

## Database Usage (and Construction)

SQL Queries and Relational Algebra  
Views

## Summary so far

- SQL is based on relational algebra.
  - Operations over relations
- Operations for:
  - Selection of rows ( $\sigma$ )
  - Projection of columns ( $\pi$ )
  - Combining tables
    - Cartesian product ( $\times$ )
    - Join, natural join ( $\bowtie$ ,  $\Join$ )

## Subqueries

- Subqueries is a term referring to a query used inside another query:

```
SELECT teacher
FROM   GivenCourses NATURAL JOIN
      (SELECT course, period
       FROM   Lectures
        WHERE weekday = 'Mon')
WHERE  period = 2;
```

What does this query mean?

```
SELECT course, period
FROM   Lectures
WHERE  weekday = 'Mon'
```

course	period	room	weekday	hour
TDA357	2	room1	Mon	8
TDA357	2	room1	Thu	8
TDA357	4	room3	Tue	8
TDA357	4	room3	Thu	13
TIN090	1	room4	Mon	8
TIN090	1	room3	Thu	13

```
SELECT course, period
FROM   Lectures
WHERE  weekday = 'Mon'
```

course	period	room	weekday	hour
TDA357	2	room1	Mon	8
TIN090	1	room4	Mon	8

```
SELECT teacher
FROM   GivenCourses NATURAL JOIN
      (SELECT course, period
       FROM   Lectures
        WHERE weekday = 'Mon')
WHERE  period = 2;
```

course	period
TDA357	2
TIN090	1

course	period	teacher	#students
TDA357	2	Niklas Broberg	130
TDA357	4	Rogardt Heldal	135
TIN090	1	Devdatt Dubashi	95

```
SELECT teacher
FROM   GivenCourses NATURAL JOIN
      (SELECT course, period
       FROM   Lectures
        WHERE weekday = 'Mon')
WHERE  period = 2;
```

course	period	teacher	#students
TDA357	2	Niklas Broberg	130
TIN090	1	Devdatt Dubashi	95

## Result

teacher
Niklas Broberg

## Renaming attributes

- Sometimes we want to give new names to attributes in the result of a query.
  - To better understand what the result models are.
  - In some cases, to simplify queries

```
SELECT *
FROM   Courses NATURAL JOIN
      (SELECT course as code, period, teacher
       FROM   GivenCourses);
```

## Renaming relations

- Name the result of a subquery to be able to refer to the attributes in it.
- Alias existing relations (tables) to make referring to it simpler, or to disambiguate.

```
SELECT L.course, weekday, hour, room
FROM   Lectures L, GivenCourses G, Rooms
WHERE  L.course = G.course
      AND L.period = G.period
      AND room = name
      AND nrSeats < nrStudents;
```

What does this query mean?

## Renaming in Relational Algebra

- Renaming = Given a relation, give a new name to it, and (possibly) to its attributes

$$\rho_{A(X)}(R)$$

- Rename R to A, and the attributes of R to the names specified by X (must match the number of attributes).
- Leaving out X means attribute names stay the same.
- Renaming the relation is only necessary for subqueries.
- $\rho$  = rho = greek letter  $\mathbf{r}$  = rename

## Quiz!

Write a query that lists all courses that are given in more than one period, with different teachers.

```
SELECT A.course
FROM   GivenCourses A, GivenCourses B
WHERE  A.course = B.course
      AND A.teacher <> B.teacher;
```

## Sequencing

- Easier to handle subqueries separately when queries become complicated.

– Example:  $\pi_x(R_1 \bowtie_c R_2)$  could be written as

```
R3 := R1 ⋈ R2
R4 := σc(R3)
R   := πx(R4)
```

– In SQL:

```
WITH
  R3 AS (SELECT * FROM R1, R2),
  R4 AS (SELECT * FROM R3 WHERE c)
SELECT x FROM R4;
```

- Example:

```
WITH DBLectures AS
  (SELECT room, hour, weekday
   FROM Lectures
   WHERE course = 'TDA357'
        AND period = 2)
SELECT weekday
FROM DBLectures
WHERE room = 'VR';
```

What does this query mean?

## Creating views

- A *view* is a "virtual table", or "persistent query" – a relation defined in the database using data contained in other tables.

```
CREATE VIEW viewname AS query
```

- For purposes of querying, a view works just like a table. The main difference is that you can't perform modifications on it – its contents is defined by other tables.

Example:

```
CREATE VIEW DBLectures AS
  SELECT room, hour, weekday
  FROM Lectures
  WHERE course = 'TDA357'
        AND period = 2;

SELECT weekday
FROM DBLectures
WHERE room = 'VR';
```

## The WHERE clause

- Specify conditions *over rows*.
- Can involve
  - constants
  - attributes in the row
  - simple value functions (e.g. ABS, UPPER)
  - subqueries
- Lots of nice tests to make...

## Testing for membership

- Test whether or not a tuple is a member of some relation.

```
tuple [NOT] IN subquery {or literal set}
```

```
SELECT course
FROM GivenCourses
WHERE period IN (1, 4);
```

List all courses that take place in the first or fourth periods.

## Quiz!

List all courses given by a teacher who also gives the Databases course (TDA357). (You must use IN...)

```
SELECT course
FROM   GivenCourses
WHERE  teacher IN
      (SELECT teacher
       FROM   GivenCourses
       WHERE  course = 'TDA357');
```

## Testing for existence

- Test whether or not a relation is empty.

`[NOT] EXISTS subquery`

e.g. List all courses that have lectures.

```
SELECT code, name
FROM   Courses
WHERE  EXISTS
      (SELECT *
       FROM   Lectures
       WHERE  course = code);
```

Note that code is in scope here since it is an attribute in the row being tested in the outer "WHERE" clause. This is called a correlated query.

## Quiz!

List all courses that are not given in the second period. (You must use EXISTS...)

```
SELECT code
FROM   Courses
WHERE  NOT EXISTS
      (SELECT *
       FROM   GivenCourses
       WHERE  course = code
              AND period = 2);
```

## Ordinary comparisons

- Normal comparison operators like =, <, <>, but also the special BETWEEN.

`value1 BETWEEN value2 AND value3`

```
SELECT course
FROM   GivenCourses
WHERE  period BETWEEN 2 AND 3;
```

List all courses that take place in the second or third periods.

– Same thing as

`value2 <= value1 AND value1 <= value3`

## Comparisons with many rows

- Two operators that let us compare with all the values in a relation at the same time.

`tuple op ANY subquery {or literal set}`  
`tuple op ALL subquery {or literal set}`

```
SELECT course
FROM   GivenCourses
WHERE  period = ANY (1,4);
```

List all courses that take place in the first or fourth periods.

## Quiz!

List the course(s) with the fewest number of students (in any period). (You must use ANY or ALL...)

```
SELECT course
FROM   GivenCourses
WHERE  nrStudents <= ALL
      (SELECT nrStudents
       FROM   GivenCourses);
```

## String comparisons

- Normal comparison operators like < use lexicographical order.
  - 'foo' < 'fool' < 'foul'
- Searching for patterns in strings:

`string LIKE pattern`

- Two special pattern characters:
  - \_ (underscore) matches any one character.
  - % matches any (possibly empty) sequence of characters.

## Quiz!

List all courses that have anything to do with databases (i.e. have the word Database in their name).

```
SELECT *
FROM   Courses
WHERE  name LIKE '%Database%';
```

## The NULL symbol

- Special symbol NULL means either
  - we have no value, or
  - we don't know the value
- Use with care!
  - Comparisons and other operations won't work.
  - May take up unnecessary space.

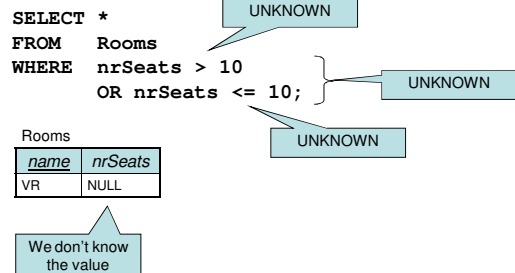
## Comparing values with NULL

- The logic of SQL is a three-valued logic – TRUE, FALSE and UNKNOWN.
- Comparing any value with NULL results in UNKNOWN.
- A row is selected if all the conditions in the WHERE clause are TRUE for that row, i.e. not FALSE or UNKNOWN.

## Three-valued logic

- Rules for logic with unknowns:
  - true AND unknown = unknown
  - false AND unknown = false
  - true OR unknown = true
  - false OR unknown = unknown
  - unknown AND/OR unknown = unknown

## Unintuitive result



## Don't expect the "usual" results

- Laws of three-valued logic are not the same as those for two-valued logic.
- Some laws hold, like commutativity of AND and OR.
- Others do not:  
p OR NOT p = true

## Arithmetic in queries

- We allow arithmetic operations in queries.

```
SELECT weekday, hour, room, course,  
       nrSeats - nrStudents as nrFreeSeats  
FROM   Rooms,  
       (Lectures NATURAL JOIN GivenCourses)  
WHERE  name = room;
```

- Not just arithmetic, but rather any operations on values.
  - Oracle has lots of pre-defined functions.

## Constants

- Constants can be used in projections.

```
SELECT code, name,  
       'Database course' as comment  
FROM   Courses  
WHERE  name LIKE '%Database%';
```

code	name	comment
TDA357	Databases	Databases course

– Beware of keywords...

## Quiz!

What will the result of this query be?

```
SELECT 1  
FROM    Courses;
```

Courses	
code	name
TDA357	Databases
TIN090	Algorithms

1
1
1

For each row in Courses that passes the test (all rows since we have no test), project the value 1.

## Aggregation

- Aggregation functions are functions that produce a single value over a relation.
  - SUM, MAX, MIN, AVG, COUNT...

```
SELECT MAX(nrSeats)  
FROM   Rooms;
```

MAX actually has Rooms as an implicit argument!

```
SELECT COUNT(*)  
FROM   Lectures  
WHERE  room = 'VR';
```

## Quiz!

List the room(s) with the highest number of seats, and its number of seats.

```
SELECT name, MAX(nrSeats)  
FROM   Rooms;
```

NOT correct!

Error when trying to execute, why is it so?

## Aggregate functions are special

- Compare the following:

```
SELECT nrSeats
FROM Rooms;
```

```
SELECT MAX(nrSeats)
FROM Rooms;
```

- The ordinary selection/projection results in a relation with a single attribute nrSeats, and one row for each row in Rooms.
- The aggregation results in a single value, not a relation.
- We can't mix both kinds in the same query!  
(almost...more on this later)

name	nrSeats
room1	10
room2	20
room3	55
room4	30
room5	34

```
SELECT nrSeats
FROM Rooms;
```

nrSeats
10
20
55
30
34

name	nrSeats
room1	10
room2	20
room3	55
room4	30
room5	34

```
SELECT MAX (nrSeats)
FROM Rooms;
```

MAX(nrSeats)
55

```
SELECT MAX (nrSeats) AS nrSeats
FROM Rooms;
```

NRSEATS
55

## Quiz! New attempt

List the room(s) with the highest number of seats, and its number of seats.

```
SELECT name,
       (SELECT MAX(nrSeats)
        FROM Rooms)
FROM Rooms;
```

Not correct either, will list all rooms, together with the highest number of seats in any room.

Let's try yet again...

name	nrSeats
room1	10
room2	20
room3	55
room4	30
room5	34

```
SELECT name,
       (SELECT MAX(nrSeats)
        FROM Rooms)
FROM Rooms;
```

name	nrSeats
room1	55
room2	55
room3	55
room4	55
room5	55

## Quiz! New attempt

List the room(s) with the highest number of seats, and its number of seats.

```
SELECT name, nrSeats
FROM Rooms
WHERE nrSeats = MAX(nrSeats);
```

Still not correct, MAX(nrSeats) is not a test over a row so it can't appear in the WHERE clause!

Let's try yet again...

## Quiz!

List the room(s) with the highest number of seats, and its number of seats.

```
SELECT name, nrSeats
FROM   Rooms
WHERE  nrSeats =
      (SELECT MAX(nrSeats)
       FROM   Rooms);
```

That's better!

## Single-value queries

- If the result of a query is known to be a single value (like for MAX), the whole query may be used as a value.

```
SELECT name, nrSeats
FROM   Rooms
WHERE  nrSeats =
      (SELECT MAX(nrSeats)
       FROM   Rooms);
```

- Dynamic verification, so be careful...

## NULL in aggregations

- NULL never contributes to a sum, average or count, and can never be the maximum or minimum value.
- If there are no non-null values, the result of the aggregation is NULL.

## Summary – aggregation

- Aggregation functions: MAX, MIN, COUNT, AVG, SUM
- Compute a single value over a whole relation.
- Can't put aggregation directly in the WHERE clause (since it's not a function on values).
- Can't mix aggregation and normal projection!  
... well, not quite true...

## Not quite true?

- Sometimes we want to compute an aggregation for every value of some other attribute.
  - Example: List the average number of students that each teacher has on his or her courses.
  - To write a query for this, we must compute the averaging aggregation *for each value of teacher*.

## Grouping

- Grouping intuitively means to partition a relation into several groups, based on the value of some attribute(s).
  - "All courses with this teacher go in this group, all courses with that teacher go in that group, ..."
- Each group is a sub-relation, and aggregations can be computed over them.
- Within each group, all rows have the same value for the attribute(s) grouped on, and therefore we can project that value as well!



## Grouping

- Grouping = given a relation R, a set of attributes X, and a set of aggregation expressions G; partition R into groups  $R_1 \dots R_n$  such that all rows in  $R_i$  have the same value on all attributes in X, and project X and G for each group.

$\gamma_{X,G}(R)$

```
SELECT  X, G
FROM    R
GROUP BY X;
```

- "For each X, compute G"
- $\gamma$  = gamma = greek letter **g** = **g**rouping

Example: List the average number of students that each teacher has on his or her courses.

course	per	teacher	nrSt.
TDA357	4	Rogardt Heldal	130
TDA590	2	Rogardt Heldal	70
TIN090	1	Devdatt Dubhashi	62

SQL?

Result?

Relational Algebra?

## Specialized renaming of attributes

- General renaming operator, rename R to A and its attributes to X :

$\rho_{A(X)}(R)$

- More convenient alternative for grouping, rename the result of expression G to B:

$\gamma_{X,G \rightarrow B}(R)$

- e.g.  $\gamma_{teacher, AVG(nrStudents) \rightarrow avgStudents}(Given\ Courses)$
- Works in normal projection ( $\pi$ ) as well.

## Summary – grouping and aggregation

- Aggregation functions: MAX, MIN, COUNT, AVG, SUM
  - Compute a single value over a whole relation, or a partition of a relation (i.e. a group).
  - If no grouping attributes are given, the aggregation affects the whole relation (and no ordinary attributes can be projected).
- Can't put aggregation directly in the WHERE clause (since it's not a function on values).
- Can't mix aggregation and normal projection!
  - If an aggregation function is used in the SELECT clause, then the only other things that may be used there are other aggregation functions, and attributes that are grouped on.

## Summary

- Complex queries, involving subqueries
  - Renaming of relations and attributes
- Creating views
- Lots and lots of tests for the WHERE clause
  - IN, EXISTS, BETWEEN, ALL, ANY, LIKE
- Arithmetic and other functions, constant values
- Aggregation functions
  - more on these next time