

Reproducibility, Proofs and Domain Specific Languages

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Who am I?

- | Patrik Jansson, Prof. of Comp. Sci., Chalmers, Sweden

- | Programming languages

- | Software technology

- | Two EU projects about “Global Systems Science (GSS)”

- | CoeGSS.eu

- | GRACeFUL-project.eu

- | a “tweet-sized definition” (from 2013):

- GSS is about developing systems, theories, languages and tools for computer-aided policy making with potentially global implications.*

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- | When is the implementation *correct* w.r.t. the spec?
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- | If we have a formal system (a logic)
 - | we may give a proof of correctness
 - | (moving trust)

Machine checked proof

- | Dependent type theory
 - | can express many specifications
 - | and implementations
 - | and proofs
- | in the same formal system.

Proof assistants: Coq, Agda, Idris, ...

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Risk: formal proofs improve 3. but damage 4.

A short story about Dee and Foo

Actors: Dee the Domain Expert & Foo the Formaliser

Starting point: Dee has an experiment + informal description

Dee + Foo work on formalisation.

Happy ending?: formal spec. and computer proof of correctness

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Drama:

- | Dee cannot understand the formal spec.
- | Foo cannot understand the domain.

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- | now Dee can “assemble” the proof herself!

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- | not merely correct programs,
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Summary

Key points:

- | experiment = implementation
- | reproducibility requires also a *specification*
- | *provability* is stronger than reproducibility
- | dependently typed languages can express spec., impl. and proof
- | formal specifications can introduce a “comprehension gap”
- | Domain-specific languages (DSLs) can be used to bridge the gap

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