On the Specification of Full Contracts

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- A contract is a binding agreement between two or more entities (enforceable by law)
- Use contracts to regulate interactions in concurrent/distributed systems
 - Components, services, etc
- Different notions (or *levels*) of contracts
 - Static interfaces
 - Behavioural interfaces
 - Design-by-contract (pre-, post-conditions, invariants, etc)
 - Quality-of-Service
 - 'Social' contracts
 - Deontic e-contracts

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- Full contracts: Normal and exceptional behaviour
 - Exceptions, compensations, tolerance to faults, penalties, etc

() Specification of CoCoME (use cases 1-8) using CL

- $\bullet \ \mathcal{CL}$ is contract language based on deontic logic
- It allows the specification of obligations, permissions and prohibitions, and the penalties in case of violations
- Or a compare suitability of operational and logical approaches to specify full contracts on a well-known case study (CoCoME)
 - Operational
 - rCOS (Relational Calculus of Object and Component Systems –CSP implementation)
 - Logical
 - Deontic-logic based language (\mathcal{CL})
 - Temporal logics

CoCoME: Common Component Modelling Example

- Trading System to handle sales and inventory of a store chain
- 8 use cases



CoCoME: Common Component Modelling Example

- Trading System to handle sales and inventory of a store chain
- 8 use cases
- Use case 1
 - How a sale is processed
- Use case 2
 - How a cash desk switches to express mode, restricting total number of customer items
- We focus on the behavioural aspects of the use cases
 - Prop1, Prop2, Prop3



Operational and Logic Specification Languages CSP and Temporal Logics

Definition (CSP)

• CSP (rCOS)

P ::= $Stop \mid a \rightarrow P \mid P[]P \mid P \sqcap P \mid X$

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Definition (CSP)

• CSP (rCOS)

$$P$$
 ::= Stop | $a \rightarrow P | P[]P | P \sqcap P | X$

Definition (Temporal Logics)

LTL

$$\varphi \quad ::= \quad p \mid \neg \varphi \mid \varphi \lor \varphi \mid \mathbf{G}\varphi \mid \mathbf{F}\varphi \mid \mathbf{X}\varphi \mid \varphi \mathbf{U}\varphi$$

• CTL

$$p ::= p | \neg \varphi | \varphi \lor \varphi | \mathbf{AG}\varphi | \mathbf{AF}\varphi | \mathbf{AX}\varphi | \varphi \mathbf{AU}\varphi | \mathbf{EG}\varphi | \mathbf{EF}\varphi \\ \mathbf{EX}\varphi | \varphi \mathbf{EU}\varphi$$

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Operational and Logic Specification Languages

 $\mathcal{CL}:$ A Deontic-Based Language for Contracts

Definition (*CL* Syntax)

$$C := C_{O} | C_{P} | C_{F} | C \land C | [\beta] C | \top | \bot$$

$$C_{O} := \mathbb{O}_{C}(\alpha) | C_{O} \oplus C_{O}$$

$$C_{P} := \mathbb{P}(\alpha) | C_{P} \oplus C_{P}$$

$$C_{F} := \mathbb{F}_{C}(\alpha) | C_{F} \lor [\alpha] C_{F}$$

$$\alpha := 0 | 1 | a | \alpha \& \alpha | \alpha; \alpha | \alpha + \alpha \quad \beta := 0 | 1 | a | \beta \& \beta | \beta; \beta | \beta + \beta | \beta^{*}$$

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Example (Specification of Prop1)

If pay with card, three allowed attempts to enter pin code; otherwise pay with cash, or return goods

$$\Box$$
[cardPay] $\mathbb{O}_{\psi_1}(correctPin)$

where $\psi_1 = \mathbb{O}_{\psi_2}(\text{correctPin})$, with $\psi_2 = \mathbb{O}_{\mathbb{O}(\text{cashPay}+\text{returnItems})}(\text{correctPin})$

Specification of CoCoME Use Cases 1 and 2 Specific clauses to be specified

Prop1: Pay by cash: obligation to swipe the card + correct pin

- Incorrect pin: two more allowed attempts
- After 3 incorrect pins: obligation to pay cash
- No cash: give up the goods

Prop1: Pay by cash: obligation to swipe the card + correct pin

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Prop2: Normal mode: the cashier may switch to express mode

- If in the last hour 50% of the sales had less than eight items (*)
- In express mode: cashier obliged to eventually go to normal mode
- If (*) holds infinitely often, then the cashier should change to express mode infinitely often

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Prop2: Normal mode: the cashier may switch to express mode

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- In express mode: cashier obliged to eventually go to normal mode
- If (*) holds infinitely often, then the cashier should change to express mode infinitely often

Prop3: In express mode: cashier obliged to service customers with less than eight items

• If customer with more than eight items: cashier decides whether to service the client



- Payment with credit card
- Three allowed attemps to enter correct pin

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- Payment with credit card
- Three allowed attemps to enter correct pin

- CSP: Specification of normal case + refinement to capture exceptional behaviour
 - Possible but intricate branching
- Can be described in both CTL and LTL
- Can be described in CL (using CTDs)



- Permission to switch from normal to express mode
- Obligation to come back to normal mode
- Fairness constraint between normal and express mode



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- Obligation to come back to normal mode
- Fairness constraint between normal and express mode

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- Fairness left underspecified in CSP
- Cannot be described in CTL (fairness)
- Cannot be described in LTL (existential branch)
- Can be described in CL (modulo semantical treatment of fairness)



- Obligation to serve customers with
 < 8 items
- Permission to service clients with > 8 items

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- Obligation to serve customers with < 8 items
- Permission to service clients with > 8 items

- Possible in CSP, but process not longer a refinement of original
- Described in CTL
- Cannot be described in LTL (existential branch)
- Can be described in CL

	LTL	CTL	CSP	\mathcal{CL}
Prop1	\checkmark	\checkmark	\checkmark	\checkmark
Prop2	-	-	-	(√)
Prop3	_	\checkmark	(√)	\checkmark

• Different specification languages suitable to different purposes

- · Contracts only for normal behaviour: temporal logic would suffice
- Composition and comparison of contracts: process calculi more flexibility
- Operational approach and temporal logic not very suitable to represent certain exceptional cases
 - $\bullet \ \mathcal{CL}$ could be encoded in CTL^*
- Deontic approach suitable to represent obligations, permissions and prohibitions (and many exceptional cases)
 - Needs to be extended to capture more complex compensations (as present in *long transactions*)