:

```
class Monad m where
  (>>=) :: m a -> (a -> m b) -> m b
```

$\gg=$

$(\gg=) :: m a \rightarrow (a \rightarrow m b) \rightarrow m b$

1. $\gg=$ TAKES
A MONAD
(LIKE **Just** 3)

2. AND A
FUNCTION THAT
RETURNS A MONAD
(LIKE **half**)

3. AND IT
RETURNS
A MONAD



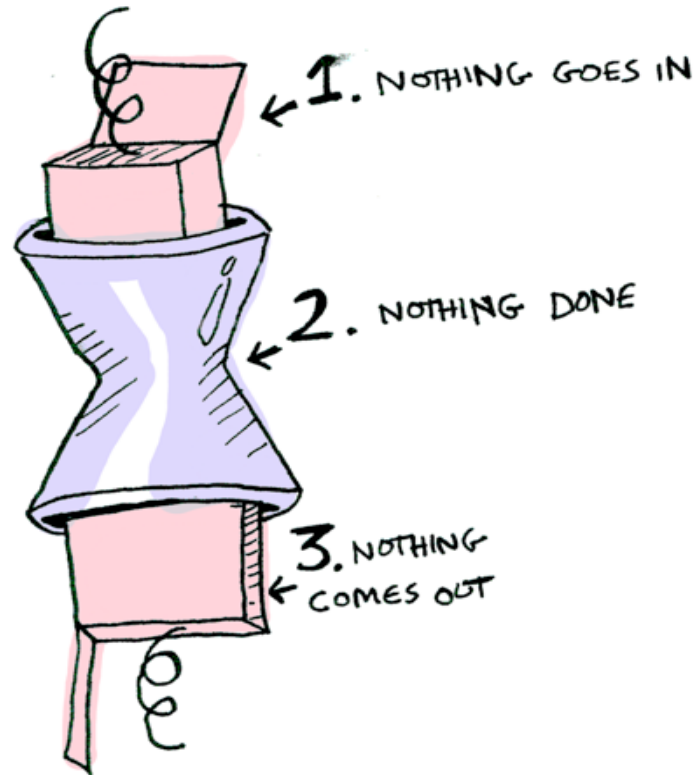
1. BIND UNWRAPS THE VALUE

2. FEEDS THE UNWRAPPED VALUE INTO THE FUNCTION



3. WRAPPED VALUE COMES OUT

>>=





Just 20 >>= half >>= half
>>= half







Instance Monad Maybe

- Maybe is a very simple monad

```
instance Monad Maybe where
  Just x  >>= k = k x
  Nothing >>= _ = Nothing

  return      = Just
  fail s      = Nothing
```

Although simple it can be useful...

Congestion Charge Billing



Vardagar
(ej dag före sön- och helgdag)

Kl	Kr
0630 - 0659	10:-
0700 - 0729	15:-
0730 - 0829	20:-
0830 - 0859	15:-
0900 - 1529	10:-
1530 - 1559	15:-
1600 - 1729	20:-
1730 - 1759	15:-
1800 - 1829	10:-

Här hamnar betalstationerna
I dag fattar kommunfullmäktige beslut om trängselskatt. Kommunstyrelsens förslag är en inre ring runt centrala Göteborg med 40 betalstationer.



Källa: Vägverket

Grafik: GP

Congestion Charge Billing

Registration number used to find the
Personnummer of the owner

`carRegister :: [(RegNr,PNr)]`

Personnummer used to find the name of the
owner

`nameRegister :: [(PNr,Name)]`

Name used to find the address of the owner

`addressRegister :: [(Name,Address)]`

Example:

Congestion Charge Billing

```
type CarReg = String ; type PNr = String
type Name = String ; type Address = String
```

```
carRegister :: [(CarReg,PNr)]
carRegister
  = [ ("JBD 007", "750408-0909"), ... ]
```

```
nameRegister :: [(PNr,Name)]
nameRegister
  = [ ("750408-0909", "Dave"), ... ]
```

```
addressRegister :: [(Name,PNr),Address]
addressRegister =
  [ ( ("Dave", "750408-0909"), "42 Streetgatan\n Askim" )
  , ... ]
```

Example:

Congestion Charge Billing

With the help of

`lookup :: Eq a => a -> [(a,b)] -> Maybe b`
we can return the address of car owners

```
billingAddress :: CarReg -> Maybe (Name, Address)
billingAddress car =
  case lookup car carRegister of
    Nothing -> Nothing
    Just pnr -> case lookup pnr nameRegister of
      Nothing -> Nothing
      Just name ->
        case lookup (name,pnr) addressRegister of
          Nothing -> Nothing
          Just addr -> Just (name,addr)
```

Example:

Congestion Charge Billing

Using the fact that Maybe is a member of class Monad we can avoid the spaghetti and write:

```
billingAddress car = do
  pnr  <- lookup car carRegister
  name <- lookup pnr nameRegister
  addr <- lookup (name,pnr) addressRegister
  return (name,addr)
```


Example:

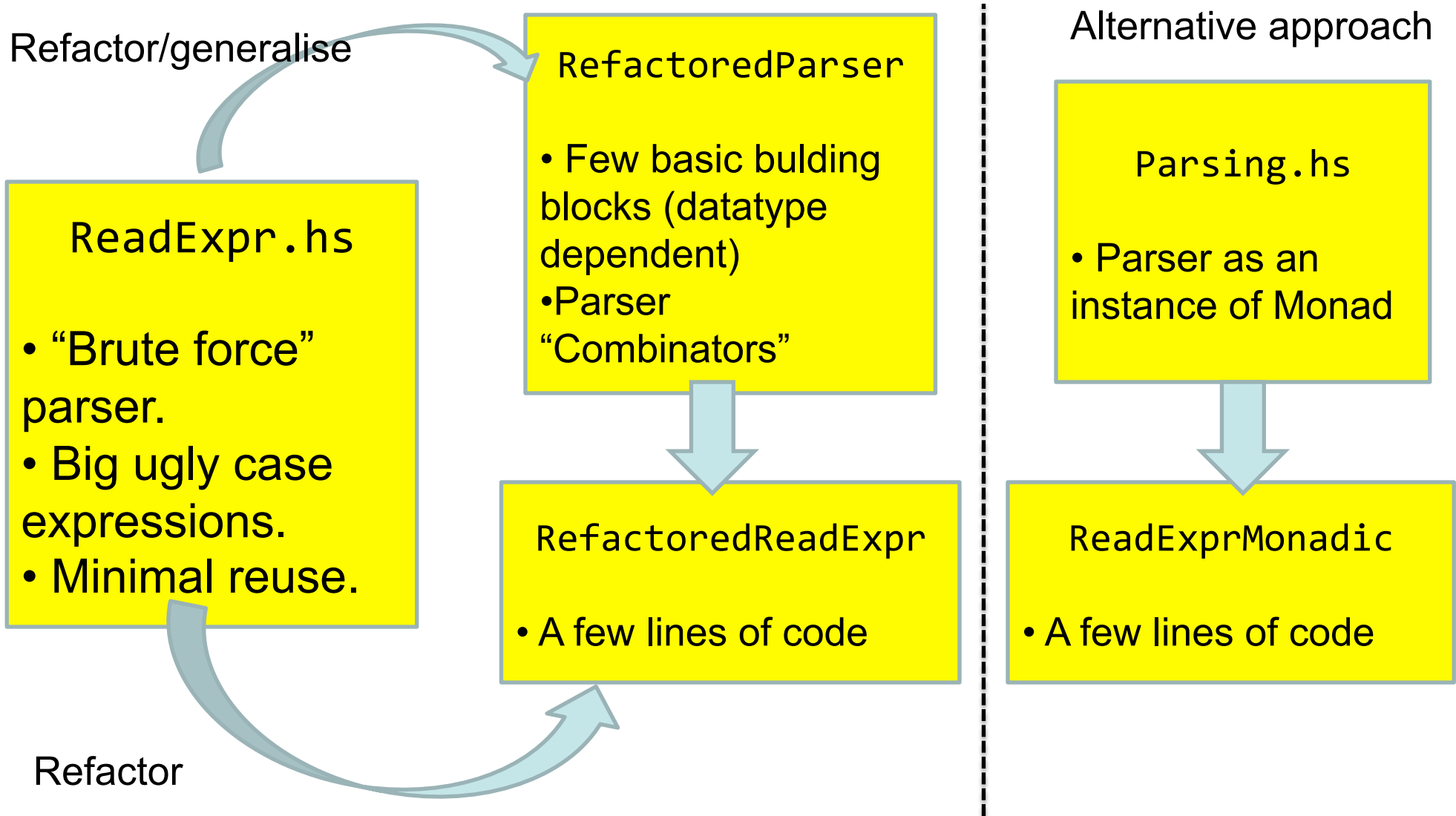
Congestion Charge Billing

Unrolling one layer of the do syntactic sugar:

```
billingAddress car ==  
  lookup car carRegister >>= \pnr ->  
  do  
    name <- lookup pnr nameRegister  
    addr <- lookup (name,pnr) addressRegister  
    return (name,addr)
```

- `lookup car carRegister` gives `Nothing`
then the definition of `>>=` ensures that the whole
result is `Nothing`
- `return` is `Just`

The Big Picture



Monad

- To be an instance of class Monad you need two operations: **>>=** and **return**

```
instance Monad Parser where
  return = succeed
  (>>=)   = (>*>)
  -- (>->) is equivalent to (>>)
```

- Why bother?

- First example of a home-grown monad
- Can understand and use do notation

Summary

- We can use higher-order functions to build Parsers from other more basic Parsers.
- Parsers can be viewed as an instance of Monad
- We can build our own Monads!
 - A lot of "plumbing" is nicely hidden away
 - The implementation of the Monad is not visible and can thus be changed or extended

Parsing

- So far: how to write

```
readExpr :: String -> Maybe Expr
```

- Key idea:

```
type Parser = String -> Maybe (a, String)
```

- This lecture: Building Parsers; Parsers as a new type of "instructions" – i.e. a monad.

Recall some key building blocks

```
succeed :: a -> Parser a
```

```
succeed a = P $ \s -> Just(a,s)
```

```
sat :: (Char -> Bool) -> Parser Char
```

```
(>->) :: Parser a -> Parser b -> Parser b
```

```
(>*>) :: Parser a -> (a -> Parser b) -> Parser b
```

```
Main> parse (digit >*> \a -> sat (==a)) "22xx"
```

```
Just ('2',"xxx")
```

```
Main> parse (digit >*> \a -> sat (==a)) "12xx"
```

```
Nothing
```

The Parser Monad

- Using these building blocks we can make Parser an instance of the class Monad
 - We get a language of “Parsing Instructions”
 - Another way to write Parsers using do notation
 - Deeper understanding of Monads

Example

- recall doTwice

```
doTwice :: Monad m => m a -> m (a,a)
doTwice cmd =
  do a <- cmd
     b <- cmd
     return (a,b)
```

```
Main> parse (doTwice digit) "9876"
Just (('9','8'), "76")
```


Example revisited: Parsing Expressions

```
expr :: Parser Expr
expr s1 = case parse num s1 of
  Just (a,s2) -> case s2 of
    '+' : s3 -> case parse expr s3 of
      Just (b,s4) -> Just (Add a b, s4)
      Nothing    -> Just (a,s2)
    _       -> Just (a,s2)
  Nothing   -> Nothing
```

modified to use the new version of Parser type. Otherwise as before

Monadic style abstracts away from implementation of the Parser type

```
expr :: Parser Expr
expr = do a <- num
       do char '+'
       b <- expr
       return (Add a b)
+++ return a
```

Parser Combinators

```
zeroOrMore, oneOrMore :: Parser a -> Parser [a]
```

```
zeroOrMore p = oneOrMore p +++ return []
```

```
oneOrMore p = do v <- p  
                 vs <- zeroOrMore p  
                 return(v:vs)
```

```
Main> parse (oneOrMore digit) "9876+"  
Just ("9876","+")
```

Combinator: a function which take functions as arguments and produces a function as a result

Parser Combinators

```
nat :: Parser Int -- Parses a non negative integer
nat = do xs <- oneOrMore number
      return (read xs)

int :: Parser Int
int = nat +++
      do char '-'
         n <- nat
         return (-n)
```

Chain

```
chain p op f = P $ \s1 ->
  case parse p s1 of
    Just (a,s2) -> case s2 of
      c:s3 | c == op -> case chain p op f s3 of
        Just (b,s4) -> Just (f a b, s4)
        Nothing    -> Just (a,s2)
      _ -> Just (a,s2)
    Nothing -> Nothing
```

Old definition (modified to work with the new type)

```
chain' p op f = do v <- p
  vs <- zeroOrMore (char op >> p)
  return (foldr1 f (v:vs))
```

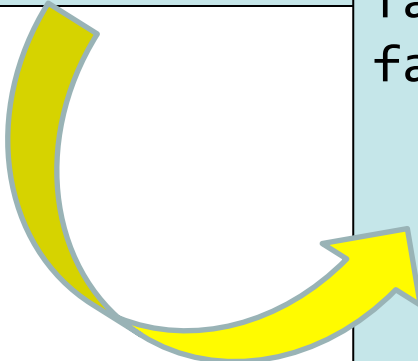
Prelude.foldr1 : fold operation for lists with at least one element (no "nil" case)

Factor

```
factor :: Parser Expr
factor ('(':s) =
  case expr s of
    Just (a, ')':s1) -> Just (a, s1)
    _                 -> Nothing

factor s = num s
```

```
factor :: Parser Expr
factor = num +++
  do char '('
  e <- expr
  char ')'
  return e
```



IO t

- Instructions for interacting with operating system
- Run by GHC runtime system produce value of type t

Gen t

- Instructions for building random values
- Run by **quickCheck** to generate random values of type t

Parser t

- Instructions for parsing
- Run by **parse** to parse a string and **Maybe** produce a value of type t

Three Monads

Code

- `Parsing.hs`
 - module containing the parser monad and simple parser combinators.
- `ReadExprMonadic.hs`
 - A reworking of `Read`

See course home page

- We can build our own Monads!
 - A lot of "plumbing" is nicely hidden away
 - A powerful pattern, used widely in Haskell
 - A pattern that can be used in other languages, but syntax support helps
 - F# computation expressions
 - Scala

More examples

- <http://adit.io/posts/2013-06-10-three-useful-monads.html>
- stack (slides/video from last year)

Another Example: A Stack

- A Stack is a stateful object
- Stack operations can push values on, pop values off, add the top elements

```
type Stack = [Int]
newtype StackOp t = StackOp (Stack -> (t,Stack))

-- the type of a stack operation that produces
-- a value of type t
pop :: StackOp Int
push :: Int -> StackOp ()
add :: StackOp ()
```

Running a StackOp

```
type Stack = [Int]
newtype StackOp t = StackOp (Stack -> (t,Stack))

run (StackOp f) = f

-- run (StackOp f) state = f state
```

Operations

```
pop :: StackOp Int
pop = StackOp $ \(x:xs) -> (x,xs) -- can fail

push :: Int -> StackOp ()
push i = StackOp $ \(s) -> ((),i:s)

add :: StackOp ()
add = StackOp $ \(x:y:xs) -> ((),x+y:xs) -- can fail
```

Building a new StackOp...

```
swap :: StackOp ()
swap = StackOp $ \s ->
    let (x,s') = run pop s
        (y,s'') = run pop s'
        (_,s''') = run (push x) s''
        (_,s''''') = run (push y) s''''
    in (_, s''''')
```

No thanks!

StackOp is a Monad

- Stack instructions for producing a value

```
-- (>>=) :: StackOp a -> (a -> StackOp b) -> StackOp b
instance Monad StackOp
  where return n = StackOp $ \s -> (n,s)
        sop >>= f  = StackOp $ \s ->
                                let (i,s') = run sop s
                                in run (f i) s'
```

So now we can write...

```
swap = do  
  a <- pop  
  b <- pop  
  push a  
  push b
```

Stack t

- Stack instructions producing a value of type t
- Run by **run**

Maybe t

- Instructions for either producing a value or nothing
- Run by ?? (not an abstract data type)

Two More Monads