

Erlang 24/7 Chalmers 2014-02-26 Cons T Åhs Keeper of The Code <u>cons@klarna.com</u> @lisztspace

Cons T Åhs

- Keeper of The Code at klarna
 - Architecture The Big Picture
 - Development getting ideas to work
 - Code Quality care about the details
 - Increase competence of developers better developers are more productive and delivers better solutions

Cons T Åhs

- Student in Uppsala 1981 85 (writing code for five years, Lisp before starting in Uppsala)
- Uppsala University (1985-2000, 2002, 2009, 2012), research & teaching; foundations, algorithms, functions, relations, objects, compilers, pragmatics, theory, theorem proving, formal program correctness..
- Consultant (1991-); online poker, low level networking, medical imaging, graphics, finance, musical notation (Lisp), speech synthesis (Lisp, Prolog), compilers (Lisp, Prolog), real time video decoding, teaching.
- Klarna from Feb. 28, 2011

Klarna - The Business

- Make shopping on the net simpler, safer, more fun.
- Pay by invoice after the goods are delivered
- Customer checks out at estore
 - Klarna identifies customer and investigates credit
 - estore sends goods and invoice
 - Klarna pays estore (Klarna takes the risk)
 - customer pays Klarna
- Identification and risk determination needs to be done fast (a few seconds)



Klarna - The Business

- Klarna makes money by taking a calculated risk
 - We're buying and selling invoices
- What do we need to be good at?
 - Identifying our customers do you exist?
 - Interesting algorithms, lots of data
 - Evaluate risk are you going to pay?
 - Interesting algorithms, lots of data
 - Bookkeeping we're essentially a bank.
 - Only lots of data, no algorithms



Klarna - The Facts

- Founded in 2005
- Revenue doubled every year from start we're growing exponentially
- Sweden, Norway, Denmark, Finland, Germany, Netherlands, Austria more to come..
- Over 800 employees
- Over 15K estores and growing

Very nice for the shareholders..

2005 2006 2007 2008 2009 2010 2011 2012 2013



Klarna - The Facts

- Currently over 2.5M transactions/month
 - average ≈ 1/s, but there are peaks both over the day and the year
 - you can't build for the average
- People shop all the time
 - Available 24/7 no downtime
 - Software upgrades with no downtime
 - Hardware upgrades and relocation with no downtime

.. we need to build and maintain a robust system

Klarna - The technology

- Use technology from a domain with similar needs
 - large amount of messages/transactions
 - high demands on availability, scalability, robustness
 - low tolerance for downtime for software or hardware upgrades
- We're using Erlang/otp
 - "functional," state is handled in processes
 - easy to distribute and communicate
 - can handle large amounts of processes
 - robust
 - soft real time



Erlang A functional language

- Dynamically typed functional language
- No side effects; variables are bound once and the value can not be changed
- Every expression computes a value
- Pattern matching provides parallel binding and compact programs (mixed blessing beware!)
- Looks very much like Prolog
- The power of higher order functions and closures
- It is easier to understand, reason about, compile and debug a functional program.



24/7 implies interesting challenges

- 24/7 means having a goal of 100% availability
- 99.99% availability translates to about 4 minutes of allowed downtime/month
- There is not too much you can do in 4 minutes
- No planned shutdown, even for upgrades
- Code and data format change without downtime
- Subsystem switching without downtime
- Hardware switching without downtime
- etc..
- There is no notion of stopping and starting the system it just runs..



Joe Armstrong says:

Each year your sequential programs will go slower.

Each year your concurrent programs will go faster.

Concurrent and distributed programming

- With concurrent programming troubles form when you have a shared <u>and</u> mutable state.
- Problem typically solved by using synchronisation with locks
 - Complicated you have to know when to lock
 - Can lead to more problems performance degradation
 - Cooperative model all parts of the program must agree
- Take away one and your on safe ground.
- Erlang takes away both!



No shared state, no mutable state

- Each process has a state of its own, or rather a sequence of states; possibly a new state after receiving a message
- Each process has a private heap
- Each process has a message queue (the implementation handles these)
- Processes can not share state, even when they live in the same VM.
 - All communication must be done with messages.
 - Asynchronous message passing messages are <u>copied</u> between processes



No shared state

- Why?
 - Background (telecom switches) with a large number of small and short lived processes
 - When a process dies there is no risk reclaiming the whole process
 - No other process can access the memory it used
 - Nothing happens if you send a message to a dead pid
 - The dead process can not reference the memory of another process
 - Leads to robustness



Erlang built for fault tolerance

- No shared state means that a crashing process will not take another process down.
- A crashing process can notify another process, which knows how to restart.
 - Processes can be linked with each other, thereby creating process and supervisor structures
- Ease of distribution and horisontal scalability also makes it easy to build redundant systems - we have immediate failover if a server dies.
 - Not trivial, but support for it exists in the language

Build it robust

- Accidents happen
 - "this can never happen"
 - unexpected input
 - missing case etc
- Be prepared!
 - Don't assume your program will never crash
 - Limit effects of a crash Erlang does this for you
- Note:
 - Exceptions are for exceptional cases

Processes everywhere

- Processes in Erlang are cheap and flexible
 - creation and destruction is fast and easy
 - initial size is small, typically just hundreds of bytes
 - they can grow surprisingly large..
- A typical system will consist of a (large) number of communicating processes
- At any point, one of them can have a mishap and die
- A dying process screams out in agony
- Catch the death and act accordingly
 - Restart a process that dies
 - Restart other parts of the system

Linking Processes

- Simple creation of processes is done using **spawn/1**
- There is also spawn_link/1
 - works like spawn/1
 - creates a <u>link</u> between the calling process (self()) and the newly created process
 - We're saying that these processes are important to each other.
 - If one dies the other dies as well.
 - Links are symmetrical
 - A process can be linked to several other processes, thus building process hierarchies



Catching the death of a process

- Instead of having process groups dying together a process can catch the death of other processes it is linked to
- Call process_flag(trap_exit, true)
 - If a linked process dies, instead of getting an exit signal that would kill the process, an ordinary message of the form {'EXIT', From, Reason} is received
 - The message can be processed by an ordinary **receive ... end** in a suitable manner

Workers vs Supervisors

- Trapping exit signals is asymmetric
 - S is linked to W1, W2, W3
 - S traps exits
 - If S dies, W1, W2 and W3 will die
 - If either of W1, W2, W3 dies, S will get to know about it
 - S is a supervisor
 - W1, W2 and W3 are workers



IDIA

- The low level mechanisms are few and simple
- Putting them together and getting it Right (tm) is tricky (and distracts you from your core task)
- Use the supervisor behaviour (very similar to an interface in Java)
- The supervisor behaviour states that you implement one function **init/1** which should return a term

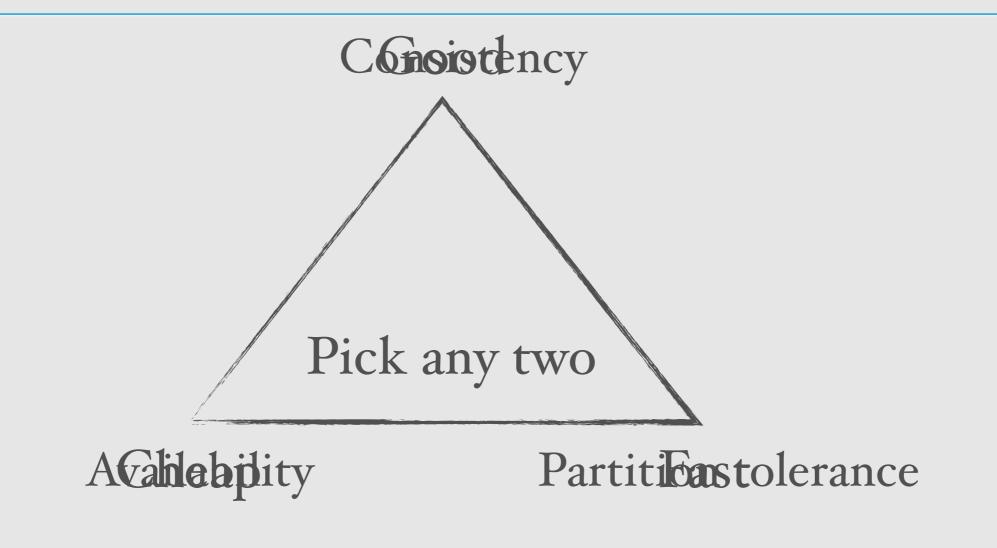
{ok,

{RestartStrategy, MaxRetries, Time},
[ChildSpec]}

Restart Strategies

- Allow a maximum of MaxRetries during Time seconds. If more are needed, the supervisor is terminated (and possibly handled by another supervisor)
- **one_for_one** restart children independently of each other
- one_for_all if one dies, restart all
- rest_for_one if one dies, restart all "after" that one
- **simple_one_for_one** like one_for_one, but all workers run the same code and can be added dynamically

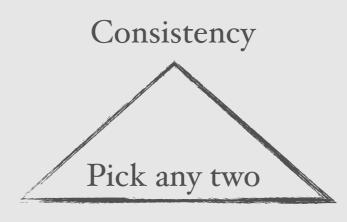
Brewer's CAP Theorem



Consistency - all nodes see the same data at the "same time" When your network partitions, your distributed Availability will be clients receive answers "immediately" System will be either consistent or available. At best. Partition tolerance - service operates despite message loss between nodes



Brewer's CAP Theorem



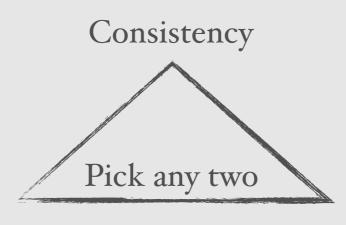
Availability

Partition tolerance

- Let the business decide which one to sacrifice
- Which two traits are crucial for making money?
- Three types of systems examples?
 - Consistent and available
 - Available and partition tolerant
 - Consistent and partition tolerant



Brewer's CAP Theorem



Availability

Partition tolerance

- We handle money and risk
 - Consistency is definitely important!
- We want customers to spend money
 - We have to be available!
- Good bye, partition tolerance..?

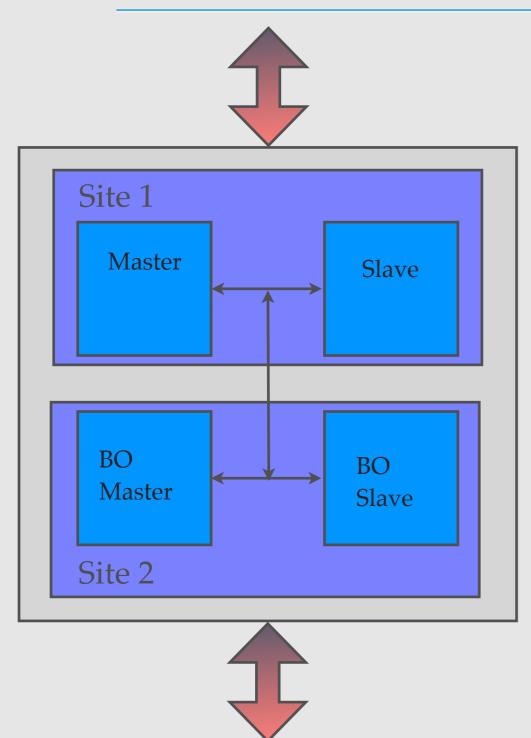


Theoretical Dilemma Meets Reality

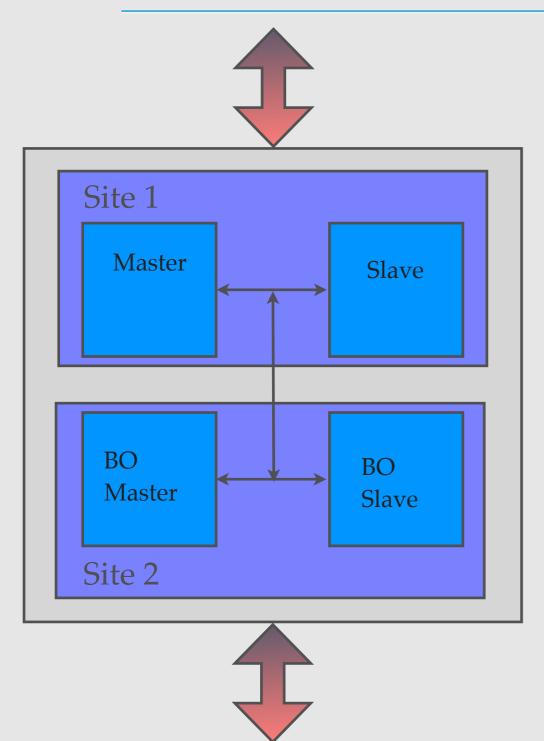
• The problem:

- Construct a system that is consistent, highly available and can scale when we grow exponentially
- Scalability is often solved by growing horisontally, i.e., by adding more nodes (thus exposing us to risk of partioning).
- The reality [time to market is important]:
 - We don't know we're going to grow exponentially focus on current problems; ignore the future.
 - This is reality solve problems when they come along with simple real world solutions. Theory..?

- Focus on consistency and availability
 - Use several nodes for fail over
 - When, not if, a node dies another one can take over
 - Replicate data between nodes
 - Data is consistent between nodes and safe
 - Know what you've done
 - Logs can be used to rebuild state after failures

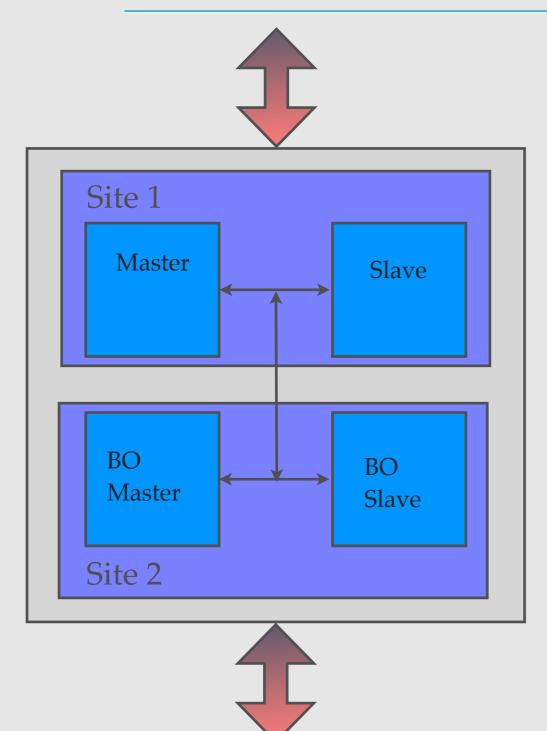


- Each node runs erlang
- Mnesia is used for persistence
- Master and slave have different main tasks
 - master handles incoming purchases
 - slave handles web traffic
- Back office (BO) handles customer service
- Data is replicated among the nodes, so all nodes have the same data
- If one node dies, the other takes over all responsibilities until the dead one comes up again (master and slave might then switch)
- Transaction logs are kept to rebuild state if needed
- Regular backups are kept



- Availability is covered by multiple nodes being able to fail over fast; this is immediate.
- Consistency is covered by letting the nodes/sites have different responsibilities and replicating data.
- Partition tolerance is handled in the same way as a dead node, i.e., a node will take over all responsibilities in a site. Data is synced upon contact again.
- We can keep availability even when external services don't answer.
- When one node has died, we're vulnerable during the time the other takes to restart.
- If both nodes die on a site, we have a serious problem.





- Software upgrades are done several times each week without downtime or stopping the system.
- Code is loaded while normal traffic is flowing.
- OS, erlang and hardware can be upgraded without downtime; stop one node, do maintenance and restart.
- Hardware can be moved without downtime; stop one node, move it and restart.
- This has served us well from the start and is still doing quite nicely.
- Growth has been exponential, but this has also been able to handle the needed scaling.
- How?



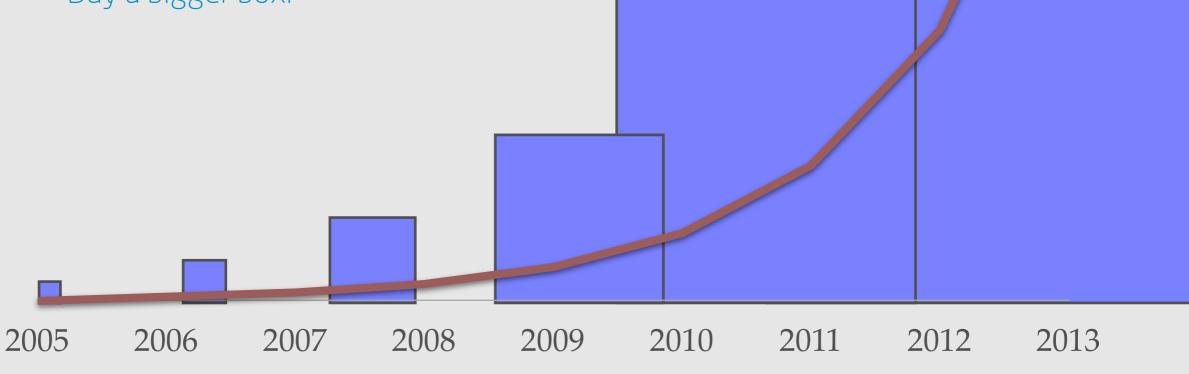
Solve Problems When They Arise

• Success means growth



- Adding new markets and customers means handling more traffic and data we need to scale up.
- Solve it by simple means and use the resources available





Theory is catching up

- We can't scale like that forever
 - Data sets will be extremely large
 - We're getting closer to the practical limits of certain aspects erlang/otp; it wasn't really designed for large scale systems
 - Traffic increase, both incoming and between nodes will expose bottlenecks
 - Moving into new markets might mean having nodes closer to the market, either for latency or regulatory reasons
- Are we going to build a second system?
 - No, the Second System Syndrome is Real.
 - Change imposes risk and we need to maintain availability
 - Remodel the existing system to be able to scale better by true horisontal scaling.
 - Rebuild and replace components incrementally



New Directions

- Scale horisontally by having multiple front end nodes, handling purchases and external web traffic.
- Move towards SOA, with stateless services for identification, risk assessment, payment handling.
- The front end has high demands for availability and scalability; high SLAs.
- Scale vertically by separating a back end to a more traditional transaction oriented bank like system.
 - Less critical.
 - We can tolerate inconsistency between front and back end. Availability (FE) is more critical.

New Directions

- Availability still high priority.
- Consistency might suffer when moving purchase nodes apart, but data will be consistent eventually.
- Fail over between purchase nodes, when a whole purchase node fails
- Partition tolerance needs to be handled within a purchase node between services.
- Messages to back end handled by a message queue and handled in order. Less critical.
- This is currently work in progress.



Interested in Klarna?

- Starting points
 - <u>engineerng.klarna.com</u>
 - <u>klarna.com/jobs</u>
 - <u>recruitment@klarna.com</u>