

Monads reference card

Jean-Philippe Bernardy

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	monad component	DSL application
Syntax	$m :: \star \rightarrow \star$	Expressions parameterized on return type
	$return :: a \rightarrow m a$	constant expression
	$(\gg=) :: m a \rightarrow (a \rightarrow m b) \rightarrow m b$	bind an a returned by the lhs into the rhs

	name	law	a DSL aspect
Laws	left identity	$return a \gg= (\lambda x. m x) \equiv m a$	inlining/factorizing a constant
	right identity	$m \gg= (\lambda x. return x) \equiv m$	removal/introduction of useless return
	associativity	$(m \gg= f) \gg= g \equiv m \gg= (\lambda x. f x \gg= g)$	extension/shrinking of scope

“do”	do	$x \leftarrow \alpha$	$\alpha \gg= \lambda x.$
		$y \leftarrow \beta$	$\beta \gg= \lambda y.$
		γ	γ

- parentheses are not needed
- x may appear in γ

Comprehensions	[γ	
		$x \leftarrow \alpha$	$\alpha \gg= \lambda x.$
	,	$y \leftarrow \beta$	$\beta \gg= \lambda y.$
]	return γ	

- $\gg=$ can be used to “flatten” levels of the monad.
- $join :: m(m a) \rightarrow m a$
- $join xs = xs \gg= id$