Intersection Testing Chapter 16

Tomas Akenine-Möller Department of Computer Engineering Chalmers University of Technology

What for?

- A tool needed for the graphics people all the time...
- Very important components:
 - Need to make them fast!
- Finding if (and where) a ray hits an object
 - Picking
 - Ray tracing and global illumination
- For speed-up techniques
- Collision detection (treated in a later lecture)





Midtown Madness 3, DICE

Some basic geometrical primitives

• Ray:

• Sphere:

• k-DOP

• Box

Axis-aligned (AABB)

- Oriented (OBB)



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Four different techniques

- Analytical
- Geometrical
- Separating axis theorem (SAT)
- Dynamic tests

- Given these, one can derive many tests quite easily
 - However, often tricks are needed to make them fast

Analytical: Ray/sphere test

- Sphere center: c, and radius r
- Ray: $\mathbf{r}(t) = \mathbf{o} + t\mathbf{d}$
- Sphere formula: ||**p**-**c**||=*r*



• Replace \mathbf{p} by $\mathbf{r}(t)$, and square it:

$$(\mathbf{r}(t) - \mathbf{c}) \cdot (\mathbf{r}(t) - \mathbf{c}) - r^2 = 0$$

$$(\mathbf{o} + t\mathbf{d} - \mathbf{c}) \cdot (\mathbf{o} + t\mathbf{d} - \mathbf{c}) - r^2 = 0$$

$$(t\mathbf{d} + (\mathbf{o} - \mathbf{c})) \cdot (t\mathbf{d} + (\mathbf{o} - \mathbf{c})) - r^2 = 0$$

$$(\mathbf{d} \cdot \mathbf{d})t^2 + 2((\mathbf{o} - \mathbf{c}) \cdot \mathbf{d})t + (\mathbf{o} - \mathbf{c}) \cdot (\mathbf{o} - \mathbf{c}) - r^2 = 0$$

$$t^{2} + 2((\mathbf{o} - \mathbf{c}) \cdot \mathbf{d})t + (\mathbf{o} - \mathbf{c}) \cdot (\mathbf{o} - \mathbf{c}) - r^{2} = 0 \quad ||\mathbf{d}|| = 1$$



Geometrical: Ray/Box Intersection

- Boxes and spheres often used as bounding volumes
- A slab is the volume between two parallell planes:

 A box is the logical intersection of three slabs (2 in 2D):



Geometrical: Ray/Box Intersection (2)

 Intersect the 2 planes of each slab with the ray



Keep max of t^{min} and min of t^{max}
If t^{min} < t^{max} then we got an intersection
Special case when ray parallell to slab

Separating Axis Theorem (SAT) Page 563 in book

- Two convex polyhedrons, A and B, are disjoint if any of the following axes separate the objects:
 - An axis orthogonal to a face of A
 - An axis orthogonal to a face of B
 - An axis formed from the cross product of one edge from each of A and B

axis

A and B overlaps on this axis

SAT example: Triangle/Box

- E.g an axis-aligned box and a triangle
- 1) test the axes that are orthogonal to the faces of the box
- That is, x,y, and z



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Triangle/Box with SAT (2)

- Assume that they overlapped on x,y,z
- Must continue testing
- 2) Axis orthogonal to face of triangle



Triangle/Box with SAT (3)

- If still no separating axis has been found...
- 3) Test axis: $t = e_{box} \times e_{triangle}$
- Example:
 - x-axis from box: $e_{box} = (1,0,0)$
 - $\mathbf{e}_{triangle} = \mathbf{v}_1 \mathbf{v}_0$
- Test all such combinations
- If there is at least one separating axis, then the objects do not collide
- Else they do overlap

Rules of Thumb for Intersection Testing

- Acceptance and rejection test
 - Try them early on to make a fast exit
- Postpone expensive calculations if possible
- Use dimension reduction
 - E.g. 3 one-dimensional tests instead of one complex 3D test, or 2D instead of 3D
- Share computations between objects if possible
- Timing!!!

Another analytical example: Ray/ **Triangle in detail** • Ray: $\mathbf{r}(t) = \mathbf{0} + t\mathbf{d}$ • Triangle vertices: \mathbf{v}_0 , \mathbf{v}_1 , \mathbf{v}_2 • A point in the triangle: • $\mathbf{t}(u,v) = \mathbf{v}_0 + u(\mathbf{v}_1 - \mathbf{v}_0) + v(\mathbf{v}_2 - \mathbf{v}_0) =$ $(1-u-v)\mathbf{v}_0 + u\mathbf{v}_1 + v\mathbf{v}_2 \quad [u,v \ge 0, u+v \le 1]$ • Set t(u,v) = r(t), and solve!

$$\begin{pmatrix} | & | & | \\ -\mathbf{d} & \mathbf{v}_1 - \mathbf{v}_0 & \mathbf{v}_2 - \mathbf{v}_0 \\ | & | & | \end{pmatrix} \begin{pmatrix} t \\ u \\ v \end{pmatrix} = \begin{pmatrix} | \\ \mathbf{o} - \mathbf{v}_0 \\ | \end{pmatrix}$$



Ray/Triangle (2)
$$\begin{pmatrix} | & | & | & | \\ -\mathbf{d} & \mathbf{v}_{1} - \mathbf{v}_{0} & \mathbf{v}_{2} - \mathbf{v}_{0} \\ | & | & | & | \end{pmatrix}\begin{pmatrix} t \\ u \\ v \end{pmatrix} = \begin{pmatrix} | \\ \mathbf{o} - \mathbf{v}_{0} \\ | \end{pmatrix}$$
$$\mathbf{e}_{1} = \mathbf{v}_{1} - \mathbf{v}_{0} \quad \mathbf{e}_{2} = \mathbf{v}_{2} - \mathbf{v}_{0} \quad \mathbf{s} = \mathbf{o} - \mathbf{v}_{0}$$
$$\mathbf{e}_{1} = \mathbf{v}_{1} - \mathbf{v}_{0} \quad \mathbf{e}_{2} = \mathbf{v}_{2} - \mathbf{v}_{0} \quad \mathbf{s} = \mathbf{o} - \mathbf{v}_{0}$$
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Use this fact: $det(a, b, c) = (a \times b) \cdot c = -(a \times c) \cdot b$

$$\begin{pmatrix} t \\ u \\ v \end{pmatrix} = \frac{1}{(\mathbf{d} \times \mathbf{e}_2) \cdot \mathbf{e}_1} \begin{pmatrix} (\mathbf{s} \times \mathbf{e}_1) \cdot \mathbf{e}_2 \\ (\mathbf{d} \times \mathbf{e}_2) \cdot \mathbf{s} \\ (\mathbf{s} \times \mathbf{e}_1) \cdot \mathbf{d} \end{pmatrix}$$

• Share factors to speed up computations

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Ray/Triangle (3) Implementation

$$\begin{pmatrix} t \\ u \\ v \end{pmatrix} = \frac{1}{(\mathbf{d} \times \mathbf{e}_2) \cdot \mathbf{e}_1} \begin{pmatrix} (\mathbf{s} \times \mathbf{e}_1) \cdot \mathbf{e}_2 \\ (\mathbf{d} \times \mathbf{e}_2) \cdot \mathbf{s} \\ (\mathbf{s} \times \mathbf{e}_1) \cdot \mathbf{d} \end{pmatrix}$$

• Be smart!

- Compute as little as possible. Then test

• Examples: $\mathbf{p} = \mathbf{d} \times \mathbf{e}_2$

$$a = \mathbf{p} \cdot \mathbf{e}_1$$
$$f = 1/a$$

• Compute $u = f(\mathbf{p} \cdot \mathbf{s})$

• Then test valid bounds

• if (u<0 or u>1) exit;

Plane: $\pi : \mathbf{n} \cdot \mathbf{p} + d = 0$ **Point/Plane** • Insert a point x into plane equation: $f(\mathbf{x}) = \mathbf{n} \cdot \mathbf{x} + d = ?$ $f(\mathbf{x}) = \mathbf{n} \cdot \mathbf{x} + d = 0$ for x's on the plane $f(\mathbf{x}) = \mathbf{n} \cdot \mathbf{x} + d < 0$ for **x**'s on one side of the plane Negative half space $f(\mathbf{x}) = \mathbf{n} \cdot \mathbf{x} + d > 0$ for x's on the other side Positive half space origin \mathbf{X}_1 n $\mathbf{n} \cdot \mathbf{x}_2 = ||\mathbf{x}_2|| \cos \gamma < 0$ $\mathbf{n} \cdot \mathbf{x}_1 = ||\mathbf{x}_1|| \cos \phi > 0$ π \mathbf{X}_{2}

Sphere/Plane AABB/Plane

Plane: $\pi : \mathbf{n} \cdot \mathbf{p} + d = 0$ Sphere: $\mathbf{c} \qquad r$ Box: $\mathbf{b}^{\min} \quad \mathbf{b}^{\max}$

Sphere: compute f(c) = n · c + d
f(c) is the signed distance (n normalized)
abs(f(c)) > r no collision
abs(f(c)) = r sphere touches the plane
abs(f(c)) < r sphere intersects plane

- Box: insert all 8 corners
- If all f's have the same sign, then all points are on the same side, and no collision

AABB/plane

Plane: $\pi : \mathbf{n} \cdot \mathbf{p} + d = 0$ Sphere: $\mathbf{c} \qquad r$ Box: $\mathbf{b}^{\min} \quad \mathbf{b}^{\max}$

The smart way (shown in 2D)
Find diagonal that is most closely aligned with plane normal



Need only test the red points

More details in book

Ray/Polygon: very briefly Intersect ray with polygon plane Project from 3D to 2D • How? • Find $\max(|n_x|, |n_v|, |n_z|)$ • Skip that coordinate! Then, count crossing in 2D





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View frustum testing

- View frustum is 6 planes:
- Near, far, right, left, top, bottom



- Create planes from projection matrix
 - Let all positive half spaces be outside frustum
 - Not dealt with here -- p. 773-774, 3rd ed.
- Sphere/frustum common approach:
 - Test sphere against 6 frustum planes
 - If in positive half space and r distances from plane
 no intersection
 - If intersecting plane or in negative half space, continue
 - If not outside after all six planes, then inside or intersecting
- Example follows...

View frustum testing example



Not exact test, but not incorrect
 A sphere that is reported to be inside, can be outside
 Not vice versa

Similarly for boxes



- Testing is often done every rendered frame, i.e., at discrete time intervals
- Therefore, you can get "quantum effects"

Frame n

Frame *n*+1

- Dynamic testing deals with this
- Is more expensive

 Deals with a time interval: time between two frames

Dynamic intersection testing Sphere/Plane t=n $s_c \& s_o$ are signed distances S_{c} S_e • No collision occur: t=n+1- If they are on the same side of the plane $(s_c s_o > 0)$ - Plus $|s_c| > r$ and $|s_o| > r$ • Otherwise, sphere can move $|s_c|-r$ • Time of collision: $t_{cd} = n + \frac{s_c - r}{s_c - s_e}$ • Response: reflect v around n, and move $(1-t_{cd})\mathbf{r}$ (r=refl vector)

BONUS

Dynamic Separating Axis Theorem SAT: tests one axis at a time for overlap





• Same with DSAT, but:

 Need to adjust the projection on the axis so that the interval moves on the axis as well

Need to test same axes as with SAT

- Same criteria for overlap/disjoint:
 - If no overlap on axis => disjoint
 - If overlap on all axes => objects overlap Moller © 20

BONUS

Dynamic Sweep-and-Prune

• http://graphics.idav.ucdavis.edu/~dcoming/papers/coming_staadt_vriphys05.pdf

Exercises

 Create a function (by writing code on paper) that tests for intersection between:

- two spheres
- a ray and a sphere
- view frustum and a sphere

Scan Line Fill

Set active edges to AB and AC For y = A.y, A.y-1,...,C.yIf $y=B.y \rightarrow$ exchange AB with BC Compute xstart and xend. Interpolate color, depth, texcoords etc for points (xstart,y) and (xend,y)

For x = xstart, xstart+1, ...,xend Compute color, depth etc for

(x,y) using interpolation.



This is the modern way to rasterize a triangle

Using Interpolation

 $C_1 C_2 C_3$ specified by glColor or by vertex shading C_4 determined by interpolating between C_1 and C_2 C_5 determined by interpolating between C_2 and C_3 interpolate between C_4 and C_5 along span



Rasterizing a Triangle

- -Convex Polygons only
- –Nonconvex polygons assumed to have been tessellated
- –Shades (colors) have been computed for vertices (Gouraud shading)
- -Combine with z-buffer algorithm
 - March across scan lines interpolating shades
 - Incremental work small

Flood Fill

- Fill can be done recursively if we know a seed point located inside (WHITE)
- Scan convert edges into buffer in edge/inside color (BLACK)

```
flood_fill(int x, int y) {
    if(read_pixel(x,y) = = WHITE) {
        write_pixel(x,y,BLACK);
        flood_fill(x-1, y);
        flood_fill(x+1, y);
        flood_fill(x, y+1);
        flood_fill(x, y-1);
    }
}
```