### Efficient Verification of Software Product Lines

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## Software Product Lines

- A product line is "a family of products designed to take advantage of their common aspects and predicted variabilities."[Weiss; 1999]
- A product line is "a set of systems sharing a common set of features that satisfy the specific needs of a particular market segment." [Clements, Northrop; 2001]

## Product Line Development



## Verification of SPL

- High configurative variability of products
- Correctness of products is crucial.
- Formal verification by theorem proving and model checking can establish critical product properties.
- But, it is not feasible to verify each product in isolation.



Verification of Software Product Lines







Verification of Software Product Lines

### Outline

- Model-based Software Product Line Engineering
- Variability Modelling using  $\Delta s$
- Implementing SPL with  $F\Delta J$
- Proof Reuse for Verification of  $F\Delta Js$

## Product Map

Example from [Batory et al., FOAL09]:

	Base	Sync	Ret	Inv	WHolder
BaseAccount	X				
SyncAccount	X	X			
AccWHolder	X				×
RetAccount	X		x		
SyncAccWH	X	X			x

## Feature Model



Master Thesis: Comparison between Product Maps and Feature Models

Verification of Software Product Lines





I. Schaefer, A. Worret, A. Poetzsch-Heffter: A Model-based Framework for Automated Product Derivation. Workshop on Model-driven Approaches to Product Line Engineering (MAPLE), August 2009



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## $\Delta$ -Modelling

A product line is represented by a core product and a set of product- $\Delta$ s.

The core product is a complete product for a valid feature configuration. It is not uniquely determined.

<u>Guideline:</u> The core product contains

- all mandatory features.
- a minimal number of required alternative features

Master Thesis: Impact of Core Product to SPL Design and Implementation

## $\Delta$ -Modelling (2)

- Product-Δs specify Additions, Modifications, Removals to the Core Product.
- Application constraint over the features in the feature model determines for which feature configuration the product- $\Delta$  has to be applied to the core product.
- Product-Δs can be partially ordered to avoid conflicting changes of the core product.

## Configuration

For a Feature Configuration:

- Determine product- $\Delta s$  with valid application condition.
- Determine linear ordering of product- $\Delta s$  compatible with partial ordering.
- Apply changes specified by product-Δs to core product in the linear order.

## Variability in System Design



## Configuring Designs

#### Product: Base & Sync

#### Core Design

class Account

implements IAccount

int balance

void update(int x)



#### **Configured Design**

class SyncAccount
implements IAccount
int lock
int balance
void update(int x)
void sync(int i)

#### Verification of Software Product Lines

## Configuring Designs

#### Product: Base & Ret & With Holder



#### Configured Design class RetAccount implements IBonus int 401balance void update(int x) void addBonus(int i)

Verification of Software Product Lines

## Variability Modelling with $\Delta s$

- Evolutionary Development by Adding Product- $\Delta s$
- Explicit Treatment of Combinations of Features by Complex Application Conditions
- Usable with Different Modelling Formalisms and Implementation Techniques
- Model Refinements are Orthogonal to Variability Modelling.

## Model Refinement



It holds that: refine(configure((Core,  $\Delta$ s),fc)) = configure(refine(Core, $\Delta$ s),fc)

<u>Master Thesis</u>: Case Study and Tool Support for MDD with  $\Delta s$ 

Implementing SPL using Deltas and Traits



## $F\Delta J - A PL for SPL$

- Extension of Java with Core and  $\Delta$ -Modules
- Core Product is implemented by Core Module.
- Product- $\Delta s$  are implemented by  $\Delta$ -Modules.
- A Product Implementation is obtained by application of Δ-Modules to Core Module.
- Type System ensures safety of  $\Delta$ -application.

## $F\Delta J - A PL for SPL$

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- A Product Implementation is obtained by application of Δ-Modules to Core Module.
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L. Bettini, V. Bono, F. Damiani, I. Schaefer: A Programming Language for Software Product Lines. Draft, December 2009

### Core Module

A core module contains a set of Java classes.

core BaseAccount {
 class Account extends Object {
 int balance;
 void update(int x) { balance += x; }
 }
}

## $\Delta$ -Modules

- Modifications on Class Level:
  - Addition, Removal and Modification of Classes
- Modifications of internal Class Structure:
  - Adding, Removing, Renaming Fields
  - Adding, Removing, Renaming Methods
- Application Condition in when clause: Boolean Constraint on Features in Feature Model
- Partial Ordering of  $\Delta$ -Modules by after clauses

## $\Delta$ -Module for Sync

delta DsyncUpdate after Dretirement, Dinvestment when Sync {
 modifies class Account {
 adds Lock lock;
 renames update to unsync\_update;
 adds void update(int x) { lock.lock(); unsync\_update(x); lock.unlock(); }
 }
}

## $\Delta$ -Application

```
core BaseAccount {
                                        Core
    class Account extends Object {
                                      Module
         int balance;
         void update(int x) { balance += x; }
```

}

**delta** DsyncUpdate **after** Dretirement, Dinvestment **when** Sync { **modifies class** Account {  $\Delta$ - Module adds Lock lock; **renames** update **to** unsync\_update; **adds void** update(**int** x) { lock.lock(); unsync\_update(x); lock.unlock(); } class Account extends Object { Product int balance; Lock lock; void unsync\_update(int x) { balance += x; } void update(int x) { lock.lock(); unsync\_update(x); lock.unlock(); } Verification of Software Product Lines

## $\Delta$ -Application

#### 

class Account extends Object {
 int balance;
 Lock lock;
 void unsync\_update(int x) { balance += x; }
 void update(int x) { lock.lock(); unsync\_update(x); lock.unlock(); }

# class Client { Account a; void payday(int x, int bonus) { a.addBonus(bonus); a.update(x); } }

## Type System for $F\Delta J$

- The core and  $\Delta$ -modules can be typed in isolation.
- If a core module and a set of  $\Delta$ -modules are type correct,  $\Delta$ -application is safe:
  - all renamed/removed fields and methods exist
  - all added fields and methods do not exist
  - removed classes exists and added classes do not exist
  - there are not conflicting modifications in a class

## Verification of $F\Delta J$

- We use the KeY System for deductive verification.
- Input: Java Program + JML Specifications
- KeY generates proof obligations in dynamic logic.
- KeY supports interactive and automatic verification of the proof obligations.

#### Specification of Base Account

We want to prove that the balance of an account is always positive.

```
/*@
@ public instance invariant balance >= 0; ← — Instance Invariant
(d*/
public class BaseAccount {
  int balance;
  /*@
   @ ensures \result.balance==0;
   @*/
  public BaseAccount(){
  balance = 0;
  }
  /*@
   @ public normal_behavior
   @ requires x > 0;
                                                 Method Contract
   @ assignable \everything;
   @ ensures balance >= \old(balance);
   (d*/
  public void update(int x){
      balance = balance + x;
  }
```

}

## Specification of SyncAccount

We want to prove that the balance of an account is always positive.

```
/*@
@ public instance invariant balance >= 0;
 (d*/
public class SyncAccount {
                                               /*@
                                                   @ public normal_behavior
  int balance;
                                                   @ requires x > 0;
  Lock lock:
                                                   @ assignable \everything;
                                                   @ ensures balance >= \old(balance);
                                                   @*/
  /*@
                                                  public void unsync_update(int x){
   @ ensures \result.balance==0;
                                                      balance = balance + x;
   (d*/
                                                  }
  public SyncAccount(){
  balance = 0;
                                                  /*@
  lock = new Lock():
                                                   @ public normal_behavior
   }
                                                   @ requires x > 0;
                                                   @ assignable \everything;
                                                   @ ensures balance >= \old(balance);
                                                   @*/
                                                  public void update(int x){
                                                  lock.lock(); unsync_update(x); lock.unlock();
                                                       }
                                               }
```

## Comparison

```
public class BaseAccount {
```

#### [...]

}

```
/*@
 @ public normal_behavior
 @ requires x > 0;
 @ assignable \everything;
 @ ensures balance >= \old(balance);
 @*/
public void update(int x){
    balance = balance + x;
}
```

public class SyncAccount {

```
[...]
```

```
/*@
  @ public normal_behavior
  @ requires x > 0;
  @ assignable \everything;
  @ ensures balance >= \old(balance);
  @*/
public void unsync_update(int x){
     balance = balance + x;
}
```

```
/*@
```

```
@ public normal_behavior
@ requires x > 0;
@ assignable \everything;
@ ensures balance >= \old(balance);
@*/
public void update(int x){
lock.lock(); unsync_update(x); lock.unlock();
}
```

```
}
```

## Comparison



## Comparison



```
public class RetAccount {
public class BaseAccount {
                                                int bbalance;
int balance;
                                              [...]
[...]
                                                /*@
  /*@
                                                 @ public normal_behavior
   @ public normal_behavior
                                                 @ requires x > 0;
   @ requires x > 0;
                                                 @ assignable \everything;
   @ assignable \everything;
                                                 @ ensures bbalance >= \old(bbalance);
   @ ensures balance >= \old(balance);
                                                 @*/
   @*/
                                                public void update(int x){
  public void update(int x){
                                                     bbalance = bbalance + x;
      balance = balance + x;
                                                 }
  }
                                                /*@
}
                                                 @ public normal_behavior
                                                 @ requires x > 0;
                                                 @ assignable \everything;
                                                 @ ensures bbalance >= \old(bbalance);
                                                 (d*/
                                                public void addBonus(int x){
                                                     bbalance = bbalance + x;
                                                }
                                              }
```







### Observations

- I7 Method Contracts in 6 Variants of the Bank Account SPL verified.
- Only 3 Contracts have to be proven from scratch.
- $\Delta$ -Modules imply Specification- $\Delta$ s.
- Structure of Δ-Modules indicates Proof Reuse
   Potential
- Proofs can be reused if only fields and methods are renamed, but internal class structure is unchanged.
- More reuse scenarios to be identifed.

#### Master Thesis: Case Study on Proof Reuse for Example SPL

## Conclusion

- Model-based Software Product Line Engineering
- Variability Modelling using  $\Delta s$
- Implementing SPL with  $F\Delta J$
- Proof Reuse for Verification of  $F\Delta Js$

## Master Thesis Proposals

- Comparison between Product Maps and Feature Models
- Impact of the Core Product in  $\Delta$ -Modelling
- Evaluation and Tool Support for Model-based
   Development using Δ-Modelling
- Case Study on Proof Reuse for FΔJ