

CHALMERS

# Graphics Hardware



Ulf Assarsson

# Graphics hardware – why?

- About 100x faster!
- Another reason: about 100x faster!
- Simple to pipeline and parallelize
  
- Current hardware based on triangle rasterization with programmable shading (e.g., OpenGL acceleration)
- Ray tracing: there are research architectures, and few commercial products
  - Renderdrive, RPU, (Gelato), NVIDIA OptiX
  - Or write your own GPU ray-tracer



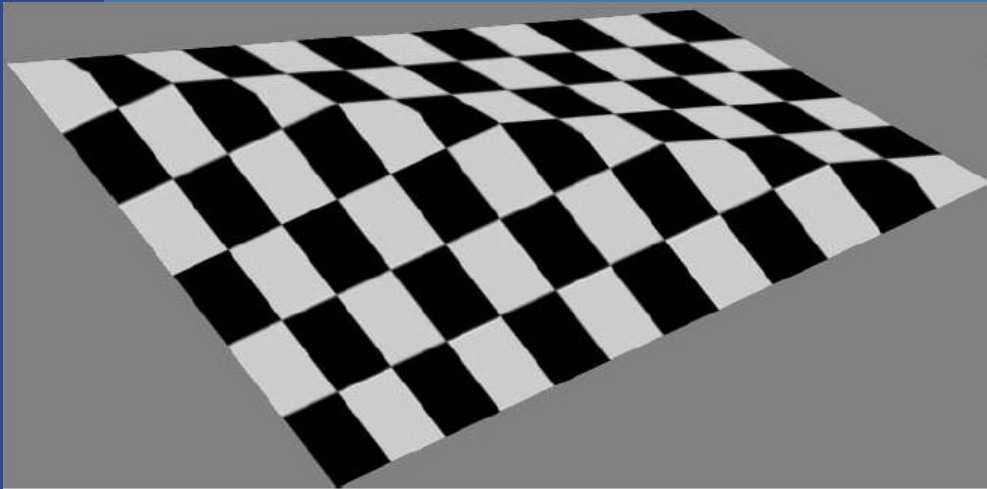
**Steel  
Monkeys**



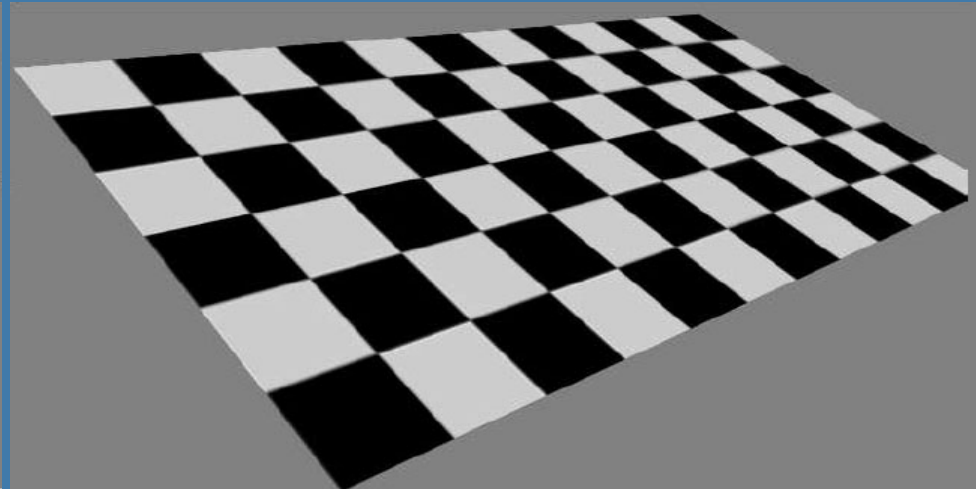
**CLASS 1  
Game product**

# Perspective-correct texturing

- How is texture coordinates interpolated over a triangle?
- Linearly?

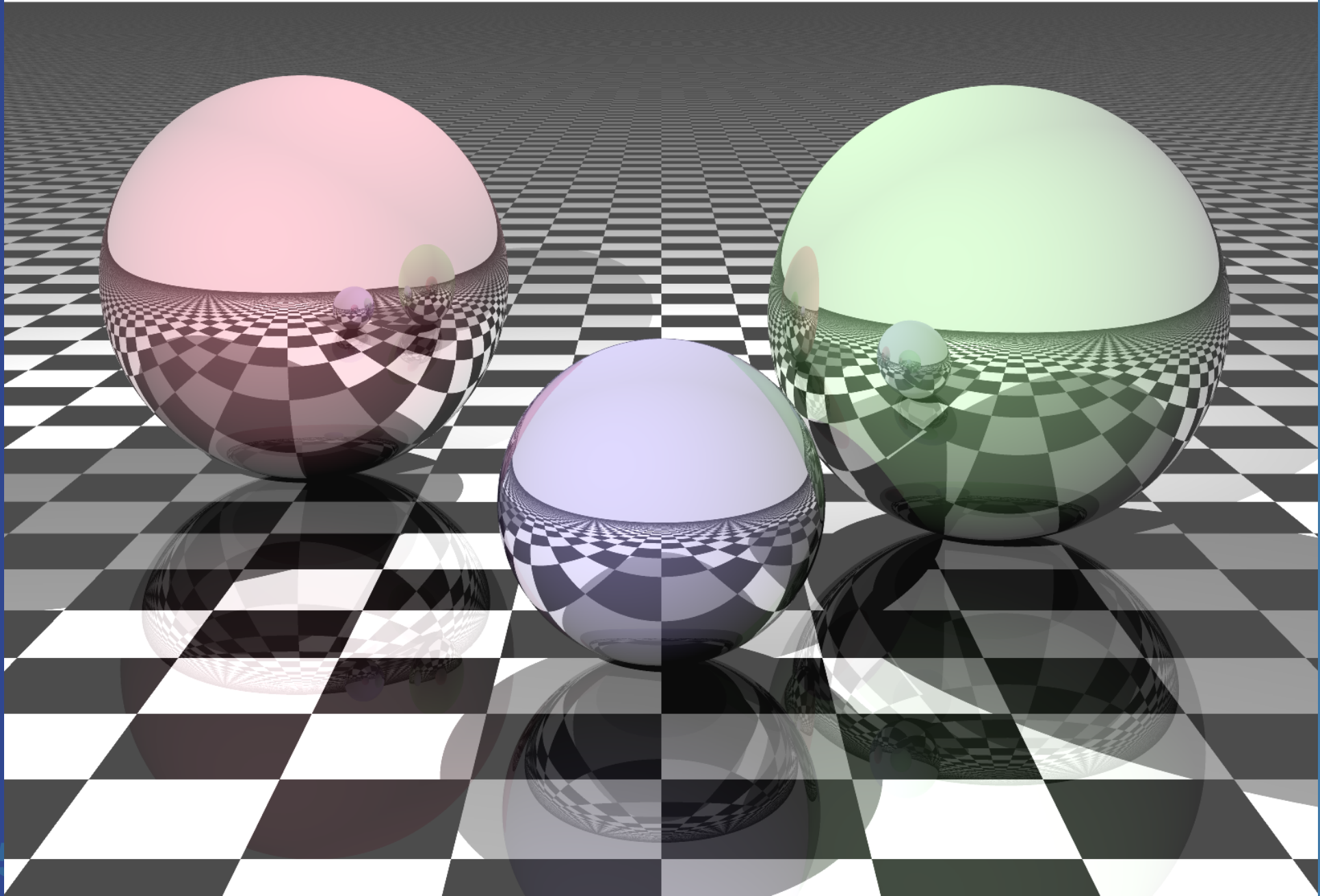


**Linear interpolation**



**Perspective-correct interpolation**

- Perspective-correct interpolation gives foreshortening effect!
- Hardware does this for you, but you need to understand this anyway!



# Recall the following

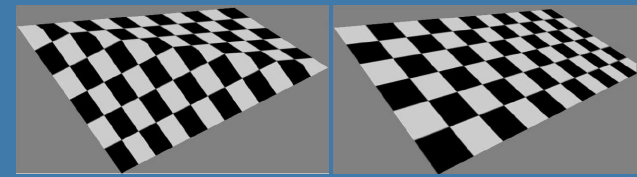
- Before projection,  $\mathbf{v}$ , and after  $\mathbf{p}$  ( $\mathbf{p}=\mathbf{M}\mathbf{v}$ )
- After projection  $p_w$  is not 1!
- Homogenization:  $(p_x/p_w, p_y/p_w, p_z/p_w, 1)$
- Gives  $(p_x', p_y', p_z', 1)$

$$\mathbf{p} = \mathbf{M}\mathbf{v} = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & -1/d & 0 \end{pmatrix} \begin{pmatrix} v_x \\ v_y \\ v_z \\ 1 \end{pmatrix} = \begin{pmatrix} v_x \\ v_y \\ v_z \\ -v_z/d \end{pmatrix}$$

# Texture coordinate interpolation

- Linear interpolation does not work
- Rational linear interpolation does:
  - $u(x) = (ax+b) / (cx+d)$  (along a scanline where  $y = \text{constant}$ )
  - $a, b, c, d$  are computed from triangle's vertices  $(x, y, z, w, u, v)$
- Not really efficient to compute  $a, b, c, d$  per scan line
- Smarter:
  - Compute  $(u/w, v/w, 1/w)$  per vertex
  - These quantities can be linearly interpolated!
  - Then at each pixel, compute  $1/(1/w) = w$
  - And obtain:  $(w * u/w, w * v/w) = (u, v)$
  - The  $(u, v)$  are perspective-correctly interpolated
- Need to interpolate shading this way too
  - Though, not as annoying as textures
- Since linear interpolation now is OK, compute, e.g.,  $\Delta(u/w) / \Delta x$ , and use this to update  $u/w$  when stepping in the  $x$ -direction (similarly for other parameters)

# Put differently:



- Linear interpolation in screen space does not work for  $u, v$
- Solution:
  - We have applied a non-linear transform to each vertex  $(x/w, y/w, z/w)$ .
    - Non-linear due to  $1/w$  – factor from the homogenisation
  - We must apply the same non-linear transform to  $u, v$ 
    - E.g.  $(u/w, v/w)$ . This can now be correctly screenspace interpolated since it follows the same non-linear  $(1/w)$  transform and then interpolation as  $(x/w, y/w, z/w)$
    - When doing the texture lookups, we still need  $(u, v)$  and not  $(u/w, v/w)$ .
    - So, multiply by  $w$ . But we don't have  $w$  at the pixel.
    - So, linearly interpolate  $(u/w, v/w, 1/w)$ , which is computed in screenspace at each vertex.
    - Then at each pixel:
      - $u = (u/w) / (1/w)$
      - $v = (v/w) / (1/w)$

For a formal proof, see Jim Blinn, "W Pleasure, W Fun", IEEE Computer Graphics and Applications, p78-82, May/June 1998

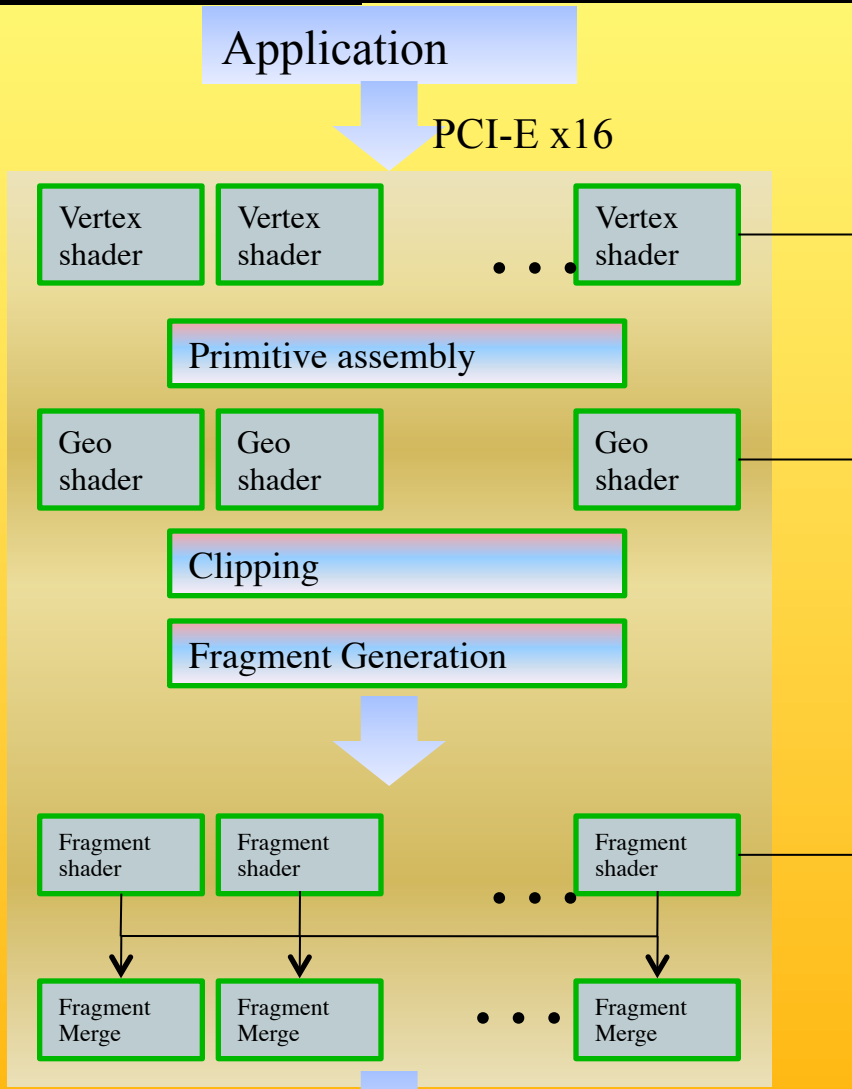
Need to interpolate shading this way too, though, not as annoying as textures



# Background:

## Graphics hardware architectures

- Evolution of graphics hardware has started from the end of the pipeline
  - Rasterizer was put into hardware first (most performance to gain from this)
  - Then the geometry stage
  - Application will not be put into GPU hardware (?)
- Two major ways of getting better performance:
  - Pipelining
  - Parallelization
  - Combinations of these are often used

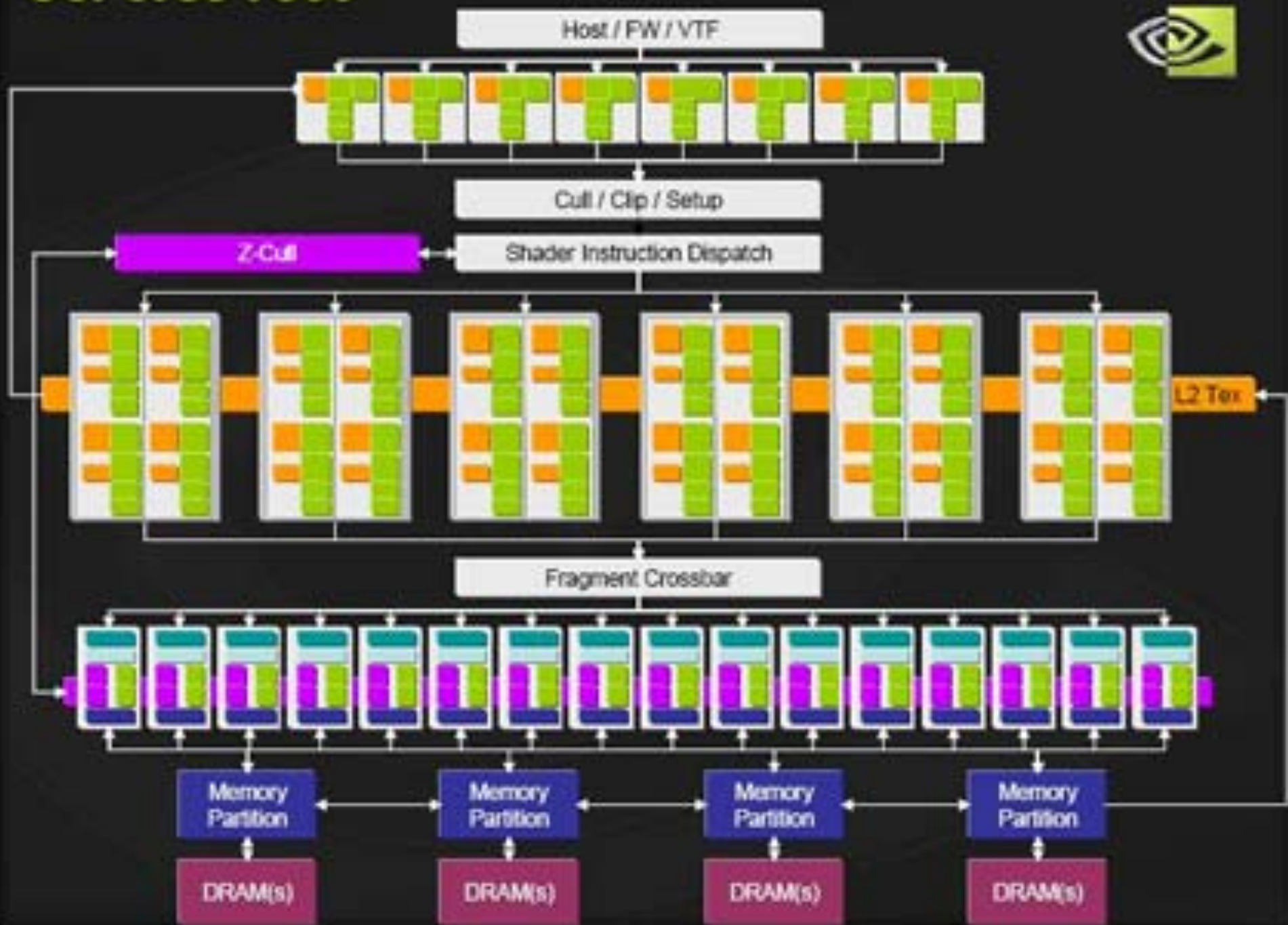


On NVIDIA  
8000/9000/200/400/500/600/700/TITAN/900-series:  
Vertex-, Geometry- and Fragment  
shaders allocated from a pool of  
128/240/480/512/1536/2880  
ALUs

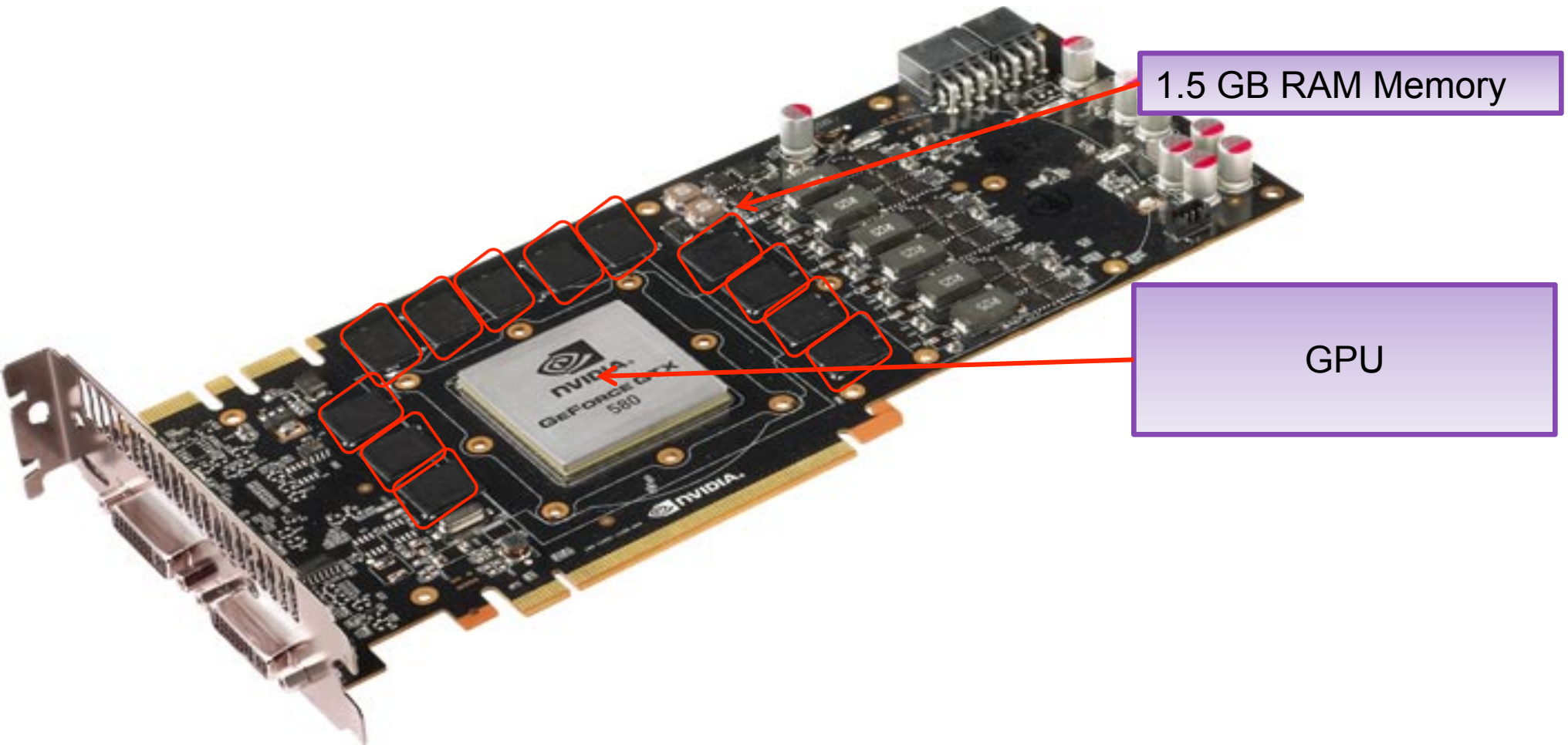


Display

# GeForce 7800

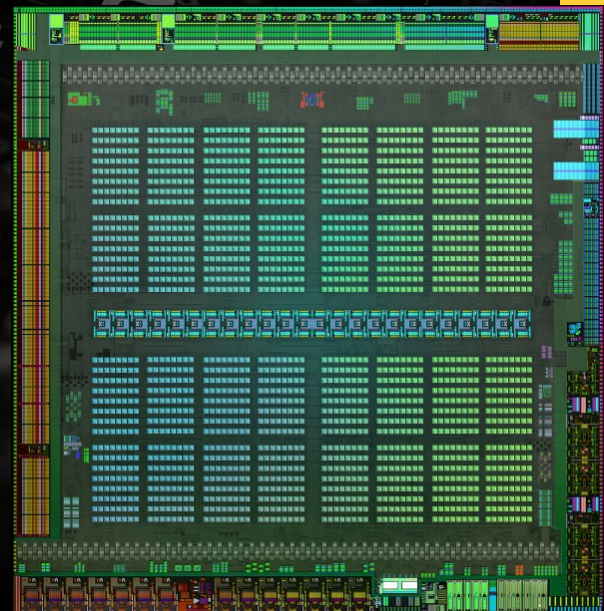
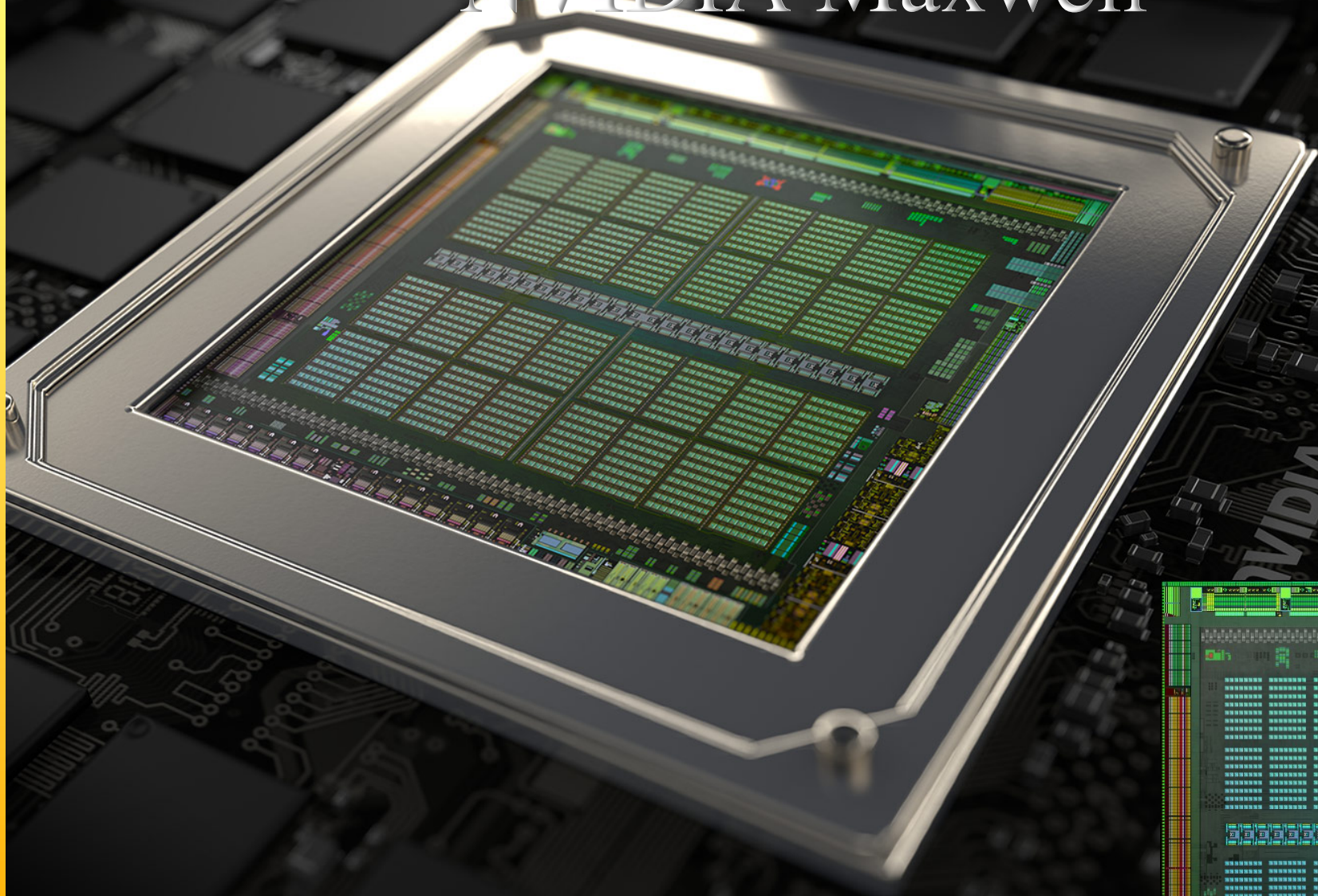


# Graphics Processing Unit - GPU



- NVIDIA Geforce GTX 580

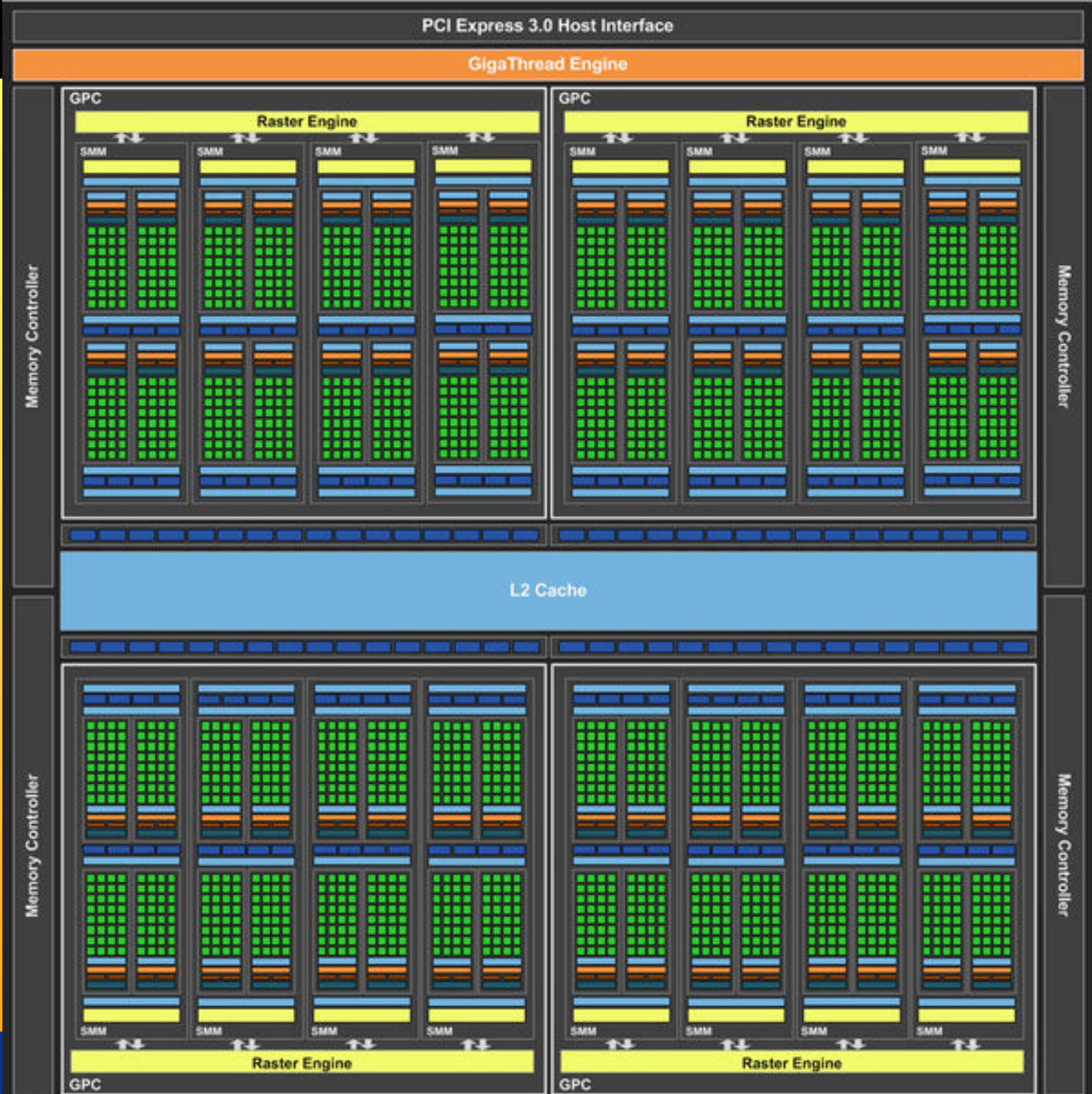
# NVIDIA Maxwell



16 SMMs (“Cores”)  
2MB L2 cache  
64 color outputs/  
clock (i.e., 64 ROPs)

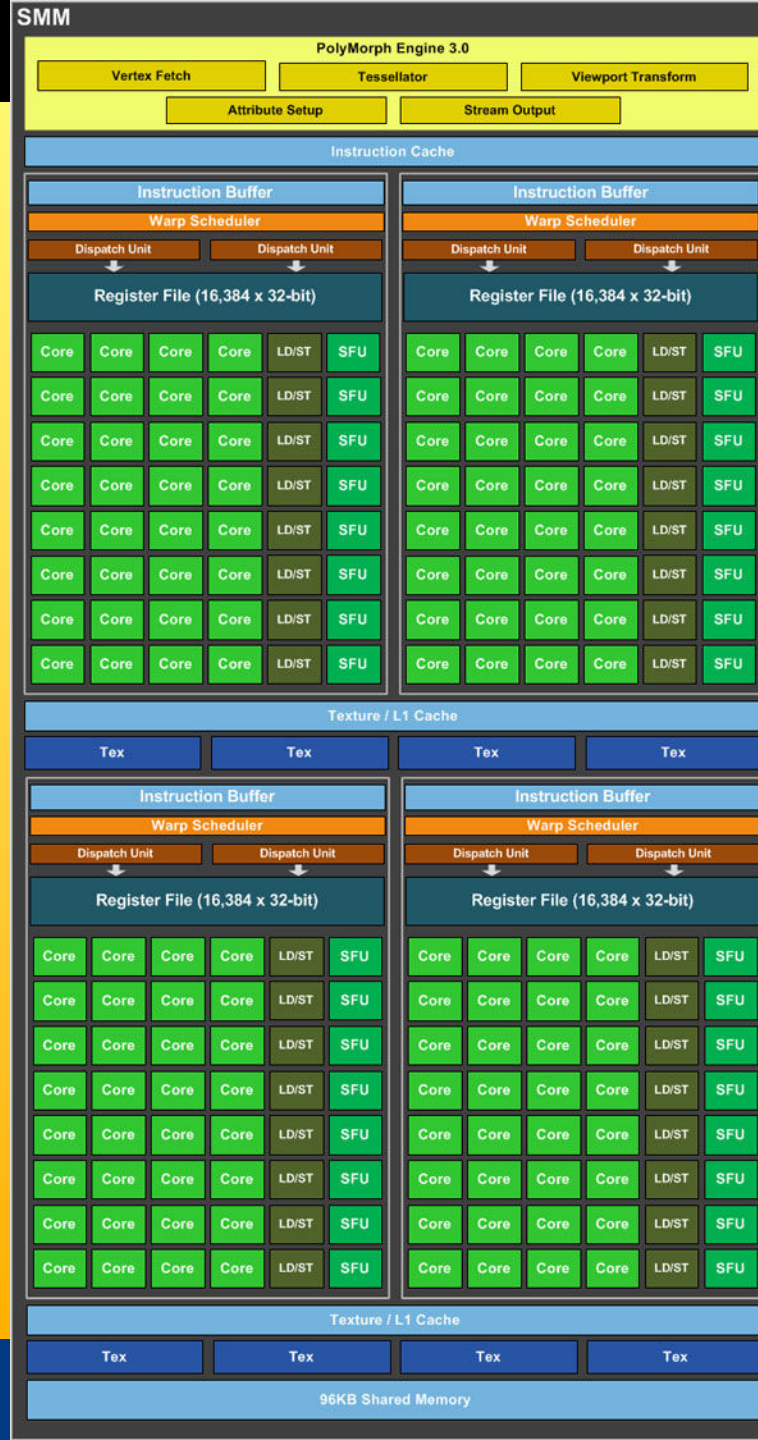
Each SMM:

- 128 ALUs
- 96KB L1 cache
- 8 TexUnits
- 32 Load/Store units for access to global memory



Each SMM:

- 128 ALUs
- 96KB L1 cache
- 8 TexUnits
- 32 Load/Store units for access to global memory



# Graphics Hardware History

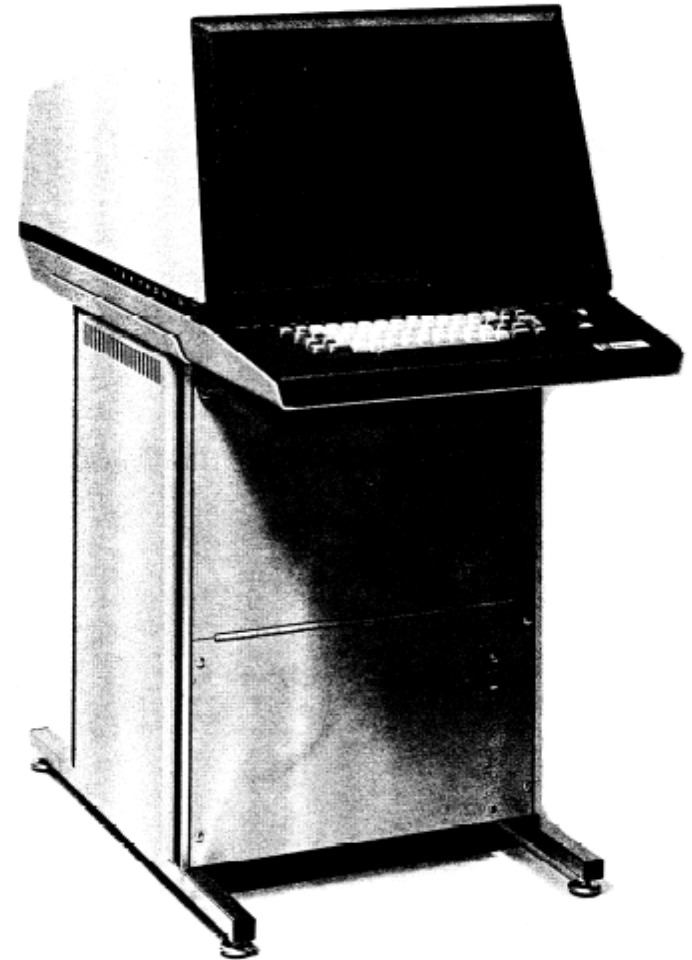
- 80's:
  - linear interpolation of color over a scanline
  - Vector graphics
- 91' Super Nintendo, Neo Geo,
  - Rasterization of 1 single 3D rectangle per frame (FZero)
- 95-96': Playstation 1, 3dfx Voodoo 1
  - Rasterization of whole triangles (Voodoo 2, 1998)
- 99' Geforce (256)
  - Transforms and Lighting (geometry stage)
- 02' 3DLabs WildCat Viper, P10
  - Pixel shaders, integers,
- 02' ATI Radion 9700, GeforceFX
  - Vertex shaders and **Pixel shaders** with floats
- 06' Geforce 8800
  - Geometry shaders, integers and floats, logical operations
- Then:
  - More general multiprocessor systems, higher SIMD-width, more cores





# Direct View Storage Tube

- Created by Tektronix
  - Did not require constant refresh
  - Standard interface to computers
    - Allowed for standard software
    - Plot3D in Fortran
  - Relatively inexpensive
    - Opened door to use of computer graphics for CAD community



**Tektronix 4014**

Fig. 1-1. 4014 Computer Display Terminal.

# Briefly about Graphics HW pipelining

2001 • In GeForce3: 600-800 pipeline stages!

- 57 million transistors
- First Pentium IV: 20 stages, 42 million transistors,

## • Evolution of cards:

2004 - X800 - 165M transistors

2005 - X1800 - 320M trans, 625 MHz, 750 Mhz mem, 10Gpixels/s, 1.25G verts/s

2004 - GeForce 6800: 222 M transistors, 400 MHz, 400 MHz core/550 MHz mem

2005 - GeForce 7800: 302M trans, 13Gpix/s, 1.1Gverts/s, bw 54GB/s, 430 MHz core, mem 650MHz(1.3GHz)

2006 - GeForce 8800: 681M trans, 39.2Gpix/s, 10.6Gverts/s, bw:103.7 GB/s, 612 MHz core (1500 for shaders), 1080 MHz mem (effective 2160 GHz)

2008 - Geforce 280 GTX: 1.4G trans, 65nm, 602/1296 MHz core, 1107(\*2)MHz mem, 142GB/s, 48Gtex/s

2007 - ATI Radeon HD 5870: 2.15G trans, 153GB/s, 40nm, 850 MHz, GDDR5, 256bit mem bus,

2010 - Geforce GTX480: 3Gtrans, 700/1401 MHz core, Mem (1.848G(\*2)GHz), 177.4GB/s, 384bit mem bus, 40Gtexels/s

2011 - GXT580: 3Gtrans, 772/1544, Mem: 2004/4008 MHz, 192.4GB/s, GDDR5, 384bit mem bus, 49.4 Gtex/s

2012 - GTX680: 3.5Gtrans (7.1 for Tesla), 1006/1058, 192.2GB/s, 6GHz GDDR5, 256-bit mem bus.

2013 - GTX780: 7.1G, core clock: 837MHz, 336 GB/s, Mem clock: 6GHz GDDR5, 384-bit mem bus

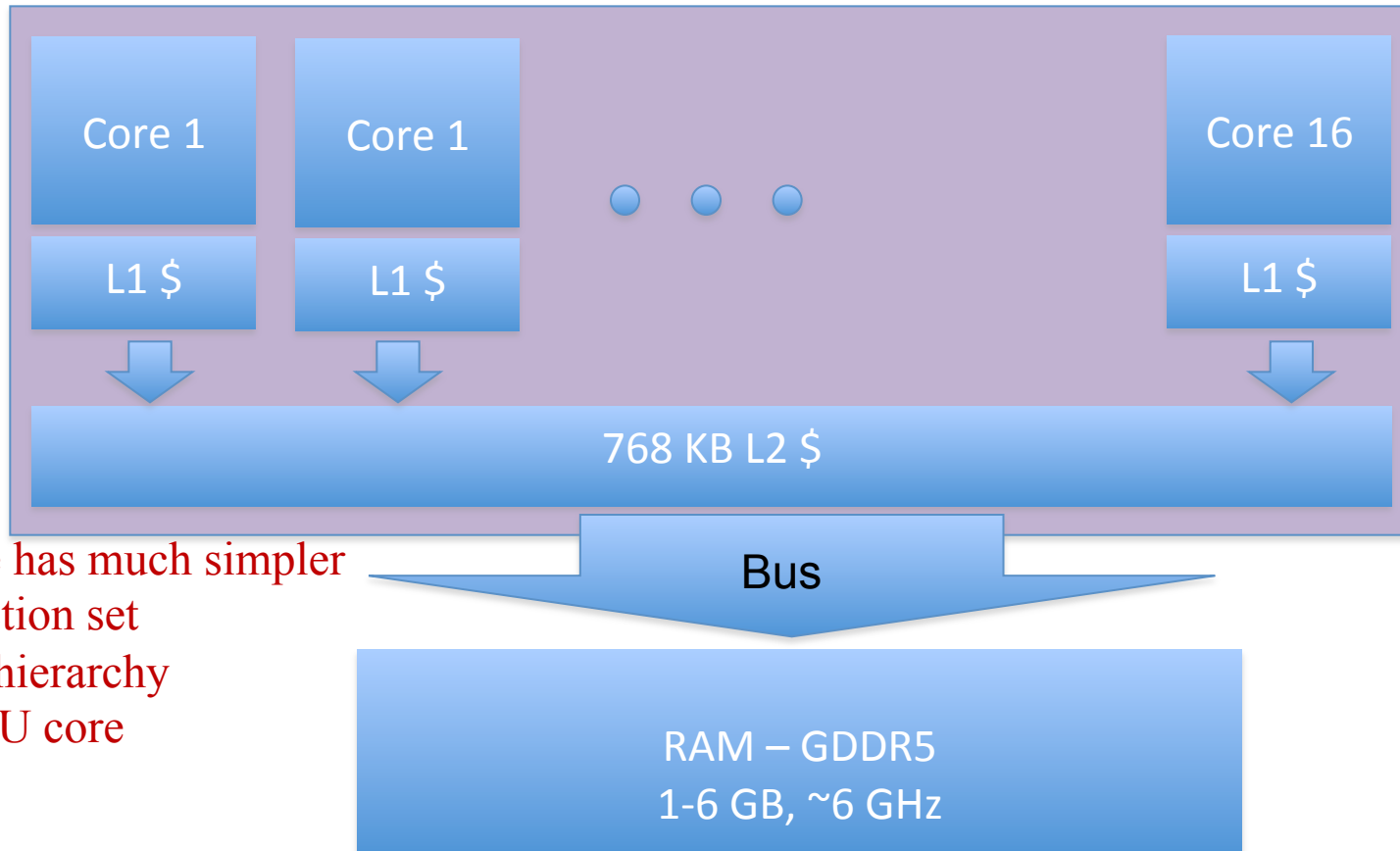
2014 - GTX980: 7.1G?, core clock: ~1200MHz, 224GB/s, Mem clock: 7GHz GDDR5, 256-bit mem bus

- Lesson learned: #trans doubles ~per 2 years. Core clock increases slowly. Mem clock -increases with new technology DDR2, DDR3, GDDR5, and with more memory busses (à 64-bit)
- We want as fast memory as possible! Why?
  - Parallelization can cover for slow clock. Parallelization more energy efficient than high clock frequency. Powerconsumption prop. to freq<sup>2</sup>.
  - Memory transfers often the bottleneck

# GPU- Nvidia's Kepler 2012



Overview:



16 cores à  
192-SIMD width  
(2\*6\*16)

16/48 KB per  
each 48 SIMD

Bandwidth  
~330 GB/s

Bus: 256/384  
bits

Compare to  
ATI 2900:

- 2x512bits

Larrabee:

- 2x512bits

GPU core has much simpler

- instruction set
  - cache hierarchy
- than a CPU core

Wish:

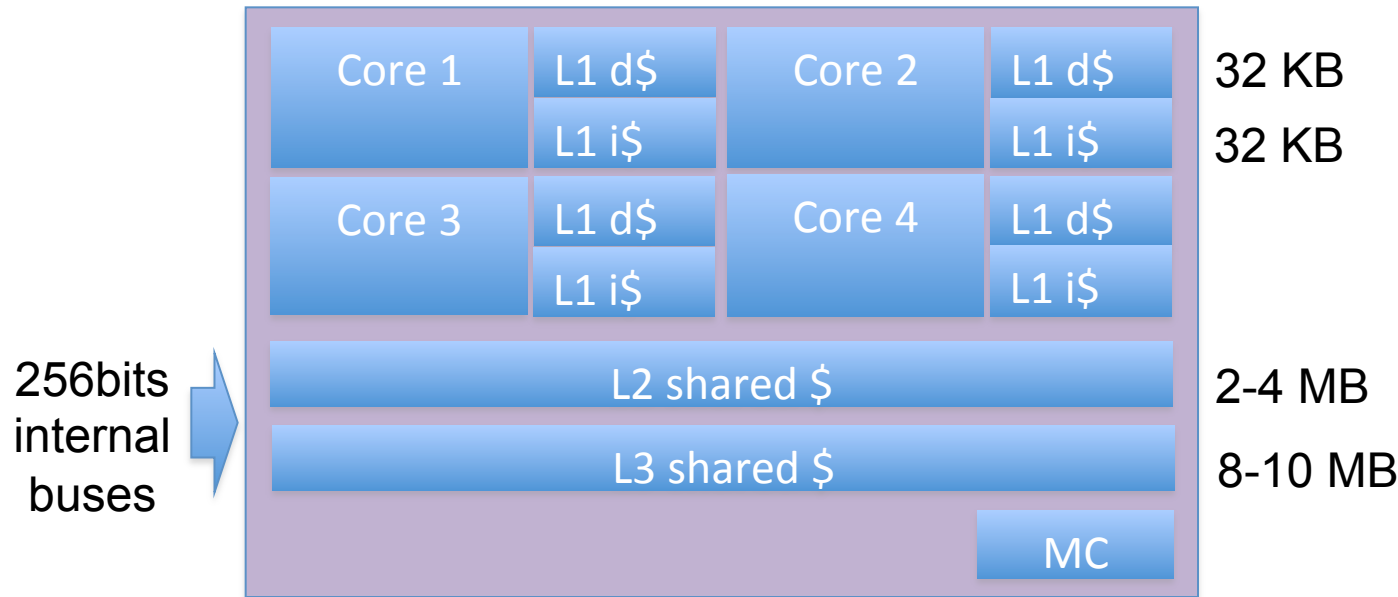
3072 ALUs à 1 float/clock => 12KB/clock

~1GHz core clock => 12000 GB/s request

We have ~330GB/s. In reality we can do 20-40 instr. between each RAM-read/write. Solved by L1\$ + L2\$ + latency hiding (warp switching)

# CPU - 2011

AVX:  
Intel's Sandybridge  
AMD's Bulldozer



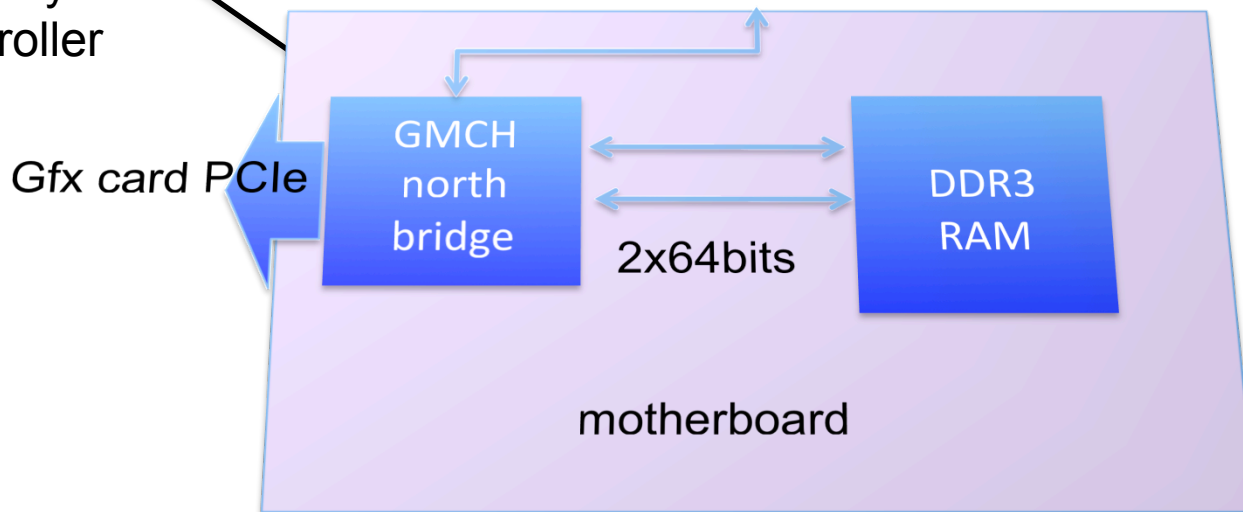
1 – 8 cores à  
4 SIMD floats  
(16 SIMD for  
bytes)

- 8 cores à 4 floats  
⇒ We want 128 bytes/clock  
(e.g. from RAM)  
⇒ 128GByte/s, 1GHz CPU

- In addition, x3, since:  
 $r1 = r2 + r3;$

In reality: 6-12GB/s

Solved by \$-hierarchy + registers



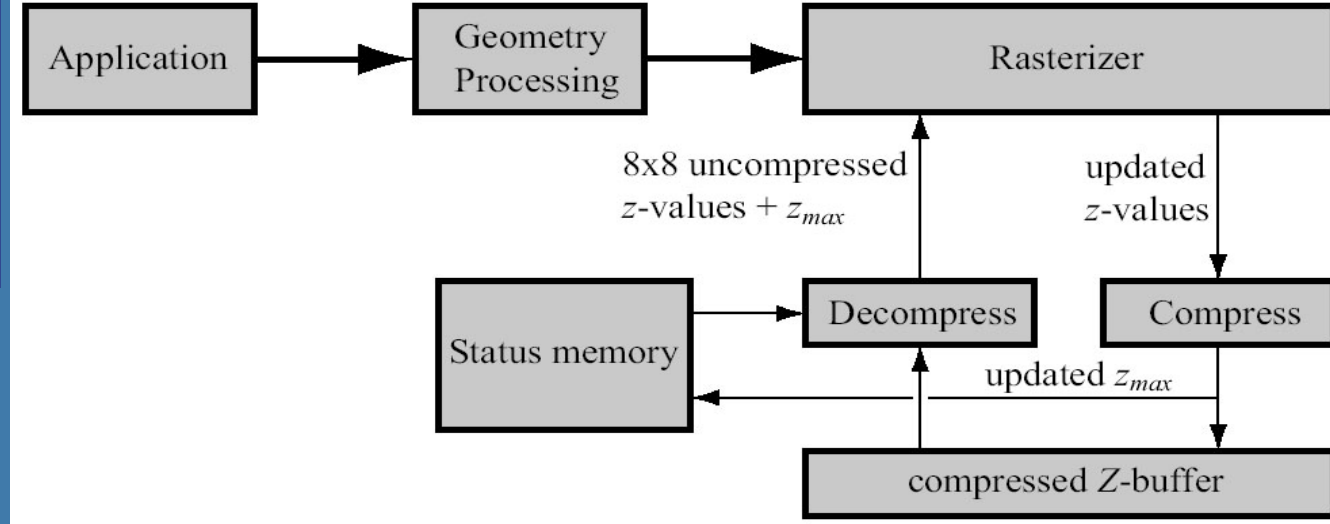
# Memory bandwidth usage is huge!!

- On top of that bandwidth usage is never 100%.
- However, there are many techniques to reduce bandwidth usage:
  - Texture caching with prefetching
  - Texture compression
  - Z-compression
  - Z-occlusion testing (HyperZ)

# Z-occlusion testing and Z-compression

- One way of reducing bandwidth
  - ATI Inc., pioneered with their HyperZ technology
- Very simple, and very effective
- Divide screen into tiles of 8x8 pixels
- Keep a status memory on-chip
  - Very fast access
  - Stores additional information that this algorithm uses
- Enables occlusion culling on triangle basis, z-compression, and fast Z-clears

# Architecture of Z-cull and Z-compress



- Store  $z_{max}$  per tile, and a flag (whether cleared, compressed/uncompressed)
- Rasterize one tile at a time
- Test if  $z_{min}$  on triangle is farther away than tile's  $z_{max}$ 
  - If so, don't do any work for that tile!!!
  - Saves texturing and z-read for entire tile – huge savings!
- Otherwise read compressed Z-buffer, & unpack
- Write to unpacked Z-buffer, and when finished compress and send back to memory, and also: update  $z_{max}$
- For fast Z-clears: just set a flag to "clear" for each tile
  - Then we don't need to read from Z-buffer, just send cleared Z for that tile

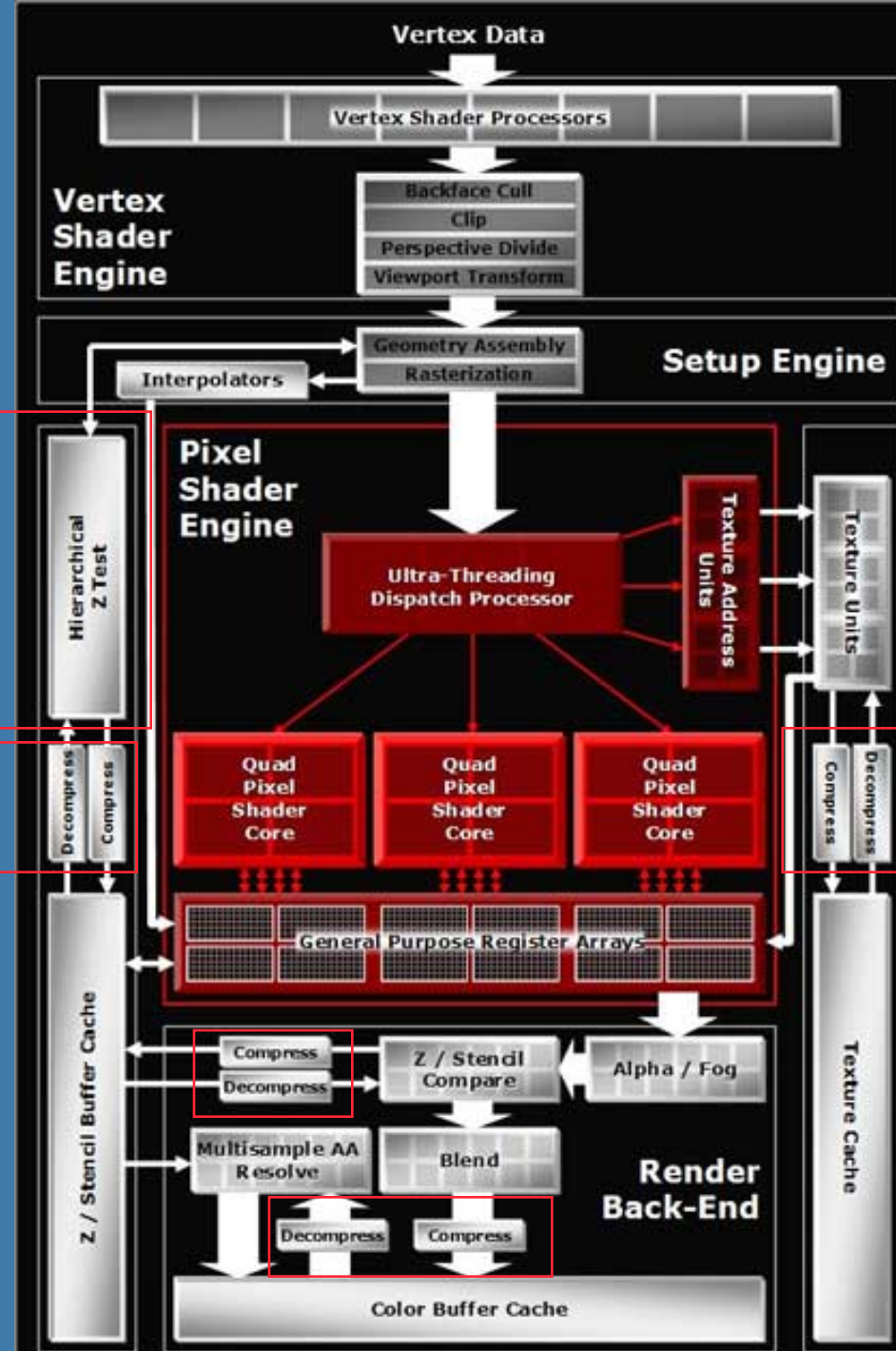
# X1800 GTO

- Real example

Z-cull

Z-compress

Also note texture compress and color compress





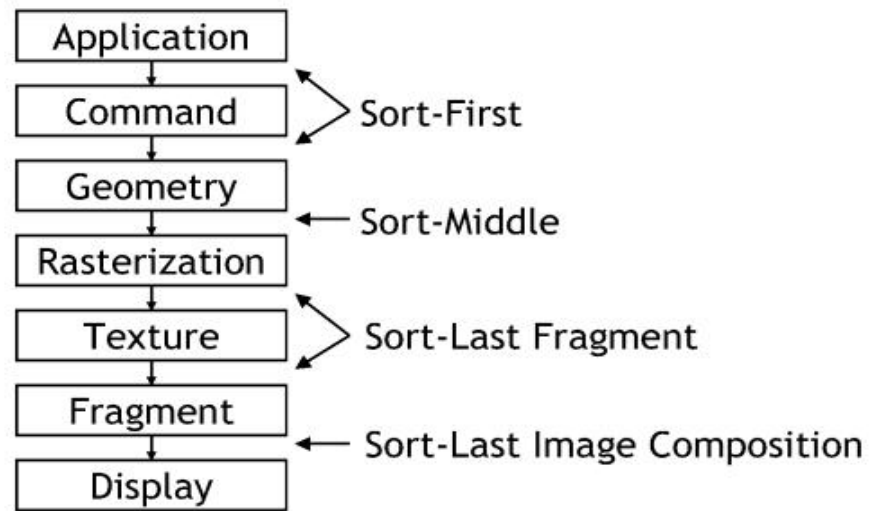
# Taxonomy of Hardware

- We can do many computations in parallel:
  - Pixel shading, vertex shading, geometry shading
    - x,y,z,w   r,g,b,a
- But results need to be sorted somewhere before reaching the screen.
  - Operations can be parallelized but result on screen must be as if each triangle were rendered one by one in their incoming order (according to OpenGL spec)
    - E.g., for blending (transparency), (z-culling, stencil test)

# Taxonomy of hardware

- Need to sort from model space to screen space
- Gives four major architectures:
  - Sort-first
  - Sort-middle
  - Sort-Last Fragment
  - Sort-Last Image

## Sorting Taxonomy



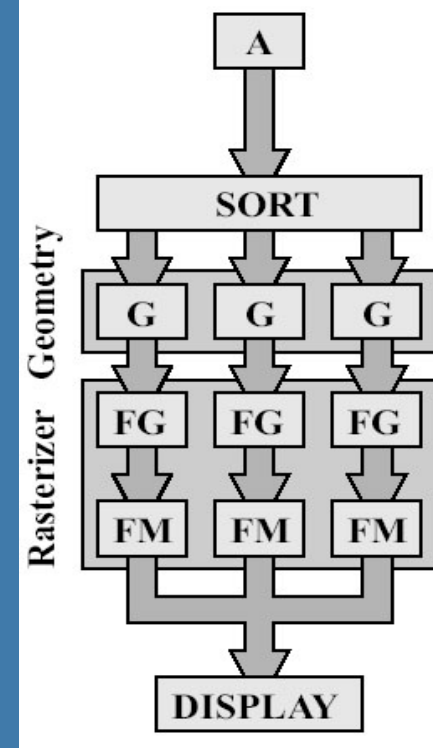
CS448 Lecture 9

Kurt Akeley, Pat Hanrahan, Fall 2001

- Will describe these briefly. Sort-last fragment (and sort middle) are most common in commercial hardware

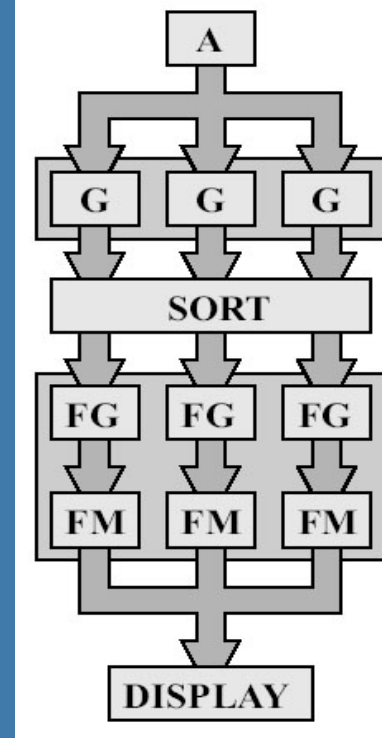
## Sort-First

- Sorts primitives before geometry stage
  - Screen is divided into large regions
  - A separate pipeline is responsible for each region (or many)
  - But vertex shader can change screen location!
- G is geometry, FG & FM is part of rasterizer
  - A fragment is all the generated information for a pixel on a triangle
  - FG is Fragment Generation (finds which pixels are inside triangle)
  - FM is Fragment Merge (merges the created fragments with various buffers (Z, color))
- Not explored much at all



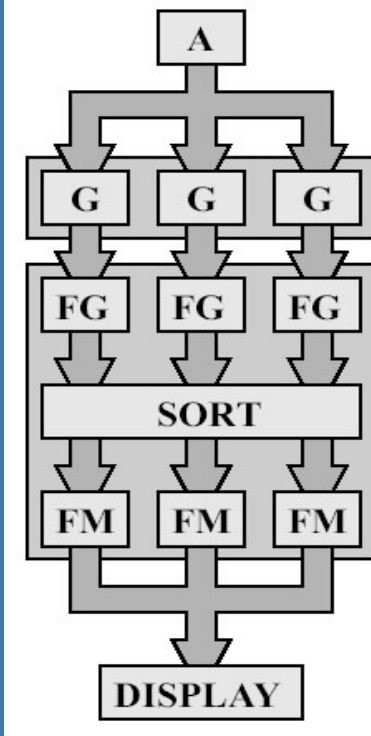
# Sort-Middle

- Sorts between G and R
- Pretty natural, since after G, we know the screen-space positions of the triangles
- Older/cheaper hardware uses this
  - Examples include InfiniteReality (from SGI) and the KYRO architecture (from Imagination)
- Spread work arbitrarily among G's
- Then depending on screen-space position, sort to different R's
  - Screen can be split into "tiles". For example:
    - Rectangular blocks (8x8 pixels)
    - Every n scanlines
- The R is responsible for rendering inside tile
- A triangle can be sent to many FG's depending on overlap (over tiles)



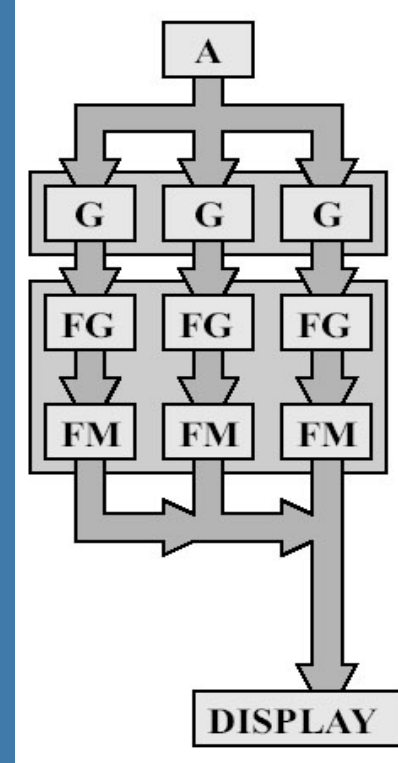
# Sort-Last Fragment

- Sorts between FG and FM
- XBOX, PS3, nVidia use this
- Again spread work among G's
- The generated work is sent to FG's
- Then sort fragments to FM's
  - An FM is responsible for a tile of pixels
- A triangle is only sent to one FG, so this avoids doing the same work twice
  - Sort-Middle: If a triangle overlaps several tiles, then the triangle is sent to all FG's responsible for these tiles
    - Results in extra work

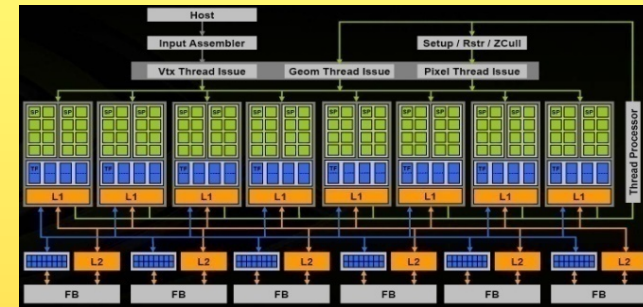
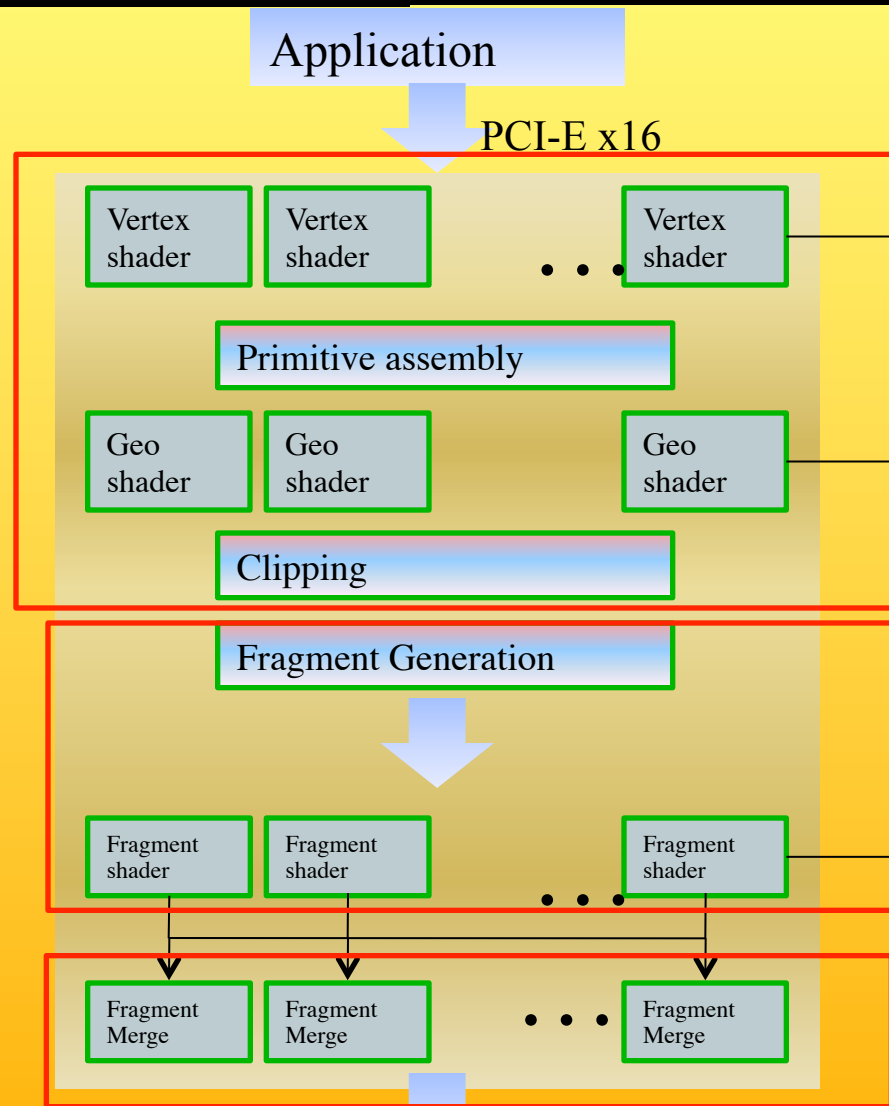


# Sort-Last Image

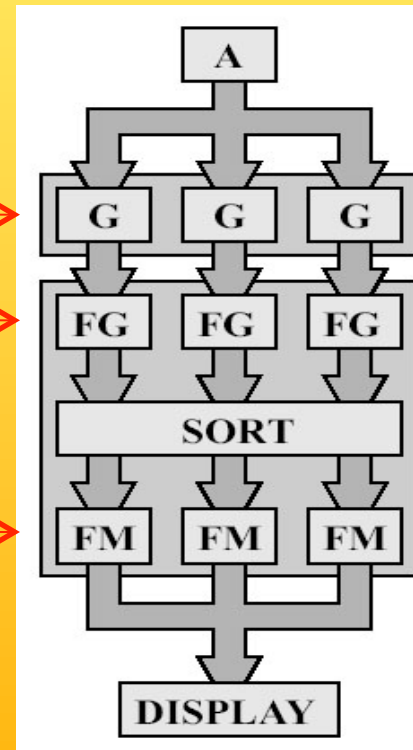
- Sorts after entire pipeline
- So each FG & FM has a separate frame buffer for entire screen (Z and color)
- After all primitives have been sent to the pipeline, the z-buffers and color buffers are merged into one color buffer
- Can be seen as a set of independent pipelines
- Huge memory requirements!
- Used in research, but probably not commercially



# Logical layout of a graphics card:



On NVIDIA  
8000/9000/200/400  
/600-series:  
→ Vertex-, Geometry- and  
Fragment shaders  
allocated from a pool of  
128/240/480/1536/3072  
processors (=ALUs)

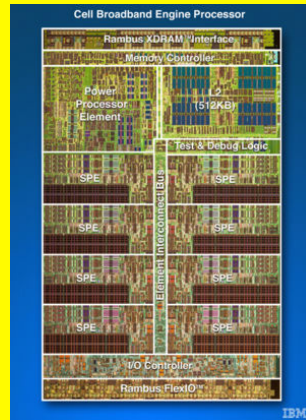


# Current and Future Multicores in Graphics

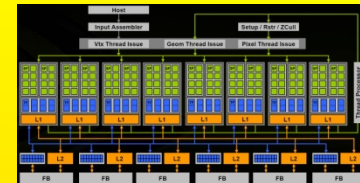
- Cell – 2005
  - 8 cores à 4-float SIMD
  - 256KB L2 cache/core
  - 128 entry register file
  - 3.2 GHz

## PowerXCell 8i Processor – 2008

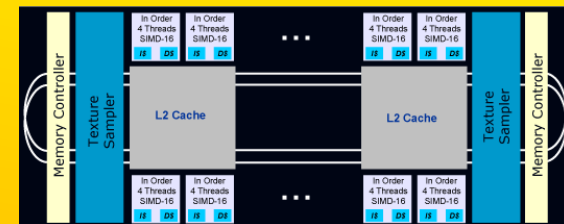
- 8 cores à 4-float SIMD
- 256KB L2 cache
- 128 entry register file
- but has better double precision support



- NVIDIA 8800 GTX – Nov 2006
  - 16 cores à 8-float SIMD (GTX 280 - 30 cores à 8, june '08)
  - 16 KB L1 cache, 64KB L2 cache (rumour)
  - 1.2-1.625 GHz



- Larrabee – "2010"
  - 16-24 cores à 16-float SIMD (Xeon Phi: 61 cores, 2012)
  - Core = 16-float SIMD (=512bit FPU) + x86 proc with loops, branches + scalar ops, 4 threads/core
  - 32KB L1cache, 256KB L2-cache (512KB/core)
  - 1.7-2.4 GHz (1.1 GHz)



- NVIDIA Fermi GF100 – 2010, (GF110 2011)
  - 16 cores à 2x16-float SIMD (1x16 double SIMD)
  - 16/48 KB L1 cache, 768 KB L2 cache

- NVIDIA Kepler 2012 - 16 cores à 2x3x16=96 float SIMD
- NVIDIA Kepler 2013 - 16 cores à 2x6x16=192 float SIMD



# Intel Xeon Phi

- Knights Corner

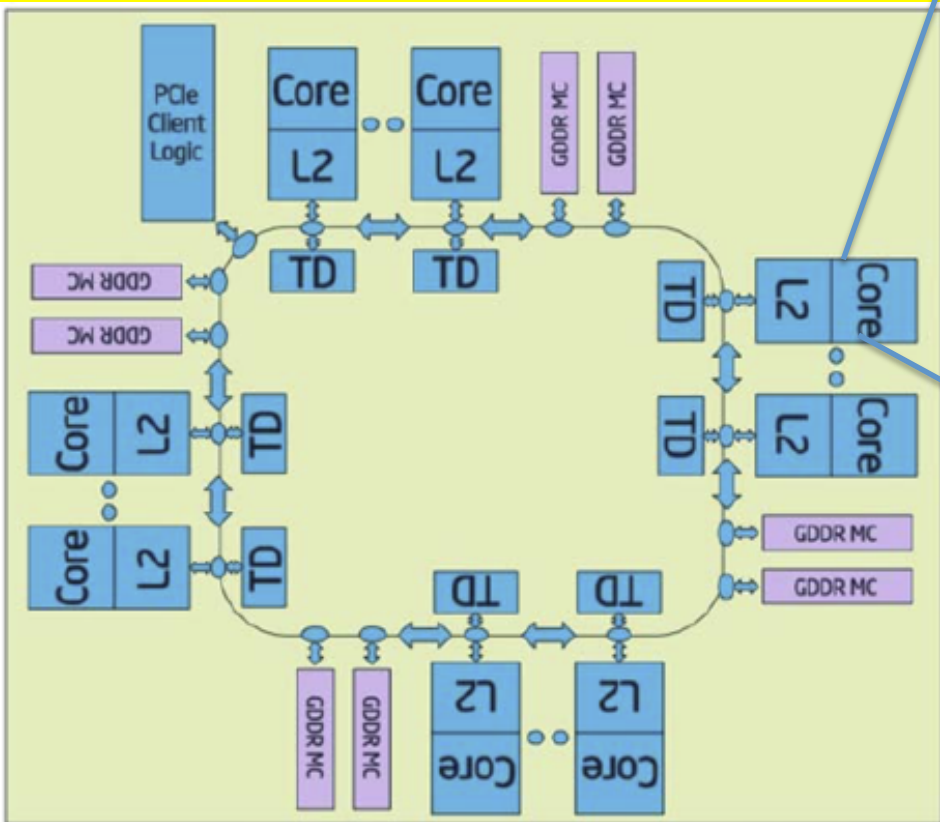


Figure 5: Knights Corner Microarchitecture

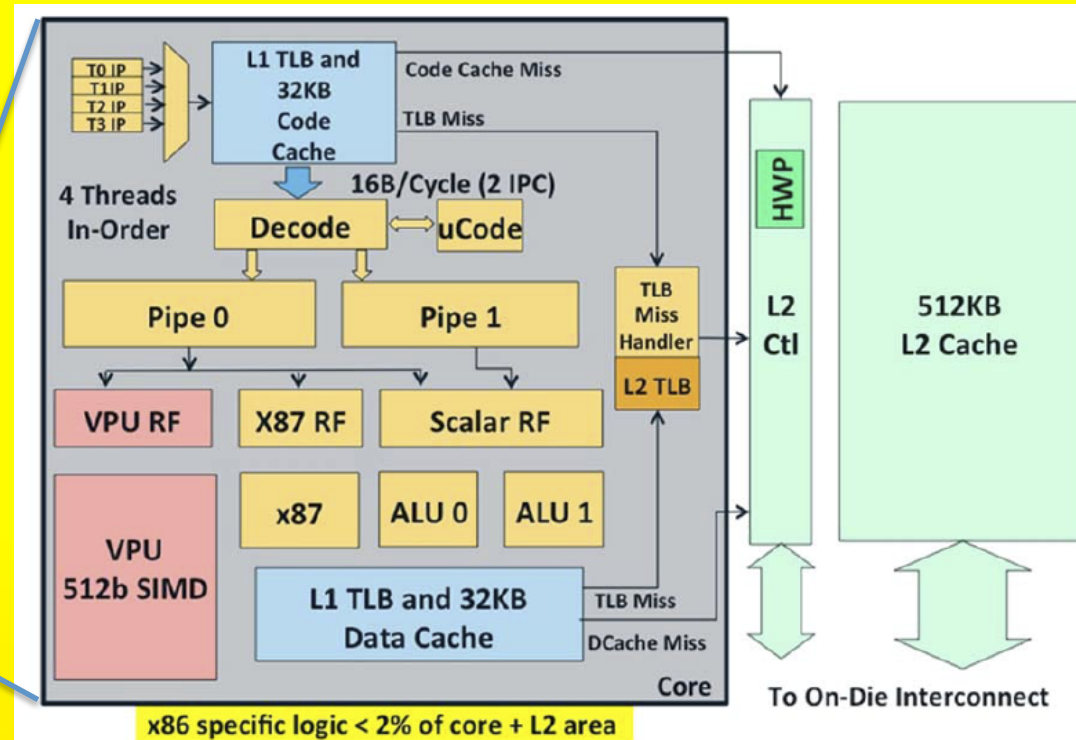


Figure 4: Knights Corner Core

<http://www.tomshardware.com/reviews/xeon-phi-larrabee-stampede-hpc,3342-3.html>

# NVIDIA year 2020

- Exaflop machine:
- Google on:  
"The Challenge of Future High-Performance Computing" Uppsala
- [http://media.medfarm.uu.se/play/video/3261#\\_utma=1.4337140.1361541635.1361541635.1361541635.1&\\_utmb=1.4.10.1361541635&\\_utmc=1&\\_utmz=1.1361541635.1.1.utmcsr=\(direct\)%7Cutmccn=\(direct\)%7Cutmcmd=\(none\)&\\_utmh=-&\\_utmk=104508928](http://media.medfarm.uu.se/play/video/3261#_utma=1.4337140.1361541635.1361541635.1361541635.1&_utmb=1.4.10.1361541635&_utmc=1&_utmz=1.1361541635.1.1.utmcsr=(direct)%7Cutmccn=(direct)%7Cutmcmd=(none)&_utmh=-&_utmk=104508928)
- Bill Dally, Chief Scientist & sr VP of Research, NVIDIA, prof. of Engineering, Stanford Univ.
- “Energy efficiency is key to performance”  
– Flops/W.

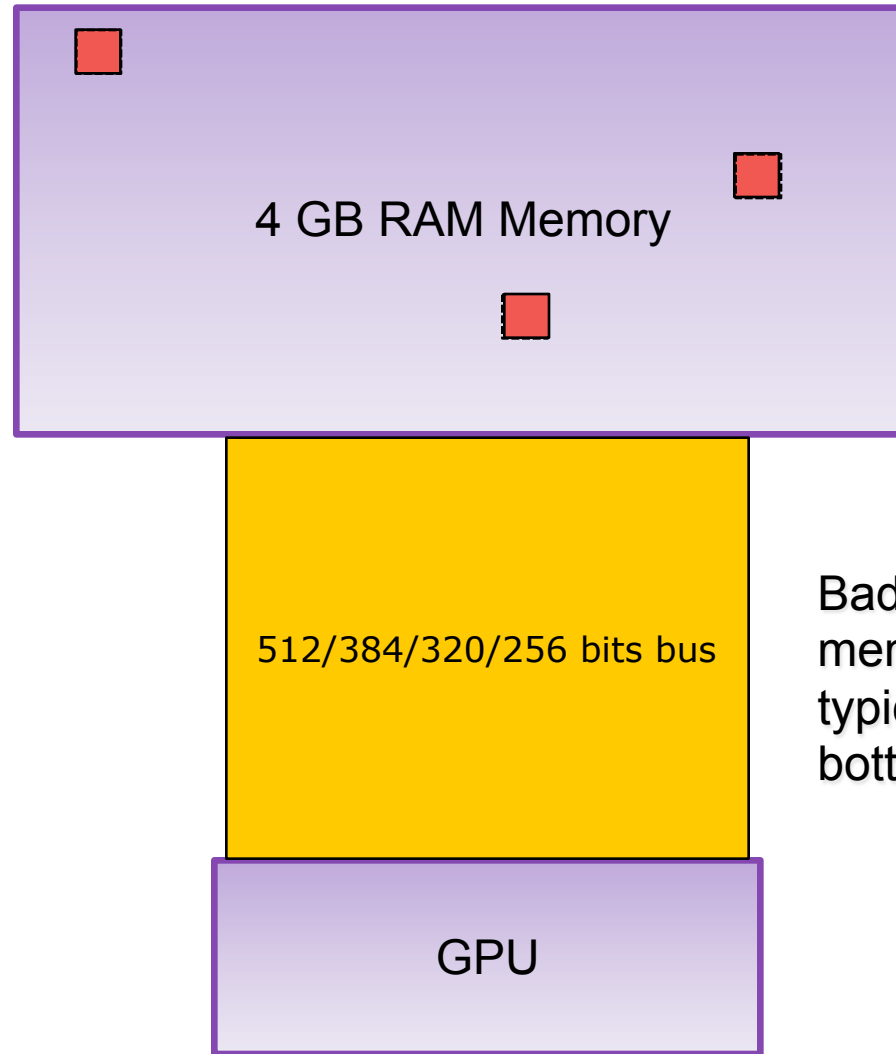
If we have time...

How create GPU algorithms  
with optimal performance?


Answer: coalesced memory  
accesses

# Graphics Processing Unit - GPU

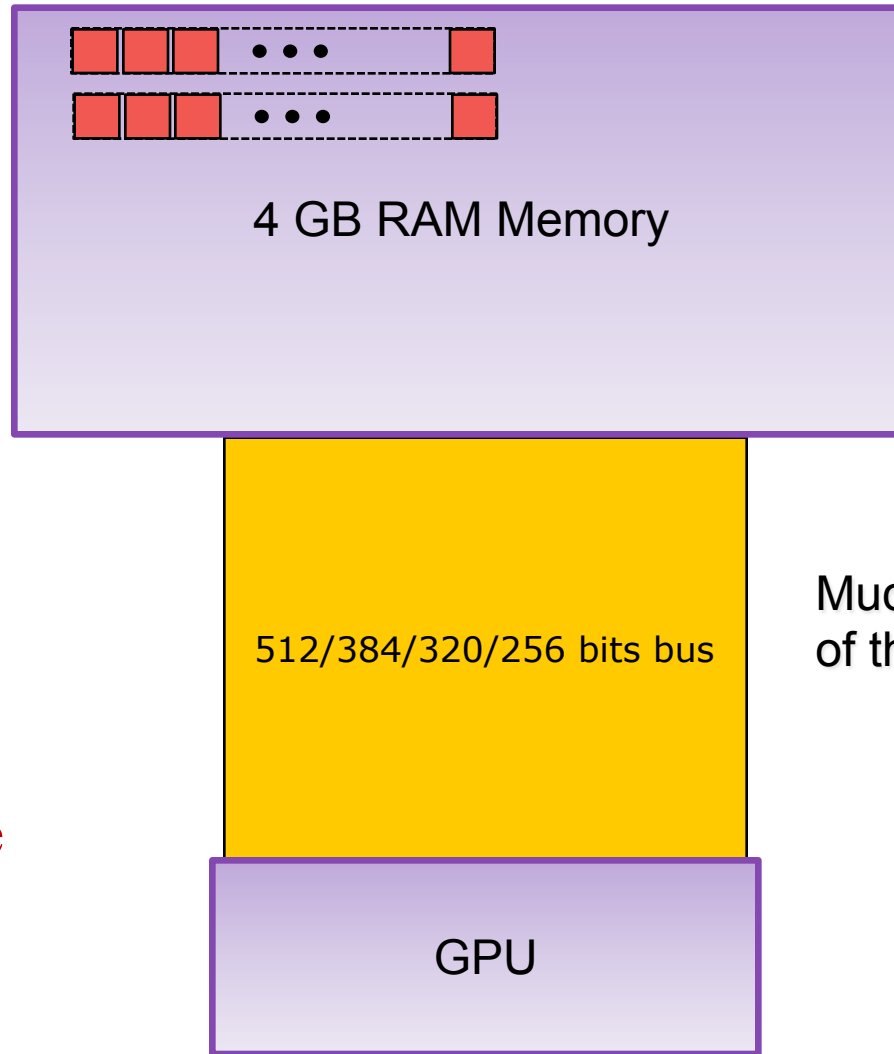
Conceptual layout:



Bad utilization of the memory bus, which typically is the bottleneck!


 = memory element (32 bits)

# Graphics Processing Unit - GPU

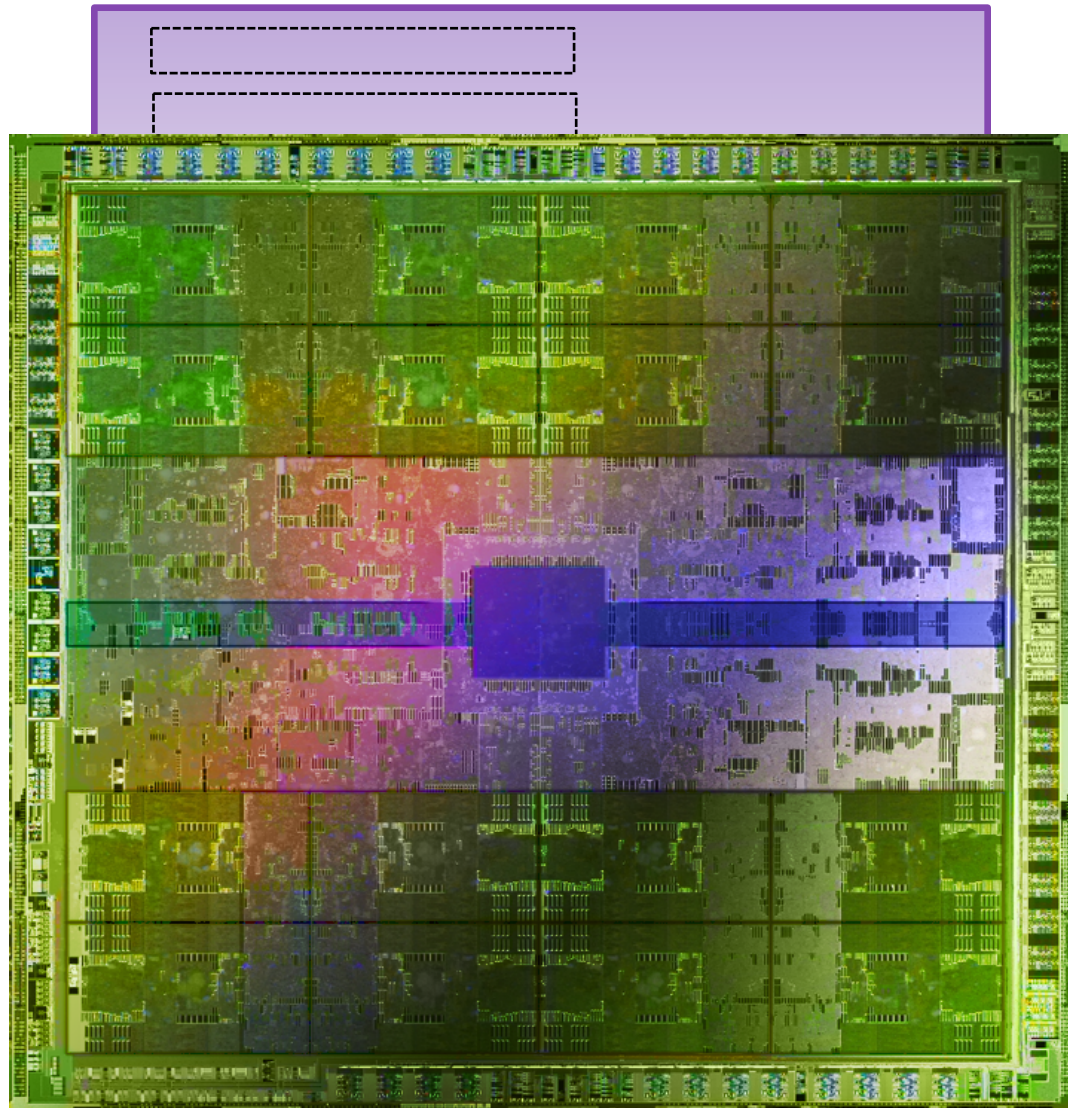


Read 32  
coalesced floats  
for max  
bandwidth usage

Much better utilization  
of the memory bus!

 = memory element (32 bits)

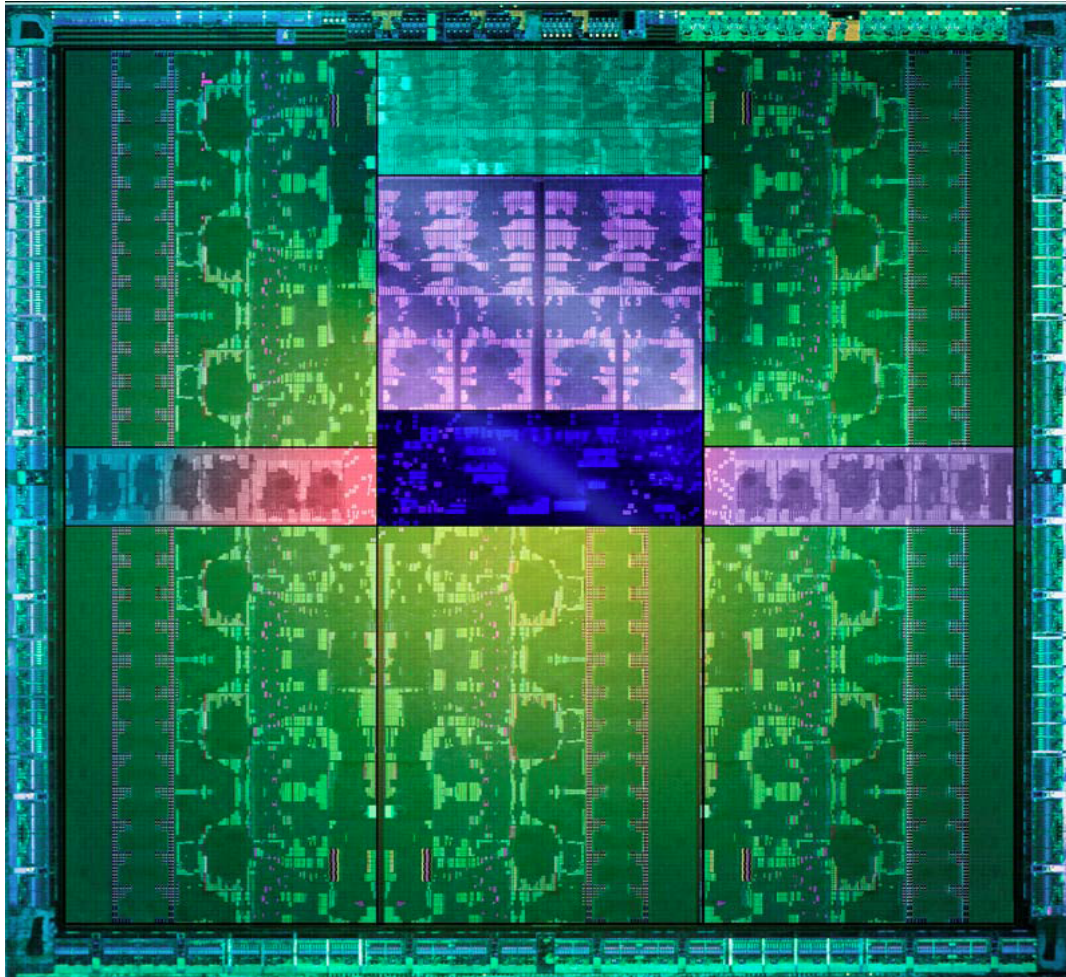
# Let's look at the GPU



Higher utilization  
memory bus!

■ = memory element (32 bits)

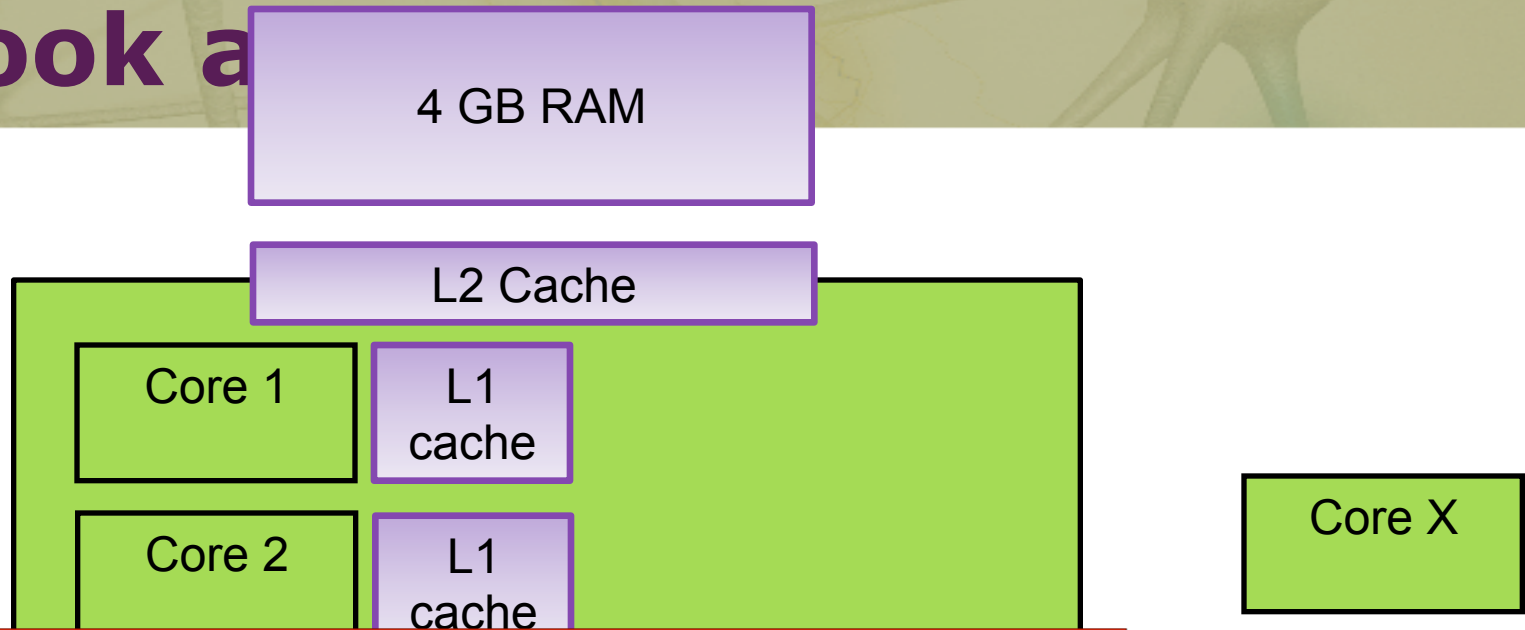
# Let's look at the GPU



Kepler: 15-16 multi-processors



# Let's look at



## Terminology

CPU:	Core	ALU (SIMD lane)
NVIDIA:	Streaming Multiprocessor	core
ATI	SIMD core	stream core

192 ALUs or "lanes"  
 (logically: 6 x 32-SIMD width)  
 6x32 mul/add per 1-2 clocks  
 (6x32 "threads")

SIMD = single instruction multiple data

Kepler: 15-16 multi-processors

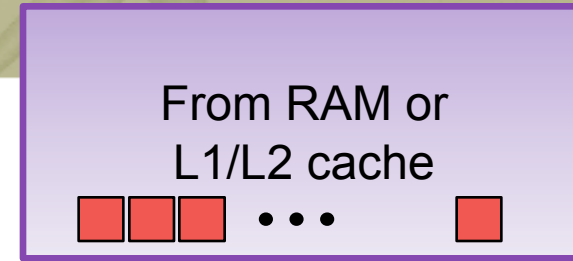
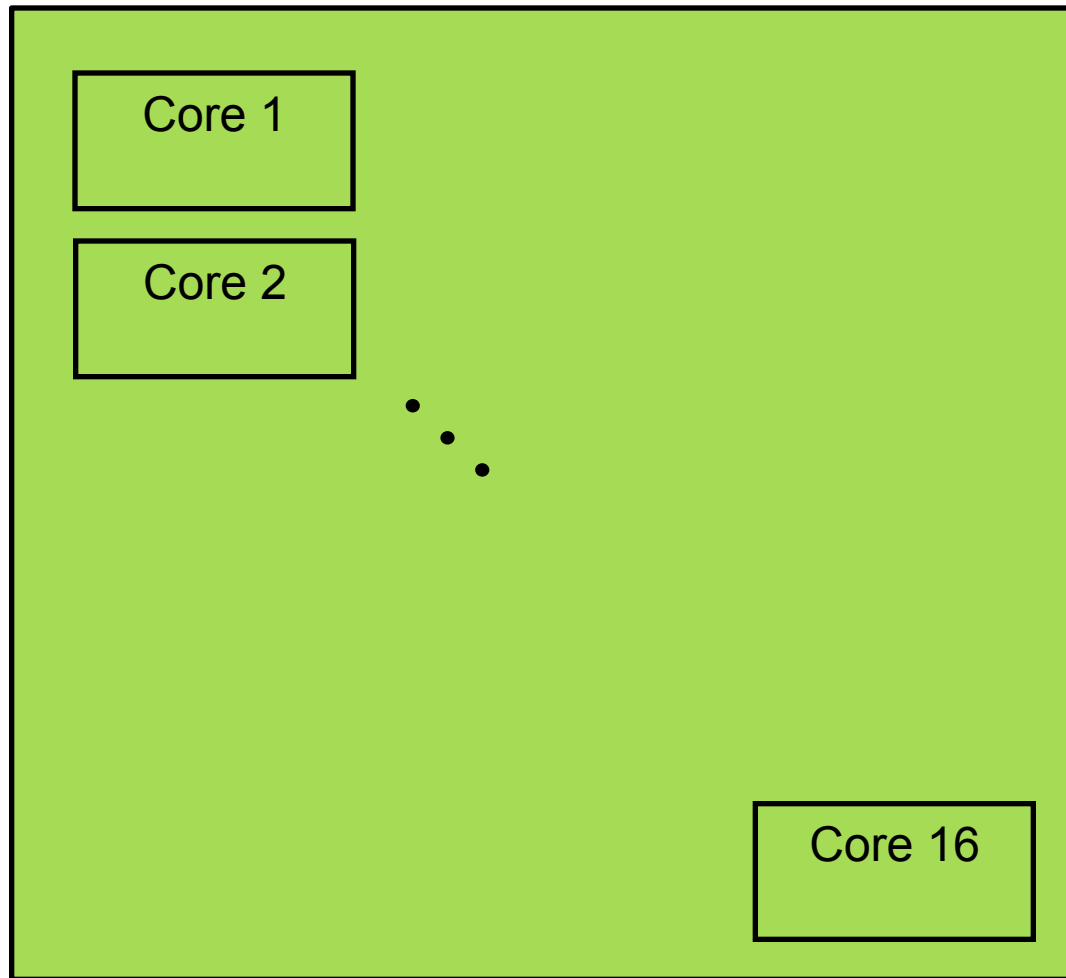
# Let's look at the GPU

## Each core:

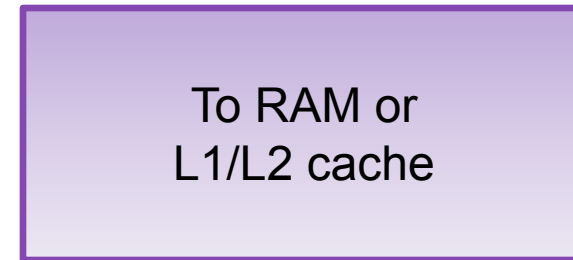
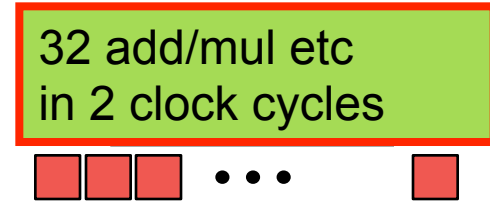
- executes one program (=shader).

## Each cycle:

- 192 flops
- 6x32 SIMD for up to 4 different instr.



6x



Kepler: 15-16 multi-processors

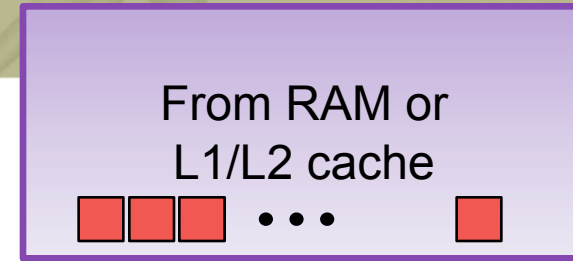
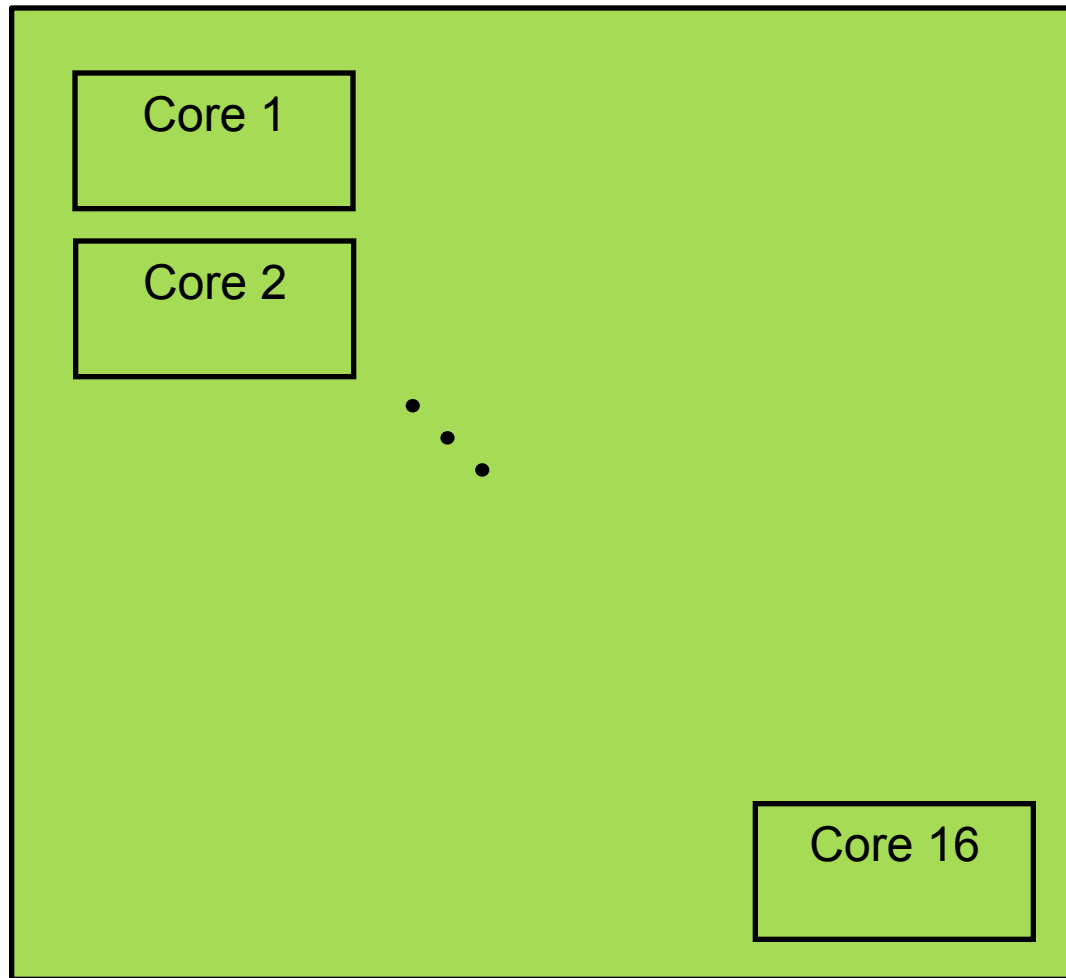
# Let's look at the GPU

## Each core:

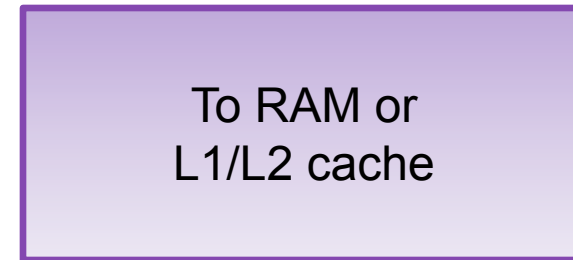
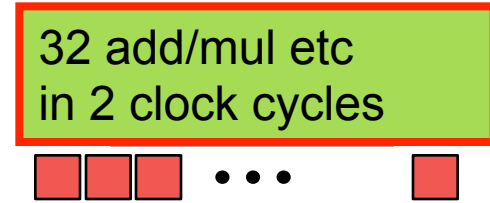
- executes one program (=shader).

## Each cycle:

- 192 flops
- 6x32 SIMD for up to 4 different instr.



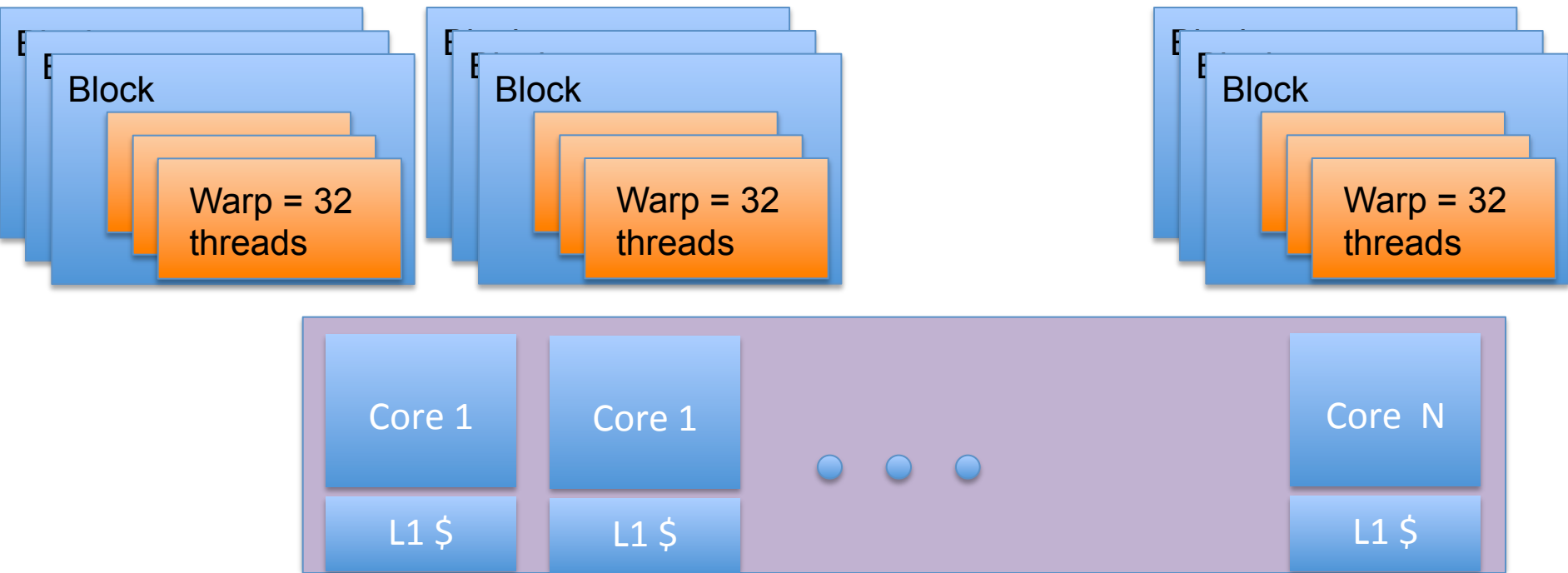
6x



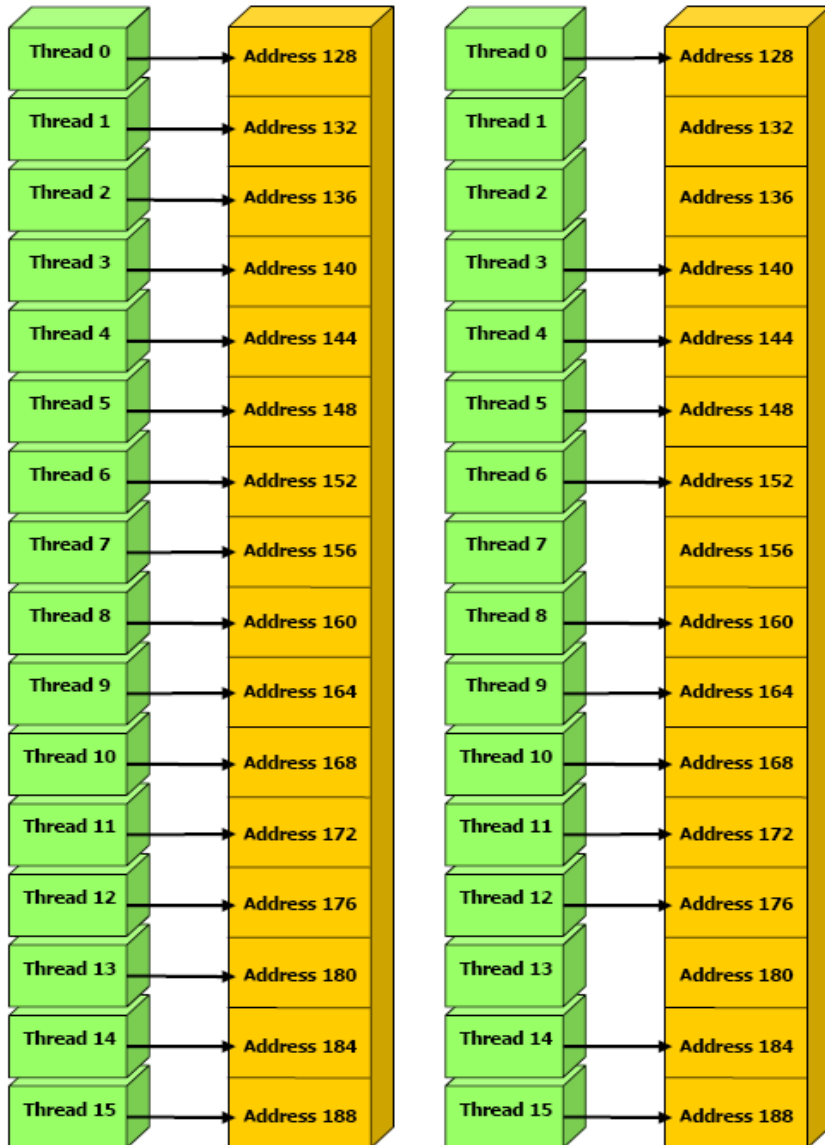
Kepler: 15-16 multi-processors

# CUDA

- A kernel (=CUDA program) is executed by 100:s-1M:s threads
  - A "warp" = 32 threads, one thread per ALU
  - Warps (one to ~32) are grouped into one block
  - Block: executed on one core
    - One to 48 warps execute on a core



# Memory Accesses – Global Memory



4 GB RAM

- Coalesced reads and writes
- For maximum performance, each thread should read from the same 16-float block (128 bytes)
  - i.e., the same cache-line

# Fermi

- Global mem accesses.
- One transaction:
- Two transactions:

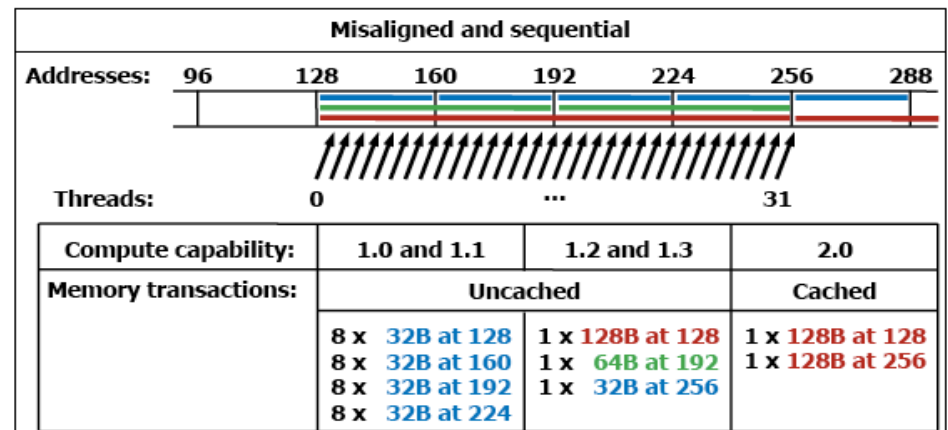
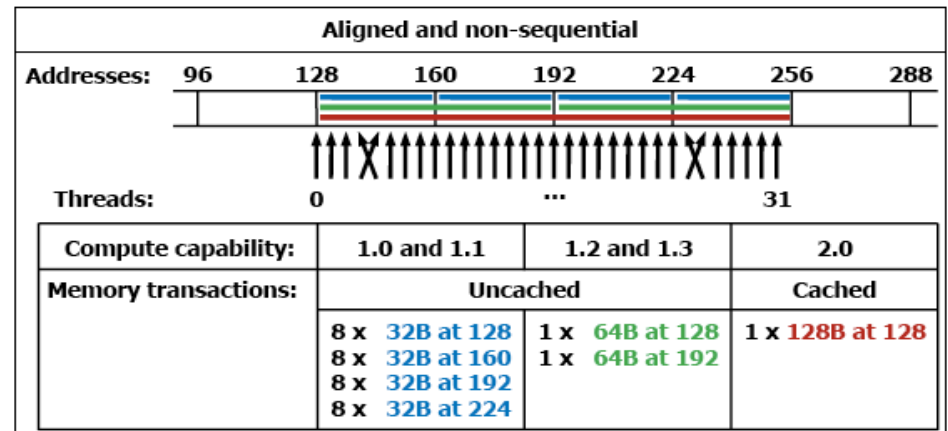
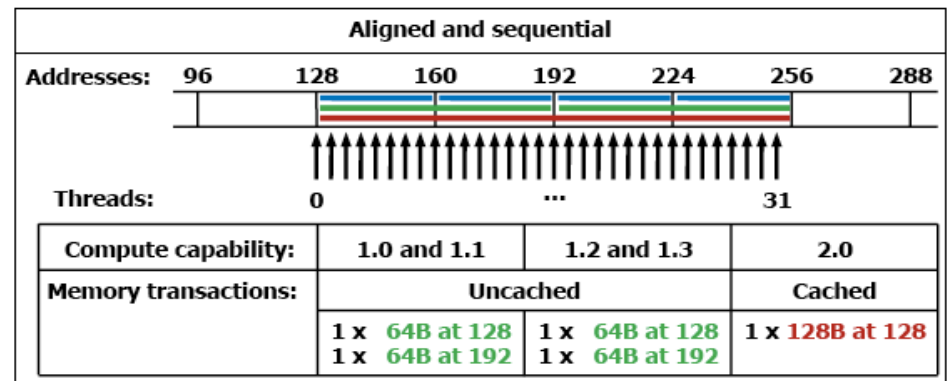
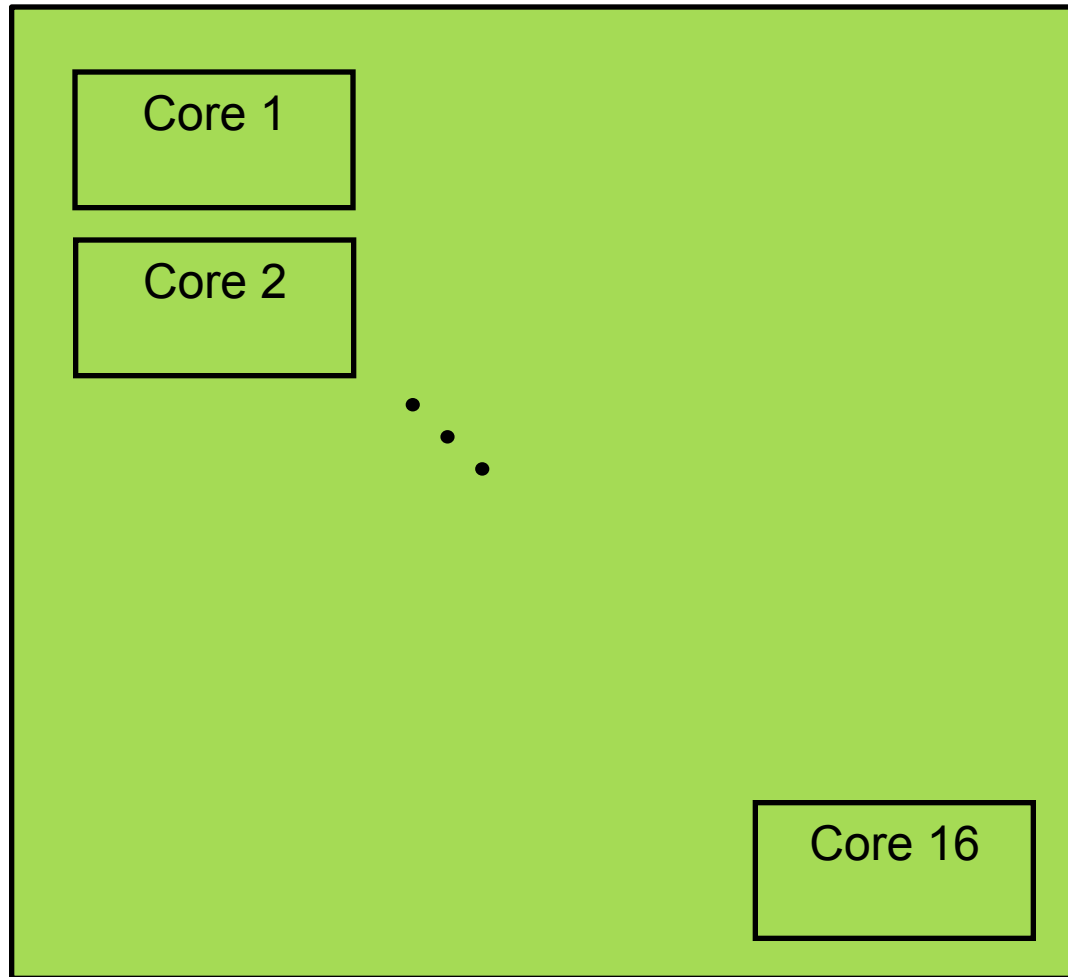


Figure G-1. Examples of Global Memory Accesses by a Warp, 4-Byte Word per Thread, and Associated Memory Transactions Based on Compute Capability

# Efficient programming



From RAM or L1/L2 cache

Four red squares are shown in a row, with three dots between the second and third squares, representing data being fetched from memory.

32 add/mul etc in 2 clock cycles

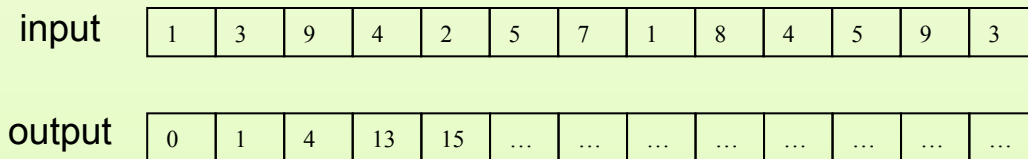
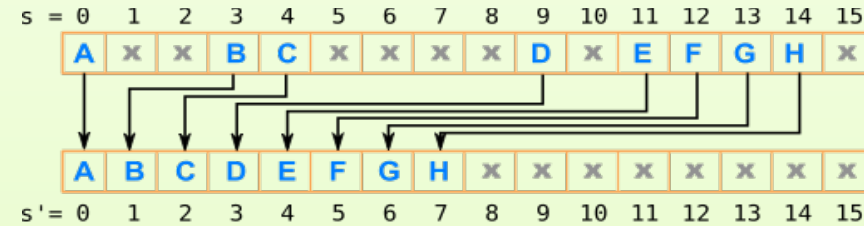
Four red squares are shown in a row, with three dots between the second and third squares, representing data being processed by the cores.

To RAM or L1/L2 cache

Fermi: 16 multi-processors à 2x16 SIMD width

# Efficient Programming

- If your program can be constructed this way, you are a winner!
- More often possible than expected
  - Stream compaction
  - Prefix sums
  - Sorting

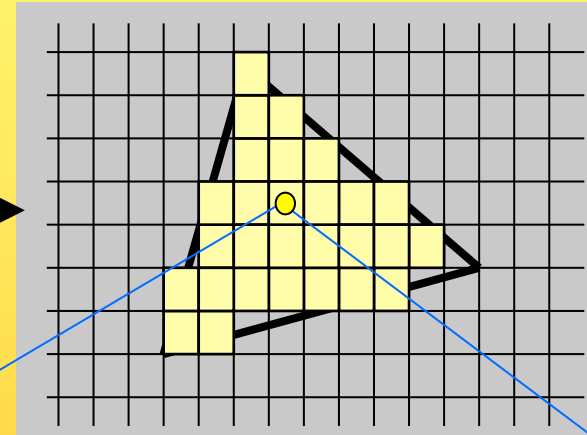
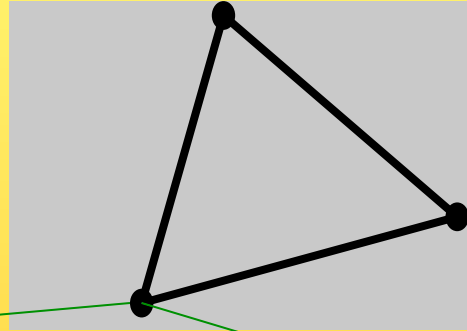
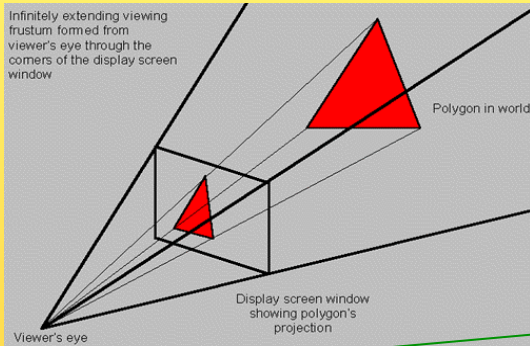


19 5 100 1 63 79  
 ↓  
 1 5 19 63 79 100

Fermi: 16 multi-processors à 2x16 SIMD width



# Shaders



```
// Vertex Shader
#version 130

in vec3 vertex;
in vec3 color;
out vec3 outColor;
uniform mat4 modelViewProjectionMatrix;

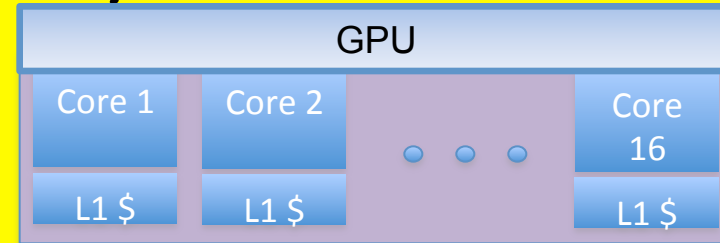
void main()
{
    gl_Position = modelViewProjectionMatrix*vec4(vertex,1);
    outColor = color;
}
```

```
// Fragment Shader:
#version 130
in vec3 outColor;
out vec4 fragColor;

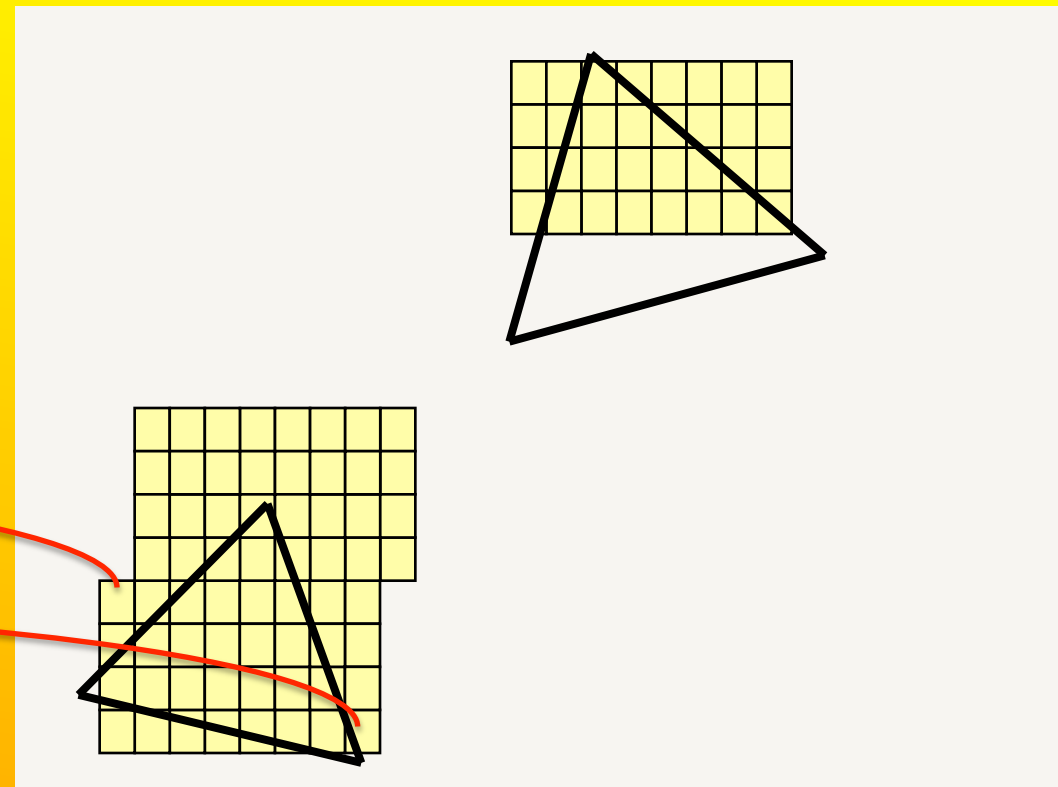
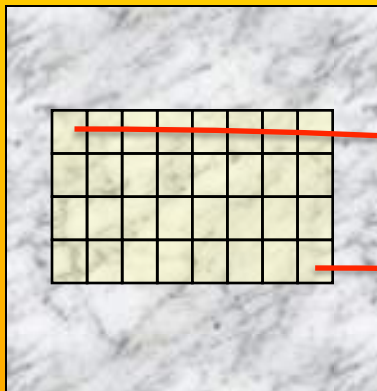
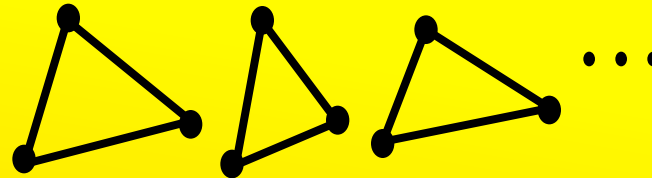
void main()
{
    fragColor = vec4(outColor,1);
}
```

# Shaders and coalesced memory accesses

- Varje core exekverar upp till 4 olika instruktioner per klockcykel för



- Vertex shader:
  - 192 vertices
- Fragment shader:
  - 192 pixlar  
i block om 32 pixlar  
(eller möjligen 16 pixlar)



# Thread utilization

- Each core executes one program (=shader)
- Each of the 192 ALUs execute one "thread" (a shader for a vertex or fragment)
- Since the core executes the same instruction for at least 32 threads (as far as the programmer is concerned)...
- If (...) ...the core must
  - Then,  $a = b + c$ ; execute both paths
  - ... if any of the 32 threads need the if and else-path.
- Else
  - $a = c + d$ ; But not if all need the same path.

# Need to know:

- Perspective correct interpolation (e.g. for textures)

- Taxonomy:

- Sort first
- sort middle
- sort last fragment
- sort last image

- Bandwidth

- Why it is a problem and how to "solve" it
  - Texture caching with prefetching
  - Texture compression
  - Z-compression
  - Z-occlusion testing (HyperZ)

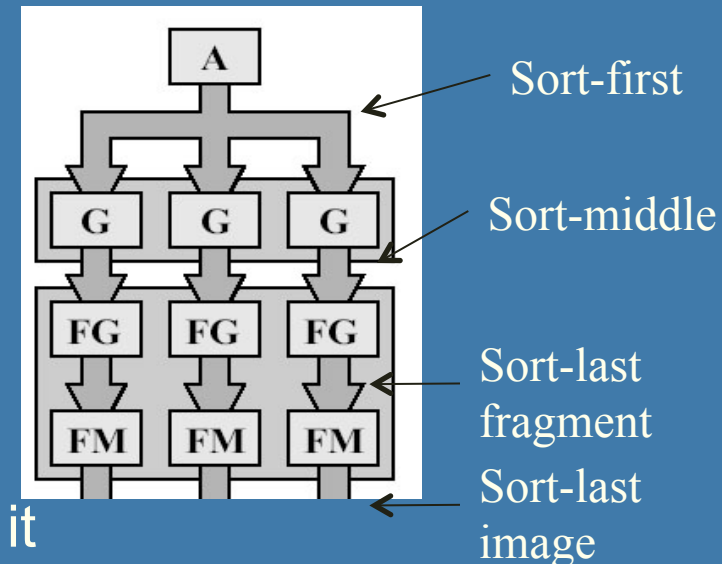
- Be able to sketch the architecture of a moder graphics card

Linearly interpolate  $(u_i/w_i, v_i/w_i, 1/w_i)$  in screenspace from each triangle vertex  $i$ .

Then at each pixel:

$$u_{ip} = (u_{ip}/w_{ip}) / (1/w_{ip})$$
$$v_{ip} = (v_{ip}/w_{ip}) / (1/w_{ip})$$

where  $ip$  = screen-space interpolated value from the triangle vertices.



# Need to know:

