Computer Communications DIT 420 EDA343

Summary

Computer Communication

Important for the exam

When/where: Friday Dec 21, 14.00-18.00 , H

You may have with you:

- English-X dictionary
- no calculators, PDAs, etc (if/where numbers matter, do rounding)

Grading

- 30-40, 41-50, 51-60 (out of 60)= 3, 4, 5 (CTH)
- 30-50, 51-60 (out of 60) = G, VG (GU)

To think during last, summary-study

Overview; critical eye; explain: why is this so? / How does it work?

Flashback

Principles, Organisation

Network Problems (in the order faced in the 1st intro):

- producer-consumer problems, flow and error control,
- manage access to shared (broadcast) transmission media ,
- routing,
- congestion,
- connecting transparently different networks,
- serving different types of traffic,
- performance,
- mobility
- security

Layering : principle, why

Highlights

- types of delay; performance
- reliable data transfer (flow, error control)
- datagram vs VC end-to-end communication, congestion control
- routing, also with mobility
- multiple access protocols (wired, wireless)
- LANs and related technologies
- network security issues
- TCP/IP protocol stack (also applications), evolution (p2p applications, overlays, NAT, streaming apps)

Types of delay; performance

Types of delay; performance

- Propagation, transmission, queueing, processing
- Throughput (effective bandwidth) ٠
- Utilization (efficiency) ٠
- Packet-switching: impact of store&forward ٠

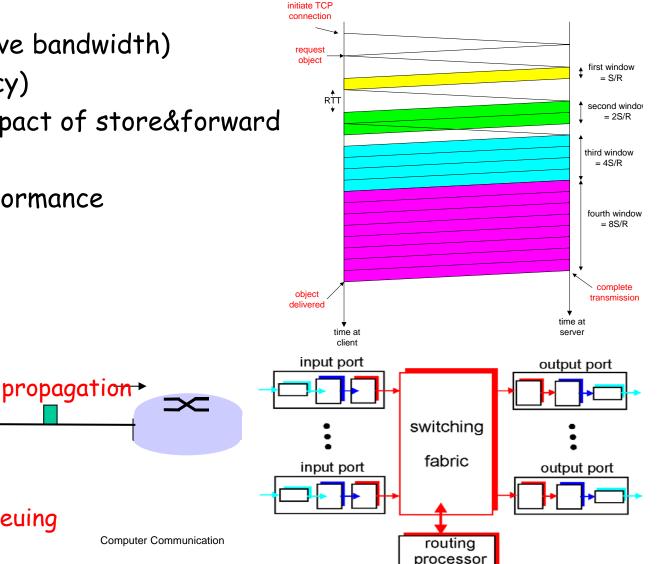
queuing

- TCP's slow start ٠
- Sliding windows performance ٠

transmission

nodal

processing

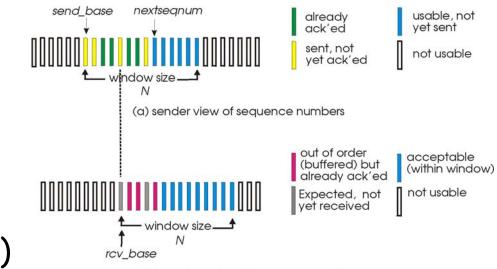


Reliable data transfer

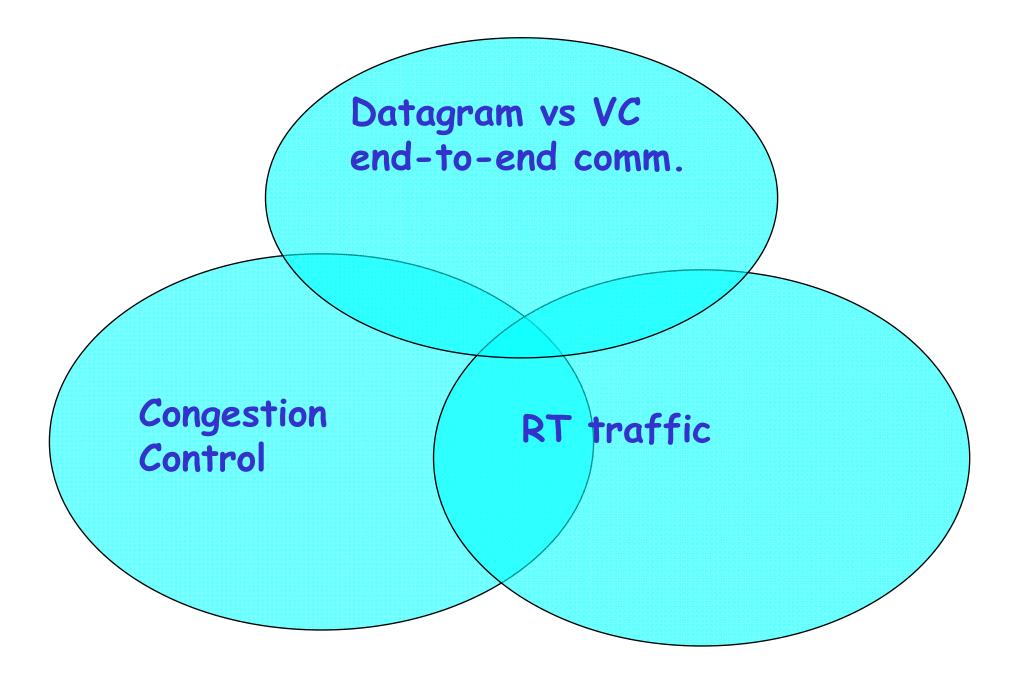
Reliable data transfer

Guaranteed, in-order, correct delivery:

- Flow control:
 - stop&wait
 - sliding windows
 - sequence numbers
 - window sizes
 - dynamic windows (TCP)
 - performance
- Error detection: checksums, CRC (*)
- Error control: go-back-n, selective repeat, FEC methods



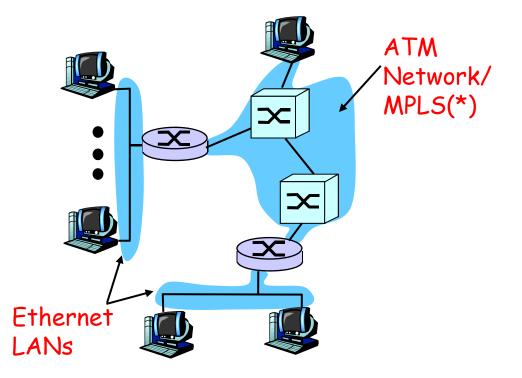
(b) receiver view of sequence numbers



Computer Communication

Datagram vs VC end-to-end communication

- Conceptual differences
- Decisions, comparison, why



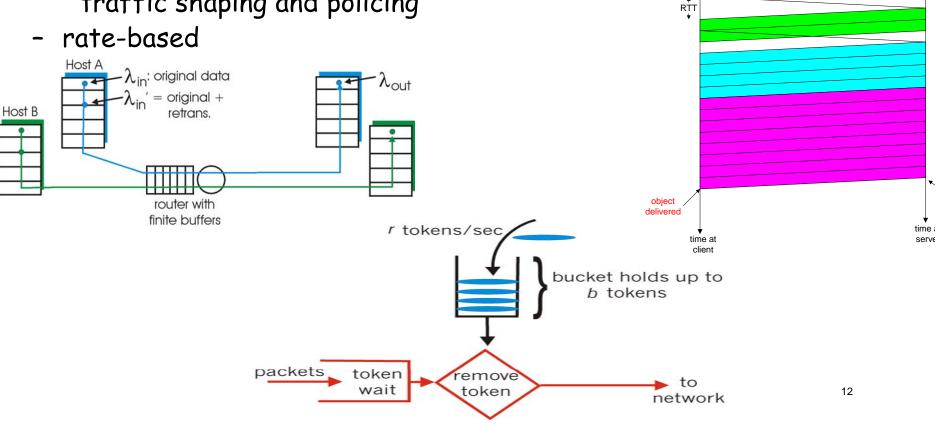
Congestion control (CC)

initiate TCP

connection

request object

- why, how congestion occurs
- CC in TCP and performance; implied weaknesses
- CC in other ways, e.g. VC-based networks
 - RT-traffic resource reservation: traffic shaping and policing



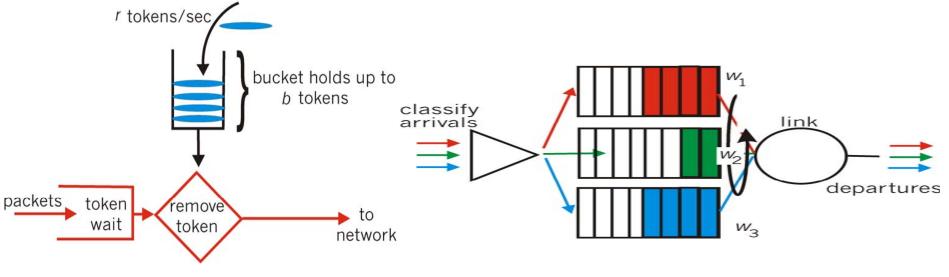
RT/streaming traffic

Conceptual needs:

- packet/flow marking
- Admission control
- Traffic shaping & policing
- Packet scheduling (switches)

Internet context

- Application-level solutions (FEC, playout delay, caching-CDN)
 - Intserv, Diffserv



Computer Communication

Routing, also with mobility

Routing, also with mobility

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B

3

1

Mobile

- Routing algorithms
- Forwarding
- Resource, policy issues
- · Addressing mobility, tunneling

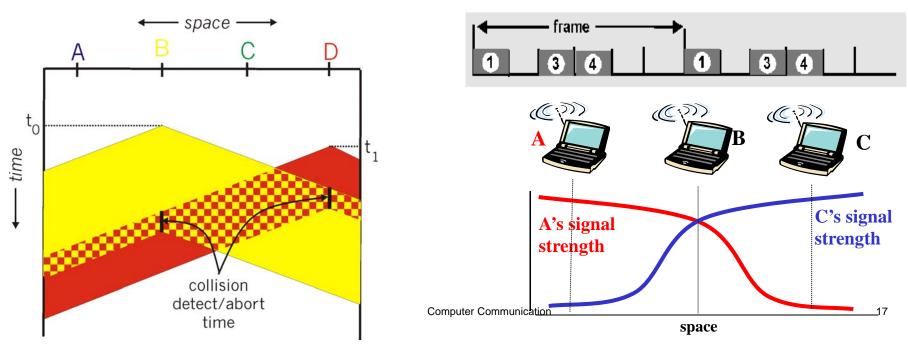
Switching Center **Public telephon** Complementary video links network, and Internet - IP addresses and subnets Mobile http://www.youtube.com/watch?v=ZTJIkjgyuZE Switching Center &list=PLE9F3F05C381ED8E8&feature=plcp - How does BGP choose its routes http://www.youtube.com/watch?v=RGe0qt9Wz4U &feature=plcp Computer Communication 15

Multiple access algorithms

Medium access: multiple access methods

Strategies: (functionality, appropriateness)

- Contention-based (random access), wired/wireless:
 - Aloha, CSMA(CD/CA) (collision-delay trade-off)
- Collision-free:
 - Channel partitioning: TDMA, FDMA, CDMA
 - Taking turns: token-passing, reservation-based



LANs & related technologies

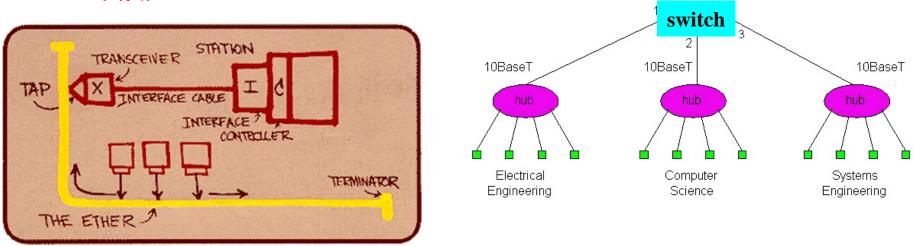
LANs & related link technologies

Protocol Examples: wired, wireless

Ethernet, 802.11 (+ 802.16 wimax), GSM:

Functionality, performance under low/high load

- Connecting devices;
 - functionalities and differences (Hubs, switches)
 - Algorithms for switch-"routing": learning& forwarding of packets
- · ARP



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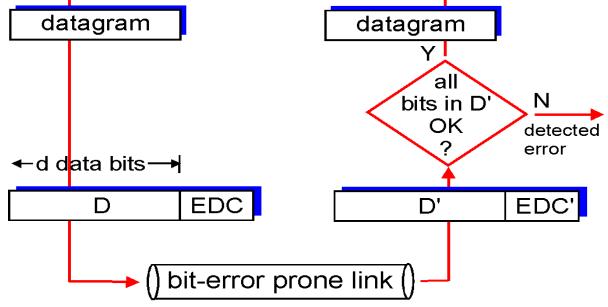
BTW: A little "backpatching"

Computer Communication

Error Detection

EDC= Error Detection and Correction bits (redundancy)

- D = Data protected by error checking, may include header fields
- Error detection not 100% reliable!
 - protocol may miss some errors, but this should happen only rarely
 - larger EDC field yields better detection and correction



Internet checksums

TCP (UDP)'s checksum:

- segment contents = sequence of 16-bit integers
- checksum: addition (1's complement sum) of segment contents
- sender puts checksum value into UDP (TCP) checksum field

<u>m:</u> <u>Cyclic redundancy check (CRC)</u>

- more powerful error-detection coding
- data bits, D = binary number
- consider r+1 bit pattern (generator), G
- □ goal: compute r CRC bits, R, such that
 - O,R> exactly divisible by G (modulo 2)
 - receiver knows G, divides <D,R> by G. If non-zero remainder: error detected!
 - can detect errors on less than r+1 bits

D: data bits to be sent R: CRC bits

bit pattern

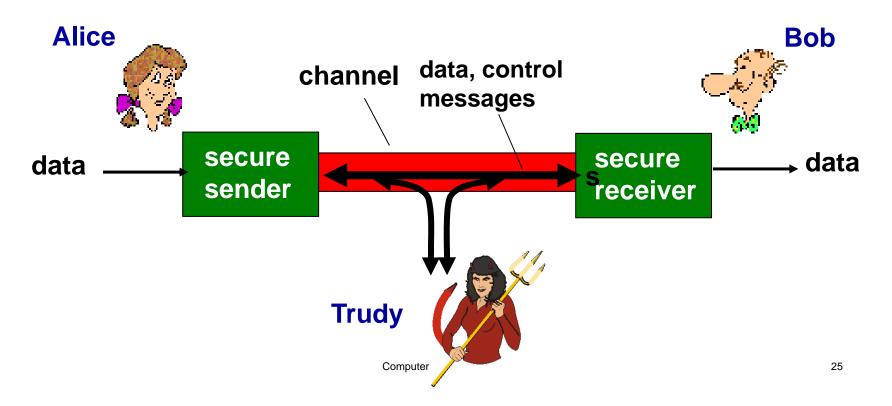
mathematical formula

CRC Example 3 Recall we want: 101011 $D \cdot 2^r XOR R = nG$ 10111 1001 0.000 equivalently: G if we divide $D \cdot 2^r$ by 10 G, want 101 remainder R 100 11 110 1001 101 R = remainder $\left[\frac{D \cdot 2^{r}}{G}\right]$ 1001 R

Security issues

Security issues

- C, I, A and methods to achieve them
- Instantiation in Internet: RSA, email PGP, authentication
- Firewalls and packet filtering



TCP/IP protocol stack (also applications), evolution

TCP/IP protocol stack (also applications), evolution

- Instantiation of networksolutions (Routing, Congestion Control, Flow & error control, applications, link layer technologies)
- Limitations, advantages, updates
- Application-layer networking
 (P2P applications, overlays, CDNs, multimedia/streaming application protocols)

| application |
|-------------|
| transport |
| network |
| link |
| physical |

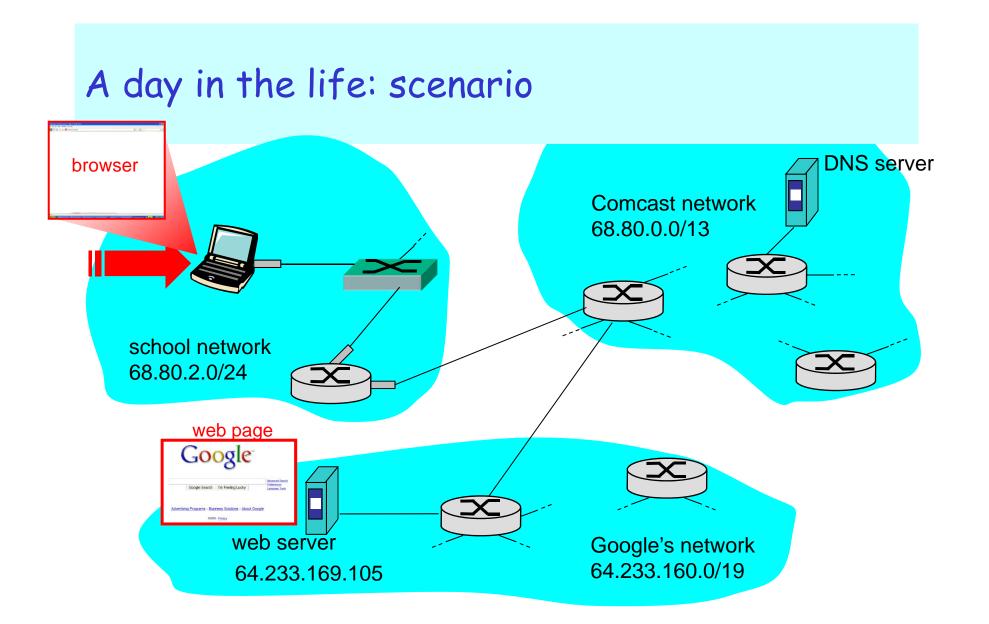
Items we discussed in less detail

- But exist in the slides you received
- Good to read (<u>remember: reading is from the book!</u> <u>The slides cannot replayce the book reading</u>) for own benefit; but need not consider for the exam next week
 - CRC's
 - Intserv, Diffserv, RSVP
 - MPLS

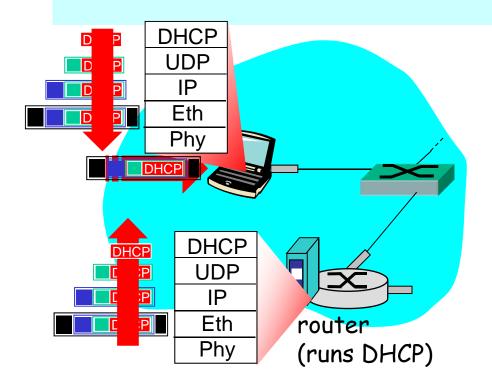
Synthesis: a day in the life of a web request

putting-it-all-together: synthesis!

- *goal:* identify, review protocols (at all layers) involved in seemingly simple scenario: requesting www page
- *scenario:* student attaches laptop to campus network, requests/receives www.google.com

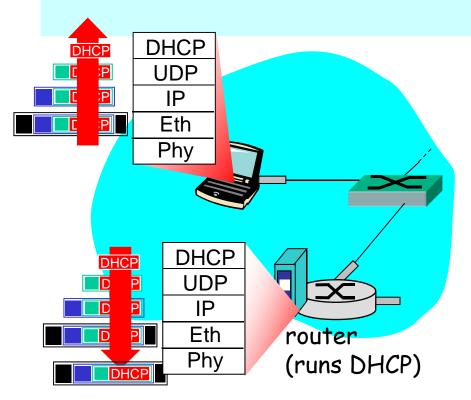


A day in the life ... connecting to the Internet



- connecting laptop needs to get its own IP address: use DHCP
- DHCP request encapsulated in UDP, encapsulated in IP, encapsulated in Ethernet
- Ethernet frame broadcast (dest: FFFFFFFFFFF) on LAN, received at router running DHCP server
- Ethernet demux'ed to IP demux'ed to UDP demux'ed to DHCP

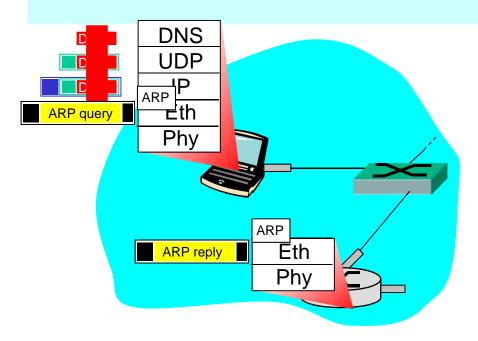
A day in the life ... connecting to the Internet



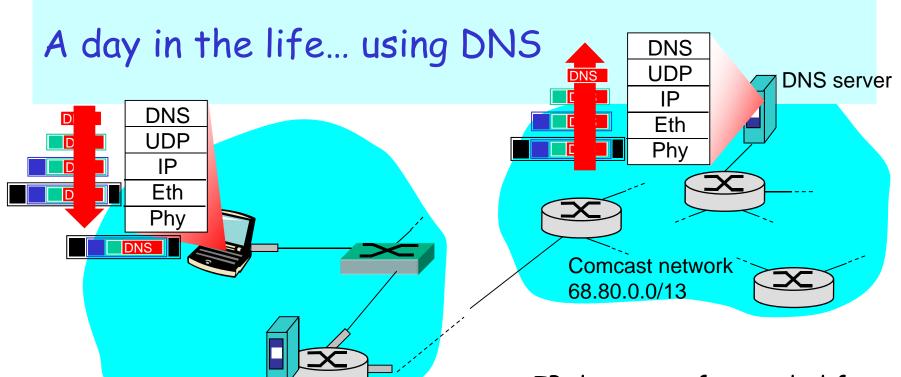
- DHCP server formulates DHCP ACK containing client's IP address (and also IP address of firsthop router for client, name & IP address of DNS server)
- frame forwarded (switch learning) through LAN, demultiplexing at client
- DHCP client receives DHCP ACK reply

Client now has IP address, knows name & addr of DNS server, IP address of its first-hop router

A day in the life... ARP (before DNS, before HTTP)

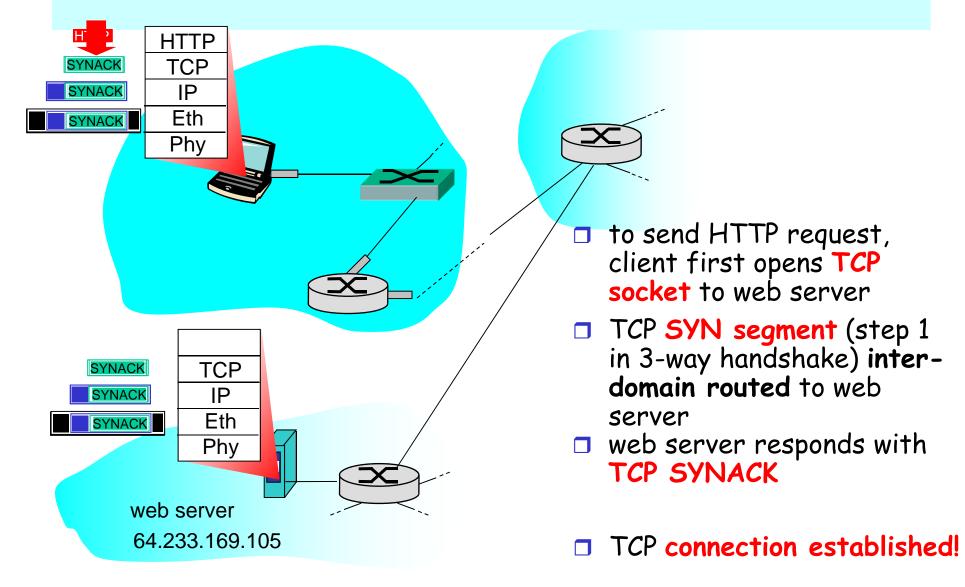


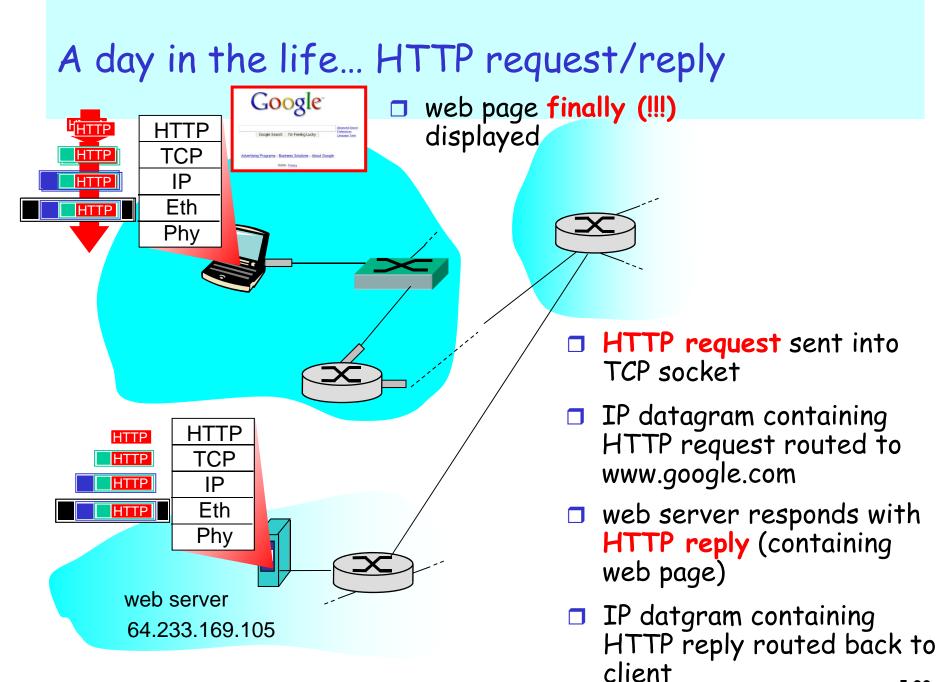
- before sending *HTTP* request, need IP address of www.google.com: *DNS*
- DNS query created, encapsulated in UDP, encapsulated in IP, encasulated in Eth. In order to send frame to router, need MAC address of router interface: ARP
- ARP query broadcast, received by router, which replies with ARP reply giving MAC address of router interface
- client now knows MAC address of first hop router, so can now send frame containing DNS query



- IP datagram containing DNS query forwarded via LAN switch from client to 1st hop router
- IP datagram forwarded from campus network to destination (DNS-server) network, routed (tables created by RIP, OSPF and BGP routing protocols) to DNS server
- demux'ed to DNS server
- DNS server replies to client with IP address of www.google.com_{5: DataLink Layer} 34

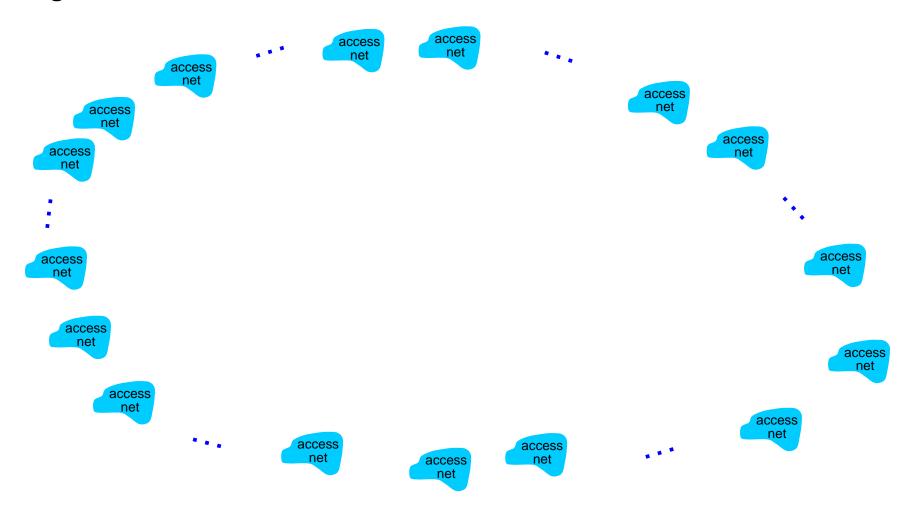
A day in the life... TCP connection carrying HTTP



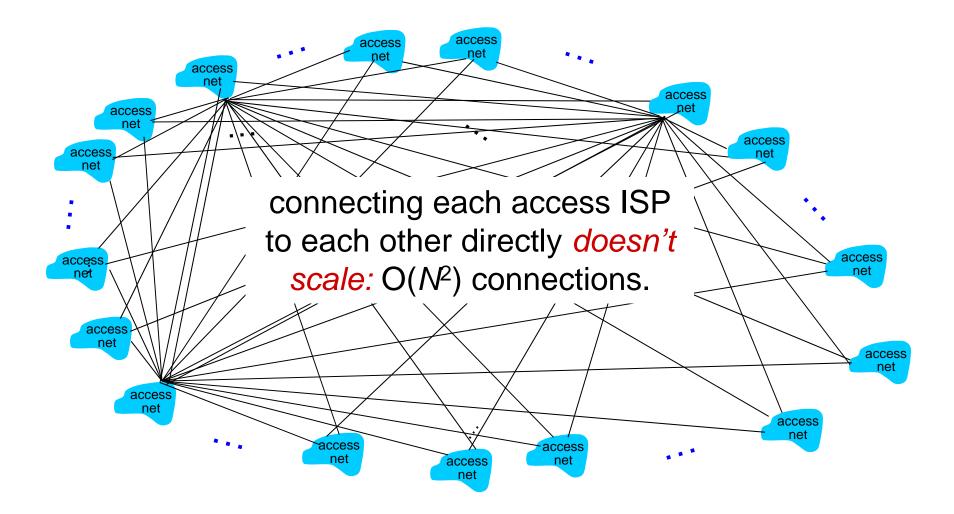


Synthesis cont.

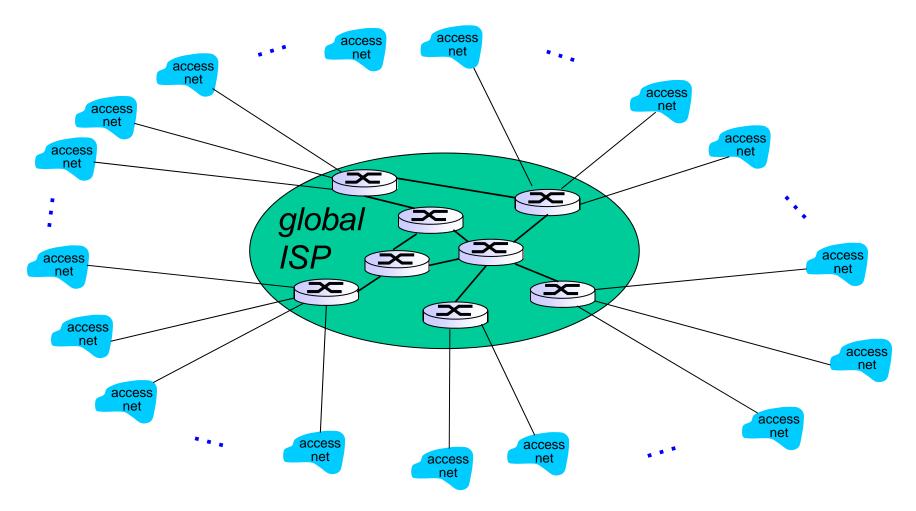
Question: given millions of access ISPs, how to connect them together?



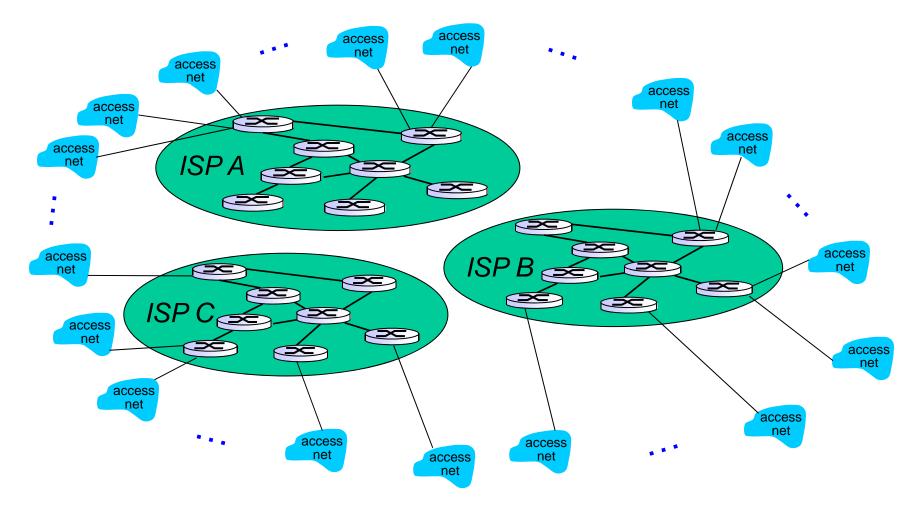
Option: connect each access ISP to every other access ISP?



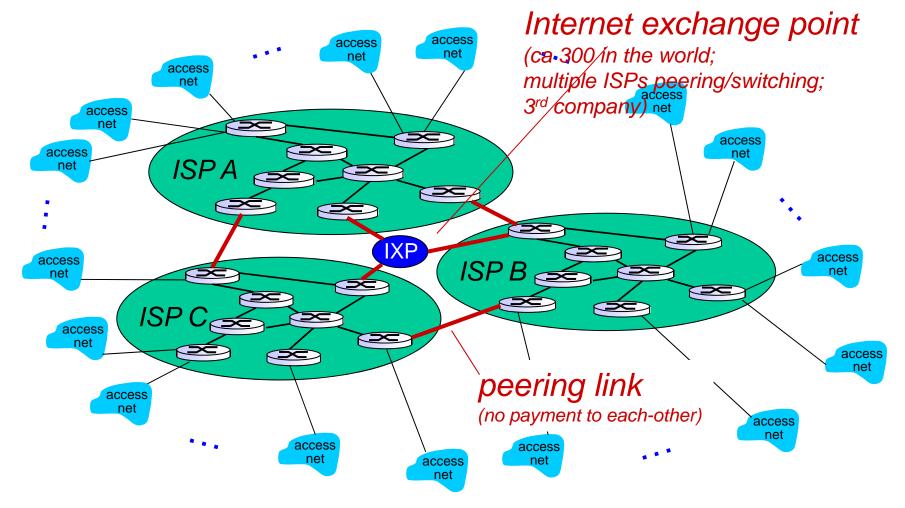
Option: connect each access ISP to a global transit ISP? **Customer** and **provider** ISPs have economic agreement.



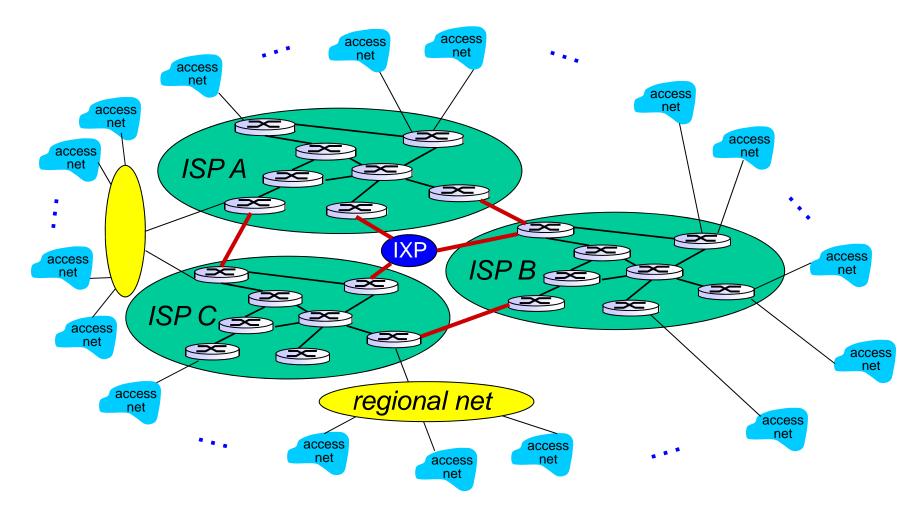
But if one global ISP is viable business, there will be competitors



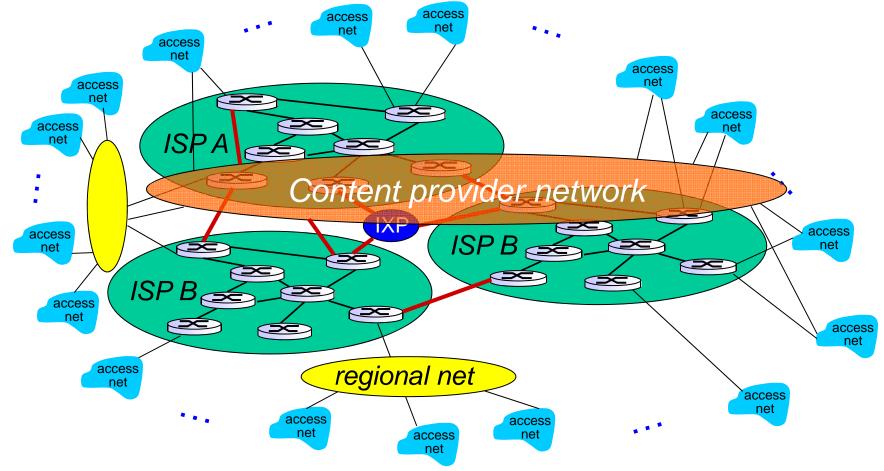
But if one global ISP is viable business, there will be competitors which must be interconnected

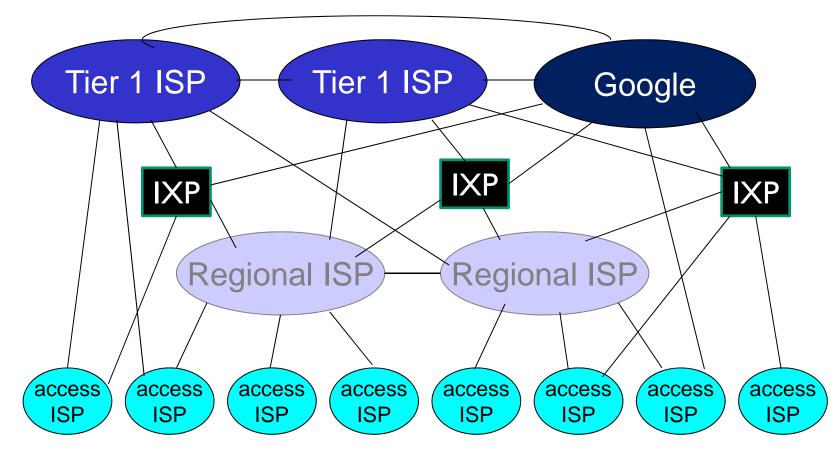


... and regional networks may arise to connect access nets to ISPS



... and content provider networks (e.g., Google, Microsoft, Akamai) may run their own network, to bring services, content close to end users





- at center: small # of well-connected large networks
 - "tier-I" commercial ISPs (e.g., Level 3, Sprint, AT&T, NTT), national & international coverage
 - A new form of content provider network (e.g, Google): private network that connects it data centers to Internet, often bypassing tier-I, regional ISPs

<u>Thank you &</u> good luck in the exam!!

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